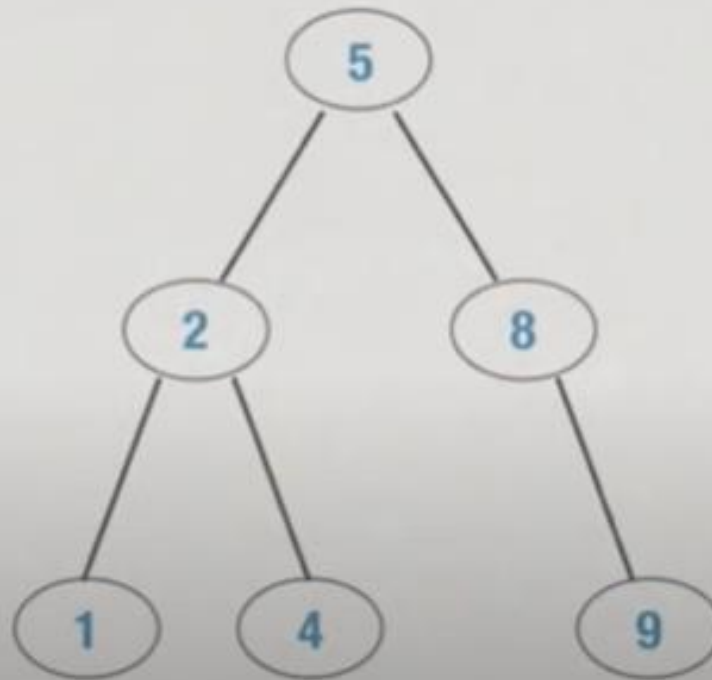


Binary search tree

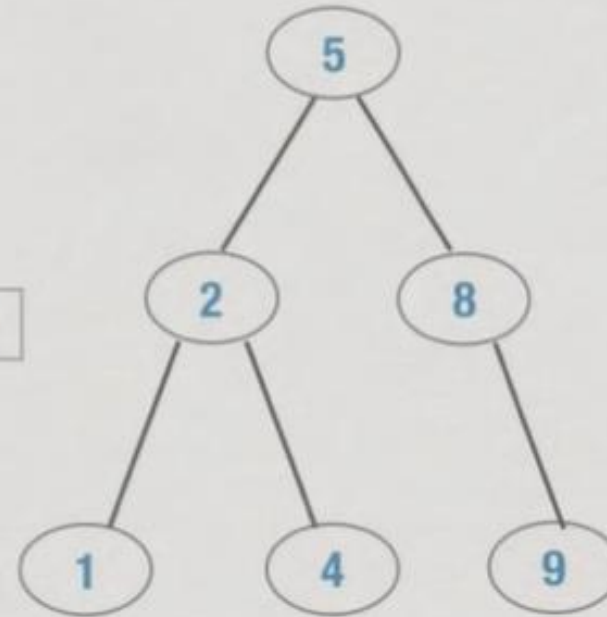
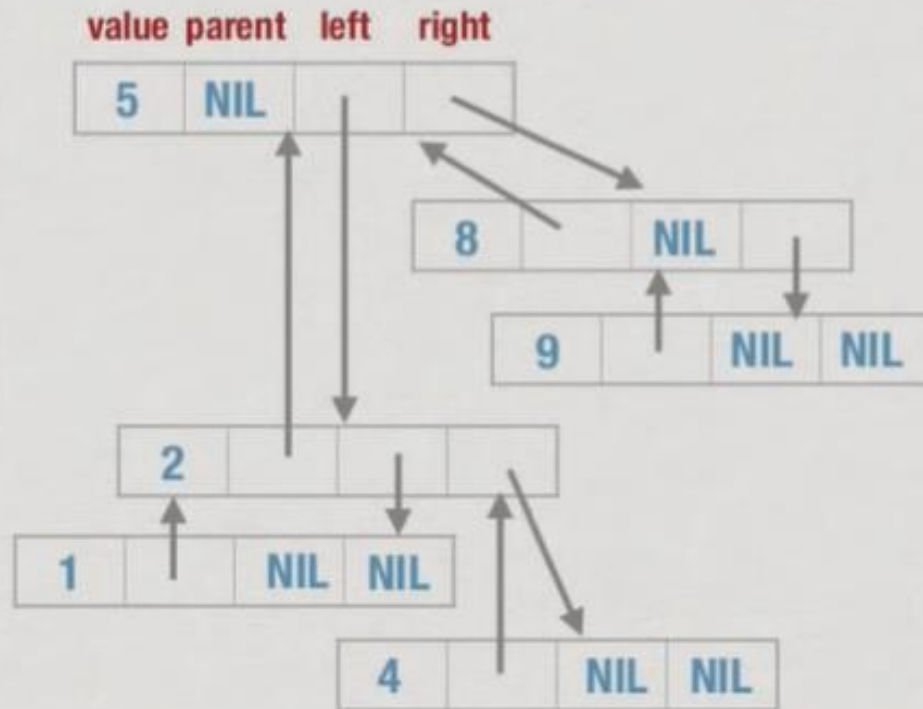
Binary search tree

