

Part II - Number Theory

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0 Introduction

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Books:

- A. Baker, *A concise introduction to the theory of numbers*, CUP 1984
- N. Koblitz, *A course in number theory & cryptography*, Springer 1994
- H. Davenport, *The higher arithmetic*, CUP 2008

Number theory studies the hidden and mysterious properties of the integers and the rational numbers.

It has always been an experimental science. Examining numerical data leads to **conjectures**, many of which are very old and still unproven today.

Example 0.1. (i) Let $N \geq 1$ be an integer of the form $8n + 5, 8n + 6$ or $8n + 7$. Does there exist a right-angled triangle of area N , all of whose sides have rational length? We don't know.

(ii) Let $\pi(x)$ be the number of primes less than or equal to x and define $\text{li}(x) = \int_2^x \frac{dt}{\log t}$. Then for all $x \geq 3$, $|\pi(x) - \text{li}(x)| \leq \sqrt{x} \log x$. This is in fact equivalent to the Riemann hypothesis.

(iii) There are infinitely many twin primes. We now know there is an integer $N \leq 246$ such that there are infinitely many pairs of primes the form $p, p + N$.

1 Euclid's algorithm and factoring

Definition 1.1 (Division algorithm). Given $a, b \in \mathbb{Z}$, with $b > 0$, there exist $q, r \in \mathbb{Z}$ such that $a = qb + r$, and $0 \leq r < b$.

Notation. If $r = 0$, then we write $b|a$, else $b \nmid a$.

Proof. Let $S = \{a - nb \mid n \in \mathbb{Z}\}$. This certainly contains integers ≥ 0 , so take the smallest one r . We claim $r < b$. Indeed, if not, then $r - b \geq 0$, contradicting minimality. \square

Given $a_1, \dots, a_n \in \mathbb{Z}$ not all zero, let $I = \{\lambda_1 a_1 + \dots + \lambda_n a_n \mid \lambda_i \in \mathbb{Z}\}$.

Lemma 1.2. $I = d\mathbb{Z}$ for some $d > 0$.

Proof. I certainly contains integers ≥ 0 . Let d be the least positive element of I . We claim it works. Take $a \in I$, then $a = qd + r$ with $0 \leq r < d$. But $r = a - qd \in I \implies r = 0$. \square

Remark. We get from this that d divides each a_i , and any common divisor of the a_i must divide d . Why?

We write $d = \gcd(a_1, \dots, a_n)$ for the **greatest common divisor** (or **highest common factor**), or just use the shorthand $d = (a_1, \dots, a_n)$.

Corollary 1.3. Let $a, b, c \in \mathbb{Z}$. Then there exist $x, y \in \mathbb{Z}$ such that $ax + by = c$ if and only if $(a, b) | c$.

The division algorithm gives a very efficient way to compute (a, b) . Assume $a > b > 0$. Apply the division algorithm recursively to get

$$\begin{array}{ll} a = q_1 b + r_1 & 0 \leq r_1 < b \\ b = q_2 r_1 + r_2 & 0 \leq r_2 < r_1 \\ r_1 = q_3 r_2 + r_3 & 0 \leq r_3 < r_2 \\ \vdots & \\ r_{k-2} = q_k r_{k-1} + r_k & 0 \leq r_k < r_{k-1}, r_k \neq 0 \\ r_{k-1} = q_{k+1} r_k + 0 & \end{array}$$

Claim. $r_k = (a, b)$. Indeed, $(a, b) = (b, r_1) = (r_1, r_2) = \dots = (r_{k-1}, r_k) = r_k$. This is called **Euclid's algorithm**.

Remark. If $d = (a, b)$, then by Lemma 1.2, there exist $r, s \in \mathbb{Z}$ such that $ra + sb = d$. Euclid's algorithm gives us a way to find r and s .

In the following table, x and y stand for 34 and 25, and we then compute remainders as linear combinations of them.

