

# Part II - Graph Theory

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## 0 Introduction

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Lecture 1

**Notation.** We write  $[n]$  for  $\{1, 2, \dots, n\}$ . For a set  $X$  and  $k \in \mathbb{N}$ , define  $X^{(k)} = \{S \subset X \mid |S| = k\}$ , i.e. the set of all subsets of size  $k$ .

## 1 Fundamentals

**Definition 1.1.** A **graph** is an object  $G = (V, E)$  where  $V$  is a set and  $E \subseteq V^{(2)}$ .

$V$  is the set of vertices, and  $E$  is the set of edges.

$V(G)$  will denote  $V$ ,  $E(G)$  will denote  $E$ , and we define  $|G| = |V(G)|$  (sometimes called the order) and  $e(G) = |E(G)|$  (sometimes called the size).

**Example 1.1.** The **complete graph** on  $n$  vertices is denoted  $K_n$ . This is the graph where  $V(K_n) = [n]$  and  $E(K_n) = [n]^{(2)}$ .

**Remark.** We assume the following:

- We don't order edges;
- We don't allow loops (an edge joining a vertex to itself);
- We don't allow multiple edges;
- Most of the time,  $V(G)$  will be finite (we will explicitly say when it's not).

**Example 1.2.** The **empty graph** on  $n$  vertices, denoted  $\overline{K_n}$ , has  $V(\overline{K_n}) = [n]$  and  $E(\overline{K_n}) = \emptyset$ .

**Example 1.3.** The path of length  $n$ , denoted  $P_n$ , is a path: it has  $V(P_n) = [n+1]$  and  $E(P_n) = \{\{i, i+1\} \mid 1 \leq i \leq n\}$ .

**Example 1.4.** The cycle of length  $n$ , denoted  $C_n$ , has  $V(C_n) = [n]$  and  $E(C_n) = \{\{i, i+1\} \mid 1 \leq i \leq n-1\} \cup \{\{1, n\}\}$ .

Let  $G$  be a graph and  $x \in V(G)$ . The **neighborhood** of  $x$  is  $N(x) = \{y \mid xy \in E(G)\}$ , i.e. all the vertices connected to  $x$ . If  $y \in N(x)$ , we write  $x \sim y$  and say  $y$  is a **neighbor** of  $x$  or that  $y$  is **adjacent** to  $x$ .

The **degree** of  $x$  is  $\deg(x) = |N(x)|$ .

Just as a formality, we define graph isomorphism: let  $G, H$  be graphs. A graph isomorphism is a bijection  $\phi : V(G) \rightarrow V(H)$  such that it maps edges to edges, i.e.  $uv \in E(G) \iff \phi(u)\phi(v) \in E(H)$ .

**Definition 1.2** (Subgraph). We say  $H$  is a **subgraph** of  $G$  if  $V(H) \subseteq V(G)$  and  $E(H) \subseteq E(G)$ .

Two subgraph types that are important enough to have their own notation:

- If  $G$  is a graph, and  $xy \in E(G)$ , define  $G - xy$  to be the graph  $(V(G), E(G) \setminus \{xy\})$ .
- For  $x, y \in V(G)$ , define  $G + xy$  to be the graph  $(V(G), E(G) \cup \{xy\})$ .

**Definition 1.3** (Path). Let  $G$  be a graph,  $x, y \in V(G)$ . A  $x - y$  **path** in  $G$  is a sequence  $x_1, \dots, x_k$  where  $x_1 = x$ ,  $x_k = y$  and  $x_i x_{i+1} \in E(G) \forall 1 \leq i \leq k - 1$  and all the  $x_i$  are distinct.

**Definition 1.4.** A graph is **connected** if  $\forall x \neq y \in V(G)$ , there exists an  $x - y$  path in  $G$ .

**Remark.** A little annoyingly, if  $P$  is a  $x - y$  path and  $P'$  is a  $y - z$  path, then the concatenation  $PP'$  may not be a path (since the vertices of the new path might not be unique).

So let an  $x - y$  **walk** in a graph  $G$  be a sequence  $x_1, \dots, x_k$  where  $x_1 = x$ ,  $x_k = y$  and  $x_i x_{i+1} \in E(G) \forall 1 \leq i \leq k - 1$ . Then a concatenation of walks is again a walk.

**Proposition 1.1.** If  $W$  is an  $xy$  walk, then  $W$  contains a  $xy$  path.

*Proof.* Let  $W' \subseteq W$  be a minimal  $xy$  walk. We claim this is a path. If not, then some vertex  $x_i$  must be visited at least twice, say  $W' = x_1 x_2 \dots x_i \dots x_i x_l \dots x_k$ . Then take  $W'' = x_1 x_2 \dots x_i x_l \dots x_k$ . This contradicts the minimality of  $W'$ , so we're done.  $\square$

**Remark.** We may define a **distance** on  $V(G)$ : for  $x, y \in V(G)$ , let  $d(x, y)$  be the length of the shortest  $xy$  path. If  $G$  is connected, then this distance defines a metric on  $V(G)$ .

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Lecture 2

## 1.1 Trees

**Definition 1.5.** A graph  $G$  is **acyclic** if it does not contain a cycle as a subgraph.

**Definition 1.6.** A graph  $G$  is a **tree** if it is acyclic and connected.

**Proposition 1.2.** The following are equivalent:

1.  $G$  is a tree;
2.  $G$  is minimally connected (  $\forall xy \in E(G)$ ,  $G - xy$  is not connected);
3.  $G$  is maximally acyclic (  $\forall xy \notin E(G)$ ,  $G + xy$  contains a cycle).

*Proof.* (a)  $\implies$  (b): A tree is connected. Assume for contradiction that  $\exists xy \in E(G)$  such that  $G - xy$  is connected. Let  $P$  be a  $xy$  path in  $G - xy$ . But then  $P$  defines a cycle in  $G$ , contradiction.

(b)  $\implies$  (a): Minimally connected implies connected. For acyclicity, assume for contradiction that  $G$  contains a cycle  $C$ . Let  $xy \in E(C)$ . We claim that  $G - xy$  is connected. Choose  $u \neq v \in V(G - xy)$ . Let  $P$  be a  $uv$  path in  $G$ . If  $P$  does not contain  $xy$ , we're done. If  $P$  does contain  $xy$ , then take paths  $u \rightarrow x$ ;  $x \rightarrow y$  along our cycle without using  $xy$ ;  $y \rightarrow v$ . The concatenation of these gives a  $uv$  walk, which contains a  $uv$  path. Hence  $G - xy$  is connected, contradiction.

(a)  $\implies$  (c): A tree is acyclic. Let  $xy \notin E(G)$ ,  $x \neq y$ . Let  $P$  be a  $xy$  path. Then  $P$  defines a cycle in  $G + xy$ .

(c)  $\implies$  (a): We have acyclicity. If  $G$  is not connected,  $\exists x \neq y \in V(G)$  with no  $xy$  path. Then  $G + xy$  is acyclic.  $\square$

**Definition 1.7.** If  $T$  is a tree and  $v \in V(T)$  with  $\deg(v) = 1$ , we call  $v$  a **leaf**.

**Definition 1.8.** Let  $G$  be a graph and  $X \subseteq V(G)$ . Then

$$G[X] = (X, E(G) \cap \{(u, v) \mid u, v \in X\}).$$

We call this **the graph induced on  $X$** .

**Definition 1.9.** If  $x \in V(G)$ , define  $G - x = G[V(G) \setminus \{x\}]$ .

**Proposition 1.3.** Let  $T$  be a tree,  $|T| \geq 2$ . Then  $T$  has a leaf.

*Proof.* Let  $P = x_1 \dots x_k$  be the a longest possible path in  $T$ . Note  $N(x_k) \subseteq \{x_1, \dots, x_{k-1}\}$ . If  $x_i \sim x_k$  for some  $1 \leq i \leq k-2$ , there is a cycle in  $T$ , contradiction. Thus  $N(x_k) = \{x_{k-1}\} \implies x_k$  is a leaf.  $\square$

**Remark.** We can show that any  $T$  has two leaves, but we can't do any better (consider a path).

**Remark.** We could have also proved this by taking a non-backtracking walk in  $G$  (i.e. go out of a different edge then you came in on). Either we return (hence have a cycle) or get stuck (hence have a leaf).

**Proposition 1.4.** Let  $T$  be a tree on  $n \geq 1$  vertices. Then  $e(G) = n - 1$ .

*Proof.* By induction.  $n = 1$  is trivial. Assume the claim holds for  $n$ . Take a tree  $T$  with  $n + 1$  vertices. Let  $x \in V(T)$  be a leaf. Then  $T - x$  is connected and acyclic, therefore a tree, thus  $e(T - x) = n - 1$ . But  $e(G) = e(G - x) + 1$  and  $|V(G)| = |V(G - x)| + 1$ , hence we're done.  $\square$

**Definition 1.10.** Let  $G$  be a connected graph. Then a subgraph  $T$  of  $G$  is a **spanning tree** if  $T$  is a tree on  $V(G)$ .

**Proposition 1.5.** Every connected graph contains a spanning tree.

*Proof.* Start with the graph  $G$ , then throw away edges of  $E(G)$  one by one subject to keeping the graph connected. At some point, the removal of any further edge will disconnect the graph, at which point we have a minimal connected subgraph of  $G$ , which by Prop. 1.2 is a tree.  $\square$

## 1.2 Bipartite graphs

**Definition 1.11.** Let  $G = (V, E)$  be a graph.  $G$  is **bipartite** if there exists a partition  $V = A \cup B$  such that  $E(G) \subseteq \{uv \mid u \in A, v \in B\}$ .

**Definition 1.12.** The **complete bipartite graph**  $K_{n,m}$  is the graph with vertex set  $A \cup B$ ,  $A = \{x_1, \dots, x_n\}$ ,  $B = \{y_1, \dots, y_m\}$  and edge set  $E(K_{n,m}) = \{x_i y_j \mid x_i \in A, y_j \in B\}$ .

**Remark.** There obviously exist non-bipartite graphs: odd cycles are not bipartite.

**Definition 1.13.** A **circuit** is a sequence  $x_1, x_2, \dots, x_l x_{l+1}$ , where  $x_i x_{i+1} \in E(G)$  and  $x_{l+1} = x_1$ . The length of this circuit is  $l$ . We say a circuit is **odd** if its length is odd.