- For each node with value v
- Values in left subtree $< v$
- Values in right subtree $> v$
- No duplicate values



Binary search tree

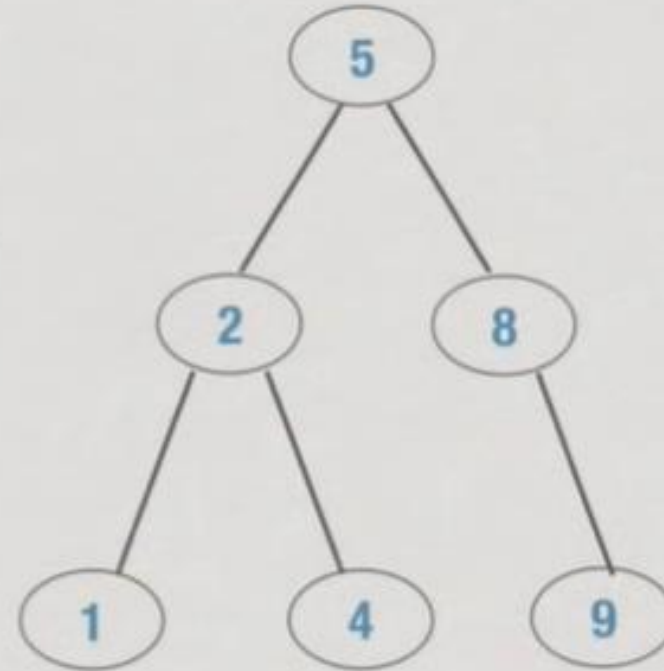
- Implement using pointers



Inorder traversal

```
function inOrder(t)
  if (t != NIL)
    inOrder(t.left)
    print(t.value)
    inOrder(t.right)
```

- Lists values in sorted order



1 2 4 5 8 9

find(v)

Recursive

```
function find(t,v)

if (t == NIL)
    return(False)

if (t.value == v)
    return(True)

if (v < t.value)
    return(find(t.left,v))
else
    return(find(t.right,v))
```

Iterative

```
function find(t,v)

while (t != NIL) {

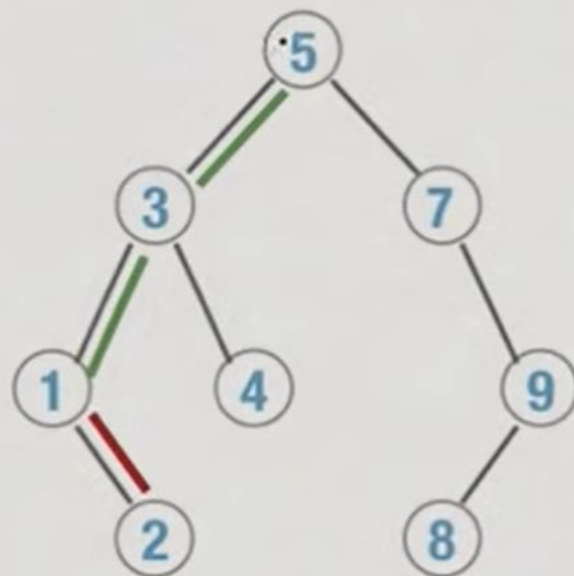
    if (t.value == v)
        return(True)

    if (v < t.value)
        t = t.left
    else
        t = t.right
}

return(False)
```

Minimum

- Left most node in the tree



Minimum

- * Left most node in the tree

Recursive

```
function minval(t)
```

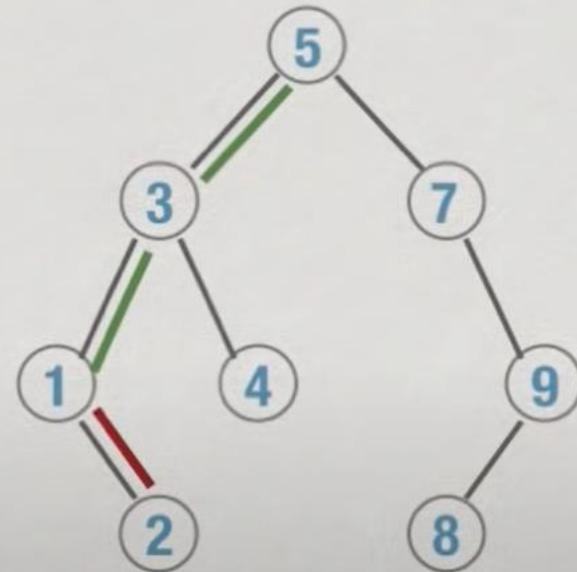
```
# Assume t is not empty
```

```
if (t.left == NIL)
```

```
    return(t.value)
```

```
else
```

```
    return(minval(t.left))
```