We can use a trick here to speed this up: find each row as $q \cdot$ the row before it + the second row before it, then figure out signs at the end. (In fact, the minus signs zigzag down).

$$\begin{array}{r|rr}
 & x & y \\
 a = 34 & 1 & 0 \\
 b = 25 & 0 & 1 \\
 34 = 1 \cdot 25 + 9 & 1 & -1 \\
 25 = 2 \cdot 9 + 7 & -2 & 3 \\
 9 = 1 \cdot 7 + 2 & 3 & -4 \\
 7 = 3 \cdot 2 + 1 & -11 & 15
 \end{array}$$

We hence get $-11 \cdot 34 + 15 \cdot 25 = 1$.

Definition 1.4. An integer $n > 1$ is **prime** if its only positive divisors are 1 and n . Otherwise n is **composite**.

Lemma 1.5. Let p be a prime, and $a, b \in \mathbb{Z}$. If $p|ab$, then $p|a$ or $p|b$.

Proof. Assume $p \nmid a$. Then $(a, p) = 1$. By Lemma 1.2, $\exists r, s \in \mathbb{Z}$ such that $ra + sp = 1 \implies rab + spb = b$. Since $p|ab$, $p|b$ follows. \square

Theorem 1.6 (Fundamental Theorem of Arithmetic). Every integer $n > 1$ can be written as a product of primes. This representation is unique up to reordering.

Proof. Existence is obvious. For uniqueness, suppose $n = p_1 p_2 \dots p_r = q_1 q_2 \dots q_s$ for p_i, q_i primes. We have $p_1 | q_1 q_2 \dots q_r$, so by Lemma 1.5, $p_1 | q_j$ for some j , so $p_1 = q_j$. Now cancel these out and induct. \square

Remark. If $m = \prod_{i=1}^k p_i^{\alpha_i}$ and $n = \prod_{i=1}^k p_i^{\beta_i}$ for p_i distinct primes and $\alpha_i, \beta_i \geq 0$, then

$$(m, n) = \prod_{i=1}^k p_i^{\min(\alpha_i, \beta_i)}.$$

However, if m and n are large, it is more efficient to compute (m, n) using Euclid's algorithm.

Suppose we have some large positive integer N . An obvious algorithm for factoring N is to trial divide by 2 and the odd integers up to \sqrt{N} .

Definition 1.7. An algorithm with input a positive integer N is **polynomial** or a **polynomial time** algorithm if it takes $\leq c(\log N)^b$ **elementary operations** for some constants b and c .

Remark. An elementary operation is just adding/multiplying two numbers in $\{0, 1, \dots, 9\}$.

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Remark. "Polynomial" makes sense here as it takes $\log N$ digits to write N .

Polynomial algorithms are known for:

- Adding and multiplying integers (the usual way);
- Computing gcd's (via Euclid's algorithm);
- Detecting n^{th} powers (compute $\sqrt[n]{}$ numerically and round)
- More remarkably, primality testing (Agrawal, Kayal, Saxena in 2002)

But trial division up to \sqrt{N} is not polynomial.

Fundamental question: Is there a polynomial time algorithm for factoring? This is unknown.

Later in this course we study the distribution of the prime numbers, in particular the function $\pi(x)$, the number of primes $\leq x$.

Theorem 1.8. There are infinitely many prime numbers, i.e. $\lim_{x \rightarrow \infty} \pi(x) \rightarrow \infty$.

Proof. Suppose there are only finitely many, say p_1, \dots, p_k . Consider $N = \prod_{i=1}^k p_i + 1$. Then N must be divisible by some prime other than the p_i , so we're done. \square

All the largest known primes are of the form $2^n - 1$ for n a prime. These are called **Mersenne primes**. 51 of them are known, the largest being $2^{82589933} - 1$.

2 Congruences

Fix a positive integer n (the modulus).

Definition 2.1. We say $a \equiv b \pmod{n}$, or that a is congruent to $b \pmod{n}$ if n divides $a - b$.