**Proposition 1.6.** Let  $C$  be an odd circuit in a graph  $G$ . Then  $C$  contains an odd cycle.

*Proof.* Let  $x_1 x_2 \dots x_i x_{i+1} \dots x_i x_k \dots x_l x_1$  be an odd circuit. Consider the circuits  $C_1 = x_1 \dots x_i x_k \dots x_l x_1$  and  $C_2 = x_i x_{i+1} \dots x_k x_{k+1} \dots x_i$ . Then one of  $C_1, C_2$  has odd length and is strictly shorter, so we're done by induction.  $\square$

**Theorem 1.7.** Let  $G$  be a graph. Then

$$G \text{ is bipartite} \iff G \text{ does not contain an odd cycle.}$$

*Proof.* ( $\implies$ ): If  $G$  contains an odd cycle, then as odd cycles are not bipartite,  $G$  cannot be bipartite.

( $\impliedby$ ): We may assume that  $G$  is connected. Let us fix  $x_0 \in V(G)$ . Let

$$V_0 = \{x \in V(G) \mid d(x, x_0) \equiv 0 \pmod{2}\}$$

$$V_1 = \{x \in V(G) \mid d(x, x_0) \equiv 1 \pmod{2}\}.$$

We claim this is a bipartition of  $G$ . Assume for contradiction that  $\exists u, v \in V_0$  s.t.  $uv \in E(G)$ . But there is an even  $ux_0$  path and an even  $vx_0$  path, thus putting these three paths together gives an odd circuit in  $G$ . By Prop 1.6,  $G$  contains an odd cycle, contradiction. (Analogous proof for  $V_1$ ).  $\square$

### 1.3 Planar graphs

**Definition 1.14.** A **planar graph** is a graph that can be drawn in the plane with no edge crossings.

**Example 1.5.**  $K_4$  is planar. A path  $P_n$  is planar.

**Definition 1.15.** A **plane graph** is one such drawing of a graph in the plane without edge crossings.

Note that this is important, since we can draw  $K_4$  in a way that it does have edges crossing.

**Example 1.6.**  $K_{2,3}$  is planar.  $K_{3,3}$  is not planar.  $K_5$  is not planar (we don't prove this right now).

**Question.** What graphs are planar? Is there a (simple) method to decide if a graph is planar?

**Definition 1.16.** Let  $G$  be a plane graph. Consider  $\mathbb{R}^2 \setminus G$ . This is broken into finitely many regions. These are called the **faces** of the plane graph.

**Definition 1.17.** The **boundary** of a face  $F$  is the collection of vertices and edges on the topological boundary.

**Remark.** The boundary of a face need not be a cycle. It need also not be connected. In fact, the boundary need not contain a cycle at all (consider a tree).

**Remark.** We also note that two different drawings of a graph in the plane can be fundamentally different.

**Theorem 1.8** (Euler). Let  $G$  be a connected plane graph with  $n$  vertices,  $m$  edges and  $f$  faces. Then  $n - m + f = 2$ .

*Proof.* We induct on  $m$ .  $m = 1$  is clear. If  $G$  is acyclic, then  $G$  is a tree, so  $m = n - 1$ ,  $f = 1$  and we're done.

So assume  $G$  contains a cycle and let  $e$  be an edge on this cycle. Delete  $e$ . Then  $n$  stays fixed,  $m$  decreases by 1, and  $f$  decreases by 1, so by induction,  $n - (m - 1) + (f - 1) = 2$  and we're done.  $\square$

**Remark.** We really do need the graph to be connected, consider  $t$  triangles in the plane as a counterexample.

**Corollary 1.9.** Let  $G$  be a planar graph,  $|G| \geq 3$ . Then  $e(G) \leq 3|G| - 6$ .

*Proof.* Draw  $G$  in the plane. We may assume that  $G$  is connected. Let  $F$  be a face, let  $\deg(F)$  = the number of edges in  $G$  that touch  $F$ . Note  $\deg(F) \geq 3$ . Now note that since every edge touches at most two faces, we get

$$3f \leq \sum_{F \text{ a face}} \deg(F) \leq 2e(G) \implies f \leq \frac{2}{3}e(G).$$

Put this into Euler's formula to get

$$n - \frac{1}{3}e(G) \geq n - e(G) + f = 2 \implies 3(n - 2) \geq e(G).$$

□

**Remarks.** (i): This is a statement about planar graphs only.

(ii): This is quite restrictive.  $K_n$  has  $\binom{n}{2} \approx n^2/2$  edges, while our above corollary says the number of edges of a planar graph is linear in  $n$ .

**Corollary 1.10.**  $K_5$  is not planar.

*Proof.* We have  $e(K_5) = 10, n = 5$ , so  $10e(G) \not\leq 3|G| - 6 = 9$ , so we're done by the above corollary. □

But  $K_{3,3}$  does not fail this test. So we need to improve our argument:

**Corollary 1.11.** Let  $G$  be a planar graph,  $|G| \geq 4$  and  $G$  has no cycles of length 3. Then  $e(G) \leq 2|G| - 4$ .

*Proof.* Repeat the proof of Corollary 1.9, but use  $\deg(F) \geq 4$  for every face. □

Now we can see that  $K_{3,3}$  is not planar.  $K_{3,3}$  has no cycle of length 3 by definition,  $n = 6, e(G) = 9$ , so  $9 = e(G) \not\leq 2 \cdot (6 - 2) = 8$ .

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**Definition 1.18.** A **subdivision** of a graph  $G$  is a subgraph where we replace some of the edges of  $G$  with disjoint paths.

**Observation.** A subdivision of a **non-planar** graph is non-planar.

**Observation.** If  $G$  contains a  $K_{3,3}$  or  $K_5$  subdivision as a subgraph, then  $G$  is non-planar.

**Theorem 1.12** (Kuratowski's theorem).  $G$  is planar  $\iff G$  does not contain a subdivided  $K_{3,3}$  or  $K_5$ .

We do not prove this, but the proof is actually not too hard.

## 2 Connectivity & matching

### 2.1 Matching in bipartite graphs

Let  $G = (X \sqcup Y, E)$  be bipartite with bipartition  $X, Y$ .

**Definition 2.1.** A **matching from  $X$  to  $Y$**  is a set of edges  $\{xy_x \mid x \in X, y_x \in Y\}$  and  $x \rightarrow y_x$  is an injection.

**Question.** When does a bipartite graph have a  $X$  to  $Y$  matching?

We can first think about examples where we do not have a matching. For example, we clearly have no matching if  $|X| > |Y|$ .

**Definition 2.2.** Let  $G$  be a graph,  $A \subseteq V(G)$ . Define  $N_G(A) = \bigcup_{x \in A} N(x)$ .

Then we clearly also don't have a matching if we have  $A \subset X$  such that  $|N(A)| < |A|$ . But this is actually the only obstruction:

**Theorem 2.1** (Hall's Marriage Theorem). Let  $G$  be a bipartite graph  $G = (X \sqcup Y, E)$ . Then

$$G \text{ has a matching from } X \text{ to } Y \iff \forall A \subseteq X, |N(A)| \geq |A|.$$

The right-hand side is called Hall's criterion.

*Proof.* ( $\implies$ ) is the easy direction.

Now let  $M$  be a matching and let  $A \subseteq X$ . Then if  $\{y_1, \dots, y_{|A|}\}$  are matched to  $A$ , we show  $|N(A)| \geq |\{y_1, \dots, y_{|A|}\}| \geq |A|$ .

( $\impliedby$ ): Apply induction on  $|X|$ . If  $|X| = 1$ , we're done. For the induction step, consider the following question: is there  $\emptyset \neq A \subsetneq X$  such that  $|N(A)| = |A|$ ?

If the answer is no, then  $\forall A \subsetneq X$  we have  $|N(A)| \geq |A| + 1$ . Let  $xy \in E(G)$  and let  $G' = G[X \setminus \{x\} \cup Y \setminus \{y\}]$ . We now check Hall's criterion for  $G'$ . If  $B \subseteq X \setminus \{x\}$ , then  $|N_{G'}(B)| \geq |N_G(B)| - 1 \geq |B|$ , so done by induction.

If the answer is yes, then let  $G_1 = G[A \cup N(A)]$  and  $G_2 = G[X \setminus A \cup Y \setminus N(A)]$ .

Claim 1:  $G_1$  satisfies Hall's criterion. Let  $B \subseteq A$ , then

$$|N_{G_1}(B)| = |N_G(B)| \geq |B|.$$

Claim 2:  $G_2$  satisfies Hall's criterion. Let  $B \subset X \setminus A$ . Consider  $N_G(A \cup B)$ . On the one hand,  $|N_G(A \cup B)| \geq |A| + |B|$ . On the other hand,  $|N_G(A)| + |N_{G_2}(B)| = |N_G(A \cup B)|$ . As  $|N(A)| = |A|$ , we get  $|N_{G_2}(B)| \geq |B|$ .

From claims 1 and 2 we can apply induction in  $G_1, G_2$  to get a matching in these graphs, and then put them together to get a matching in  $G$ .  $\square$



**Definition 2.3.** A matching of deficiency of  $d$  from  $X$  to  $Y$  is a matching from  $X'$  to  $Y$  where  $X' \subseteq X$ ,  $|X| - d = |X'|$ .

**Theorem 2.2** (Defect Hall's Theorem).

$$G \text{ contains a matching of deficiency } d \iff \forall A \subseteq X, |N(A)| \geq |A| - d.$$

*Proof.* ( $\implies$ ) : easy.

( $\impliedby$ ) : Add  $d$  phantom vertices to  $Y$ , which we join to all vertices in  $X$ , so we now satisfy Hall's condition. Apply Hall to get a matching, and then remove the  $d$  vertices we added, which removes at most  $d$  elements of  $X$ .  $\square$

**Definition 2.4.** Let  $G$  be a graph. The **minimum degree** in  $G$  is  $\delta(G) = \min_{x \in V(G)} d(x)$ , and the **maximal degree** in  $G$  is  $\Delta(G) = \max_{x \in V(G)} d(x)$ .

**Definition 2.5.** A graph is **regular** if  $\delta(G) = \Delta(G)$ . It is  **$k$ -regular** if  $k = \delta(G) = \Delta(G)$ .

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**Corollary 2.3.** For  $k \geq 1$ , if  $G = (X \sqcup Y, E)$  is a  $k$ -regular bipartite graph, then there exists a matching from  $X$  to  $Y$ .

