Path-based depth-first search for strong and biconnected components

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May 27, 2017

Outline

Several Questions

- One-pass or two-pass?
- LOWPOINT?
- Ear decomposition?
- Compele version?
- Robbin's Theorem?

Outline

1. Introduction

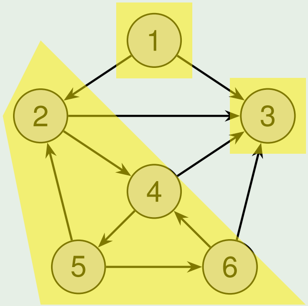
2. Strong Components
   - Thinking about Strong Components
   - Purdom and Munro's high-level algorithm
   - Contribution

Review: What have we learned from the textbook?

Concepts of Strong Components

- Two mutually reachable vertices are in the same *strong component*.

Example



Review: What have we learned from the textbook?

Algorithms to Find Strong Components

- Idea: Run DFS twice. Once on the original graph G, once on the *tranposition* of G^T.
- Trick: Using *finishing times* of each vertex computed by the first DFS.
- Linear time complexity: $O(V + E)$

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Pseudo-Code

```
H = G;
while H still has a vertex v
   start a new path P = (v);
   while P is not empty
      if the last vertex v_k of P has an edge (v_k, w)
         if w belongs to P
            contract the cycle v_i(w), ... , v_k, both
                in H and in P; /* w and v_i are
                identical. */
         else
            add w to P, as the new last vertex of P;
         end if
```

- Note that *contracting* means selecting one vertex as a representation and merging others rather than deleting them.

Pseudo-Code (Continued)

```
        else
            output v_k as a vertex of the strong component
                graph;
            delete v_k from both H and P;
        end if
    end
end
```

Assessment

- The time consumption of each statement in the pseudo-code is clear. Total time complexity is linear. except the statement in line 7:

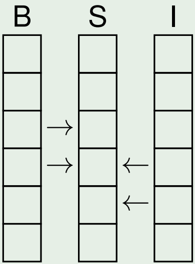
```
7 contract the cycle v_i(w), ... , v_k, both in H and in
      P; /* w and v_i are identical. */
```

- Problem is how to merge in linear time while keeping the next time accessing this vertex still in constant time.
- Therefore, a good data structure for disjoint set merging is needed.

Examples

Block B

The second line.



Block C

The third line.

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His Contribution

- He gave a simple list-based implementation that achieves linear time.
- Use only stacks and arrays as data structure.

New algorithm to discover strong components

Procedure 1: STRONG(G)

empty stacks S and B;
for $v \in V$ **do**
  $I[v] = 0$;
$c = n$;
for $v \in V$ **do**
  if $I[v] = 0$ **then**
    DFS(v);

New algorithm to discover strong components

Procedure 2: DFS(v)

```
PUSH(v,S);   I[v]=TOP(S);   PUSH(I[v],B);
/* add v to the end of P                                   */
for egdes(v, w) ∈ E do
    if I[v] = 0 then
        DFS(w);
    else /* contract if necessary                          */
        while I[w] < B[TOP[B]] do
            POP(B);
    if I[v] = B[TOP(B)] then
        /* number vertices of the next strong component    */
        POP(B);
        increase c by 1;
        while I[v] ≤ TOP[S] do
            I[POP(S)]=c;
```

Summary

- The first main message of your talk in one or two lines.
- The second main message of your talk in one or two lines.
- Perhaps a third message, but not more than that.

- Outlook
  - Something you haven't solved.
  - Something else you haven't solved.

For Further Reading I

A. Author.
Handbook of Everything.
Some Press, 1990.

S. Someone.
On this and that.
Journal of This and That, 2(1):50–100, 2000.