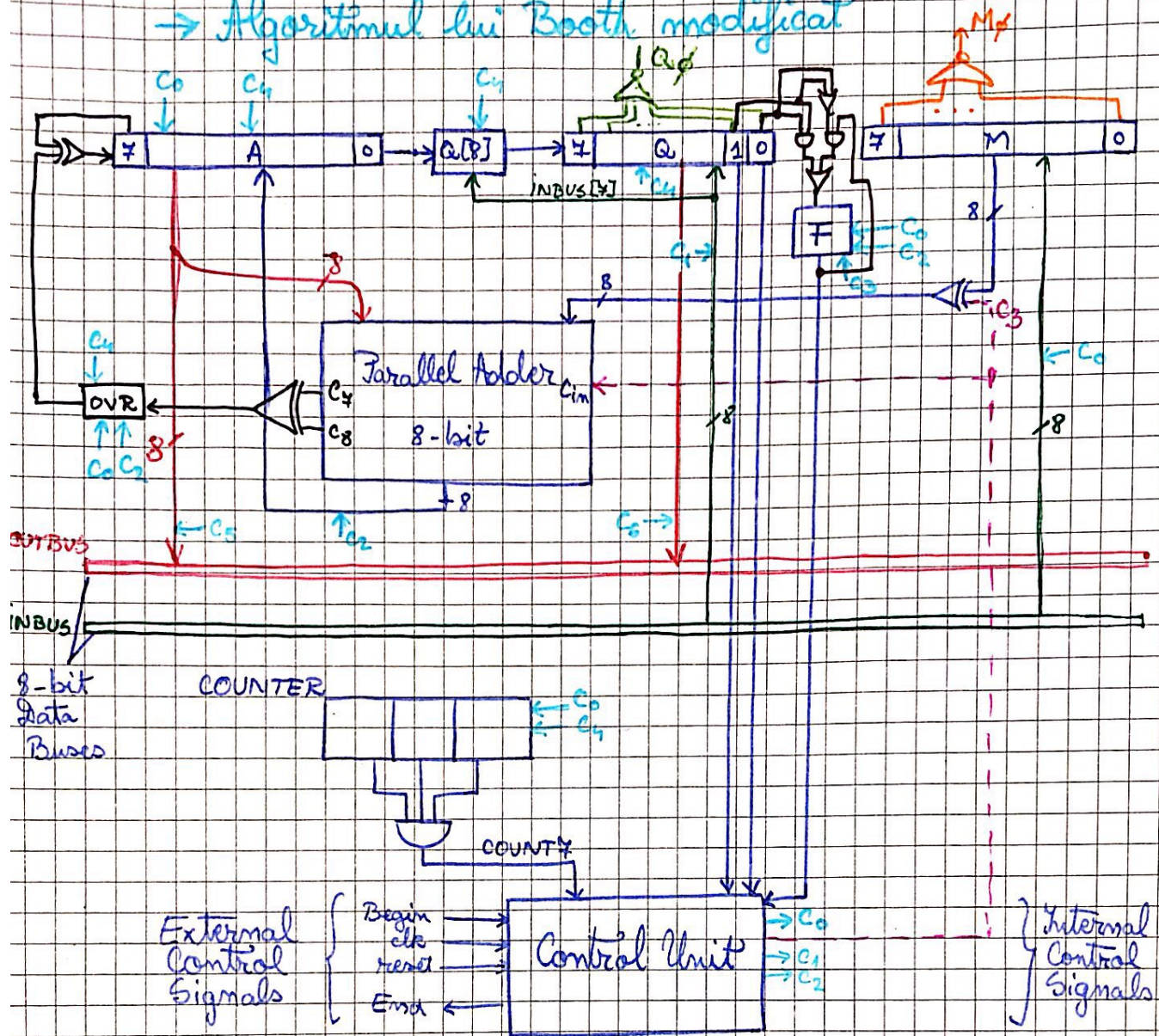
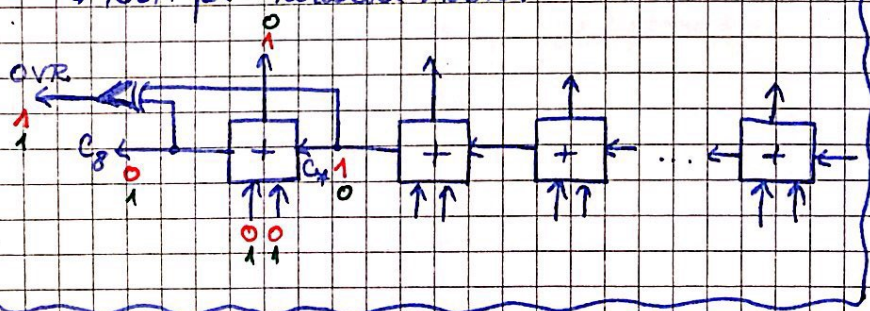