*Proof.* We check Hall's criterion. Let  $A \subseteq X$ . On the one hand,

$$e(G[A \cup N(A)]) = \sum_{v \in A} \deg(v) = k|A|.$$

On the other hand,

$$e(G[A \cup N(A)]) \leq \sum_{v \in N(A)} \deg(v) \leq k|N(A)|.$$

Hence  $|N(A)| \geq |A|$  and we're done.  $\square$

Let  $\Gamma$  be a finite group and let  $H$  be a subgroup of  $\Gamma$ . Let  $L_1, \dots, L_n$  be the set of left cosets and  $R_1, \dots, R_n$  be the right cosets (of the forms  $gH$  and  $Hg$  respectively).

**Question.** Is there  $g_1, \dots, g_n \in \Gamma$  such that  $g_1H, \dots, g_nH$  are the left cosets and  $Hg_1, \dots, Hg_n$  are the right cosets?

**Corollary 2.4.** There exist  $g_1, \dots, g_n \in \Gamma \setminus H$  such that  $g_1H, \dots, g_nH$  are the left cosets and  $Hg_1, \dots, Hg_n$  are the right cosets.

*Proof.* It is enough to find a pairing  $L_i \leftrightarrow R_{\sigma(i)}$  such that  $L_i \cap R_{\sigma(i)} \neq \emptyset \forall i$ . Then choose  $g_i \in L_i \cap R_{\sigma(i)}$  and we have  $g_iH = L_i$ ,  $Hg_i = R_{\sigma(i)}$ .

Define  $X = \{R_1, \dots, R_n\}$  and  $Y = \{L_1, \dots, L_n\}$ , and define  $R_i \sim L_j$  when  $R_i \cap L_j \neq \emptyset \forall i, j$ . Let  $A = \{r_{i_1}, \dots, r_{i_k}\}$ . Note

$$\left| \bigcup_{j=1}^k R_{i_j} \right| = k|H|.$$

But  $L_1, \dots, L_n$  partition  $\Gamma$  and  $|L_i| = |H|$ , so at least  $k$  left cosets must intersect  $\bigcup R_{i_j}$ . Thus Hall's criterion is satisfied and we're done.  $\square$

## 2.2 Connectivity

For a tree,  $G - x$  (where  $x$  is any non-leaf) is disconnected. On the other hand, remove any 2 vertices from the Petersen graph and it stays connected (but if you remove any 3, you disconnect it).

**Notation.** Let  $S \subseteq V(G)$ , and let  $G - S = G[V(G) \setminus S]$ .

**Definition 2.6.** Let  $G$  be a graph,  $|G| \geq 1$ . Define

$$\kappa(G) = \min\{|S| \mid \exists S \subseteq V(G), G - S \text{ is disconnected or a single vertex}\}.$$

We say a graph  $G$  is  **$k$ -connected** if  $\kappa(G) \geq k$ .

In other words,  $G$  is  $k$ -connected if and only if  $G - S$  is connected for all  $S \subseteq V(G)$ ,  $|S| \leq k - 1$ .

**Example 2.1.** •  $\kappa(\text{Tree}) = 1$ .

- $\kappa(\text{Petersen graph}) = 3$ , so we can say the Petersen graph is 3-connected.
- $\kappa(\text{Cycle}) = 2$ .
- $\kappa(K_n) = n - 1$ .

We have another natural definition of connectivity.

**Definition 2.7.** Let  $G$  be a graph and let  $a, b \in V(G)$ . Say that  $ab$  paths  $P_1, \dots, P_k$  are **disjoint** if  $V(P_i) \cap V(P_j) = \{a, b\} \forall i \neq j$ .

Amazingly, we have Menger's theorem: These two notions of connectivity ( $\#$  of disjoint paths and  $\kappa(G)$ ) are equivalent.

**Remarks:**

- We have  $\delta(G) \geq \kappa(G)$ . To see this, delete  $N(x)$  for  $x \in V(G)$  of minimal degree, then  $G - N(x)$  is disconnected (or a single vertex).
- We have  $\kappa(G - x) \geq \kappa(G) - 1$ . This is clear: if  $S \subset V(G - x)$  disconnects  $G - x$  with  $|S| \leq \kappa(G) - 2$ , then  $S \cup \{x\}$  disconnects  $G$ , contradiction.

- We can have  $\kappa(G - x) > \kappa(G)$ . For example, a cycle is 2-connected, but a cycle with one protruding edge is 1-connected.

**Definition 2.8.** A **component** in  $G$  is a maximal connected subgraph.

**Definition 2.9.** Let  $G$  be a graph, let  $a, b \in V(G), a \neq b, a \not\sim b$ . Say  $S \subseteq V(G) \setminus \{a, b\}$  is a  $a - b$  **separator** if  $G - S$  disconnects  $a$  from  $b$  (i.e.  $a, b$  are in different components of  $G - S$ ).

**Theorem 2.5** (Menger's theorem, form 1). Let  $G$  be a connected graph and fix  $a, b \in V(G), a \neq b, a \not\sim b$ . Then the minimum size of an  $a - b$  separator is equal to the maximal number of disjoint paths from  $a$  to  $b$ .

In other words, if all  $a - b$  separators have size  $\geq k$ , then there exist  $P_1, \dots, P_k$ , disjoint paths between  $a$  and  $b$ .

**Note.** Define  $\kappa_{a,b}(G)$  be the size of the minimal  $a - b$  separator.

**Note.** Recall  $\kappa(G - x) \geq \kappa(G) - 1$ , and also  $\kappa(G - xy) \geq \kappa(G) - 1$ . We also have  $\kappa_{a,b}(G - x) \geq \kappa_{a,b}(G) - 1$  and  $\kappa_{a,b}(G - xy) \geq \kappa_{a,b}(G) - 1$  (exercise, not hard).

*Proof.* Assume for contradiction that the statement of the theorem is false. Let  $G$  be a minimal counterexample to the theorem that

- minimizes  $k$ ;
- subject to (a), choose  $G$  to minimize  $e(G)$ .

Now let  $S$  be a minimal  $a, b$  separator in  $G$ . We have  $|S| = k$ . Note that the theorem is true for  $k = 1$ , so assume  $k \geq 2$ .

If  $S \neq N(a)$  and  $S \neq N(b)$ , consider  $G - S$  and let  $A$  be the component containing  $a$  and  $B$  be the component containing  $b$ .

Define  $G_a = G[A \cup S]$  along with a vertex  $c$  joined to each vertex in  $S$ , and  $G_b = G[B \cup S]$  along with a vertex  $c$  joined to each vertex in  $S$ . Note that  $\kappa_{a,c}(G_a) \geq k$ , since any  $a - c$  separator in  $G_a$  is a  $a, b$  separator in  $G$ . Likewise,  $\kappa_{b,c}(G_b) \geq k$ .

Note that  $e(G_a) < e(G), e(G_b) < e(G)$  since  $N(a) \not\subseteq S, N(b) \not\subseteq S$ . So there exists a neighbor  $x$  of  $b$  in  $B$  with  $\deg(x) \geq 2$ , else we can remove  $x$  and apply minimality.

So by minimality of  $G$ , we can find  $k$  disjoint  $a, c$  paths, say  $P_1, \dots, P_k$  in  $G_a$ , and likewise we can find  $k$  disjoint  $b, c$  paths  $Q_1, \dots, Q_k$  in  $G_b$ . We can put these paths together to get paths  $P_1 Q_{\sigma(1)}, \dots, P_k Q_{\sigma(k)}$ , which are  $k$  disjoint  $a, b$  paths, contradiction, done.

Let us now assume WLOG that  $S = N(a)$ .

**Claim:**  $N(a) \cap N(b) = \emptyset$ .

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Indeed, if  $\exists x \in N(a) \cap N(b)$ , consider  $G - x$ . We have  $\kappa_{a,b}(G - x) \geq k - 1$ . Thus, by minimality, we can find  $k - 1$  disjoint  $ab$  paths in  $G - x$ , so all of these, plus  $axb$ , gives us  $k$  disjoint  $ab$  paths in  $G$ , contradiction.

Let  $ax_1 \dots x_l b$  be a shortest  $ab$  path. Note that  $l \geq 2$  and  $x_2 \neq b$ . Consider  $G - x_1 x_2$ . We must have  $\kappa_{a,b}(G - x_1 x_2) \leq k - 1$  by minimality, so  $\kappa_{a,b}(G - x_1 x_2) = k - 1$ . So there is a  $a, b$  separator  $\tilde{S}$ ,  $|\tilde{S}| = k - 1$  in  $G - x_1 x_2$ . We see that  $\tilde{S} \cup \{x_1\}$  and  $\tilde{S} \cup \{x_2\}$  are  $a, b$  separators in  $G$  of size at most  $k$ . Now either  $\tilde{S} \cup \{x_1\} \neq N(a), N(b)$  or  $\tilde{S} \cup \{x_2\} \neq N(a), N(b)$ , so we're done.  $\square$

**Corollary 2.6** (Menger's theorem, form 2). Let  $G$  be a connected graph,  $|G| \geq 2$ . Then  $G$  is  $k$ -connected  $\iff \forall a, b \in V(G), a \neq b$ , there exist  $k$  disjoint  $ab$  paths in  $G$ .

*Proof.*  $\Leftarrow$  is the easy direction. Say  $G - S$  is disconnected and let  $a, b$  be in different components of  $G - S$ . Note  $a \not\sim b$ . Then  $\exists k$  disjoint  $a - b$  paths and  $S$  must intersect each of these, so  $|S| \geq k$ .

$\Rightarrow$ . Let  $a, b \in V(G), a \neq b$ . If  $a \not\sim b$ , then just apply Menger form 1 and we're done. If  $a \sim b$ , then consider  $G - ab$ . We have  $\kappa_{a,b}(G - ab) \geq k - 1$ , so apply Menger form 1 to get  $k - 1$  disjoint paths and add back  $ab$  as a  $k^{\text{th}}$  path.  $\square$

### 2.2.1 Edge connectivity

Let  $G$  be a graph. Let  $\lambda(G) = \min\{|W| \mid W \subseteq E(G), G - W \text{ is disconnected}\}$ . We say that  $G$  is  **$k$ -edge connected** if  $\lambda(G) \geq k$ .

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**Example 2.2.** • A cycle has  $\kappa(C_n) = 2, \lambda(C_n) = 2$ .

- A "bowtie graph" has  $\kappa(G) = 1, \lambda(G) = 2$ . We can generalize this and take two copies of  $K_n$  which intersect in one vertex, then  $\kappa(G) = 1$  and  $\lambda(G) = n - 1$ .

