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The vertex cover problem is an NP-Complete problem, which means that it cannot be solved in polynomial time. However, we can use approximate solutions to find a vertex cover with a size no more than twice the

optimal vertex cover.

Initialize an empty set C and consider a collection E of all the edges in the graph. While E is not empty, we pick an edge, add its two endpoints to the result set C, and remove all edges that are incident on either endpoint. We repeat this process until there are no edges left in E. Finally, we print the result set C.

The time complexity of this algorithm is O(V + E), where V is the number of vertices and E is the number of edges in the graph. The space complexity is O(V), as we use a boolean array to keep track of the already visited vertices. Overall, this approach provides a near-optimal solution to the vertex cover problem, which is useful when an exact solution is not required or feasible.