This defines an equivalence relation on \mathbb{Z} , and we write $\mathbb{Z}/n\mathbb{Z}$ for the set of equivalence classes. We can denote these by $a + n\mathbb{Z}$, or (more commonly) by $a \pmod{n}$. We can check that addition and multiplication are well-defined.

Remark. $n\mathbb{Z}$ is a subgroup/ideal of \mathbb{Z} and $\mathbb{Z}/n\mathbb{Z}$ is the quotient group/ring.

Lemma 2.2. Let $a \in \mathbb{Z}/n\mathbb{Z}$. Then the following are equivalent:

- (i) $(a, n) = 1$
- (ii) $\exists b \in \mathbb{Z}$ such that $ab \equiv 1 \pmod{n}$
- (iii) a is a generator for $\mathbb{Z}/n\mathbb{Z}$.

Proof. (i) \implies (ii): $(a, n) = 1 \implies \exists r, s \in \mathbb{Z}$ such that $ra + sn = 1$, so $ra \equiv 1 \pmod{n}$.

(ii) \implies (i): $ab \equiv 1 \pmod{n} \implies ab + kn = 1$ for some $k \in \mathbb{Z} \implies (a, b) = 1$.

(ii) \iff (iii): $\exists b \in \mathbb{Z}$ s.t. $ab \equiv 1 \pmod{n} \iff 1$ belongs to the subgroup of $\mathbb{Z}/n\mathbb{Z}$ generated by a . \square

Notation. $(\mathbb{Z}/n\mathbb{Z})^\times$ is the group of **units** in $\mathbb{Z}/n\mathbb{Z}$, i.e. the elements with an inverse under multiplication.

Definition 2.3. $\phi(n) = |(\mathbb{Z}/n\mathbb{Z})^\times|$ is called the **Euler totient function**. We also have $\phi(n) = |\{1 \leq a \leq n \mid (a, n) = 1\}|$.

Remark. $\mathbb{Z}/n\mathbb{Z}$ is a field $\iff \phi(n) = n - 1 \iff n$ is prime.

Theorem 2.4 (Euler-Fermat theorem). If $(a, n) = 1$, then $a^{\phi(n)} \equiv 1 \pmod{n}$.

Proof. Apply Lagrange's theorem to the group $G = (\mathbb{Z}/n\mathbb{Z})^\times$. Then for $a \in G$, its order divides $|G| = \phi(n)$. \square

As a corollary:

Theorem 2.5 (Fermat's little theorem). If $p \nmid a$, then $a^{p-1} \equiv 1 \pmod{p}$.

Lemma 2.6. Let G be a cyclic group of order n . We have

$$|\{g \in G \mid \text{order}(g) = d\}| = \begin{cases} \phi(d) & \text{if } d \mid n \\ 0 & \text{otherwise} \end{cases}$$

In particular, $\sum_{d \mid n} \phi(d) = n$.

Proof. WLOG let $G = (\mathbb{Z}/n\mathbb{Z}, +)$. We have $|\{g \in G \mid \text{order}(g) = n\}| \stackrel{(*)}{=} \phi(n)$ by Lemma 2.2. If $d \mid n$, say $n = dk$, then the elements of order dividing d are the classes $0, k, 2k, \dots, (d-1)k \pmod{n}$. These form a cyclic subgroup of order d . Applying $(*)$ to this cyclic subgroup shows that there are $\phi(d)$ elements of order d . \square

Example 2.7. Consider the simultaneous linear congruences $x \equiv 7 \pmod{10}$ and $x \equiv 3 \pmod{13}$. Suppose we can find $u, v \in \mathbb{Z}$ such that

$$\begin{cases} u \equiv 1 \pmod{10} \\ u \equiv 0 \pmod{13} \end{cases}, \begin{cases} v \equiv 0 \pmod{10} \\ v \equiv 1 \pmod{13} \end{cases}.$$

Then $x = 7u + 3v$ is a solution. But $(10, 13) = 1 \implies \exists r, s \in \mathbb{Z}$ such that $10r + 13s = 1$, and we can just take $u = 13s, v = 10r$. To find r, s , we can use Euclid's algorithm to get $r = 4, s = -3$, so $u = -39, v = 40$, and so $x \equiv 7 \cdot (-39) + 3 \cdot 40 \equiv 107 \pmod{130}$.