Minimum

- * Left most node in the tree

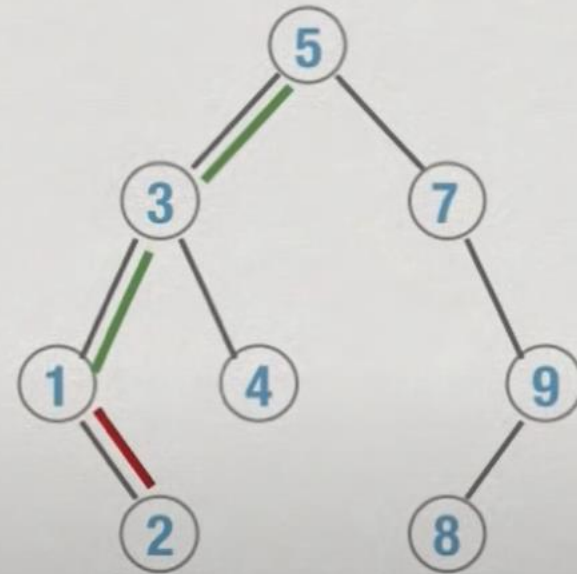
Iterative

```
function minval(t)
```

```
# Assume t is not empty
```

```
while (t.left != NIL)  
    t = t.left
```

```
return(t.value)
```

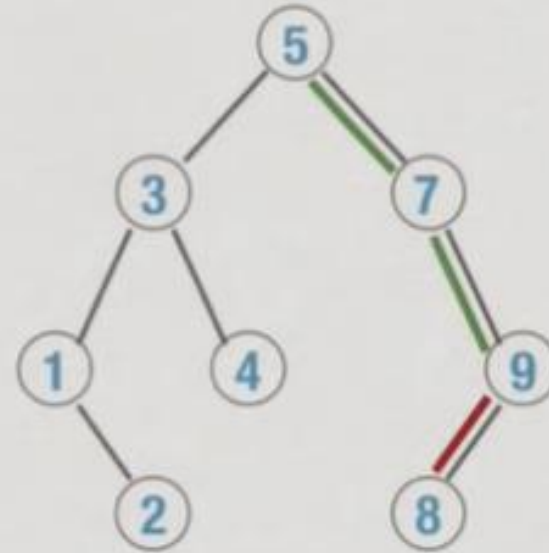


Maximum

- Right most node in the tree

Recursive

```
function maxval(t)  
# Assume t is not empty  
if (t.right == NIL)  
    return(t.value)  
else  
    return(maxval(t.right))
```

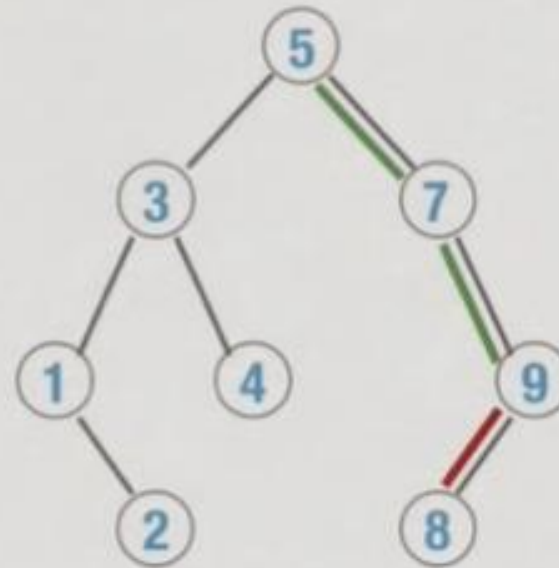


Maximum

- Right most node in the tree

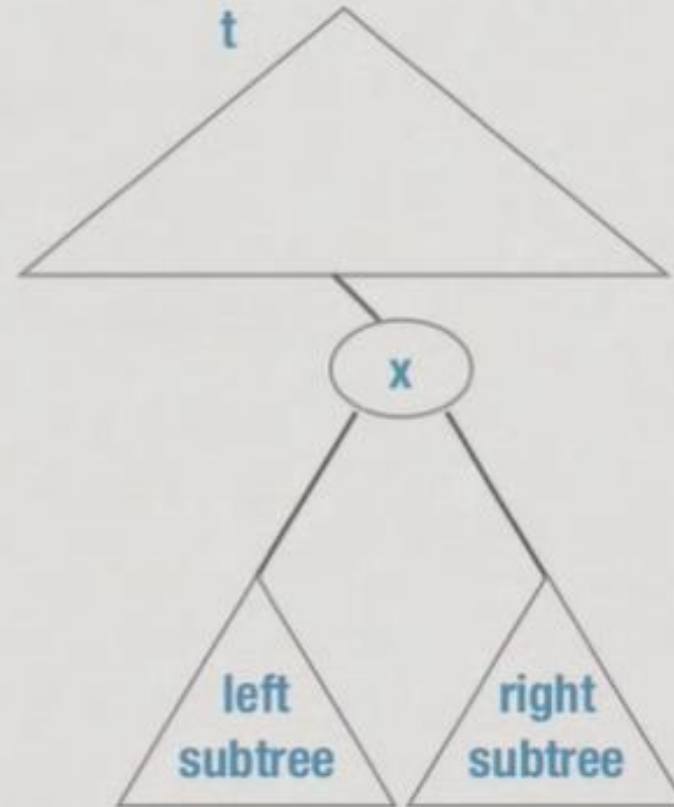
Iterative

```
function maxval(t)  
# Assume t is not empty  
while (t.right != NIL)  
    t = t.right  
return(t.value)
```