**Definition 2.10.** We say paths  $P_1, \dots, P_k$  are **edge disjoint** if

$$E(P_i) \cap E(P_j) = \emptyset \quad \forall i \neq j.$$

**Theorem 2.7** (Menger, edge version). Let  $G$  be a connected graph and  $a, b \in V(G), a \neq b$ . Then, every  $W \subseteq E(G)$  that separates  $a$  from  $b$  having size  $\geq k \implies \exists k$  edge disjoint  $a - b$  paths  $P_1, \dots, P_k$ .

**Definition 2.11.** Let  $G$  be a graph. The **line graph** of  $G$ , denoted  $L(G)$ , is defined to be the graph

$$V(L(G)) = E(G);$$

If  $e, f \in E(G)$ , then  $e \sim f$  if they share a vertex.

*Proof of Theorem 2.7.* Given  $G$ , define a new graph  $G'$  by taking the line graph of  $G$  and adding a vertex  $a'$ , which we join to all edges incident to  $a \in G$ , and similarly adding a vertex  $b'$ , which we join to all edges incident to  $b \in G$ .

Note that there is a  $ab$  path in  $G$  if and only if there is an  $a'b'$  path in  $G'$ . Thus  $W \subseteq V(G') \setminus \{a, b\}$  is a  $a'b'$  separator if and only if  $W \subseteq E(G)$  is an  $ab$  separator. Hence  $\kappa_{a,b}(G') \geq k$ .

Now apply Menger (form 1) to find  $k$  disjoint  $ab$  paths  $P_1, \dots, P_k$  in  $G'$ . These describe edge disjoint walks in  $G$  from  $a$  to  $b$ . Thus there are disjoint paths  $\tilde{P}_1 \subseteq P_1, \dots, \tilde{P}_k \subseteq P_k$  and we're done.  $\square$

**Theorem 2.8** (Menger, edge version 2). Let  $G$  be a connected graph. Then  $\lambda(G) \geq k \iff \forall a, b \in V(G)$  with  $a \neq b, \exists k$  edge disjoint  $ab$  paths.

*Proof.*  $\Leftarrow$  is the easy direction. To separate any two vertices, say  $a, b$ , we must remove an edge from each of the  $ab$  paths, so  $\lambda(G) \geq k$ .  $\Rightarrow$  follows from Menger, edge version 1.  $\square$

### 3 Graph coloring

**Definition 3.1.** We say that  $c : V(G) \rightarrow \{1, \dots, k\}$  is a  $k$ -**coloring** (or a proper  $k$ -coloring) if  $c(x) \neq c(y) \forall x \sim y$ .

**Definition 3.2.** The **chromatic number** of  $G$  is

$$\chi(G) = \min\{k \mid \exists \text{ a } k\text{-coloring of } G\}.$$

**Example 3.1.** • A path has chromatic number 2.  $\chi(P_n) = 2$ .

- For a cycle,  $\chi(G) = \begin{cases} 2 & \text{if } n \text{ is even.} \\ 3 & \text{if } n \text{ is odd.} \end{cases}$
- A tree has chromatic number 2 by induction.
- A complete graph has  $\chi(K_n) = n$ .
- A bipartite graph has  $\chi(K_{m,n}) = 2$ . In fact, a graph  $G$  is bipartite if and only if  $\chi(G) = 2$ .

**Proposition 3.1.** Let  $G$  be a graph. Then  $\chi(G) \leq \Delta(G) + 1$ .

*Proof.* Let  $x_1, \dots, x_n$  be an ordering of  $V(G)$ . We color the  $(x_i)$  one at a time in this order. When we come to vertex  $x_i$ , at most  $\Delta$  colors have been used in  $N(x_i)$ , so there is a free color for  $x_i$ .  $\square$

**Remark.** This proposition is sharp (e.g. on  $K_n$ ).

**Remark.** This is sometimes called a greedy coloring. But a greedy coloring may produce a coloring that is not optimal!

### 3.1 Coloring planar graphs

**Observation.** Let  $G$  be a planar graph. Then  $\delta(G) \leq 5$ .

*Proof.* The average degree of  $G$  is

$$\frac{1}{n} \sum_{v \in V(G)} \deg(v) = \frac{2e(G)}{n} \leq \frac{2(3n-6)}{n} = 6 - \frac{12}{n} < 6.$$

But all degrees of vertices are integers, so the minimal degree is  $\leq 5$ .  $\square$

**Proposition 3.2.** If  $G$  is planar, then  $\chi(G) \leq 6$ .

*Proof.* We induct. Base step: if  $|G| \leq 6$ , then we're clearly done.

Induction step: Given a graph  $G$ , let  $x$  be a vertex with  $\deg(x) \leq 5$ . Apply induction to  $G - x$ , which gives a coloring of  $G - x$  with 6 colors. But  $x$  has degree  $\leq 5$ , thus there is a free color to color  $x$  with.  $\square$

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Lecture 8

**Theorem 3.3.** If  $G$  is planar, then  $\chi(G) \leq 5$ .

*Proof.* We apply induction on  $|G|$ . If  $|G| \leq 5$ , we're done.

Now let  $G$  be a planar graph and  $x \in V(G)$  with  $\deg(x) \leq 5$ . Apply induction to  $G - x$ . Let neighbor  $x_i$  of  $x$  have color  $i$ , ordered clockwise around  $x$ .

Question: Can we get from  $x_1$  to  $x_3$  only walking along vertices colored 1 and 3?

If no, let  $C$  be the component of  $G$  of vertices colored 1 or 3 that contains  $x_1$ , so  $x_3 \notin C$ . "Swap" the colours 1 and 3 on  $C$ . This is a proper coloring of  $G - x$ , so we can color  $x$  with color 1 and we're done.

If yes, ask the same question for  $x_2$  and  $x_4$  - if the answer is no, swap colors on the 2-4 component containing  $x_2$ , and we can color  $x$  with color 2.

If the answer is again yes, then we get a contradiction to planarity since the 1-3 path and the 2-4 path have to intersect somewhere and we're done.  $\square$

**Theorem 3.4** (Four color theorem, non-examinable). If  $G$  is planar, then  $\chi(G) \leq 4$ .

To see this is equivalent to the map version, take the dual of our graph (place a vertex inside each face (and one for the infinite face) and connect two vertices by an edge if the two faces have any common boundary).

Kempe "proved" the four colour theorem in 1879, but his proof had a mistake. It was then proved in 1976 by reducing the problem to about 2000 configurations and checking them by computer.

Also, this is the best we can do, since  $K_4$  is planar - there is no "three color theorem".

### 3.2 Coloring graphs

**Proposition 3.5.** Let  $G$  be connected and  $\delta(G) < \Delta(G)$ . Then  $\chi(G) \leq \Delta(G)$ .

*Proof.* Let  $x_n$  have  $\deg(x_n) < \Delta(G)$ . Then choose  $x_{n-1}$  to be adjacent to  $x_n$ ,  $x_{n-2}$  to be adjacent to one of  $\{x_n, x_{n-1}\}$ , etc. Since  $G$  is connected, we eventually order everything and the ordering has the property that all vertices have less than  $\Delta(G)$  edges going to vertices with smaller index, so color greedily and we're done.  $\square$

**Theorem 3.6** (Brooks). Let  $G$  be a connected graph which is not an odd cycle or a complete graph. Then  $\chi(G) \leq \Delta(G)$ .

*Proof.* Apply induction on  $|G|$ . The claim is clearly true for  $|G| \leq 3$ .

Note that we may assume that  $\Delta \geq 3$  (where  $\Delta = \Delta(G)$ ).

**Claim 1.** If  $G$  is 3-connected, then we're done.

*Proof.* Define an ordering of  $G$  as in the proposition above. Let  $x_n$  have  $\deg(x_n) = \Delta$  and choose  $x_1, x_2 \in N(x_n)$  with  $x_1 \neq x_2, x_1 \not\sim x_2$ . This is possible, since  $G$  is not  $K_{\Delta+1}$ .

Now consider  $G \setminus \{x_1, x_2\}$ . We order this in the same way as above: connect  $x_{n-1}$  to  $x_n$ ,  $x_{n-2}$  to  $x_{n-1}$  or  $x_n$ , etc. Since  $G \setminus \{x_1, x_2\}$  is connected, we eventually order every vertex. Hence we're done by coloring greedily.

**Claim 2.** If  $\kappa(G) = 1$ , we're done.

*Proof.* Let  $x$  be a cut vertex and let  $C_1, \dots, C_k$  be the components of  $G - x$ . By induction, we can color each  $G[C_i \cup x]$ . Then permute colors to have  $x$  always be the same color, and put everything back together to get a coloring of  $G$ .

**Claim 3.** If  $\kappa(G) = 2$ , we're done.

*Proof.* Let  $S \subset V(G)$ ,  $S = \{x, y\}, x \neq y$  be our separator. Let  $C_1, \dots, C_k$  be the components of  $G - S$  and let  $G_i = G[C_i \cup S] + xy$  for each  $i$ .

- Case 1:  $\delta(G_i) < \Delta(G) \forall i$ . In this case we can always color  $C_1, \dots, C_k$  by Proposition 3.5. Note that in each of these colorings,  $x, y$  have different colors, so we can permute colors in a way that  $x, y$  always have the same color, and put everything back together to get a coloring of  $G$ .
- Case 2:  $\delta(G_1) = \Delta(G)$  (wlog assume  $i = 1$ ). In this case,  $k = 2$  and  $|N(x) \cap C_1| = \Delta - 1$ . Also,  $|N(x) \cap C_2| = 1$  and  $|N(y) \cap C_2| = 1$ . Thus let  $x', y' \in C_2$  be such that  $x' \sim x, y' \sim y$ . Now observe that  $\tilde{S} = \{x, y'\}$  is not of this bad form, so we're done by Case 1.

$\square$

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Lecture 9

### 3.3 Chromatic polynomials

**Definition 3.3.** For a graph  $G$ , define the **chromatic polynomial** of  $G$  to be the function  $P_G : \mathbb{Z}_{\geq 0} \rightarrow \mathbb{Z}_{\geq 0}$  by  $P_G(t)$  = the number of  $t$ -colorings of  $G$ .

In particular, the minimal value of  $t$  for which  $P_G(t) > 0$  is  $t = \chi(G)$ .

