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Laboratory practice No. 2: Big O Notation

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1) ONLINE EXERCISES (CODINGBAT)

1.a. Array II

```
i.
                                                               // c0
          public int[] zeroFront(int[] nums) {
               boolean [] used = new boolean [nums.length]; // c1
                                                              // c2
               int cont = 0;
               for (int i = 0; i < nums.length; i++) {</pre>
                                                              // c3 + n
                 if(nums[i] == 0) {
                                                              // c4 + n
                   if (i != cont) {
                                                              // c5 + n
                     nums[i] = nums[cont];
                                                              // c6 + n
                                                              // c7 + n
                     nums[cont] = 0;
                   }
                                                               // c8 + n
                   cont++;
                 }
               }
                                                               // c9
               return nums;
             }
```

Therefore, zeroFront is $O(c_0 + c_1 + c_2 + c_3 + c_4 + c_5 + c_6 + c_7 + c_8 + c_9 + 6n)$. Applying the sum and product properties, zeroFont is O(n).

```
public int[] notAlone(int[] nums, int val) {
                                                               // c0
ii.
             for(int i = 1; i < nums.length-1; i++) {</pre>
                                                               // c1 + n
               if(nums[i] == val && nums[i-1] != val
                 && nums[i+1] != val) {
                                                               // c2 + n
                 if (nums[i-1] > nums[i+1])
                                                              // c3 + n
                                                               // c4 + n
                   nums[i] = nums[i-1];
                                                              // c5 + n
                 else
                   nums[i] = nums[i+1];
                                                               // c6 + n
```



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```
}
return nums;  // c7
```

Therefore, notAlone is $O(c_1 + c_2 + c_3 + c_4 + c_5 + c_6 + c_7 + 6n)$. Applying the sum and product properties, notAlone is O(n).

tripleUp is $O(c_0 + c_1 + c_2 + c_3 + n - 2)$. When we apply the product and sum properties, tripleUp is O(n).

```
iv.
        public int[] tenRun(int[] nums) {
                                                                 // c0
          int tempMult = 0;
                                                                 // c1
          boolean used = false;
                                                                 // c2
          for(int i = 0; i < nums.length; i++) {</pre>
                                                                 // c3 + n
            if (nums[i] % 10 == 0) {
                                                                 // c4 + n
              used = true;
                                                                 // c5 + n
                                                                 // c6 + n
              tempMult = nums[i];
            }
                                                                 // c7 + n
            if (used)
              nums[i] = tempMult;
          }
          return nums;
        }
       public int[] shiftLeft(int[] nums) {
\mathbf{v}.
          int [] mod = new int[nums.length];
          if (nums.length==1) return nums;
          for (int i=1; i<nums.length; i++) {</pre>
            mod[nums.length-1]=nums[0];
            mod[i-1] = nums[i];
          }
          return mod;
        }
```



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1.b. Array III

```
i.
         public int[] seriesUp(int n) {
            int no = n*(n+1)/2;
            int [] nums = new int [no];
            int a = 0;
            for (int i = 1; i <= n; i++) {
              for (int j = 1; j \le i; j++) {
                nums[a] = j;
                a++;
              }
            }
            return nums;
ii.
         public int countClumps(int[] nums) {
            int c = 0;
            for (int i = 0; i < nums.length-1; i++) {
              if (nums[i] == nums[i+1]) {
                for (int j = i; j < nums.length; j++) {
                  if (nums[j] != nums[i]) {
                    i = j;
                    c++;
                  }
                  if (c == 0 \&\& j == nums.length-1) {
                    C++;
                }
              }
            }
            return c;
iii.
          public boolean linearIn(int[] outer, int[] inner) {
            int j = 0;
            int c = 0;
            if (inner.length == 0) {
              return true;
            }
            for (int i = 0; i < outer.length; i++) {</pre>
              if (inner[j] == outer[i]) {
                j++;
                if (j==inner.length) {
                  return true;
```

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```
}
             }
           }
           return false;
         }
iv.
         public int[] fix45(int[] nums) {
           boolean [] arr = new boolean[nums.length];
           for (int i = 0; i < nums.length-1; i++) {
             if (nums[i] == 4 && nums[i+1] == 5) {
               arr[i+1] = true;
             } else if (nums[i] == 4 && nums[i+1] != 5) {
               for (int j = 0; j < nums.length; j++) {
                  if (nums[j] == 5 && arr[j] == false) {
                   nums[j] = nums[i+1];
                   nums[i+1] = 5;
                    arr[i+1] = true;
                    break;
                 }
               }
           }
           return nums;
         }
```

2) SIMULATION OF PROJECT PRESENTATION QUESTIONS

2.a. ArrayMax

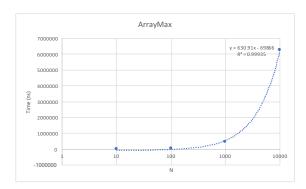


Figure 1: Time vs. N for ArrayMax

N	Time (ns)
10	6000
100	27000
1000	346000
10000	6717000

Table 1: ArrayMax's data.

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3) EXAM SIMULATION

- i. start + 1, nums, target
- ii. a) T(n) = T(n/2) + C
- iii. n-a, a, b, cres, solucionar(n-b, a, b, c)+1res, solucionar(n-c, a, b, c)+1
- iv. e) La suma de los elementos de a y es O(n).



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References