Estrutura de Dados I

Merge-Sort

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Material compilado de: Cormen.

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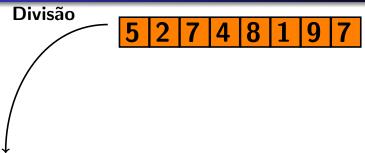
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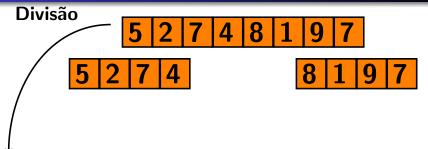
O algoritmo Merge-Sort foi inventado em 1945 por John Von Neumann. O algoritmo segue o paradigma de divisão e conquista (divide and conquer):

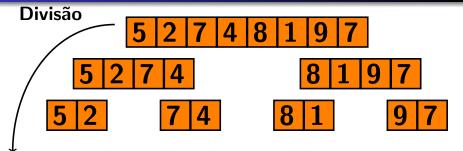
- Dividir: descubra o ponto médio do sub-arranjo (tempo constante).
- Conquistar: resolva recursivamente dois subproblemas de tamanho n/2.
- Combinar: combine os dois sub-arranjos em um único conjunto ordenado (tempo de n).

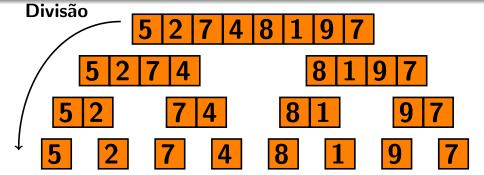
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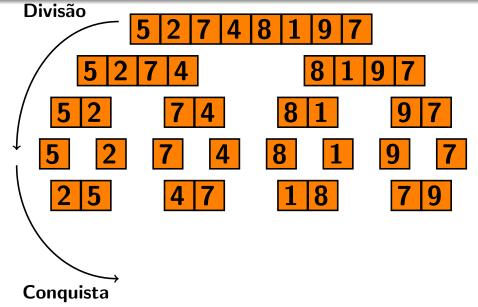
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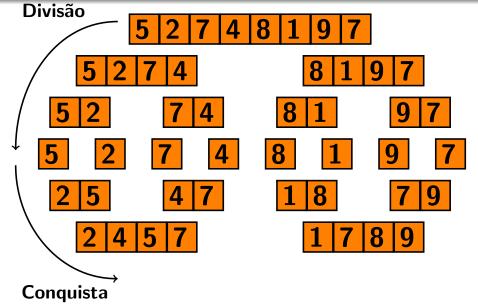


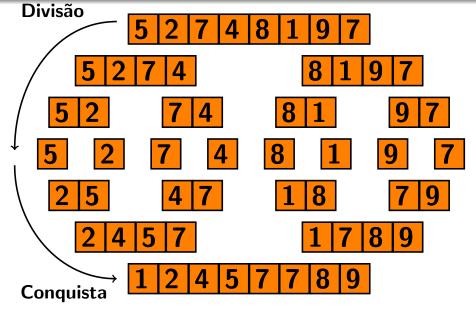












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Merge-Sort
$$(A, \mathbf{e}, \mathbf{d}, O)$$

1. se $\mathbf{e} < \mathbf{d}$ então

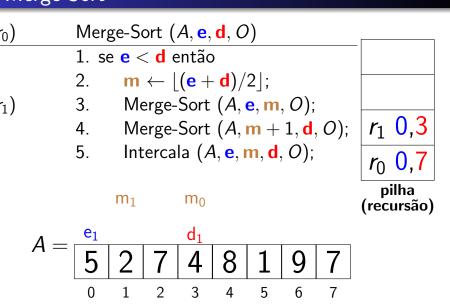
2. $\mathbf{m} \leftarrow \lfloor (\mathbf{e} + \mathbf{d})/2 \rfloor$;

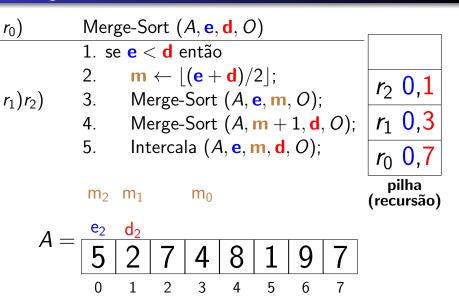
3. Merge-Sort $(A, \mathbf{e}, \mathbf{m}, O)$;

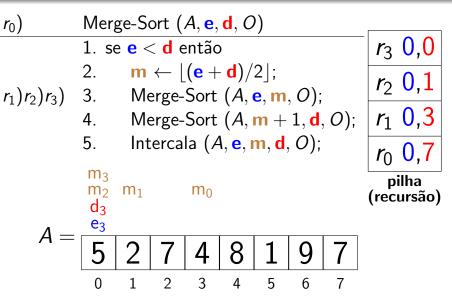
4. Merge-Sort $(A, \mathbf{m} + 1, \mathbf{d}, O)$;