Theorem 2.8 (Chinese Remainder Theorem). Let m_1, \dots, m_k be pairwise coprime integers greater than 1. Let $a_1, \dots, a_k \in \mathbb{Z}$. Let $M = m_1 m_2 \dots m_k$. Then $\exists x \in \mathbb{Z}$ satisfying

$$\begin{cases} x \equiv a_1 \pmod{m_1} \\ \vdots \\ x \equiv a_k \pmod{m_k} \end{cases}.$$

Moreover, the solution is unique mod M .

Proof. Uniqueness: Suppose $x \equiv x' \pmod{m_i} \forall i$. Then by considering the prime factorization of $x - x'$ and using the fact that the m_i are pairwise coprime, we get $x \equiv x' \pmod{M}$.

Existence: Put $M_i = \frac{M}{m_i}$, so $(M_i, m_i) = 1 \forall i$. Hence we can find $u_i \in \mathbb{Z}$ such that $u_i M_i \equiv 1 \pmod{m_i} \forall i$. Let $x = \sum_{j=1}^k a_j u_j M_j$. Then $x \equiv a_i u_i M_i \equiv a_i \pmod{m_i}$. \square

We can write this theorem in one line using ring theory.

Definition 2.9. Let $R_i = \mathbb{Z}/m_i\mathbb{Z}$, and define $R_1 \times \dots \times R_k = \{(r_1, \dots, r_k) \mid r_i \in R_i\}$ with addition and multiplication defined componentwise. This is a ring.

Theorem 2.10 (CRT, ring-theoretic version). Let m_1, \dots, m_k be pairwise coprime integers greater than 1 and put $M = m_1 \dots m_k$. Then the map

$$\begin{aligned} \theta : \mathbb{Z}/M\mathbb{Z} &\rightarrow \mathbb{Z}/m_1\mathbb{Z} \times \dots \times \mathbb{Z}/m_k\mathbb{Z} \\ a + M\mathbb{Z} &\mapsto (a + m_1\mathbb{Z}, \dots, a + m_k\mathbb{Z}) \end{aligned}$$

is an isomorphism of rings.

Proof. θ is a well defined ring homomorphism since $m_i \mid M \forall i$. Injectivity of θ follows from uniqueness in CRT, and surjectivity of θ follows from existence in CRT. \square

Corollary 2.11. θ induces an isomorphism of groups under multiplication

$$\begin{aligned} (\mathbb{Z}/M\mathbb{Z})^\times &\cong (\mathbb{Z}/m_1\mathbb{Z})^\times \times \dots \times (\mathbb{Z}/m_k\mathbb{Z})^\times \\ a + M\mathbb{Z} &\mapsto (a + m_1\mathbb{Z}, \dots, a + m_k\mathbb{Z}). \end{aligned}$$

Remark. If $a \in \mathbb{Z}$, then $(a, M) = 1 \iff (a, m_i) = 1 \forall i$.

In particular, by looking at orders of the LHS and the RHS above, we get $\phi(M) = \phi(m_1) \dots \phi(m_k)$, i.e. the Euler phi function is multiplicative.

Definition 2.12. A function $f : \mathbb{Z}^+ \rightarrow \mathbb{C}$ is **multiplicative** if $f(mn) = f(m)f(n)$ whenever $(m, n) = 1$.

Examples:

- $\phi(n) = |(\mathbb{Z}/n\mathbb{Z})^\times|$;
- $\tau(n) = \sum_{d|n} 1$, the number of divisors of n ;
- $\sigma(n) = \sum_{d|n} d$, the sum of divisors of n ;
- more generally, $\sigma_k(n) = \sum_{d|n} d^k$, so $\sigma_0 = \tau$ and $\sigma_1 = \sigma$.

