UNIVERSIDAD EAFIT SCHOOL OF ENGINEERING DEPARTMENT OF SYSTEMS AND INFORMATICS

 $\begin{array}{c} \text{Page 1 de 2} \\ \text{ST245} \\ \text{Data Structures} \end{array}$

Laboratory practice No. 5: Binary Trees

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1) CODE FOR DELIVERING ON GITHUB

The source code can be found in ParentsTree.java inside the codigo folder; everything is tested in the main method.

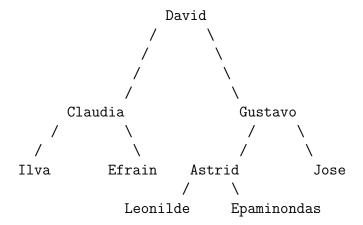
2) ONLINE EXERCISES

The source code can be found in Code. java inside the codigo folder.

3) SIMULATION OF PROYECT PRESENTATION QUESTIONS

3.a. Parents Tree???????????

You can find binary tree in ParentsTree.java, for David's family. Here's the original tree, you can find it implemented in the .java file.



UNIVERSIDAD EAFIT SCHOOL OF ENGINEERING DEPARTMENT OF SYSTEMS AND INFORMATICS

Page 2 de 2 ST245 Data Structures

3.b. Is it possible to implement a better binary tree for parentsTree?

3.c. How does exercise 2.1 work?

This exercise is really straight-forward because, when we write a tree in pre-order, we immediately know that the first integer is the root, the next integers are the left ones and, then the right ones. So, first we insert the first integer which is the root; then, as the tree is a binary search tree (BST) we know that if we insert them in another BST, it will insert them in the right order if and only if we insert them in the order they are given. So, just inserting them as the user passes them solves this problem.

3.d. What's the complexity of exercise 2.1?

```
public class Code {
   public static void main(String[] args) {
        Scanner sc = new Scanner(System.in); //c1
        System.out.println("Write quit to end the program"); //c2
        String input = sc.next(); //c3
        BinaryTree bt = new BinaryTree(); //c4
        while(!input.equals("quit")) { // c5*n
            int node = Integer.parseInt(input); // c6*n
            bt.insert(node); // c7*n*logn
            input = sc.next(); //c8*n
        }
        bt.posOrder(); // O(n)
   }
}
```

Therefore, exercise 2.1 is $O(k_0 + k_1 n + k n \log n)$. When the sum and product properties of the Big-O notation are applied, exercise 2.1 is $O(n \log n)$.

4) TEST SIMULATION

```
i. c) T(n) = 2T(n/2) + C
      i. altura(raiz.izq)+1
                                               iv.
     ii. altura(raiz.der)+1
                                                     ii. a) O(n)
ii. c) 3
                                                    iii. d)
iii.
      i. 1 == 2
                                                    iv. a)
     ii. 0
                                                      i. p.data == toInsert
                                                \mathbf{v}.
     iii. a.izq, suma-a.dato
     iv. a.der, suma-a.dato
                                                     ii. p.data > toInsert
```