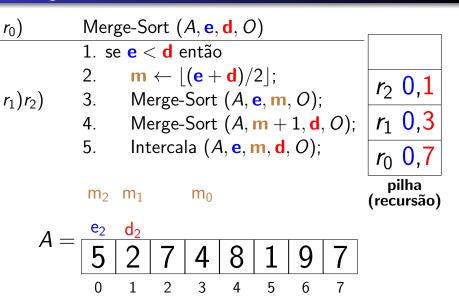
5. Intercala $(A, \mathbf{e}, \mathbf{m}, \mathbf{d}, O)$;

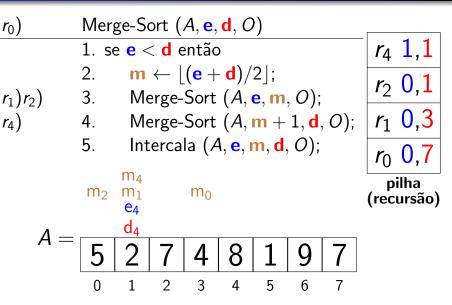
 \mathbf{m}_0
 \mathbf{m}_0

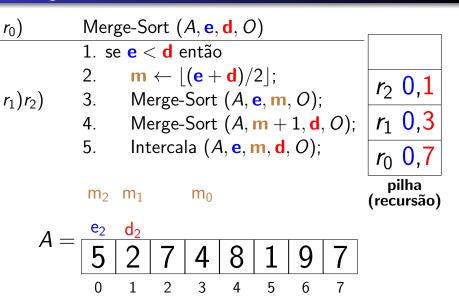


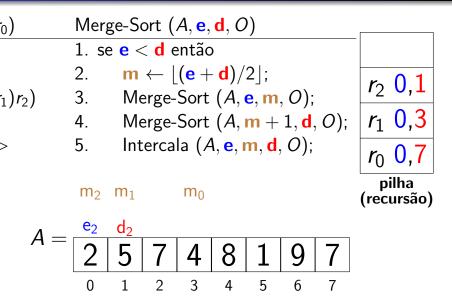


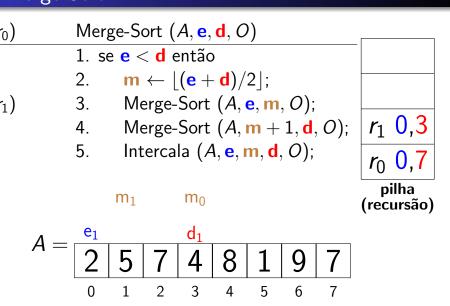


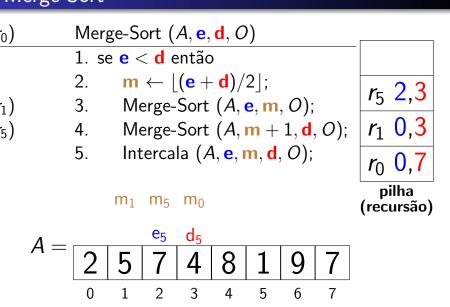


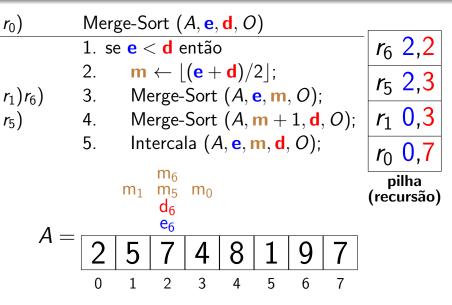


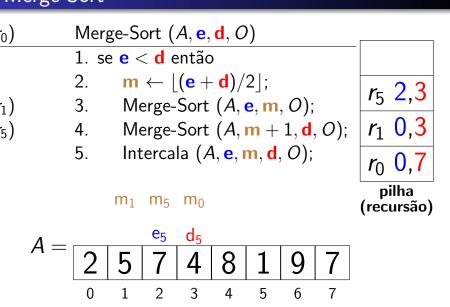


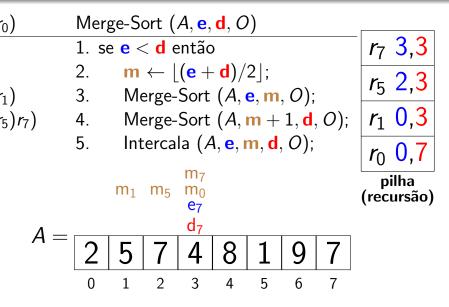


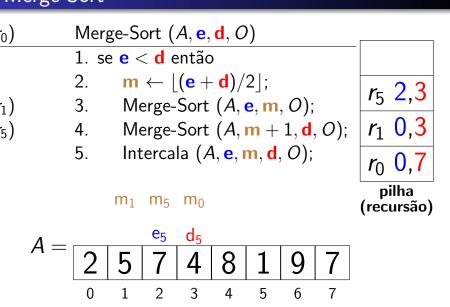


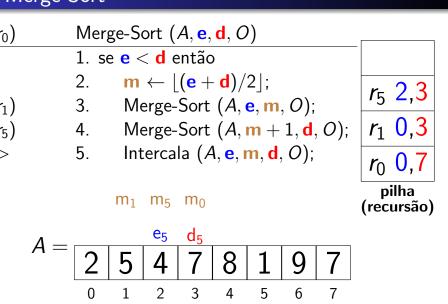


