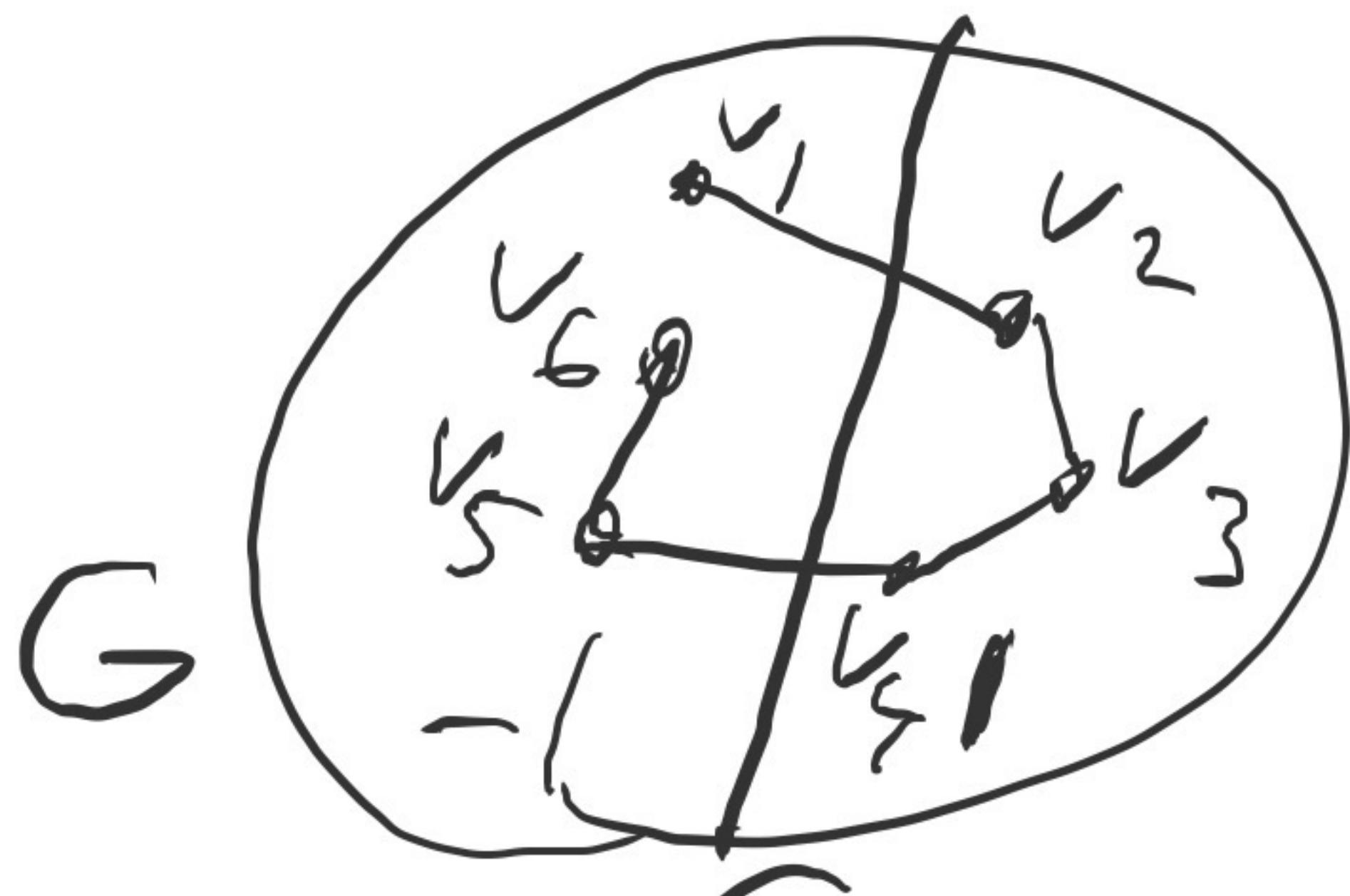


Expander Random Walks: A Fourier Analytic Approach

Gil Cohen, Noam Peri, and Amnon Ta-Shma

Overall question: Let's say we have a d -regular expander G and we take a pseudo-random sequence as follows:

1. Label half of the vertices of G $-l$ and label the other half l
2. Take a $(t-1)$ step random walk v_1, v_2, \dots, v_t and output the labels of the vertices



-1, 1, 1, -1, -1

Q: For which functions does this behave like a random $x \in \{-1, 1\}^t$.