**Example 3.2.** •  $P_{\overline{K_n}}(t) = t^n$ .

- $P_{K_n}(t) = t(t-1) \dots (t-(n-1))$ .
- $P_{P_{n-1}}(t) = t(t-1)^{n-1}$ .
- $P_T(t) = t(t-1)^{n-1}$  where  $T$  is a tree with  $n$  vertices (proof: take a leaf, induct).

**Definition 3.4.** Let  $G$  be a graph and  $e \in E(G)$ . Define  $G/e$  to be the **contraction** of  $G$  along  $e$ , which is the graph

$$\begin{aligned} V(G/e) &= V(G) \setminus \{x, y\} \cup \{xy\}, \\ E(G/e) &= E(G[V \setminus \{x, y\}]) \cup \{ez \mid x \sim z\} \cup \{ez \mid y \sim z\}. \end{aligned}$$

(In other words, you just "squish" the edge  $e$  and its endpoints become a single vertex.)

**Proposition 3.7.** Let  $G$  be a graph. We have  $P_G(t) = P_{G-e}(t) - P_{G/e}(t)$ .

*Proof.* Let  $e = xy$ . A  $t$ -coloring of  $G - e$  where  $x, y$  get different colors corresponds exactly to a  $t$ -coloring of  $G$ . On the other hand, a  $t$ -coloring of  $G - e$  where  $x, y$  get the same color corresponds exactly to a  $t$ -coloring of  $G/e$ .  $\square$

**Proposition 3.8.** For a graph  $G$ ,  $P_G$  is a polynomial.

*Proof.* Apply induction on  $e(G)$ . For the base case  $e(G) = 0$ ,  $P_{\overline{K_n}}(t) = t^n$ . Induction step: For  $e \in E(G)$ ,  $P_G(t) = P_{G-e}(t) - P_{G/e}(t)$  is a polynomial (as it is the difference of two polynomials), so we're done by induction.  $\square$

**Proposition 3.9.** Let  $G$  be a graph with  $n$  vertices and  $m$  edges. Then

$$P_G(t) = t^n - mt^{n-1} + \dots$$

*Proof.* Induction on the number of edges. If  $e(G) = 0$ , then we're done.

Induction step: Let  $e \in E(G)$  and apply the deletion-contraction relation,  $P_G(t) = P_{G-e}(t) - P_{G/e}(t)$ . By induction, the leading terms on the RHS are

$$(t^n - (m-1)t^{n-1} + \dots) - (t^{n-1} - e(G/e)t^{n-2} + \dots) = t^n - mt^{n-1} + \dots$$

$\square$



**Remarks.** We're going to stop here, but chromatic polynomials have a lot of other properties. For example:

- Other coefficients of  $P_G$  contain other information about  $G$ , e.g.

$$P_G = t^n - mt^{n-1} + \left( \binom{m}{2} - \text{the number of triangles in } G \right) t^{n-2} + \dots$$

- If  $G$  is planar, then  $P_G(2 + \frac{1+\sqrt{5}}{2}) \neq 0$ .
- The coefficients of  $c_0, \dots, c_n$  of  $P_G(t)$  are log-concave, i.e.  $c_i^2 \geq c_{i-1}c_{i+1}$ .  
June Huh won a Fields medal for this in 2021.

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Lecture 10

### 3.4 Edge coloring

**Definition 3.5.** For a graph  $G$ , a  $k$ -edge coloring is a function  $c : E(G) \rightarrow [k]$  such that  $c(e) \neq c(f)$  whenever  $e, f \in E(G)$  share an endpoint.

Define the **edge-chromatic number** (sometimes called the **chromatic index**) to be

$$\chi'(G) = \min\{k \mid \exists k\text{-edge coloring of } G\}.$$

Note that an edge coloring of  $G$  corresponds exactly to a vertex coloring of the line graph of  $G$ , and so  $\chi'(G) = \chi(L(G))$ .

**Remarks.**

- $\chi'(G) = \begin{cases} 2 & \text{if } n \text{ even.} \\ 3 & \text{if } n \text{ odd.} \end{cases}$
- $\Delta(G) \leq \chi'(G)$ .
- We can have  $\Delta(G) < \chi'(G)$ . For example,  $\chi'(\text{Petersen}) = 4$ .
- $\chi'(G) \leq 2\Delta(G) - 1$  by coloring greedily.
- $\chi$  and  $\chi'$  can be very different. For example,  $\chi'(K_{t,1}) = t, \chi(K_{t,1}) = 2$ .

We now prove Vizing's theorem. First, some setup:

**Definition 3.6.** Given an edge-coloring  $c : E(G) \rightarrow [k]$ , define the **color classes** of  $c$  to be

$$C_i = \{e : c(e) = i\}.$$

**Question.** What does  $(V(G), C_i \cup C_j)$  look like? It is the union of disjoint paths, isolated vertices and even cycles.

We say the components of this graph are  $\{i, j\}$ -components.

**Definition 3.7.** Let  $y \in V(G)$ ,  $c \in [\Delta + 1]$ . We say that  $c$  is **missing** at  $y$  if none of the edges incident at  $y$  have color  $c$ .

**Theorem 3.10** (Vizing's theorem). Let  $G$  be a graph. Then

$$\chi'(G) \in \{\Delta(G), \Delta(G) + 1\}.$$

*Proof.* Induction on  $e(G)$ . If  $e(G) = 0$ , we're done.

Induction step: given  $G$ ,  $e(G) > 0$ , let  $xv \in e(G)$ . Apply induction to the graph  $G - xv$  to get a  $\Delta + 1$  edge coloring. We want to extend this coloring to  $G$ .

Note that there is a color missing at every vertex and let  $c_0$  be the color missing at  $x$ . Define a sequence of  $v_1, \dots, v_k \in N(x)$  and corresponding colors  $c_1, \dots, c_k$  such that  $c_i$  is missing at  $v_i$  as follows:

First set  $v_1 = v$  and let  $c_1$  be any color missing at  $v_1$ . Now say we have defined  $v_i$  and  $c_i$ . Define  $v_{i+1}$  such that  $xv_{i+1}$  is colored  $c_i$ . Continue this way until one of two things happens:

- Case 1: Say we have defined  $v_k, c_k$  and  $c_k$  is missing at  $x$ . Then recolor  $xv_k$  with color  $c_k$ , so now  $c_{k-1}$  is missing at  $x$ . Hence recolor  $xv_{k-1}$  with color  $c_{k-1}$ , so now  $c_{k-2}$  is missing at  $x$ , etc, until we color  $xv_i$  with color  $c_i$  for all  $i \leq k$ . This includes  $i = 1$ , so we're done.
- Case 2: Say  $c_k = c_i$  for some  $i < k$ . Note that we may assume  $i = 1$  by recoloring as in Case 1 (i.e. uncolor  $xv_i$  and color  $xv_j \rightarrow c_j$  for all  $j < i$ ). The question we ask is whether  $v_1$  is in the same  $\{c_0, c_1\}$ -component as  $x$ .
  - If not, swap colors on the  $\{c_0, c_1\}$ -component containing  $v_1$ . Then  $c_0$  is missing at both  $x$  and  $v_1$ , and so we can color  $xv_1$  with  $c_0$  and we're done.
  - If yes, we need another question: Is  $x$  in the same  $\{c_0, c_1\}$ -component as  $v_k$ ?
    - \* If the answer is no, then swap colors in the  $\{c_0, c_1\}$ -component containing  $v_k$ . So  $c_0$  is missing at  $v_k$  and  $x$ , so we are done by Case 1.
    - \* If yes, then  $x, v_1, v_k$  are all in the same  $\{c_0, c_1\}$ -component. But one of  $c_0, c_1$  is missing at each of  $x, v_1, v_k$ . Thus  $x, v_1, v_k$  all have to be endpoints of a path in the same  $\{c_0, c_1\}$ -component, which is a contradiction.

□

### 3.5 Coloring graphs on surfaces

31 Oct 2022,

Lecture 11

If  $G$  is planar, we know that  $\chi(G) \leq 5$ . (And really  $\chi(G) \leq 4$ ). If  $G$  is drawn on a surface  $S$ , can we say anything about  $\chi(G)$ ?

$K_7$  can be drawn on a torus without edge crossings, so we definitely don't have a 4 color theorem for the torus:



**Question.** Is there a "7 color theorem" for the torus?

**Definition 3.8.** We define the **surface of genus  $g$**  (more formally, a compact orientable surface of genus  $g$ ) to be the sphere with  $g$  handles attached (or equivalently, the sphere with  $g$  holes).

Do we have an Euler-type theorem in this setting? Different cycles on the torus give different values for  $V - E + F$ , we get all of 2, 1, 0. \*picture\*

**Proposition 3.11.** Let  $G$  be a graph drawn on a surface of genus  $g$  with  $n$  vertices,  $m$  edges, and  $f$  faces. Then

$$n - m + f \geq 2 - 2g.$$

**Remarks.**

- $E = 2 - 2g$  is the **Euler characteristic** of the surface of genus  $g$ .
- This is the direction of Euler's theorem ( $\geq$ , not  $\leq$ ) that we need to bound the number of edges of  $G$ .
- We do not assume that the graph is connected here.

**Corollary 3.12.** Let  $G$  be a graph with  $|G| \geq 3$ , drawn on a surface of Euler characteristic  $E$ . Then

$$e(G) \leq 3(|G| - E).$$

*Proof.* If  $e(G) \leq 2$ , we can check the claim holds. When  $e(G) \geq 3$ , then every face in  $G$  has length at least 3, so

$$3f \leq \sum_{F \text{ a face}} \deg(F) \leq 2e(G).$$

Thus  $f \leq \frac{2}{3}e(G)$ , so  $e(G) \leq n + f - E \implies e(G) \leq 3(n - E)$ .  $\square$

**Theorem 3.13** (Heawood's theorem). Let  $G$  be a graph drawn on a surface of Euler characteristic  $E \leq 0$ . Then

$$\chi(G) \leq \left\lfloor \frac{7 + \sqrt{49 - 24E}}{2} \right\rfloor.$$

*Proof.* Let  $G$  be a graph drawn on a surface of Euler characteristic  $E$  and let  $k = \chi(G)$ . We may assume that  $G$  is minimal, i.e. all proper subgraphs of  $G$  have  $\chi \leq k - 1$ .

We estimate  $\delta(G)$  in two ways. Note  $\delta(G) \geq k - 1$ : if not, let  $x$  be a vertex with  $\deg(x) \leq k - 2$ . We have  $\chi(G - x) \leq k - 1$  by minimality, so we can put  $x$  back and color it  $\implies \chi(G) \leq k - 1$ , contradiction.