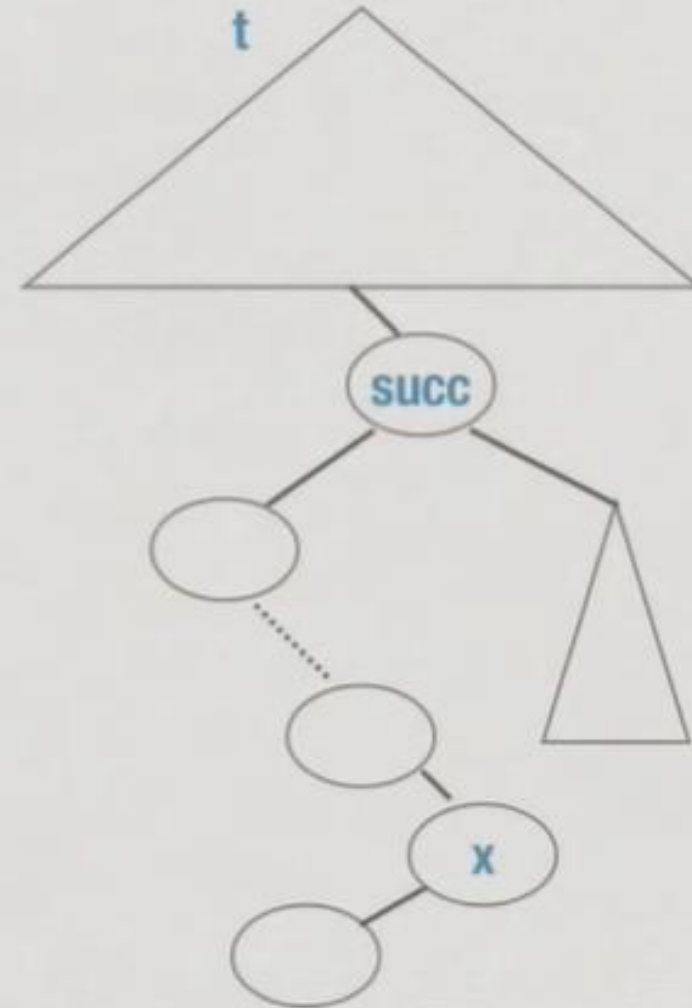
Successor

- $\text{succ}(x)$ is what $\text{inorder}(t)$ prints after x
- If x has a right subtree, $\text{min}(\text{right subtree})$