This paper uses Fourier analysis to show that

Symmetric functions, functions in AC^0 , and functions computable by a low width read once branching program are fooled.

Preliminaries: λ -spectral expanders

For a d -regular graph G , we'll also use G to denote the normalized adjacency matrix.

$$G_{ij} = \begin{cases} \frac{1}{d} & \text{if } (i,j) \in E(G) \\ 0 & \text{otherwise} \end{cases}$$

Definition: We say that G is a λ -spectral expander if for all eigenvectors except for $\vec{1}$, the magnitude of its eigenvalue is at most λ .

Definition: Define J to be $\frac{1}{n} \cdot (\text{all ones matrix})$ $\forall i, j \quad J_{ij} = \frac{1}{n}$.
Key property: $G = J + E$ where $\|E\| \leq \lambda$ and $EJ = JE = 0$.

Preliminaries: Fourier Analysis Over $\{-1/3\}^+$

Fourier characters: $X_S = \prod_{j \in S} X_j$

Fourier Coefficients: $\hat{f}_S = E_{x \in \{-1/3\}^+} [f(x) X_S(x)]$

Fourier decomposition: $f = \sum_{S \subseteq \mathbb{Z}} \hat{f}_S X_S$

Parseval's Theorem: $\sum_{S \subseteq \mathbb{Z}} \hat{f}_S^2 = E_{x \in \{-1/3\}^+} [f^2]$

We'll be analyzing ± 1 -valued functions
so this will be 1.

Preliminaries: Difference from Random

Def: Given a graph G and a labeling $\text{val}: V(G) \rightarrow \{-1, 1\}$
(where half of the vertices are labeled 1)

define $RW_{G, \text{val}, t}$ to be the distribution obtained by
taking the labels of a $(t-1)$ step random walk

Definition: Given a function $f: \{-1, 1\}^+ \rightarrow \{-1, 1\}$,

define $\mathcal{E}_{G, \text{val}}(f) = E_{x \sim RW_{G, \text{val}, t}} [f(x)] - E_{x \sim \{-1, 1\}^+} [f(x)]$

Define $\mathcal{E}_t(f) = \max_{\substack{\text{δ-spectral expanders G} \\ \text{val}}} |\mathcal{E}_{G, \text{val}}(f)|$