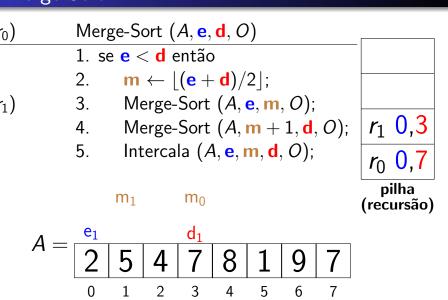


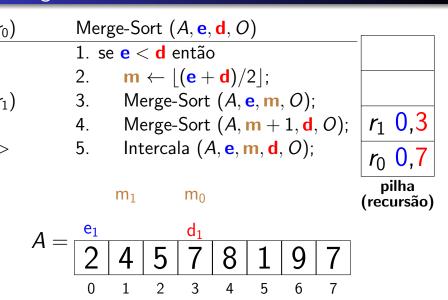


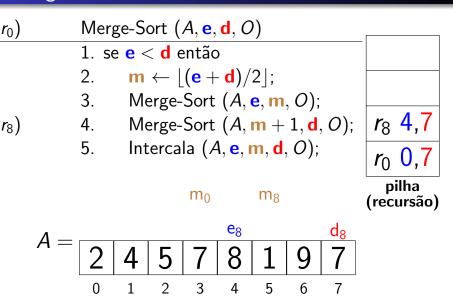










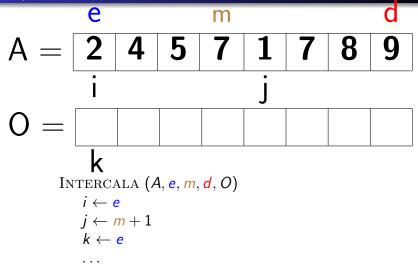


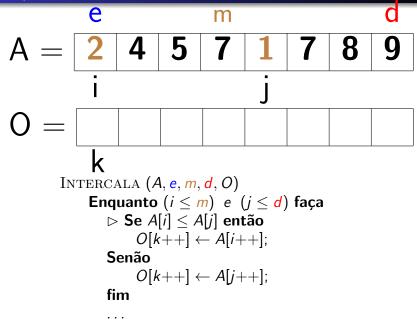
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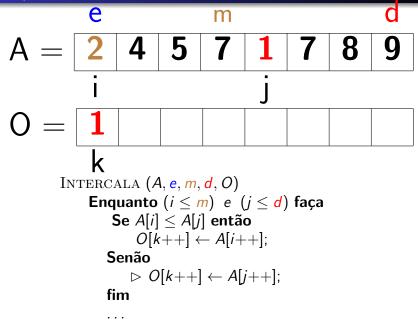
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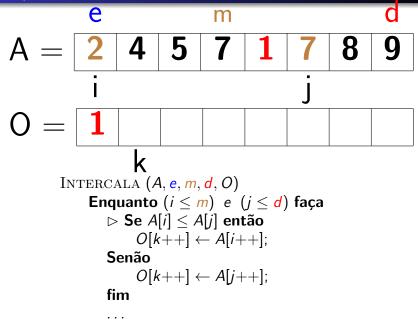
$$\mathsf{O} = igcup | \mathsf{O} = \mathsf{O}$$

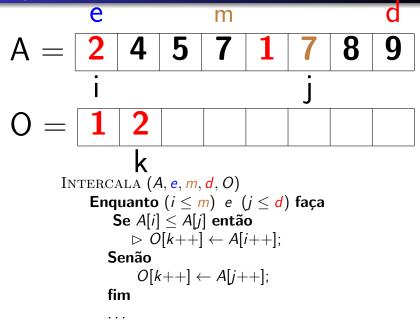
Intercala (A, e, m, d, O)

