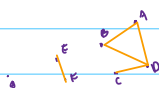




A Graph $G = (V, E)$ is V -vertices and edges $- E$



$V = \{A, B, C, D, E, F, G\}$

$E = \{ \{A, B\}, \{A, C\}, \{A, D\}, \{B, C\}, \{C, D\}, \{D, E\}, \{E, F\}, \{F, G\}, \{G, A\} \}$

A **Path** in graph G is a list

$V_0 - V_1 - V_2 - \dots - V_k$ st $V_i - V_{i+1}$

has an edge ; **length of path** = k

ex. $B - A - D$

$D - A - D - C$

A **cycle** is a path where $V_0 = V_k$

Cycle length = 0 ; it's a loop

The Degree of vertex v is the # of edges connected to v .

$d(A) = 2$ $d(E) = 1$ $d(G) = 0$

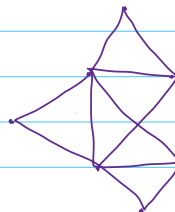
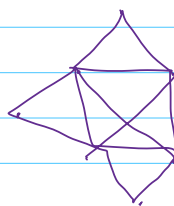
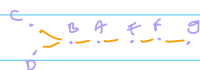
Vertex w/ degree 0 is **"isolated"**

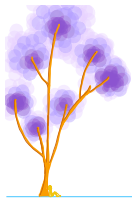
A graph is **"connected"** if for every pair of (diff) vertices there is a path

SIMPLE GRAPHS

A **simple** graph has no loops and no double edges

A **tree** is a connected graph and has no cycles



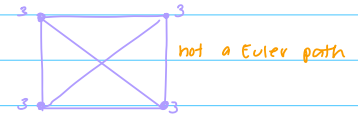
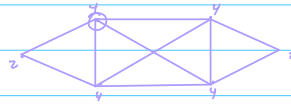


EULER PATHS & CYCLE

Path - graph that goes over every edge, exactly once

Cycle - Start & end at same point

A graph G has an Euler path iff all degrees are even, except (maybe) 2 vertices

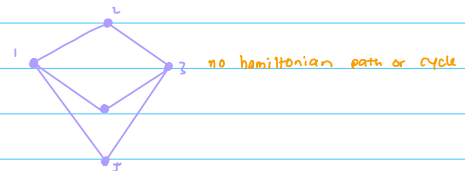
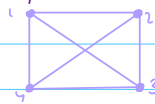


HAMILTONIAN PATHS & CYCLES

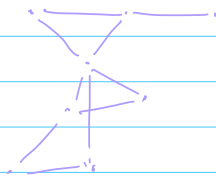
Path/cycle that visits every vertex once

ex: 2-1-4-3-2

2-4-1-3-2



Traveling sales person problem: difficult



REPRESENTING GRAPHS

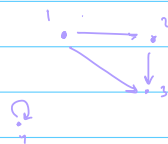
Simple: NO loops

NO parallel edges

ADJACENCY MATRIX

Let R be a relation on $\{1, 2, 3, 4\}$

	1	2	3	4
1	.	1	1	.
2	.	.	1	.
3
4	.	.	.	1



Can tell if a graph is symmetric

How many reflexive relations can we find?

	1	2	3	4
1	1	.	.	.
2	.	1	.	.
3	.	.	1	.
4	.	.	.	1

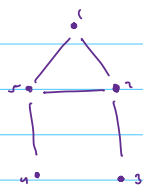
by multiplication principle:

$$2 \cdot 2 \cdot 2 \cdot 2 \cdot 2 \cdots 2$$

$\underbrace{\hspace{10em}}_{12}$

$$= 2^{12}$$

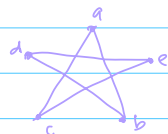
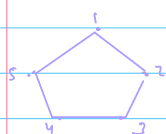
GRAPH ISOMORPHISM