Note

$$\begin{aligned} \delta(G) &\leq \text{the average degree of the graph} = \frac{1}{n} \sum_{x \in V(G)} \deg(x) = \\ &= \frac{2e(G)}{n} \leq \frac{6(n - E)}{n} = 6 - \frac{6E}{n} \leq 6 - \frac{6E}{k} \end{aligned}$$

because  $n \geq k$  and  $E \leq 0$ . So

$$k - 1 \leq 6 - \frac{6E}{k} \implies k^2 - 7k + 6E \leq 0 \implies k \leq \frac{7 + \sqrt{49 - 24E}}{2}.$$

As  $k$  is an integer, we're done.  $\square$

**Remarks.**

- Define  $H(E) = \left\lfloor \frac{7 + \sqrt{49 - 24E}}{2} \right\rfloor$ .
- Did we just prove the four color theorem, since the above gives  $H(2) = 4$ ? No, since we assumed  $E \leq 0$ .
- Amazingly, Heawood's theorem is sharp. In particular,  $K_{H(E)}$  can be drawn in the surface of Euler characteristic  $E$  (but this is very hard!).

## 4 Extremal graph theory

**Question.** What is the minimum number of edges a graph must have before it is forced to have a triangle? In other words, what is the value of  $k$  such that if  $e(G) > k$ , then  $K_3 \subset G$ ?

**Question.** How large does  $\delta(G)$  have to be before  $G$  is forced to have a cycle of length  $|G|$ ?

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**Definition 4.1.** A graph is said to be **Hamiltonian** if it contains a cycle that contains all vertices. Such a cycle is called a **Hamilton** cycle.

**Theorem 4.1** (Dirac). Let  $G$  be a graph on  $n$  vertices,  $n \geq 3$ . If  $\delta(G) \geq \frac{n}{2}$ , then  $G$  is Hamiltonian.

**Remark.** This theorem is sharp. If  $n$  is even, take two disjoint  $K_{n/2}$ 's. If  $n$  is odd, take two copies of  $K_{(n+1)/2}$  that overlap at exactly one vertex.

*Proof.* First note that  $G$  is connected. If  $x \not\sim y$ , then  $N(x), N(y) \subset G - x - y$ , but  $|N(x)|, |N(y)| \geq \frac{n}{2}$ , so  $N(x) \cap N(y) \neq \emptyset$  by the pigeonhole principle, so  $x$  and  $y$  have a common neighbor.

Now take a path of maximal length  $x_1 \dots x_l$ . The claim is that there is no cycle in  $G$  of length  $l$ . Indeed, if such a cycle did exist, then either  $l = n$  and we're done, or  $l < n$  and we can use this cycle to find a path of length  $l + 1$  in  $G$ , contradicting maximality.

Observe that  $N(x_1) \subset \{x_2, \dots, x_{l-1}\}$  and  $N(x_l) \subset \{x_2, \dots, x_{l-1}\}$ . Now define  $N^-(x_1) = \{x_i \mid x_{i+1} \in N(x_1)\}$ . Note

$$|N^-(x_1) \cup N^-(x_l)| \leq \{x_1, x_2, \dots, x_{l-1}\} \leq l - 1 \leq n - 1,$$

but  $|N^-(x_1)|, |N^-(x_l)| \geq \frac{n}{2}$ . So there must exist  $x_i$  in their intersection.

So we have  $x_1 \dots x_i x_l x_{l-1} \dots x_{i+1} x_1$ , a cycle of length  $l$ , a contradiction and we're done.  $\square$

**Remark.** Note that there's no interesting theorem of the form "if  $e(G) > K$ , then  $G$  is Hamiltonian", since  $K_{n-1} + x$  where  $x$  has degree 1 is not Hamiltonian.

**Theorem 4.2.** Let  $G$  be a graph on  $n$  vertices. Then  $e(G) > \frac{n(k-1)}{2} \implies G$  contains a path of length  $k$ .

**Remark.** If  $k \mid n$ , then this is sharp: take the disjoint union of  $\frac{n}{k}$  copies of  $K_k$ .

**Lemma 4.3.** Let  $G$  be a graph on  $n$  vertices,  $n \geq 3$ , and let  $k < n$ . If  $G$  is connected and  $\delta(G) \geq \frac{k}{2}$ , then  $G$  contains a path of length  $k$ .

**Remarks.**

- We need the assumption that  $k < n$ , otherwise take  $K_k$ .
- We need the assumption that  $G$  is connected, otherwise take a bunch of disjoint  $K_k$ .
- The  $\frac{k}{2}$  is sharp. Take a bunch of  $K_k$ 's that overlap at a single vertex.

*Proof.* Let  $x_1 \dots x_l$  be a path of maximum length. We have that there is no cycle of length  $l$ : If  $l = n$ , then we're done since  $k < n$ , and if  $l < n$ , then use a cycle of length  $l$  to build a path of length  $l + 1$  (since  $G$  is connected), a contradiction.

So we get  $N(x_1) \subset \{x_2, \dots, x_{l-1}\}$ ,  $N(x_l) \subset \{x_2, \dots, x_{l-1}\}$ , and  $N^-(x_1) = \{x_i \mid x_{i+1} \in N(x_1)\} \subset \{x_1, \dots, x_{l-2}\}$ . Thus

$$|N^-(x_1) \cup N(x_l)| \leq |\{x_1, \dots, x_{l-1}\}| \leq l - 1 \stackrel{(\star)}{\leq} k - 1$$

(where  $(\star)$  follows since else we're immediately done), but  $|N^-(x_1)|, |N(x_l)| \geq \frac{k}{2}$ , so  $N^-(x_1) \cap N(x_l) \neq \emptyset$ . So as before, we can build a cycle of length  $l$ , contradiction, done.  $\square$

*Proof of Theorem 4.2.* If  $k = 1$ , we're done. Now assume  $k \geq 2$  and apply induction on  $n$ ,  $n = 2$  being clear.

For the induction step, given a graph  $G$  on  $n \geq 3$  vertices, first note that  $\frac{n(k-1)}{2} < e(G) \leq \frac{n(n-1)}{2} \implies k \leq n$ .

We may assume that  $G$  is connected: say  $C_1, \dots, C_r$  are the components of  $G$  and  $|C_i| = n_i$ . Then  $\sum_{i=1}^r e(G[C_i]) = e(G) > \frac{n(k-1)}{2}$ , so since  $\sum_{i=1}^r n_i = n$ ,

$$\sum_{i=1}^r \left( e(G[C_i]) - \frac{n_i(k-1)}{2} \right) > 0.$$

So there exists some component  $C_i$  that has  $e(G[C_i]) > \frac{n_i(k-1)}{2}$ , so apply induction to this component and we're done.

If  $\delta(G) \geq \frac{k}{2}$ , then apply Lemma 4.3 to see that  $G$  contains a path of length  $k$ . Otherwise, there exists a vertex  $x \in V(G)$  with  $\deg(x) \leq \frac{k-1}{2}$ . Note that we have

$$e(G - x) > \frac{n(k-1)}{2} - \frac{k-1}{2} = \frac{(n-1)(k-1)}{2},$$

hence  $G$  contains a path of length  $k$  and we're done.  $\square$

**Question.** How many edges in a graph with  $n$  vertices do we need before we are forced to have a triangle?

**Note.** For  $n$  even,  $K_{n/2, n/2}$  contains no triangle and has  $\frac{n^2}{4}$  edges.

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**Theorem 4.4.** Let  $G$  be a graph on  $n$  vertices. If  $e(G) > \frac{n^2}{4}$ , then  $G$  contains  $K_3$ .

**Fact** (Jensen). Let  $a < b$ ,  $f : [a, b] \rightarrow \mathbb{R}$  convex,  $x_1, \dots, x_n \in [a, b]$ . Then

$$\frac{1}{n} \sum_{i=1}^n f(x_i) \geq f\left(\frac{1}{n} \sum_{i=1}^n x_i\right).$$

*Proof.* Assume for contradiction that  $K_3 \not\subset G$ . We may assume that  $n \geq 3$ . Consider  $x \sim y$ . We have  $d(x) + d(y) \leq n$ , else we have a triangle (since  $x, y$  have a common neighbor by pigeonhole). Thus

$$\begin{aligned} ne(G) &= \sum_{x \sim y} (d(x) + d(y)) = \frac{1}{2} \sum_x \sum_y (d(x) + d(y)) \mathbb{1}(x \sim y) = \\ &= \sum_x \sum_y d(x) \mathbb{1}(x \sim y) = \sum_x d(x) \sum_y \mathbb{1}(x \sim y) = \sum_x d(x)^2. \end{aligned}$$

Now by Jensen,

$$\begin{aligned} ne(G) &= n \left( \frac{1}{n} \sum_x d(x)^2 \right) \geq n \left( \frac{1}{n} \sum_x d(x) \right)^2 = n \left( \frac{2e(G)}{n} \right)^2 \\ &\implies e(G) \leq \frac{n^2}{4}. \end{aligned}$$

□

**Definition 4.2.** We say that a graph  $G = (V, E)$  is  $r$ -**partite** if there is a partition of  $V = V_1 \sqcup V_2 \sqcup \dots \sqcup V_r$  such that  $\forall xy \in E, \exists V_i, V_j, i \neq j$  s.t.  $x \in V_i$  and  $y \in V_j$ . This is the same as saying  $\chi(G) \leq r$ .

**Definition 4.3.** Given  $n_1, \dots, n_r \in \mathbb{Z}_{\geq 0}$ , define  $K_{n_1, \dots, n_r}$  to be the **complete  $r$ -partite graph** where  $V = V_1 \sqcup V_2 \sqcup \dots \sqcup V_r$ ,  $|V_i| = n_i \ \forall i$ , and  $\forall x \in V_i, y \in V_j, i \neq j, xy \in E(K_{n_1, \dots, n_r})$ .

Note that if  $r \mid n$ , then

$$e(K_{\frac{n}{r}, \dots, \frac{n}{r}}) = \frac{n^2}{r^2} \binom{r}{2} = \left(1 - \frac{1}{r}\right) \frac{n^2}{2}.$$

**Theorem 4.5** (Turan). Let  $G$  be a graph on  $n$  vertices,  $r \geq 1$ . Then

$$e(G) > \left(1 - \frac{1}{r}\right) \frac{n^2}{2} \implies K_{r+1} \subset G.$$

*Proof.* We prove the contrapositive. Apply induction first on  $r \geq 1$  and then on  $n$ . If  $n \leq r$ , we're done. Note that the theorem is trivial if  $r = 1$ , so assume  $r \geq 2$ . Now let  $G$  be a graph,  $K_{r+1} \not\subset G$ .