1.1. Introduction

→ Algoritmul lui Booth modificat

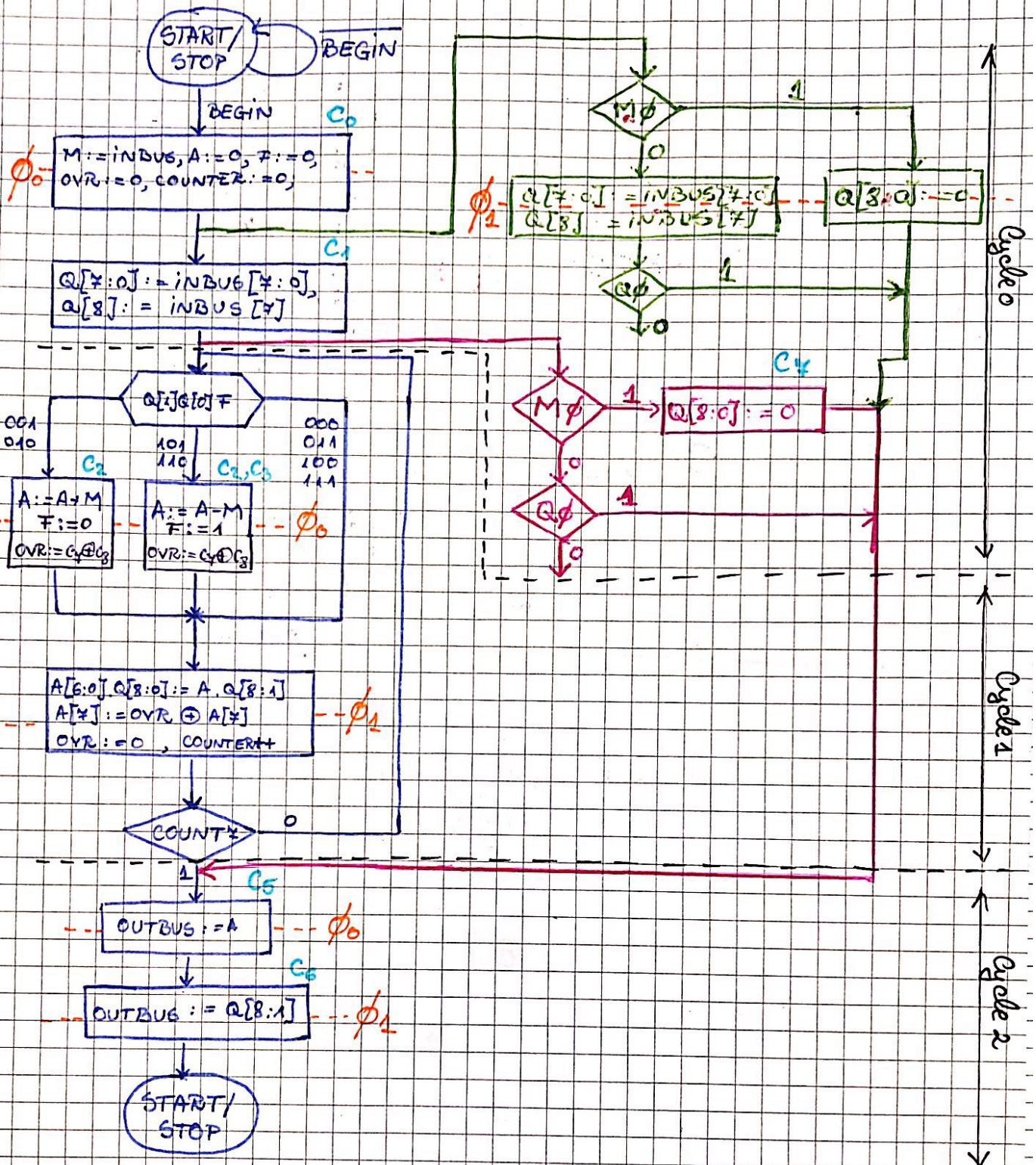
RCA pt^o Parallel Adder

$Q[1]$	$Q[0]$	F	F'	OP
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	0
1	0	1	1	1
1	1	0	1	1
1	1	1	1	0