Idea: Bound $\mathcal{E}_t(f)$ by analyzing $\mathcal{E}_{G, \text{val}}(X_S)$

Results on Symmetric Functions

Definition: We say that a function

$f: \{-1, 1\}^+ \rightarrow \{-1, 1\}$ is symmetric if

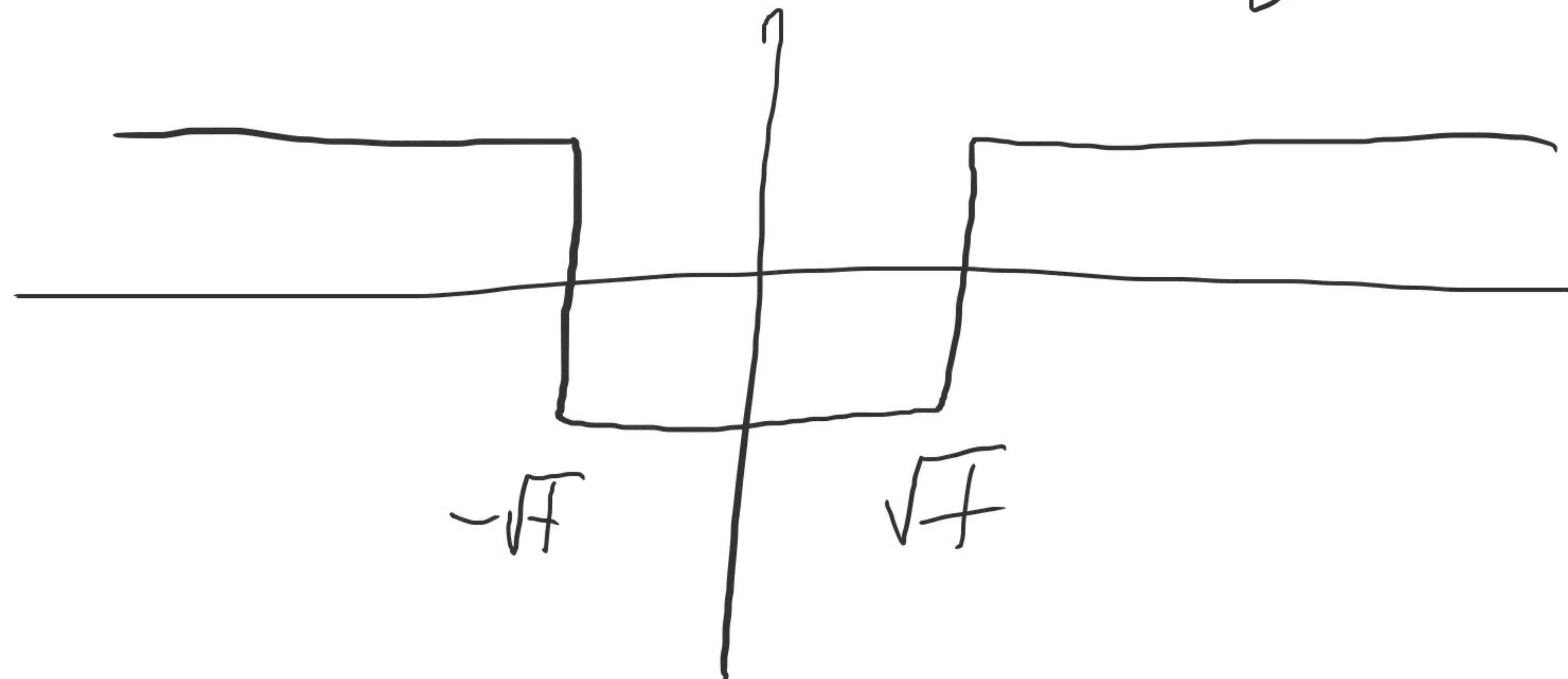
$f(x_1, x_2, \dots, x_t)$ only depends on the number of -1s and 1s in x_1, \dots, x_t .

Theorem 1.1': For all symmetric functions f ,
 $\sum_x f(x)$ is $O(\lambda)$.

Restatement: If we only consider the number of 1s and -1s in x , the total variation distance between $RW_{G, \text{val}, t}$ and $U_t = \text{uniform dist on } \{-1, 1\}^t$ is $O(\lambda)$