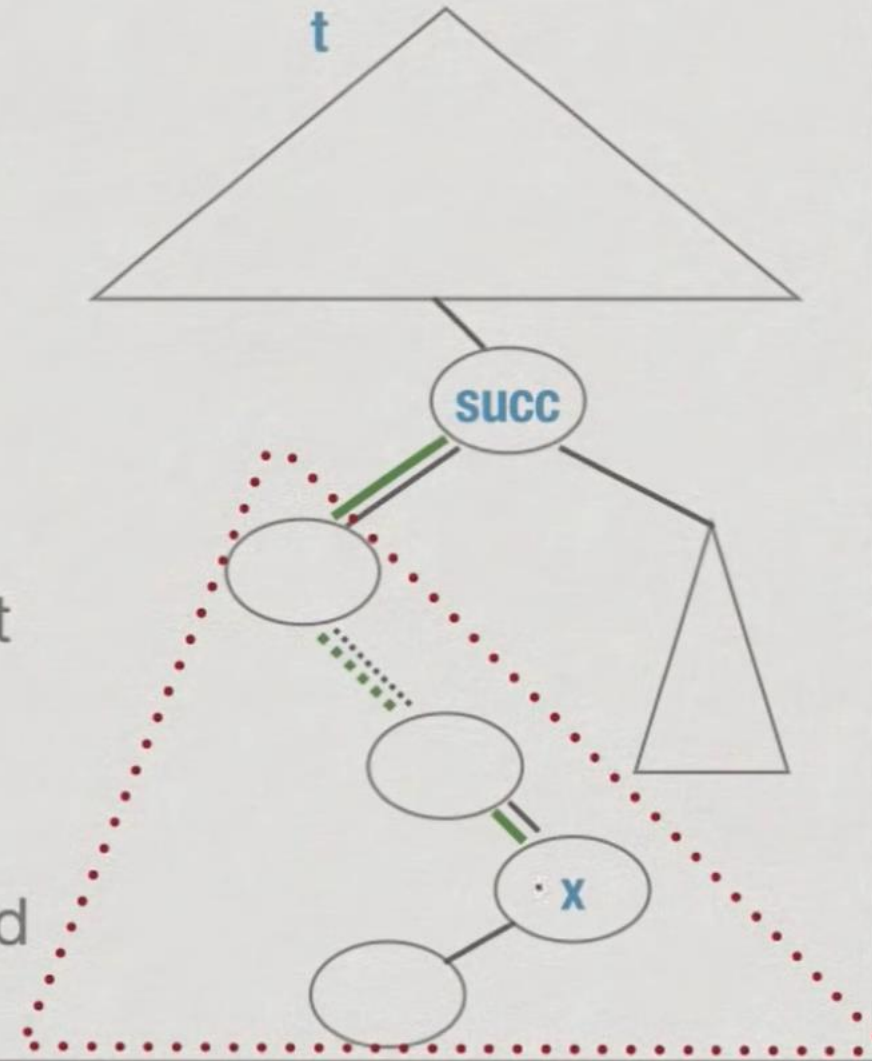
Successor

- $\text{succ}(x)$ is what $\text{inorder}(t)$ prints after x
- If x has a right subtree, $\text{min}(\text{right subtree})$
- if x has no right subtree
 - x is max of the subtree it belongs to
 - walk up to find where this subtree is connected



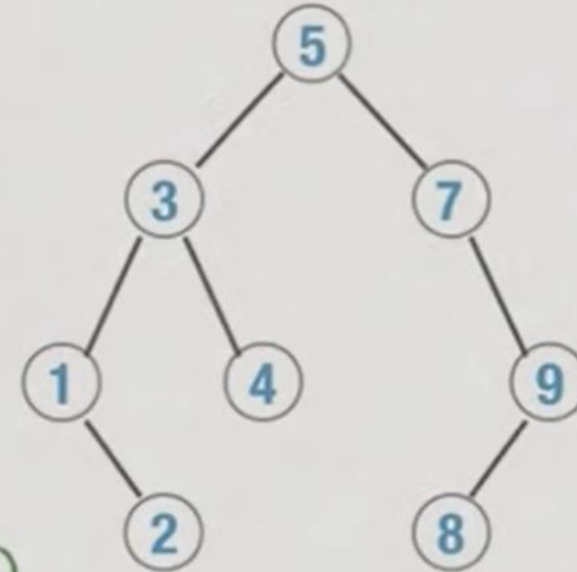
Successor

- * $\text{succ}(x)$ is what $\text{inorder}(t)$ prints after x
- * If x has a right subtree, $\text{min}(\text{right subtree})$
- * if x has no right subtree
 - * x is max of the subtree it belongs to
 - * walk up to find where this subtree is connected



Successor

```
function succ(t)
  if (t.right != NIL)
    return(minval(t.right))
  y = t.parent
  while (y != NIL and t == y.right)
    t = y
    y = y.parent
  return(y)
```



if no right sub tree then

2- 3

4-5

8-9

9- no ?

3 -4

1-2

7-8

Predecessor

- * Symmetric

```
function pred(t)
```

```
  if (t.left != NIL)
    return(maxval(t.left))
```

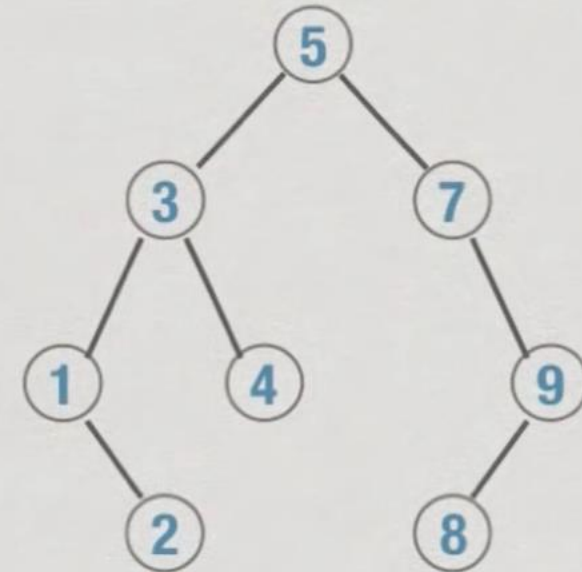
```
  y = t.parent
```

```
  while (y != NIL and t == y.left)
```

```
    t = y
```

```
    y = y.parent
```

```
  return(y)
```



Left side

5-4

3-2

9-8

1-no?