To prove this:

Lemma 2.13. If $f : \mathbb{Z}^+ \rightarrow \mathbb{C}$ is multiplicative, then so is $g : \mathbb{Z}^+ \rightarrow \mathbb{C}$, defined by $g(n) = \sum_{d|n} f(d)$.

Proof. Let m, n be coprime. Note that every divisor d of mn can be written as $d = d_1 d_2$, where $d_1 | m$, $d_2 | n$ and $(d_1, d_2) = 1$. Thus

$$g(mn) = \sum_{d|mn} f(d) = \sum_{d_1|m} \sum_{d_2|n} f(d_1 d_2) = \sum_{d_1|m} \sum_{d_2|n} f(d_1) f(d_2) = g(m)g(n).$$

□

Lemma 2.14. (i) For p a prime, $\phi(p^k) = p^{k-1}(p-1) = p^k(1 - \frac{1}{p})$.

(ii) $\phi(n) = n \prod_{p|n} (1 - \frac{1}{p})$.

Proof. (i): $\phi(p^k)$ counts the number of integers a between 1 and p^k such that $(p^k, a) = (p, a) = 1$. So we have p^a numbers, and we don't count the multiples of p , so $\phi(p^k) = p^k - p^{k-1}$.

(ii): Follows from the fact that ϕ is multiplicative.

□

Alternative proof that $\sum_{d|n} \phi(d) = n$ (cf Lemma 2.6).

Proof. Obviously the RHS is multiplicative. Since $\phi(n)$ is multiplicative, the LHS is multiplicative by Lemma 2.13, so it suffices to check for n a prime power, say $n = p^k$. To this end, compute

$$\sum_{d|p^k} \phi(d) = \phi(1) + \phi(p) + \dots + \phi(p^k) = 1 + (p-1) + (p^2-p) + \dots + (p^k - p^{k-1}) = p^k.$$

□

2.1 Polynomial congruences

Let $R = \mathbb{Z}, \mathbb{Q}, \mathbb{Z}/n\mathbb{Z}$ (or more generally any commutative ring). Set $R[X] = \{\text{polynomials with coefficients in } R\}$, i.e. $a_n X^n + a_{n-1} X^{n-1} + \dots + a_1 X + a_0$ for $a_i \in R$.

By definition, two polynomials are equal if and only if they have the same coefficients. We can check that $R[X]$ is a ring (with usual $+$ and \times).

Warning. The map $R[X] \rightarrow \{\text{functions } R \rightarrow R\}$ by $f \mapsto (\alpha \mapsto f(\alpha))$ is not always injective. For example, if $R = \mathbb{Z}/p\mathbb{Z}$ for p a prime, and $f(X) = X^p - X$, then $f(\alpha) = 0 \forall \alpha \in R$, but f is not the zero function.

Question. Can we show that if $f \in R[X]$ has degree n , then f has at most n roots in R ?

Answer. No. For example, take $R = \mathbb{Z}/8\mathbb{Z}$, then $f(X) = X^2 - 1$ has 4 solutions in $\mathbb{Z}/8\mathbb{Z}$.

Let $R = \mathbb{Z}, \mathbb{Q}, \mathbb{Z}/n\mathbb{Z}$ (or any commutative ring).

We have a **division algorithm** on $R[X]$:

Let $f, g \in R[X]$ and suppose the leading coefficient of g is a unit. Then $\exists q, r \in R[X]$ such that $f(X) = Q(X)g(X) + r(X)$ and $\deg(r) < \deg(g)$.

Proof. By induction on $\deg(f)$. If $\deg(f) < \deg(g)$, take $q = 0, r = f$. Otherwise, let $f(X) = aX^m + \dots$ and $g(X) = bX^n + \dots$ with $m \geq n$ and b a unit.