If there is no  $K_r$  in  $G$ , then apply induction (on  $r$ ) to get an even stronger bound  $e(G) \leq (1 - \frac{1}{r-1}) \frac{n^2}{2}$ .

Hence let  $K$  be a  $K_r$  in  $G$ . Each vertex in  $V(G) \setminus K$  has at most  $r - 1$  neighbors in  $K$  (else we'd have  $K_{r+1} \subset G$ ), So

$$\begin{aligned} e(G) &\leq \binom{k}{2} + (r-1)(n-r) + e(G \setminus K) \leq \\ &\leq \binom{k}{2} + (r-1)(n-r) + \left(1 - \frac{1}{r}\right) \frac{(n-1)^2}{2} = \left(1 - \frac{1}{r}\right) \frac{n^2}{2}. \end{aligned}$$

□

**Definition 4.4.** The **Turan graph**  $T_{r,n}$  is the graph  $K_{n_1, \dots, n_r}$  where  $\sum n_i = n$  and the  $\{n_i\}$  are as equal as possible, i.e.  $|n_i - n_j| \leq 1 \forall i, j$ .

**Lemma 4.6.** Let  $G$  be an  $r$ -partite graph on  $n$  vertices. Then  $e(G) \leq e(T_{r,n})$ .

*Proof.* Let  $G$  be a  $r$ -partite graph with maximal number of edges. First note  $G$  is a complete  $r$ -partite graph  $K_{n_1, \dots, n_r}$  for some  $n_1, \dots, n_r$ . If  $\exists i, j$  such that  $n_i \geq n_j + 2$ , move a vertex from part  $i$  to part  $j$ . Then we gain  $n_i - 1$  edges and lose  $n_j$  edges, so we get a graph with more edges, contradiction. Hence  $|n_i - n_j| \leq 1 \forall i, j$ , so  $G = T_{r,n}$ . □

**Theorem 4.7** (Turan, version 2). Let  $G$  be a graph on  $n$  vertices and  $r \geq 2$ . If  $G$  does not contain a  $K_{r+1}$ , then  $e(G) \leq e(T_{r,n})$ .

*Idea for proof.* Given  $G$ , we "transform"  $G$  into a complete  $r$ -partite graph without decreasing the number of edges or creating any  $K_{r+1}$ 's. □

*Proof.* Say  $V(G) = [n]$ . Let  $\alpha_1, \dots, \alpha_n > 0$  be linearly independent over  $\mathbb{Q}$  and define  $\mu(S) = \sum_{i \in S} \alpha_i$  for  $S \subseteq [n]$ .

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If  $H$  is a graph on  $n$  vertices, define  $T(H)$  as follows: Let  $x, y$  be a pair maximizing  $\mu(\{x, y\})$  such that  $N(x) \neq N(y)$  and  $x \not\sim y$ . If  $\deg(x) > \deg(y)$  or  $\deg(x) = \deg(y)$  but  $\mu(N(x)) > \mu(N(y))$ , define  $T(H)$  to be  $H - y$  along with a vertex  $x'$  with  $N(x') = N(x)$ .

**Claim.** If  $K_{r+1} \not\subset H$ , then  $K_{r+1} \not\subset T(H)$ .

*Proof.* Assume for contradiction that  $T(H)$  contains  $K_{r+1}$ , call it  $K$ . We must have  $x' \in K$ , because all other vertices stayed the same. But  $x \notin K$  since  $x \not\sim x'$ . But then  $K \setminus \{x'\} \cup x$  is a  $K_{r+1}$  in  $H$ , contradiction. □



Now consider the sequence,  $G, T(G), T(T(G)), \dots$

**Claim.** This sequence eventually stabilizes.

*Proof.* Note  $e(T(H)) \geq e(H)$ , so  $e(T^{(i)}(H))$  is an increasing sequence. It is also bounded above, so eventually we reach a maximum. Now note

$$\sum_{x \in [n]} \mu(N_{T^{(i)}(G)}(x))$$

is always increasing past that point, so eventually it stabilizes.  $\square$

So at some point, our transformation does nothing anymore. Let  $G_\infty$  be the graph where we stop.

**Claim.**  $G_\infty$  is a complete  $k$ -partite graph for some  $k \in \mathbb{N}$ .

*Proof.* Let  $k = \chi(G_\infty)$  and let  $c$  be a  $k$ -coloring of  $G_\infty$  with  $V(G_\infty) = C_1 \cup \dots \cup C_k$ , where  $C_i = \{x \mid c(x) = i\}$ .

Note that if  $x, y \in c(i)$ , then  $N(x) = N(y)$ , otherwise  $T(G_\infty) \neq G_\infty$ .

Now let  $x \in C_i$  and  $y \in C_j$  for  $i \neq j$ . If  $x \not\sim y$ , then  $x' \not\sim y \forall x' \in C_i$ . It follows that  $\forall x' \in C_i$  and  $\forall y' \in C_j$ ,  $x' \not\sim y'$ . We can now recolor  $i, j$  to be the same color to get a  $k-1$ -coloring, contradiction. So  $x \sim y \forall x \in C_i, \forall y \in C_j, \forall i \neq j$ .  $\square$

So  $G_\infty$  is complete  $k$ -partite for some  $k$ . Since  $G_\infty$  does not contain  $K_{r+1}$  by our first claim, we have  $k \leq r$ . By Lemma 4.6,  $e(G_\infty) \leq e(T_{r,n})$  and also  $e(G) \leq e(G_\infty)$ , so we're done.  $\square$

### The Zarankiewicz problem.

Let  $Z(n, t)$  be the maximal number of edges in a bipartite graph  $G = (X \cup Y, E)$ , with  $|X| = |Y| = n$  and  $K_{t,t} \not\subset G$ .

**Theorem 4.8.** Let  $t \geq 2$ . Then

$$Z(n, t) \leq t^{\frac{1}{t}} n^{2 - \frac{1}{t}} + tn.$$

In particular, if  $n$  is large compared to  $t$ , we have  $Z(n, t) \leq 2n^{2 - \frac{1}{t}}$ .

**Lemma 4.9.** Let  $t \geq 2$  be an integer. The function

$$f_t(x) = \frac{x(x-1)\dots(x-t+1)}{t!}$$

is convex for  $x \geq t-1$ .

*Proof.* Let  $s = x - t + 1$ , so  $f_t(s) = \frac{(s+t-1)(s+t-2)\dots(s)}{t!}$  which is a polynomial with non-negative coefficients. So  $f''(s) \geq 0$  if  $s \geq 0 \iff x \geq t - 1$ .  $\square$

*Proof of Theorem 4.8.* We may assume that  $d(y) \geq t - 1 \forall y \in Y$ , since if  $d(y) < t - 1$ , we could add an edge incident to  $y$  without creating a  $K_{t,t}$ . Note

$$|N(x_1) \cap \dots \cap N(x_t)| \leq t - 1.$$

Hence

$$(t-1) \binom{n}{t} \geq \sum_{x_1, \dots, x_t \in X \text{ distinct}} |N(x_1) \cap \dots \cap N(x_t)|.$$

Rewrite the RHS as

$$\begin{aligned} & \sum_{x_1, \dots, x_t \text{ distinct}} \sum_y \mathbb{1}(y \sim x_1) \mathbb{1}(y \sim x_2) \dots \mathbb{1}(y \sim x_t) = \\ & = \sum_y \sum_{x_1, \dots, x_t} \mathbb{1}(y \sim x_1) \dots \mathbb{1}(y \sim x_t) = \sum_y \binom{d(y)}{t}. \end{aligned}$$

Let  $\bar{d} = \frac{e(G)}{n}$ . As  $d(y) \geq t - 1$ , our function is convex by Lemma 4.9, so by Jensen's inequality,

$$\sum_y \binom{d(y)}{t} \geq n \binom{\bar{d}}{t}.$$

So

$$\begin{aligned} (n-1) \binom{n}{t} \geq n \binom{\bar{d}}{t} & \implies \frac{tn^t}{t!} \geq n \frac{(\bar{d}-t)^t}{t!} \\ & \implies t^{\frac{1}{t}} n^{1-\frac{1}{t}} \geq \bar{d} - t \\ & \implies t^{\frac{1}{t}} n^{1-\frac{1}{t}} \geq \frac{e(G)}{n} - t \\ & \implies e(G) \leq t^{\frac{1}{t}} n^{2-\frac{1}{t}} + tn. \end{aligned}$$

$\square$

**Remarks.**

- If  $t = 2$ , then it is known that  $Z(n, t) \geq cn^{\frac{3}{2}}$  for some constant  $c > 0$ .
- If  $t = 3$ , it is known that  $Z(n, t) \geq cn^{\frac{5}{3}}$  for some constant  $c > 0$ .
- $t = 4$  is an open problem.

Non-examinable aside:

**Theorem 4.10** (Non-examinable aside).  $\exists c > 0$  such that for infinitely many  $n$ ,  $Z(n, 2) \geq cn^{\frac{3}{2}}$ .

*Proof.* Think of  $X$  as lines in a finite geometry and  $Y$  as points, and we have an edge in our graph if a line contains a given point.

For  $p$  a prime, let  $\mathcal{L} = \mathbb{F}_p^2 \setminus \{(0, 0)\}$  and  $\mathcal{P} = \mathbb{F}_p^2 \setminus \{(0, 0)\}$ . For  $(a, b) \in \mathcal{L}$ , define  $N((a, b)) = \{(x, ax + b) \mid x \in \mathbb{F}_p\} \setminus \{(0, 0)\}$ .

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Note that if we have two lines  $(a, b), (a', b') \in \mathcal{L}$ , then  $|N(a, b) \cap N(a', b')| \leq 1$ , since if  $(x, y)$  is in the intersection, then  $y = ax + b, y = a'x + b'$ . If  $a' = a$ , then  $b = b'$  and we have the same line. Otherwise  $x = \frac{-(b-b')}{a-a'} \bmod p$ , so  $x$  (and hence  $y$ ) is uniquely determined.