No left side

2-1

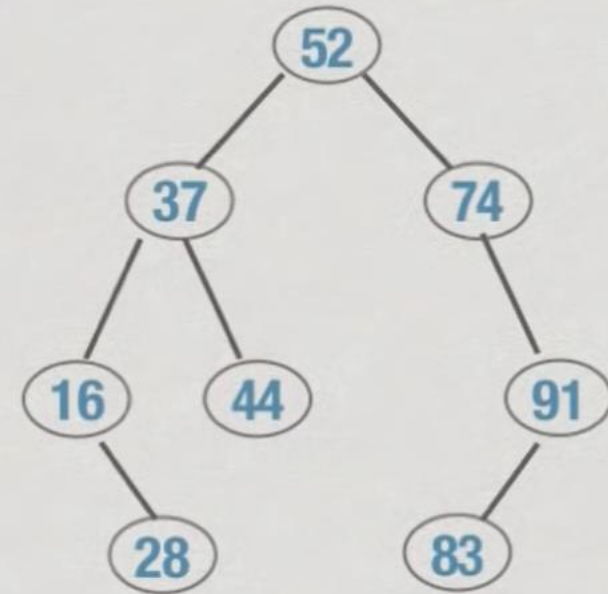
4-3

7-5

8-7

Insert

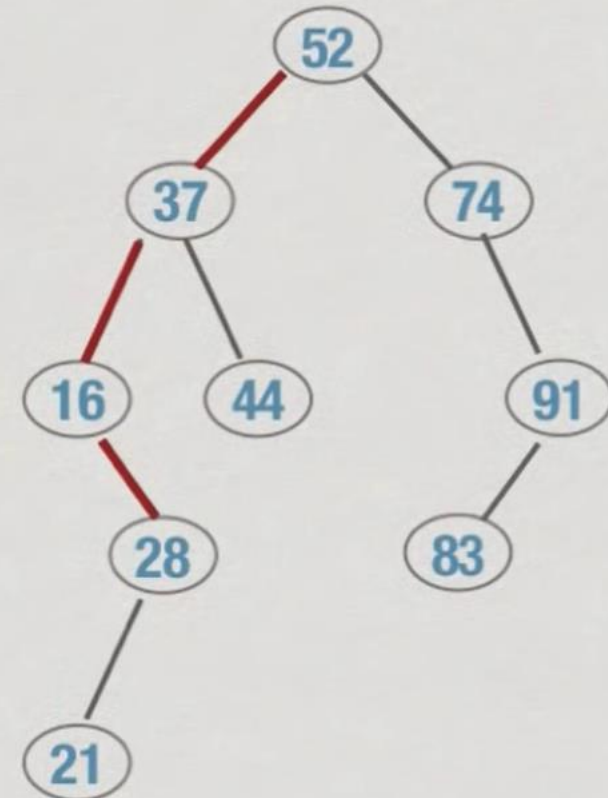
- * Try to find v
- * If it is not present, add it where the search fails