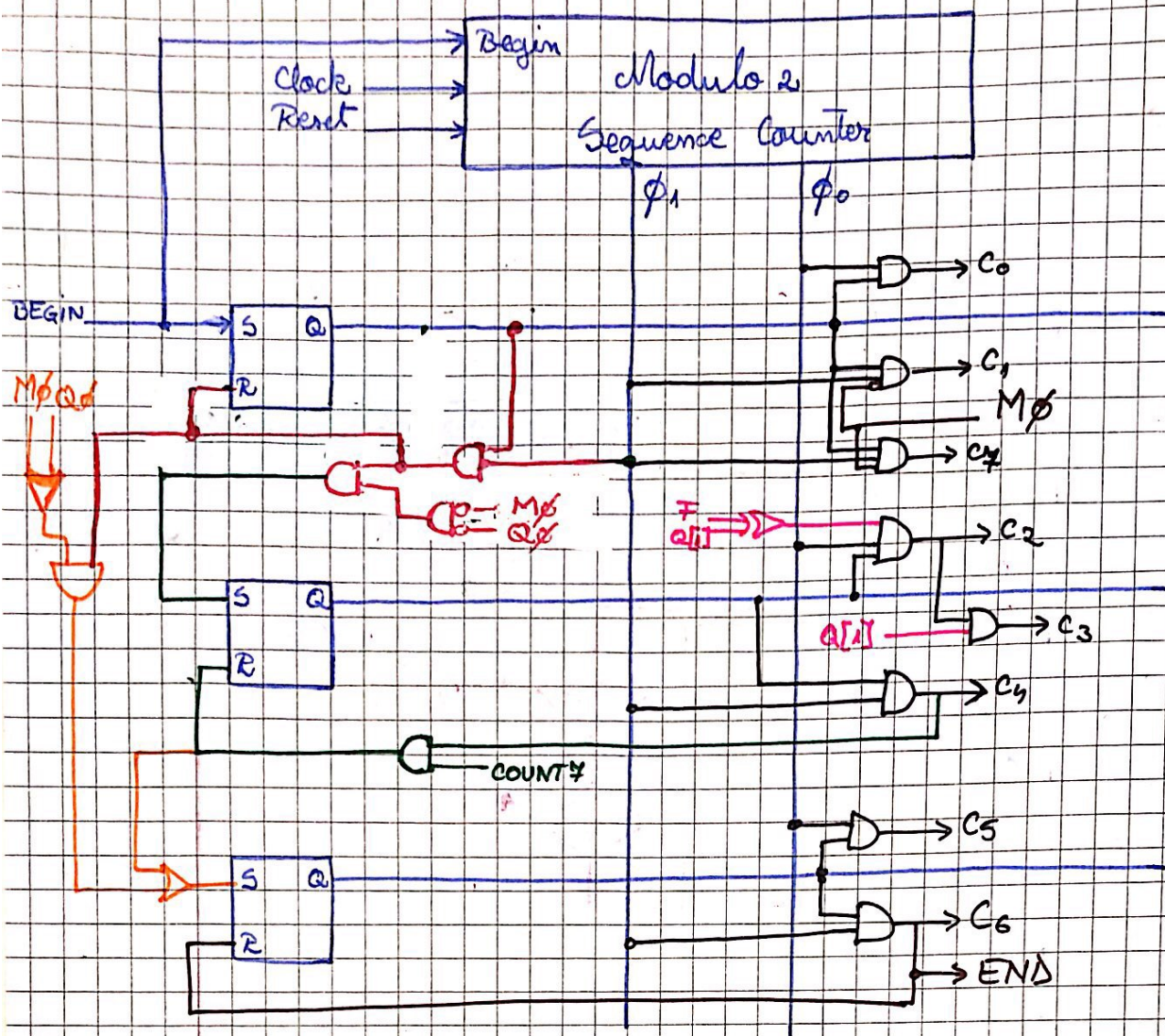
$Q[1]Q[0]$	00	01	11	10
0			1	
1		1	1	1

$$\Rightarrow F' = Q_1 Q_0 + F(Q_0 + Q_1)$$

- Q_{ϕ} este 1 dacă $Q = 0$
 - M_{ϕ} este 1 dacă $M = 0$
- } toate procesele alea la să evităm înmulțirea cu 0