Note  $\forall (a, b) \in \mathcal{L}$ ,  $|N(a, b)| \geq p - 1$ , so  $V(G) = 2(p^2 - 1) = 2n$  and  $e(G) \geq (p - 1)(p^2 - 1) \geq cp^3 = cn^{\frac{3}{2}}$  for an absolute constant  $c$ .  $\square$

End of aside.

## The Erdős–Stone theorem.

**Definition 4.5.** Let  $H$  be a (fixed) graph and  $n \in \mathbb{N}$ . Define

$$\text{ex}(n, H) = \max\{e(G) \mid V(G) = n \text{ and } H \not\subseteq G\}.$$

**Example 4.1.** • We have  $\text{ex}(n, K_{r+1}) = e(T_{r,n}) \leq (1 - \frac{1}{r}) \frac{n^2}{2}$ .

• We have  $\text{ex}(P_k, n) = \frac{n(k-1)}{2}$ .

• We've basically shown  $\text{ex}(K_{t,t}, n) \leq 2n^{2-\frac{1}{t}} + tn$ .

**Theorem 4.11** (Erdős–Stone). Let  $H$  be a fixed graph,  $H$  nonempty. Then

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, H)}{\binom{n}{2}} = \left(1 - \frac{1}{\chi(H) - 1}\right).$$

### Remarks.

- If  $\chi(H) \geq 3$ , then this determines the leading order term in  $\text{ex}(n, H)$  for large  $n$ .
- If  $\chi(H) = 2$ , then this just says  $\frac{\text{ex}(n, H)}{n^2} \rightarrow 0$ . But we have a stronger result already:  $H \subseteq K_{t,t}$  for  $t = |H|$ , so we (almost) know  $\text{ex}(H, n) \leq cn^{2-\frac{1}{t}}$ .
- It is easy to see that  $\text{ex}(n, H) \geq (1 - \frac{1}{\chi(H)-1}) \frac{n^2}{2}$ , since  $H \not\subseteq T_{(\chi(H)-1), n}$ .

We sketch the proof of this. It is non-examinable.

*Sketch of proof, non-examinable.* We only consider  $\chi(H) = 3$ . Let  $K_3(t)$  be the complete 3-partite graph with  $t$  vertices in each part (i.e.  $T_{3,3t}$ ). Note that it is enough to prove our theorem for  $H = K_3(t) \forall t$ , since  $H \subseteq K_3(|H|)$ .

We prove the following:

**Theorem 4.12.** Let  $\epsilon > 0$ ,  $t \in \mathbb{N}$ ,  $G$  a graph with  $|G| = n$  and

$$e(G) \geq \left(\frac{1}{2} + \epsilon\right) \frac{n^2}{2}$$

where  $n \geq n_0(\epsilon, t)$ . Then  $G \supset K_3(t)$ .

**Lemma 4.13.** Let  $G$  be a graph on  $n$  vertices with average degree  $\geq (\frac{1}{2} + \epsilon)n$  (in other words,  $e(G) \geq (\frac{1}{2} + \epsilon) \frac{n^2}{2}$ ). Then  $\exists \tilde{G} \subset G$  with minimum degree at least  $(\frac{1}{2} + \frac{\epsilon}{2})m$ , where  $m = |\tilde{G}| \geq \frac{\epsilon \frac{1}{2} n}{8}$ .

*Proof.* Omitted, but the idea is to remove edges with degrees smaller than the number stated above and eventually we're left with a subgraph that works.  $\square$

*Proof of theorem.* Let  $\tilde{G} \subset G$  be a subgraph with  $\delta(\tilde{G}) \geq (\frac{1+\epsilon}{2})m$ ,  $|\tilde{G}| = m$ . Find a  $K_{s,s}$  in  $\tilde{G}$  with  $s \approx \frac{10t}{\epsilon}$ , call this graph  $K$ . Consider

$$e(Y, K) = \sum_{x \in K} |N(x) \cap Y| = \sum_{x \in K} (\deg(x) - s) \geq \sum_{x \in K} \left(\frac{1}{2} + \frac{\epsilon}{4}\right) m = \left(\frac{1}{2} + \frac{\epsilon}{4}\right) m |K|.$$

But also  $e(Y, K) = \sum_{y \in Y} |N(y) \cap K|$ , so

$$\left(\frac{1}{2} + \frac{\epsilon}{4}\right) |K| < \frac{1}{m} \sum_{y \in Y} |N(y) \cap K| = (\star).$$

Let  $\mathcal{B} = \{y \in Y \mid |N(y) \cap K| \geq (\frac{1}{2} + \frac{\epsilon}{8})|K|\}$ . We claim  $|\mathcal{B}| \geq \frac{\epsilon m}{8}$ . Indeed, suppose not. Then

$$(\star) \leq \frac{1}{m} \frac{\epsilon m}{8} |K| + \left(\frac{1}{2} + \frac{\epsilon}{8}\right) |K| = \left(\frac{1}{2} + \frac{\epsilon}{4}\right) |K|,$$

a contradiction.  $\square$

and now finish, but we're out of time.  $\square$

## 5 Ramsey Theory

11 Nov 2022,  
Lecture 16

**Fact.** Let  $c$  be a 2-edge coloring of  $K_6$ . Then there exists a monochromatic  $K_3$  (i.e. a triangle  $T$ ) for which  $c$  is constant on  $E(T)$ .

This is just saying "if we color the edges of a  $K_6$ , we get a monochromatic triangle".

*Proof.* Let  $x \in V(K_6)$ .  $x$  has 5 neighbors, so WLOG it has 3 red neighbors  $y_1, y_2, y_3$ . If there is no red triangle, then  $y_1y_2, y_2y_3, y_3y_1$  are all blue, but then  $y_1y_2y_3$  is a blue triangle.  $\square$

**Definition 5.1.** For  $s \geq 2$ , define

$$R(s) = \min\{n \mid \text{every 2-edge coloring of } K_n \text{ contains a monochromatic } K_s\}.$$

**Definition 5.2.** Let  $s, t \geq 2$ . Define

$$R(s, t) = \min\{n \mid \text{every 2-coloring of } K_n \text{ contains either a red } K_s \text{ or a blue } K_t\}.$$

Note that  $R(s, t) = R(t, s)$  and  $R(2, t) = t, R(s, 2) = s$ . We just showed that  $R(3, 3) = R(3)$  (in fact,  $R(3) = 6$ ).

**Theorem 5.1** (Ramsey). For all  $s$ ,  $R(s)$  exists.

This is implied by:

**Theorem 5.2** (Ramsey, version 2). For all  $s, t$ ,  $R(s, t)$  exists, and  $\forall s, t \geq 2$ ,

$$R(s, t) \leq R(s-1, t) + R(s, t-1).$$

*Proof.* We induct on  $s+t$ . If  $s, t \leq 2$ , we're done.

Given  $s, t$ , fix  $n = \underbrace{R(s-1, t)}_{=a} + \underbrace{R(s, t-1)}_{=b}$ . Let  $c : E(K_n) \rightarrow \{\text{Red, Blue}\}$  be given. Pick  $x \in V(K_n)$  and let  $N_{\text{red}}(x), N_{\text{blue}}(x)$  be the red and blue neighborhoods of  $x$ .

If  $|N_{\text{red}}(x)| \geq a$ , then  $N_{\text{red}}(x)$  contains either a red  $K_{s-1}$ , so this  $K_{s-1}$  along with  $x$  is a red  $K_s$  in  $K_n$ , or a blue  $K_t$ .

Similarly, if  $|N_{\text{blue}}(x)| \geq b$ , then we get either a blue  $K_{t-1}$ , which extends to a blue  $K_t$ , or a red  $K_s$ .

If  $|N_{\text{red}}(x)| < a$  and  $|N_{\text{blue}}(x)| < b$ , then  $n \leq (a-1) + (b-1) + 1 < a+b$ , contradiction.  $\square$

**Definition 5.3.** For  $k \in \mathbb{N}, s_1, \dots, s_k \geq 2$ , define

$$R_k(s_1, \dots, s_k) = \min\{n \mid \text{every } k\text{-coloring of } K_n \text{ contains a monochromatic } K_{s_i} \text{ in color } i \text{ for some } i\}$$

**Theorem 5.3** (Multicolor Ramsey theorem). For  $s_1, \dots, s_k$  and  $k \geq 2$ , the number  $R_k(s_1, \dots, s_k)$  exists.

*Proof.* Induct on  $k$ . We will show that

$$R_k(s_1, \dots, s_k) \leq R_2(s_1, R_{k-1}(s_2, \dots, s_k)) = (\star).$$

Say  $(\star) = n$ . Let  $c$  be a  $k$ -coloring of  $E(K_n)$ . Now apply the 2-color Ramsey theorem to get either a subgraph  $K_{s_i}$  which is monochromatic in  $s_i$ , or we get a  $K_{R_{k-1}(s_2, \dots, s_k)}$  on which color 1 does not appear.

In the second case, induct on this complete graph to get a  $K_{s_i}$  which is monochromatic in color  $i$ .  $\square$

**Theorem 5.4.** Let  $c : \mathbb{N}^{(2)} \rightarrow \{\text{Red}, \text{Blue}\}$ . Then there exists an infinite monochromatic subset of  $\mathbb{N}$ .

**Remark.** This does not follow from the previous theorems, which only give proofs for larger and larger finite cliques.

*Proof.* Given  $c : \mathbb{N}^{(2)} \rightarrow \{\text{Red}, \text{Blue}\}$ , we construct a sequence  $x_1, x_2, x_3, \dots$  inductively as follows.

Let  $x_1 \in \mathbb{N}$ .  $x_1$  has either an infinite blue neighborhood or an infinite red neighborhood. Define  $S_1$  to be either  $N_{\text{red}}(x_1)$  or  $N_{\text{blue}}(x_1)$ , depending on which one is infinite. Choose  $x_2 \in S_1$  and let  $S_2 = S_1 \cap N_{\text{red}}(x_2)$  or  $S_2 = S_1 \cap N_{\text{blue}}(x_2)$  depending on which one is infinite. Continue this indefinitely.

Note that the sets  $F_i = \{x_i x_j \mid j > i\}$  are all monochromatic (but possibly some are blue and some are red). Call  $x_i$  red if  $F_i$  is monochromatic in red and call  $x_i$  blue otherwise.

As every  $x_i$  is either red or blue, we must have infinitely many red or infinitely many blue vertices, WLOG assume we have infinitely many red vertices. Then we're done, as  $X = \{x_i \mid x_i \text{ is red}\}$  works.  $\square$