Let $G_1 = (V_1, E_1)$

$G_2 = (V_2, E_2)$ then G_1 is isomorphic to G_2

if we can find a bijection $b: V_1 \rightarrow V_2$ and
 $(v_i, v_j) \in E_1$ iff $(b(v_i), b(v_j)) \in E_2$



$1 \rightarrow a$

$2 \rightarrow b$

$3 \rightarrow d$

$4 \rightarrow e$

$5 \rightarrow c$

all have edges

✓

SEQUENCES & STRINGS - RECURSIVE

Sequence - a function $A \rightarrow Y$ where $A \subseteq \mathbb{Z}$

$\{a_i\}_{i=0}^{\infty}$

$a_0 = 1$

$a_i = 2^i$

$a_1 = 2$

$a_3 = 8$

$a_3 = 4$

$a_4 = 16$

⋮

String - Sequence $A \rightarrow Y$, $A \subseteq \mathbb{Z}$ s.t. A, Y are finite sets

$|A|$ is length of string

Y alphabet

$Y = \{a, b, c\}$

empty string, length = 0

ex: a, b, ca, cca

λ

$|aaab| = 4$ $ccbaaab = c^2ba^3b$

$\alpha = abcc$ $\beta = bcabb$

Concatination: $\alpha\beta = abccbcabb$

Let X be a finite set.

$X^* = \{\text{all strings over } x\}$ $X^* = X^* \cup \lambda$

Let t be a string, then w is a substring of t iff there are string α, β
 s.t. $t = \alpha w \beta$

$$a_i = 2^i \rightarrow a: 2, 4, 8, 16, 32, \dots$$

$$b_n = 2n \rightarrow b: 2, 4, 6, 8, 10, \dots$$

$\{a_i\}$ is a subsequence of $\{b_i\}$.

$$\{S_{i,0}^1 : 0, 3, 6, 9, 12, \dots$$

$$\{S_{i,2}^5 : S_2' = 6$$

$$S_3' = 12$$

$$S_4' = 300$$

$$S_5' = 330$$

S' is a subsequence of S

Let L be a set of strings over the alphabet $\{a, b\}$

$$\textcircled{1} \lambda \in L$$

$$\textcircled{2} \text{ if } \alpha \in L \text{ then } a\alpha \in L$$

$$\textcircled{3} \text{ if } \alpha \in L \text{ then } b\alpha \in L$$

$$L = \{\lambda, aa, bb, abba,$$

$$[\text{say } \alpha = \lambda, \text{ then } \alpha \in L \text{ so } a\lambda \in L \text{ then } aa \in L], aaaa, baaaaab, \dots\}$$

RECURSIVE SEQUENCES

$$\{a_i\}_{i=1}^{\infty} \quad a_i = 2^i$$

$$\{f_n\}_{n=1}^{\infty} \quad f_n = f_{n-1} + f_{n-2} \quad \text{for } f_1 = 1, f_2 = 1$$

$$1, 1, 2, 3, 5, 8, 13, 21, \dots$$

Fibonacci sequence

Recursive Formula \rightarrow Explicit Formula

homogeneous linear recursive formulas (degree 2)

Let $\{S_n\}_{n=0}^{\infty}$ be a sequence described by:

$$S_n = C_1 S_{n-1} + C_2 S_{n-2} \quad C_1, C_2 \text{ are constants}$$

$$\text{and } S_0 = A, S_1 = B \quad A, B \text{ are constants}$$

Apply algorithm to get Explicit Formula

$$\textcircled{1} \text{ Solve for } t: t^2 - C_1 t - C_2 = 0$$

$$t = r_1 \quad t_2 = r_2$$

case 1: if $r_1 \neq r_2$

explicit formula: $S_n = br_1^n + dr_2^n$ b, d are constants

case 2: $r_1 = r_2$ then the formula:

$$S_n = br_1^n + dn_1r_1^n$$

② use A, B to find b, d