Se poate implementa sau , e mai eficient



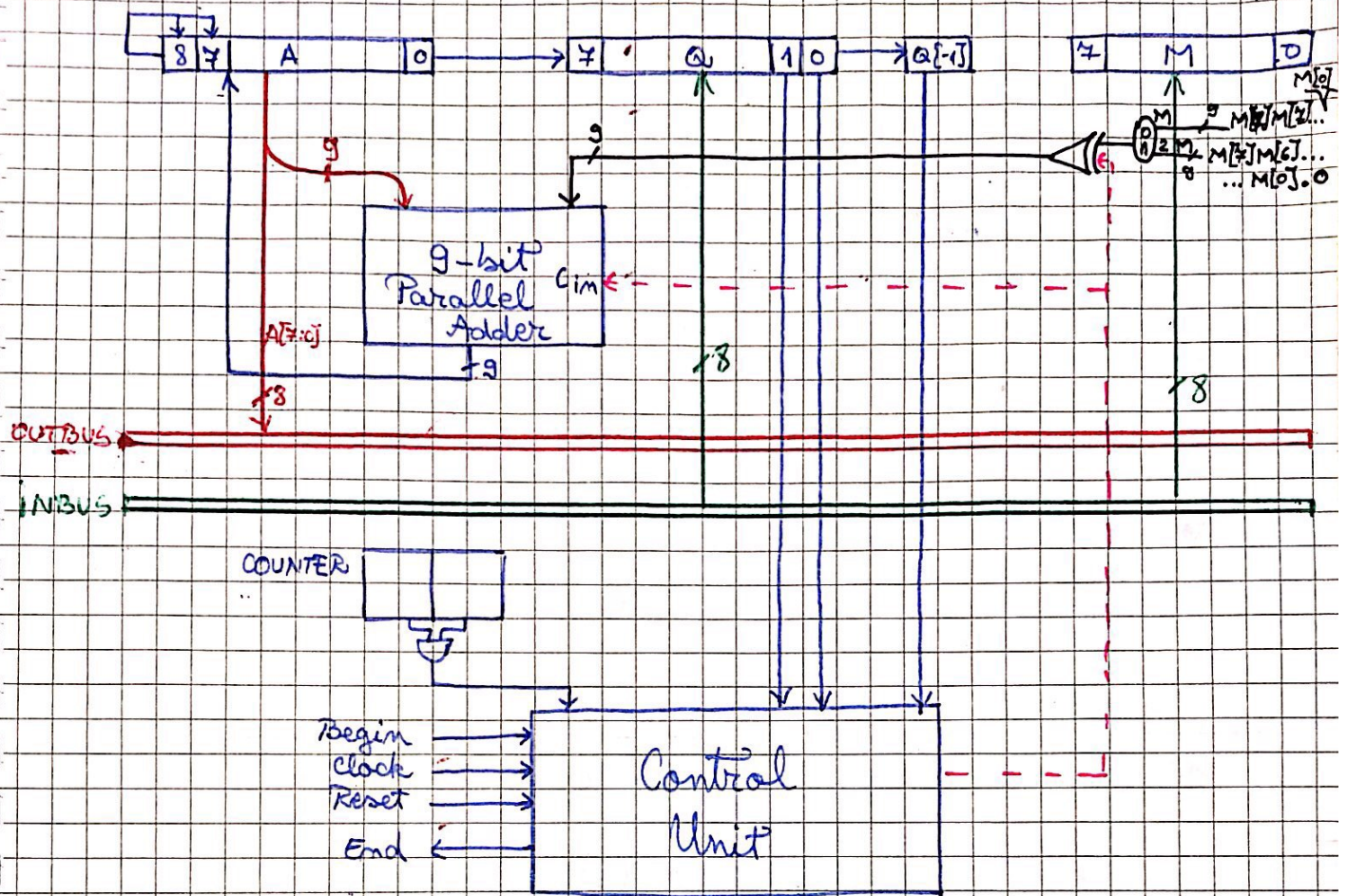
1.2. Speeding up: higher radix

Radix-4 step i with weight 2^i

b_{i+1}	b_i	b_{i-1}	OP
0	0	0	0
0	0	1	+A
0	1	0	+A
0	1	1	+2A
1	0	0	-2A
1	0	1	-A
1	1	0	-A
1	1	1	0

$$\begin{aligned}
 A \times 0 \times 2^i + A \times 0 \times 2^{i+1} &= 0 \times 1 \times 2^i \\
 + A \times 2^i + 0 \times A \times 2^{i+1} &= +A \times 2^i \\
 - A \times 2^i + A \times 2^{i+1} &= A \times 2^i \\
 0 + A \times 2^{i+1} &= +2A \times 2^i \\
 0 - A \times 2^{i+1} &= -2A \times 2^i \\
 + A \times 2^i - A \times 2^{i+1} &= -A \times 2^i \\
 + A \times 2^i + 0 &= +A \times 2^i
 \end{aligned}$$

• Pentru a salva clockul, treb. să putem face shiftări duble \Rightarrow
 \Rightarrow A pe 9 biți.



LAB 2

♥ $M = 13 = 01101_{5M}$

♥ $Q = 4 = 0111_{5M}$

Decim să facem înmulțirea în Radix 2 pt Booth modificat