Let $f_1(X) = f(X) - ab^{-1}X^{m-n}g(X)$. Then $\deg(f_1) < \deg(f)$, so by the induction hypothesis, $f_1(x) = q_1(x)g(x) + r_1(x)$ for some $q_1, r_1 \in R[X]$ and $\deg(r_1) < \deg(g)$. Now take $q(X) = ab^{-1}X^{m-n} + q_1(X)$ and $r = r_1$, so we're done. \square

Corollary 2.15. If $f \in R[X]$ and $\alpha \in R$ is such that $f(\alpha) = 0$, then $f(X) = (X - \alpha)f_1(X)$ for some $f_1 \in R[X]$.

Proof. By the division algorithm, $f(X) = (X - \alpha)f_1(X) + r$ for some $r \in R$ (as $\deg(r) < \deg(X - \alpha)$). Plug in $X = \alpha$ to get $r = 0$. \square

Definition 2.16. R is an **integral domain** if R has no zero divisors, i.e. $\alpha, \beta \in R, \alpha\beta = 0 \implies \alpha = 0$ or $\beta = 0$.

Note. Let $n > 1$. Then $\mathbb{Z}/n\mathbb{Z}$ is an integral domain $\iff n$ is prime.

Theorem 2.17. If R is an integral domain, then any polynomial $f \in R[X]$ of degree n has at most n roots.

Proof. By induction on n , the degree of f . If $n = 0$, then our polynomial is a nonzero constant and we're done. Now suppose $\exists \alpha \in R$ such that $f(\alpha) = 0$ (otherwise we're done). By Corollary 2.15, $f(X) = (X - \alpha)f_1(X)$. Since R is an integral domain, every root of f , except possibly α is a root of f_1 . By induction, f_1 has at most $n - 1$ roots, hence f has at most n roots and we're done. \square

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Corollary 2.18 (Lagrange's Theorem). Let p be a prime and $a_0, \dots, a_n \in \mathbb{Z}$ with $p \nmid a_n$. Then the congruence

$$a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0 \equiv 0 \pmod{p}$$

has at most n solutions mod p .

Proof. Take $R = \mathbb{Z}/p\mathbb{Z}$ in Theorem 2.17. □

Remark. In this course, we will refer to the above theorem as Lagrange's Theorem.

Example 2.19. Let p be a prime. We will factor $X^{p-1} - 1 \pmod{p}$. Let $f(X) = X^{p-1} - 1 - \prod_{a=1}^{p-1} (X - a)$ in $\mathbb{Z}/p\mathbb{Z}[X]$. By Fermat's Little Theorem, f has at least $p - 1$ roots mod p . But $\deg(f) < p - 1$, since the X^{p-1} terms cancel out, so by Lagrange's Theorem, $f = 0$, i.e. $X^{p-1} - 1 = \prod_{a=1}^{p-1} (X - a)$ in $\mathbb{Z}/p\mathbb{Z}[X]$. Plugging in $X = 0$ gives $(p - 1)! \equiv -1 \pmod{p}$, i.e. Wilson's Theorem.

Example 2.20. Working mod 7, the powers of 3 (starting from 0) are 1, 3, 2, 6, 4, 5. So $(\mathbb{Z}/7\mathbb{Z})^\times$ is cyclic, generated by 3.

Theorem 2.21. Let p be a prime. Then $(\mathbb{Z}/p\mathbb{Z})^\times$ is cyclic.

Proof. Let $S_d = \{a \in (\mathbb{Z}/p\mathbb{Z})^\times \mid \text{ord}(a) = d\}$. Suppose $S_d \neq \emptyset$, say $a \in S_d$. Then $1, a, a^2, \dots, a^{d-1}$ are distinct elements in $\mathbb{Z}/p\mathbb{Z}$ and they are solutions of $x^d \equiv 1 \pmod{p}$. By Lagrange's theorem, this has at most d solutions, and we found d solutions, so those are all of them, i.e. $S_d \subseteq \{1, a, a^2, \dots, a^{d-1}\}$. Note that the LHS is a cyclic group of order d , so this has $\phi(d)$ elements of order d .