Insert

- * Try to find v
- * If it is not present, add it where the search fails

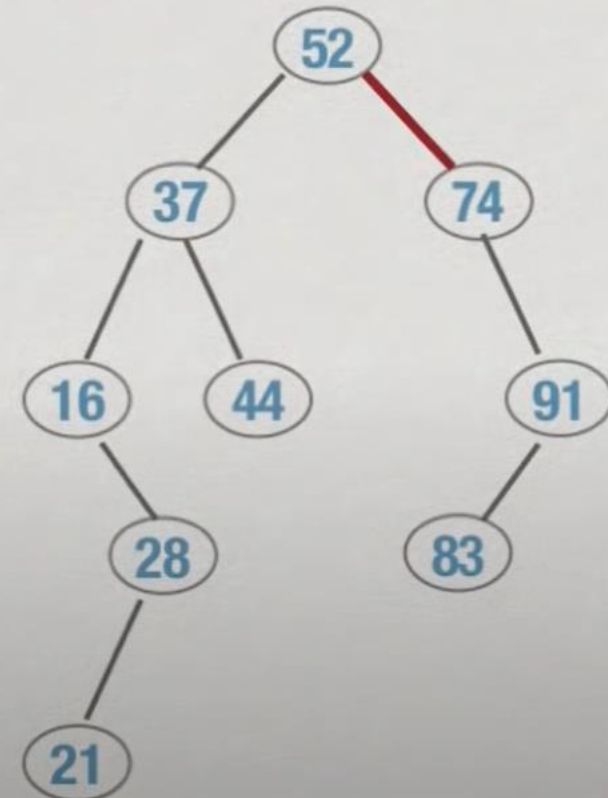
Insert 21



Insert

- * Try to find v
- * If it is not present, add it where the search fails

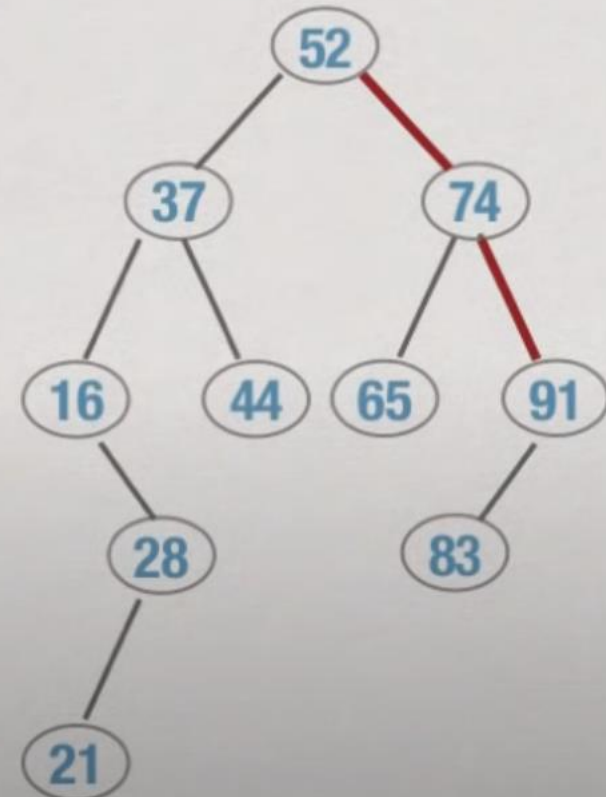
Insert 65



Insert

- * Try to find v
- * If it is not present, add it where the search fails

Insert 91



insert(v)

```
function insert(t,v)

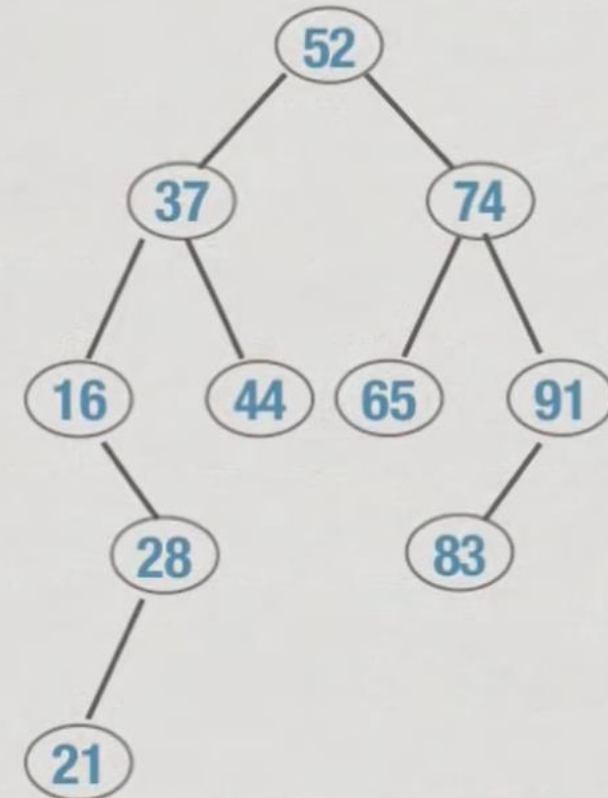
if (t == NIL)
    t = Node(v); return  # Node(v) : isolated node, value v

if (t.value == v) return

if (v < t.value)
    if (t.left == NIL) # Add a left child with value v
        t.left = Node(v); t.left.parent = t; return
    else                # Recursively insert in left subtree
        insert(t.left,v); return
else
    if (t.right == NIL) # Add a right child with value v
        t.right = Node(v); t.right.parent = t; return
    else                # Recursively insert in right subtree
        insert(t.right,v)
```

Delete

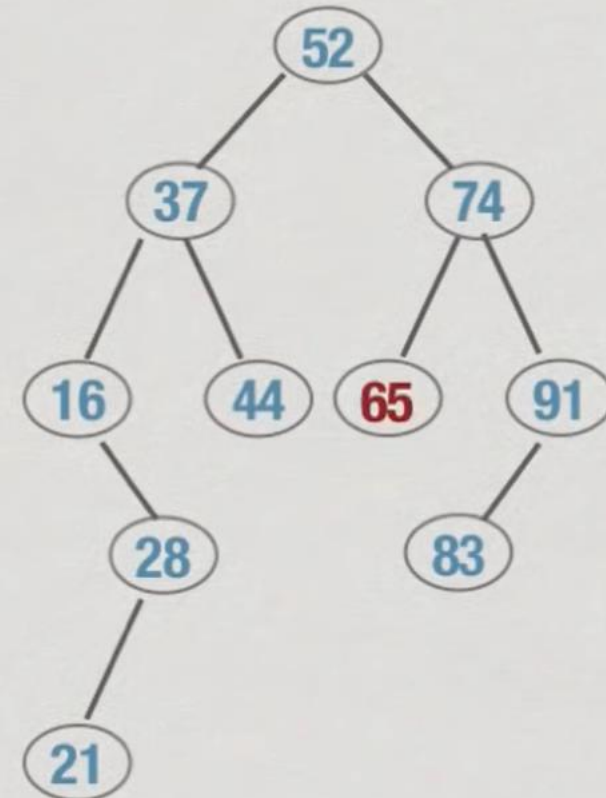
- * If v is present, delete it
- * If deleted node is a leaf, done
- * If deleted node has only one child, “promote” that child
- * If deleted node has two children, fill in the hole with $\text{pred}(v)$ or $\text{succ}(v)$
 - * Delete $\text{pred}(v)$ / $\text{succ}(v)$
 - * Either leaf or only one child



Delete

- * If v is present, delete it
- * If deleted node is a leaf, done
- * If deleted node has only one child, “promote” that child
- * If deleted node has two children, fill in the hole with $\text{pred}(v)$ or $\text{succ}(v)$
 - * Delete $\text{pred}(v)$ / $\text{succ}(v)$
 - * Either leaf or only one child

