

Final Presentatiton

COMP702 Project of Modularity in empire network

Student name: Daozheng QU, Student ID: 201518453

MSc of Data science and AI, University of Liverpool

Supervisor: Michele Zito, Keith Dures.

Abstract

What we design is an application in a professional field, that is, a combination of graph theory and modular theory and implemented by programming. We combine the known algorithms that are most suitable for this research (which will be explained in the summary), another network(empire graph), and find its relatively optimal structural division and the value of the optimal structure division represented by modularity. We convert a random planar graph into multiple empire graphs with different numbers of vertices, and then perform modular processing, then run the experiment several times and take the average value to observe the value of modularity distribution. Finally, we set the vertex growth range and interval of the random planar graph to observe the modularity distribution of namely the Markov Chain algorithm and the Fast Newman algorithm, to generate a network(random planar graph) and transform it into the random planar graph of different nodes and the difference in community division.

Acknowledgement:

as random planar graph, empire graph, modularity problems, etc. I am also very grateful for the valuable comments provided by the second supervisor. I can correct many design errors and improve the details of the proof. I also thanks to Senior Sheng Chen for giving me inspiration on the traversal (single line of code) on the opposite side of the empire graph, which will enable me to better I would like to sincerely thank the first supervisor for his help in the whole project and the meticulous answers to each question, such complete the transformation of the empire graph.

Ethical issues and data protection:

reference. It is strictly forbidden to use the project data and materials without the consent of the project supervisor and person in This project does not use any data of others, does not involve any ethical issues, and only uses publicly available resources for

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Main Text:

1.SUMMARY OF PROPOSAL:

1.1.The aims of the project.

Basic aim: Generate a small(50-100 nodes) random planar graph [1,2], and convert it into an empire graph according to the content of the empire graph conversion material given by the project, and find the empire graph whose outline structure is based on vertices), and show its details about calculating the modularity value

Challenge: Generate a random plan graph with 900-9000 nodes at an interval of 900 nodes, and repeat the above basic experiment 20 times. In the case of r = 2,3, 5, take the average Q value and analyze the distribution of the Q value



1.2.A summary of the design.

We have made adjustments to some details of the original plan over time and actual needs, such as changing the number of runs of the experiment. According to the author's Newman's proof[3,9], the formula was revised in the design. We have not changed the algorithm selection. The algorithm selected by the original design is the most appropriate and accurate. The reason why we chose the original design algorithm will be shown below

1.2.1. The design method and function of the program:

The design of basic aim: We generate a random planar graph according to the Markov chain algorithm, and then generate an empire graph generated empire graph is modularization, and the community is divided into the largest value of modularity. At this time, we have completed (node merging) based on the data of the generated random planar graph. Then according to the modularization algorithm (Fast the design of program of basic aim, and we prove the reliability and feasibility of the design in the experiment. The design of challenge aim: On the basis of the basic target program that we have completed, we designed a random planar graph of large graph generates three empire graphs with different vertices, and modularization the empire graph. There is an empire graph with three vertices Each under the random planar graph of each large node, and the average modularity of each empire graph is taken after multiple runs nodes (continuously increasing the number of nodes) and the value of r converted into an empire graph is also different.

1.2.2.Reasons for choosing the original design algorithm:

memoryless walk in the state space when time is homogeneous. According to the algorithm of the Markov chain, all qualified (under the Random planar graph(Markov chain):We still use the Markov chain algorithm[5,6]. Firstly, the Markov chain algorithm performs a random condition of keeping the planarity of the random planar graph) random planar sub-graphs can reach a uniform limit steady-state distribution, which is irreplaceable by other methods and is currently the best algorithm

algorithms (GN[9], Louvain method[10], etc.), this algorithm is faster and the accuracy almost same, so it is the best algorithm choice in this there is no good algorithm that can be directly obtained. Therefore, Newman proposed an idea based on greedy calculation, searching for all **Modularity(Fast Newman):**First of all, this algorithm was designed by the author Newman himself and proposed the concept of modularity for local optimizations until completion, and treating the results as global optimizations to obtain an approximate modularity. Compared with other the first time. Second, it is very difficult for us to really get the real modularity in a large network, it is very time-consuming and laborious, research

1.2.3.Pseudo-code of each algorithm

Empire graph: Create n/r vertices of the empire graph, for each vertex v in planar graph: list all connections {v,u}add edge to empire graph between the empire of v and that of u.