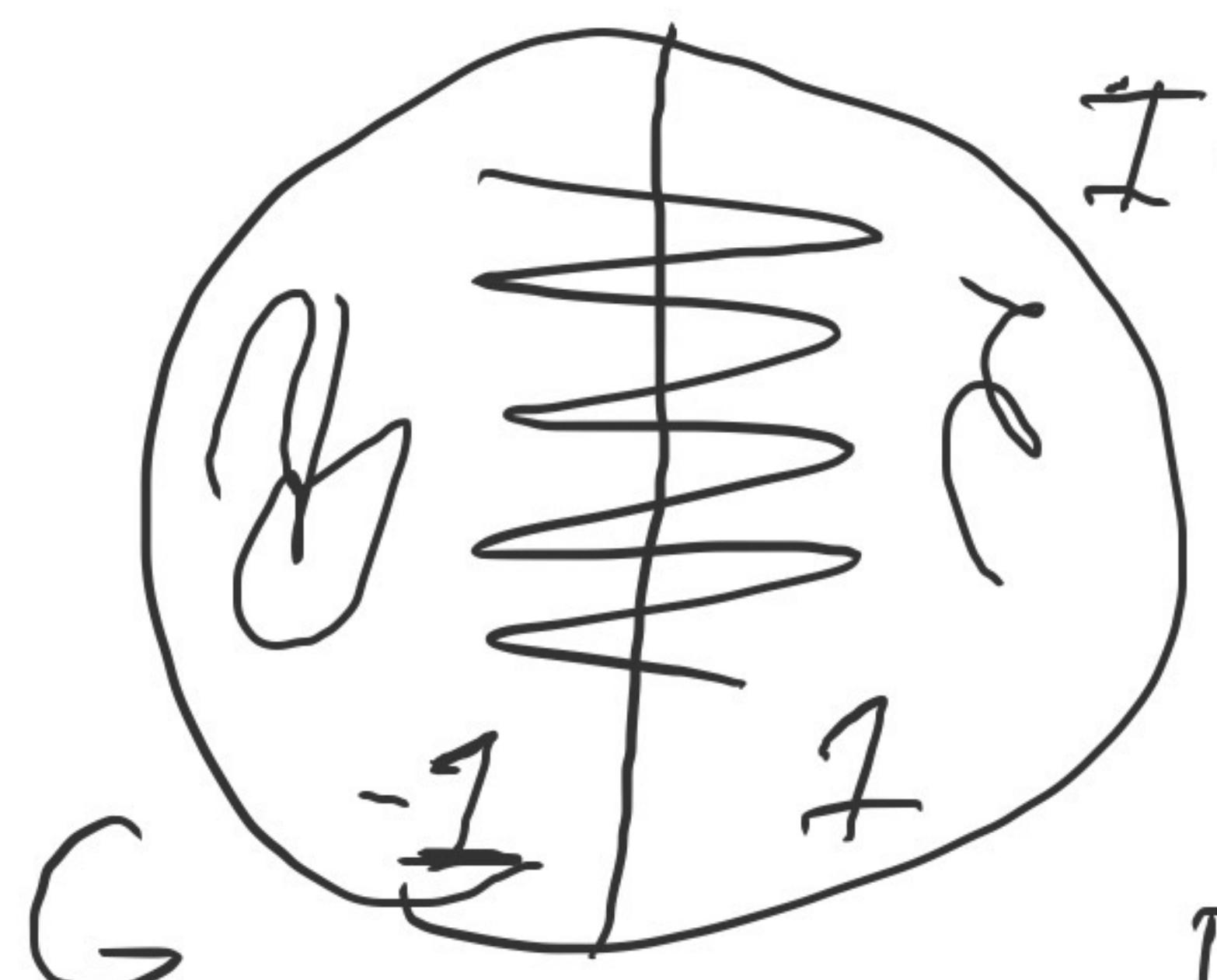
Theorem 1.4: $\sum_x MAJ_t(x)$ is $O\left(\frac{\lambda^2}{\sqrt{t}}\right)$

I think $O(\lambda)$ is tight. I think



is an example.

Note: $RW_{G, \text{val}, t}$ does not necessarily fool all functions



If $fr(\text{crossing}) = \frac{1}{2} + \Omega(\lambda)$ then

for $RW_{G, \text{val}, t}$, $E[\# \text{ of crossings}]$

is $\frac{t}{2} + \Omega(\lambda t)$

For U_t , # of crossings is $\frac{t}{2} \pm \Omega(\sqrt{F})$.