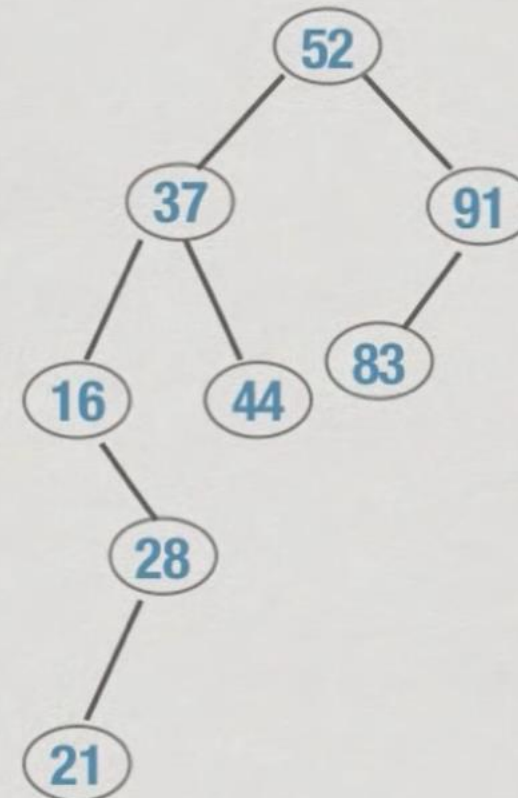
Delete 65



Delete

- * If v is present, delete it
- * If deleted node is a leaf, done
- * If deleted node has only one child, “promote” that child
- * If deleted node has two children, fill in the hole with $\text{pred}(v)$ or $\text{succ}(v)$
 - * Delete $\text{pred}(v)$ / $\text{succ}(v)$
 - * Either leaf or only one child

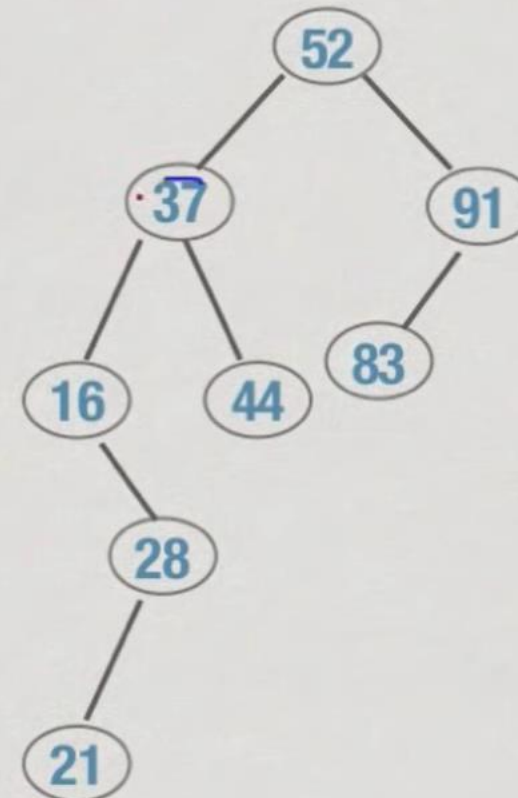
Delete 74



Delete

- * If v is present, delete it
- * If deleted node is a leaf, done
- * If deleted node has only one child, “promote” that child
- * If deleted node has two children, fill in the hole with $\text{pred}(v)$ or $\text{succ}(v)$
 - * Delete $\text{pred}(v)$ / $\text{succ}(v)$
 - * Either leaf or only one child

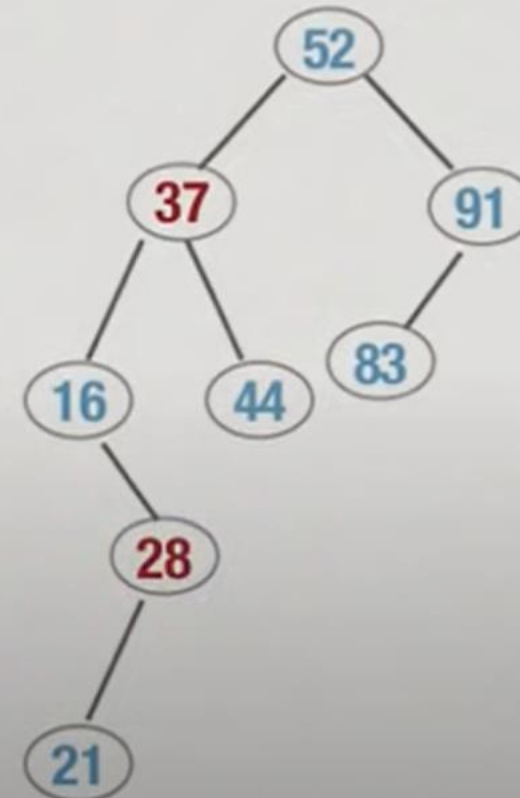
Delete 37



Delete

- If v is present, delete it
- If deleted node is a leaf, done
- If deleted node has only one child, “promote” that child
- If deleted node has two children, fill in the hole with $\text{pred}(v)$ or $\text{succ}(v)$
 - Delete $\text{pred}(v)$ / $\text{succ}(v)$
 - Either leaf or only one child

Delete 37



delete(v)

```
# t.value == v, delete here
```

```
# Delete root
```

```
if (t.parent == NIL)
    t = NIL
    return
```

```
# Delete leaf
```

```
if (t.left == NIL and t.right == NIL)
    if (t == t.parent.left)
        t.parent.left = NIL
    else
        t.parent.right = NIL
    return
```