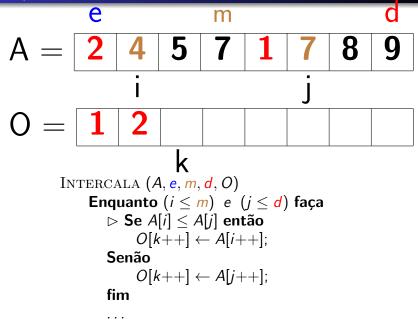


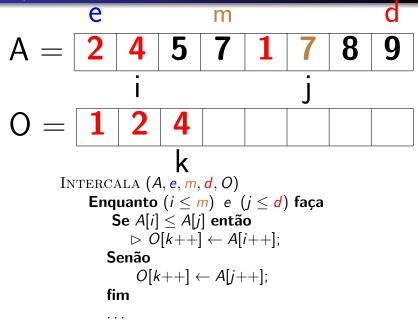


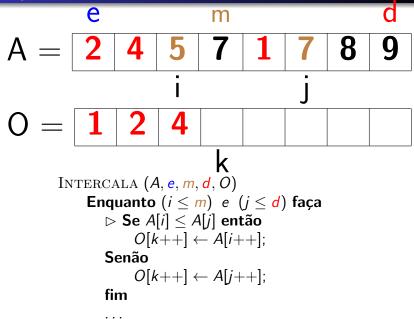


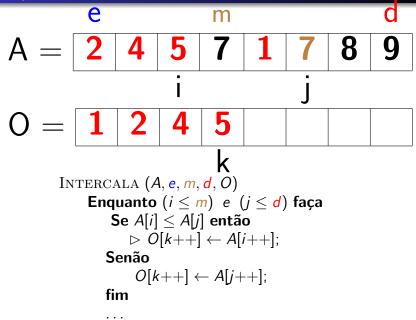




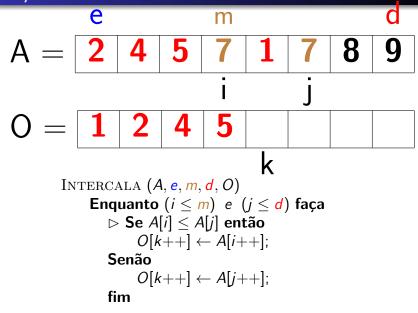


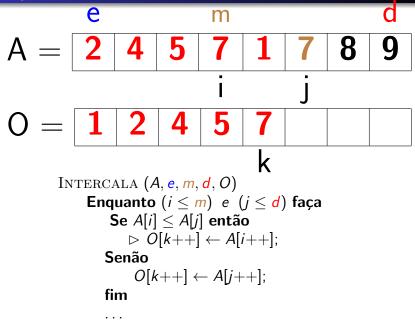


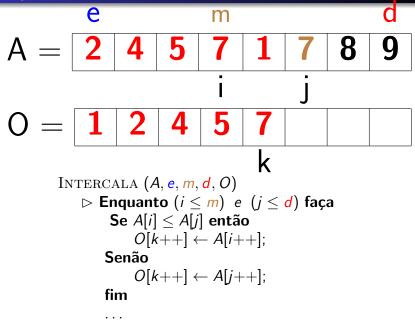


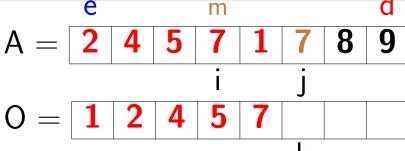


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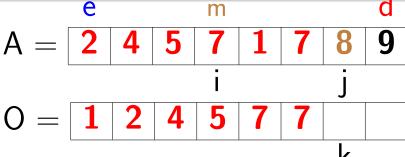




INTERCALA (A, e, m, d, O) \triangleright Enquanto $(j \le d)$ faça $O[k++] \leftarrow A[j++]$;



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INTERCALA (A, e, m, d, O) **Enquanto** $(j \le d)$ faça $\triangleright O[k++] \leftarrow A[j++];$

$$O = | 1 | 2 | 4 | 5 | 7 | 7 | 8$$

INTERCALA (A, e, m, d, O) \triangleright **Enquanto** $(j \le d)$ **faça** $O[k++] \leftarrow A[j++];$

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Complexidade

O algoritmo do Merge-Sort é estável e sua equação de recorrência é dado por:

$$T(\mathbf{n}) = \left\{ egin{array}{ll} \Theta(1) & ext{se} & \mathbf{n} = 1, \\ T(\lceil \mathbf{n}/2 \rceil) + T(\lfloor \mathbf{n}/2 \rfloor) + \Theta(\mathbf{n}) & ext{se} & \mathbf{n} > 1. \end{array}
ight.$$

Complexidade de tempo: $\Theta(n \log n)$.

Complexidade de espaço: $\Theta(n)$.