Community detection and modularity

#set a temporal empire graph for fast newman algorithm, which will # set adjacency matrix equal to one at [node_i_index, node_j_index] # find node i and node j index in adjacency matrix # Iterate each edge to create adjacency matrix be changed during the calculation # node index in adjacency matrix # create zeros adjacency matrix define an inititial function of G define a class of Empire_Graph

The community number of each node # The partitions for empire graph # every node will be a division

define a function Fast_Newman_Algo(empire_graph) define a function for calculating the value Q define a function for calculating delta Q #set a while loop(if division >1) set a Temp_Max_Delta_Q #set a division rotation set an All_Delta_Q set a Max_Merge #set a max _q

If multiple delta Q are equal, join communities together in pairs with

Find maximal Delta Q

communitiy j

join communities together in pairs with maximal delta Q

maximal Q

calcualte delta Q before and after joining communitiy i and

Will not be considered, if node i and node j are in same

Get communitiy number for node i and node j

set division, max_Q = Fast_Newman_Algo(empire_graph) #set the empire_graph = Empire_Graph(G) # Find the devisions with maximal Q #set G=create_empire(graph_ori, r) # remove the edge after joining # choosing at each step the join that results in the greatest increase # Iterator all edges in temporal graph, repeatedly join communities (or smallest decrease) in Q. together in pairs,

2.1.A description of what was produced in the project.

We have completed the program design for the basic goals and the challenge-level goals and successfully run two programs. A small node application and a large node application. Due to time reasons, we only performed 40% (900-3600 nodes) of the total amount of the experiment, and produced 80 simulated graphs, 80 random planar graphs, 240 empire graphs, 240 community divisions and maximum Q values, and 12 average maximum Q values. The different number of nodes, different r values and average maximum Q value of the random planar graph are shown in the following figure.

average Q value of 20 times

| 900/2 0.442 |
|-------------------------|
| 0.475 0.491 0.504 |

0.355

0.364

0.327

0.372

| average Q value of 20 times | 0.231 | 0.257 | 0.266 | 0.271 |
|--------------------------------|-------|--------|--------|--------|
| nodes/r | 900/2 | 1800/5 | 2700/5 | 3600/5 |



2.2. Any interesting aspects of the implementation.

empire graph with a large number of nodes but the higher the density of its topology, sometimes it looks like a sphere. With the With the increase in the number of random planar graph nodes, under a fixed number of simulations, the larger the Q value of the nodes is like a hospital tester that fluctuates within a certain level, and the large nodes seem to be sitting on a rocket close to increase in the number of nodes in the random planar graph, under a fixed number of simulations, the trend of the number of straight line.

3.EVALUATION:

3.1.An evaluation of what has been achieved.

Our basic design and the selected algorithm are feasible, effective, correct, which enables us to complete the program design of the only yielded 40% of the results, but even this can reflect the conclusions and assumptions of this project. Generally speaking, the project has been completed, but it takes time to run more results. It is expected to continue to be described in detail in the basic goals and the more difficult goals, and run the program smoothly. However, due to time reasons, our relatively difficult program dissertation

3.2. Any shortfall from the original proposal.

It is no design alternative case, and many problems are not well mastered. Many algorithm problems and details have not been sorted out in place, and the basic materials and extracurricular materials are insufficient. The time management does not take into account that the time required for random planar graph and modularization increases sharply with the increase of the number of nodes, which results in the experiment with a larger number of nodes that has not been completed.

3.3.Suggestions for future development.

- 3.3.1. I hope to complete more random plane graph experiments with large numbers of nodes, and convert into more empire graphs graphs under different r values. And explore the limit value of r, and finally observe the distribution of modularity of random planar with different numbers of nodes and empire graphs with different vertices numbers, and get more experimental data
- global optimization on large networks without resorting to the idea of greedy algorithms to complete all local optimizations to simulate 3.3.2. I hope that with the development of hardware and the introduction of new algorithms in the near future, we can directly perform global optimization[7,8]



4.CONCLUSION:

With the increase of the number of nodes in the random planar graph, the longer the number of iterations is required, the topological structure, which makes its modularity higher and more community groups. As the r value increases, more topology is lost, and the modularity value also decreases. The value of modularity is positively correlated with the number of iterations to generate the random graph, the number of nodes in the random plane graph, and the value of r used to Based on the results of the two programs that have been completed, we analyzed and summarized the following conclusions. number of vertices for the program to run, and the number of iterations must be multiples of the number of vertices to have good structure. Thus, when random planar graph transformed into an empire graph, the empire graph can have a must be graph also increases exponentially. The number of iterations time to generate the random planar transform the empire graph

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