We conclude that for every d , $|S_d| = 0$ or $|S_d| = \phi(d)$. In particular, $|S_d| \leq \phi(d)$. Hence

$$p - 1 \stackrel{(\star)}{=} \sum_{d \mid (p-1)} |S_d| \leq \sum_{d \mid (p-1)} \phi(d) = p - 1,$$

where (\star) follows since we just count all the elements in $(\mathbb{Z}/p\mathbb{Z})^\times$. Hence $|S_d| = \phi(d) \forall d \mid (p - 1)$. In particular, $S_{p-1} \neq \emptyset$, i.e. $(\mathbb{Z}/p\mathbb{Z})^\times$ contains elements of order $p - 1$, i.e. $(\mathbb{Z}/p\mathbb{Z})^\times$ is cyclic. □

Remark. The same argument shows that any finite subgroup of the multiplicative group of a field is cyclic.

Definition 2.22. An integer a such that $a \pmod{n}$ generates $(\mathbb{Z}/n\mathbb{Z})^\times$ is called a **primitive root** mod n .

Theorem 2.21 showed that primitive roots exist mod p .

Example 2.23. Let $p = 19$. Let d be the order of 2 in $(\mathbb{Z}/19\mathbb{Z})^\times$. We know $d \mid 18$, so we work out

$$\begin{aligned} 2^3 &\equiv 8 \pmod{19} \\ 2^6 &\equiv 7 \not\equiv 1 \pmod{19} \implies d \nmid 6 \\ 2^9 &\equiv -1 \not\equiv 1 \pmod{19} \implies d \nmid 9, \end{aligned}$$

so $d = 18$ and hence 2 is a primitive root mod 19.

In general, $g \in \mathbb{Z}$ (coprime to p) is a primitive root mod p if and only if $g^{\frac{p-1}{q}} \not\equiv 1 \pmod{p} \quad \forall \text{ primes } q \mid (p-1)$.

Remark. The number of primitive roots mod p is $\phi(p-1) = \phi(\phi(p))$.

Here are some (open) problems concerning primitive roots:

- (i) Artin's conjecture (1927) – Let $a > 1$ be an integer which is not a square. Then a is a primitive root mod p for infinitely many primes p . This is unknown for $a = 2$. Hooley (1967) proved this assuming GRH. Heath-Brown (1986) proved that Artin's conjecture holds for at least one of 2, 3 or 5. In fact, he proved something stronger: he proved the conjecture fails for at most 2 prime values of a .
- (ii) How large is the smallest primitive root mod p ? Burgess (1962) showed it is $\leq cp^{1/4+\epsilon} \quad \forall \epsilon > 0$ and some constant $c = c(\epsilon)$. Shoup (1992) showed it is $\leq c(\log p)^6$ assuming GRH.

We now consider $\mathbb{Z}/p^n\mathbb{Z}$ for $n > 1$. For $n \geq 3$, there is a surjective group homomorphism from $(\mathbb{Z}/2^n\mathbb{Z})^\times \rightarrow (\mathbb{Z}/8\mathbb{Z})^\times = \{\pm 1, \pm 3\} \cong C_2 \times C_2$, so $(\mathbb{Z}/2^n\mathbb{Z})^\times$ is not cyclic (since generators map to generators).

Theorem 2.24. Let p be an odd prime. Then $(\mathbb{Z}/p^n\mathbb{Z})^\times$ is cyclic $\forall n \geq 1$.

We divide the proof into 3 lemmas.

Lemma 2.25. Let $n \geq 2$. Then g is a primitive root mod p^n if and only if the following two conditions hold:

$$\begin{cases} g \text{ is a primitive root mod } p \\ g^{p^{n-2}(p-1)} \not\equiv 1 \pmod{p^n} \end{cases}.$$

Proof. (\implies) is clear, as $\phi(p^n) = p^{n-1}(p-1)$.

(\impliedby): Let d be the order of g in $(\mathbb{Z}/p^n\mathbb{Z})^\times$. Then $d \mid \phi(p^n) = p^{n-1}(p-1)$. Since $g^d \equiv 1 \pmod{p^n}$, we have $g^d \equiv 1 \pmod{p}$. Hence by assumption 1, we have $(p-1) \mid d$. Say $d = p^j(p-1)$ for some $0 \leq j \leq n-1$. If $j \leq n-2$, then this contradicts assumption 2. Hence $j = n-1$, so $d = \phi(p^n)$ is a primitive root mod p^n . \square

Next we show $\exists g \in \mathbb{Z}$ satisfying conditions 1 and 2 in the case $n = 2$.

Lemma 2.26. $\exists g \in \mathbb{Z}$ a primitive root mod p such that $g^{p-1} \not\equiv 1 \pmod{p^2}$.

Proof. Let g be a primitive root mod p . If $g^{p-1} \equiv 1 \pmod{p^2}$, then consider $g + p$, which is still a primitive root mod p , but

$$(g + p)^{p-1} = g^{p-1} + (p-1)g^{p-2}p + \dots \equiv 1 + (p-1)g^{p-2}p \pmod{p^2},$$

where the second term is not divisible by p^2 , so $(g + p)^{p-1} \not\equiv 1 \pmod{p^2}$. \square

Next we show that if g is a primitive root mod p^2 , then it is a primitive root mod $p^n \forall n \geq 2$.

Lemma 2.27. If $g^{p-1} \not\equiv 1 \pmod{p^2}$, then $g^{p^{n-2}(p-1)} \not\equiv 1 \pmod{p^n} \forall n \geq 2$.

Proof. By induction on n , the case $n = 2$ being given. Suppose the result is true for n . By Euler-Fermat, $g^{p^{n-2}(p-1)} \equiv 1 \pmod{p^{n-1}}$, so $g^{p^{n-2}(p-1)} = 1 + bp^{n-1}$ for some $b \in \mathbb{Z}$, where $p \nmid b$ by the induction hypothesis. Taking p^{th} powers gives

$$\begin{aligned} g^{p^{n-1}(p-1)} &= (1 + bp^{n-1})^p = 1 + bp^n + \binom{p}{2}b^2p^{2(n-1)} + \dots \equiv \\ &1 + bp^n + \binom{p}{2}b^2p^{2(n-1)} \star 1 + bp^n \pmod{p^{n+1}}, \end{aligned}$$

where \star follows since p is odd, so $p \mid \binom{p}{2}$ (and also we use $3(n-1) \geq n+1$ and $2(n-1)+1 \geq n+1$). Thus $g^{p^{n-1}(p-1)} \equiv 1 + bp^n \not\equiv 1 \pmod{p^{n+1}}$, so the result follows for $n+1$. \square

This completes the proof of Theorem 2.24.

Example 2.28. We saw 3 is a primitive root mod 7. We calculate $3^3 = -1 + 4 \cdot 7$, so $3^6 \equiv 1 - 8 \cdot 7 \not\equiv 1 \pmod{7^2}$. Hence 3 is a primitive root mod $7^n \forall n$.

For the case $p = 2$, let $G = \{a \in (\mathbb{Z}/2^n\mathbb{Z})^\times \mid a \equiv 1 \pmod{4}\}$. Then $(\mathbb{Z}/2^n\mathbb{Z})^\times \cong \{\pm 1\} \times G$ by $a + 2^n\mathbb{Z} \mapsto \begin{cases} (1, a + 2^n\mathbb{Z}) & \text{if } a \equiv 1 \pmod{4} \\ (-1, -a + 2^n\mathbb{Z}) & \text{if } a \equiv 3 \pmod{4}. \end{cases}$

Exercise. Show that G is cyclic (and generated by 5).

Exercise. For which n is $(\mathbb{Z}/n\mathbb{Z})^\times$ cyclic?