Results on General Functions

Definition: Given a function $f: \{0,1\}^+ \rightarrow \{0,1\}$,
define $L_{1,k}(f) = \sum_{\substack{S \subseteq \{1\}^+ \\ |S|=k}} |f_S|$

Definition: Define $L_i^+(b)$ to be the family
of functions $f: \{0,1\}^+ \rightarrow \{0,1\}$
such that

$$\forall k \quad L_{1,k}(f) \leq b^k$$

Claim 5.2': For any function $f \in L_i^+(b)$,
 $\sum_x f(x)$ is $O(b^2)$

Avishay Tal

Tight Bounds on the Fourier spectrum of AC^0

Theorem: Any function computable by a size s depth d AC^0 circuit is in $L_1^{+}(b)$ for $b = O((\log s)^{(d-1)})$
 $\mathbb{E}_x(f)$ is $O(\lambda (\log s)^{2(d-1)})$

Eshan Chattopadhyay, Pooya Hatami, Omer Reingold,
and Avishay Tal

Improved Pseudorandomness for
Unordered Branching Programs Through Local Monotonicity

Theorem: If f can be computed by a width w
read once branching program, $b = O((\log t)^w)$
 $\mathbb{E}_x(f)$ is $O(\lambda ((\log t)^w))$

$b \leq DT(f)$

$E_X(f)$ is $\mathcal{O}(\lambda DT(f)^2)$.

Analysis:

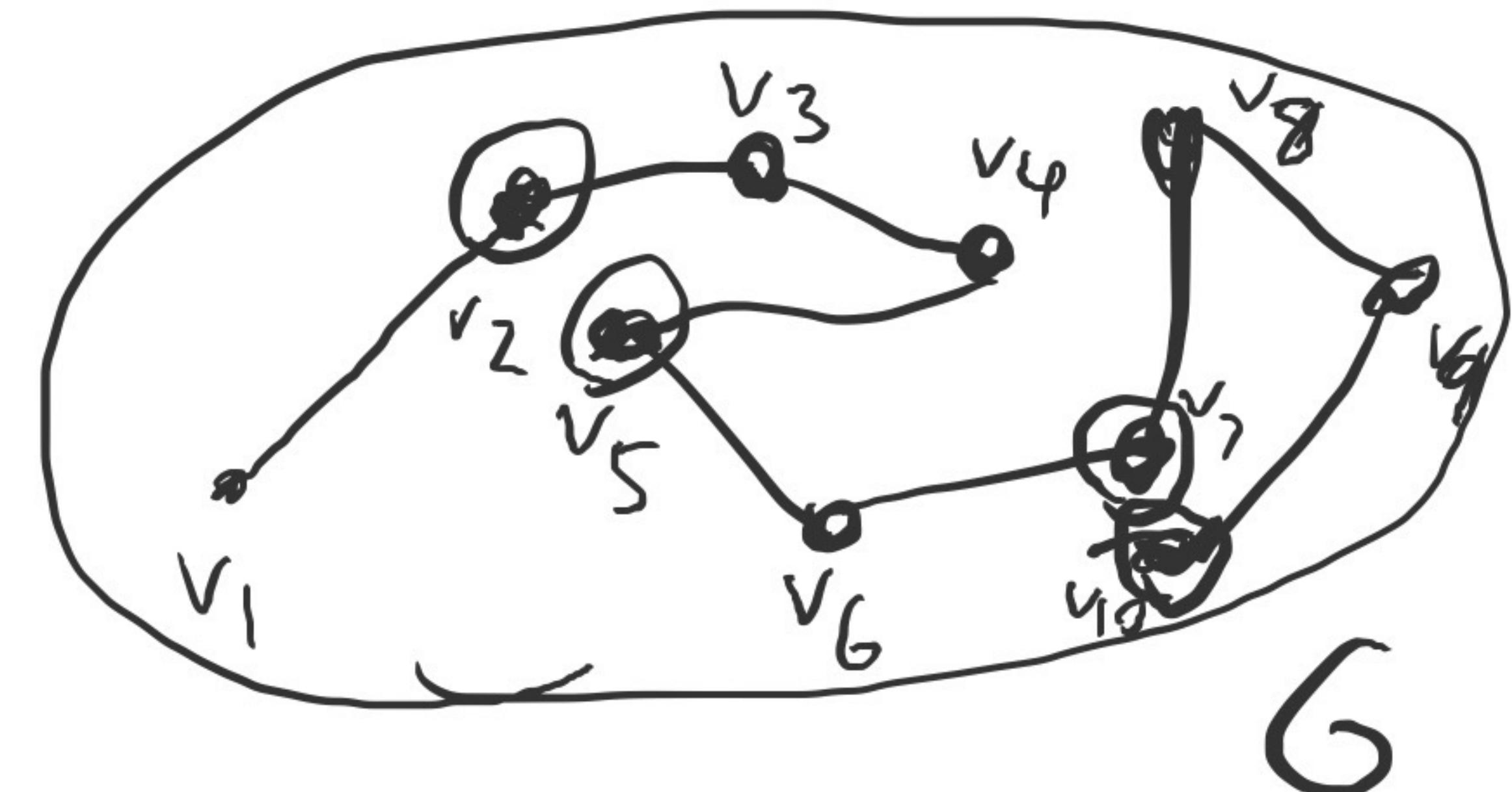
Consider $E_{G_{\text{val}}}(X_S)$

$$= \underset{\substack{\text{RW} \\ G_{\text{val}}, t}}{E}[X_S(x)]$$

Define P to be the diagonal matrix where $P_{ii} = \text{val}(i)$.

$$E_{G_{\text{val}}}(X_S) = \vec{1}^T (P G^{\Delta_{t-1}}) \dots (P G^{\Delta_1}) (P G^{\Delta_0}) \frac{\vec{1}}{n}$$

Idea: We'll analyze this by breaking G^{Δ_j} up into $G^{\Delta_j} = J + E^{\Delta_j}$ (recall $G = J + E$. when $\|E\| \leq \lambda$)



$$S = \{2, 5, 7, 10\}$$

$$\Delta_0 = 1 \quad \Delta_1 = 3 \quad \Delta_2 = 2 \quad \Delta_3 = 3$$

$$\sum_{\text{Gral}}(x_s) = \vec{1}^T (PG^{\Delta_{t-1}}) \dots (PG^{\Delta_2})(PG^{\Delta_1}) P \frac{\vec{1}}{n}$$

Idea: Consider which E^{Δ_j} terms we have.

Proposition: $\vec{1}^T P \vec{1} = 0$. Since $J = \frac{1}{n} \vec{1} \cdot \vec{1}^T$,

$$J P \vec{1} = 0 \quad \vec{1}^T P J = 0 \quad J P J = 0.$$

Claim: If we don't have E^{Δ_1} or $E^{\Delta_{t-1}}$, this part is 0.

If we don't have E^{Δ_j} and $E^{\Delta_{j+1}}$ for some j , this part is 0.

Definition: Define F_k to be the set of $I \subseteq [k-1]$ such that

$$1. I \in F_k, k-1 \in I,$$

$$2. \forall j \in [k-2] \quad j \in I \text{ or } j+1 \in I$$

Proposition 4.2: $|\sum_{\text{Gral}}(x_s)| \leq \sum_{I \in F_k} \lambda^{\sum_{j \in I} \Delta_j(s)}$

Proposition 4.2: $|\sum_{G, \text{val}}(x_S)| \leq \sum_{I \in \mathcal{F}_k} \lambda \sum_{j \in I} \Delta_j(S)$

Examples:	$ S =2$	$\leq \lambda^{\Delta_1}$	$F_2 = \{\{1\}\}$
	$ S =3$	$\leq \lambda^{\Delta_1 + \Delta_2}$	$F_3 = \{\{1, 2\}, \{1, 3\}, \{2, 3\}\}$
	$ S =4$	$\leq \lambda^{\Delta_1 + \Delta_3} + \lambda^{\Delta_1 + \Delta_2 + \Delta_3}$	$F_4 = \{\{1, 3\}, \{1, 2, 3\}\}$

$$|S|=6 \quad F_6 = \{\{1, 3, 5\}, \{1, 2, 4, 5\}, \{1, 2, 3, 5\}, \{1, 3, 4, 5\}, \{1, 2, 3, 4, 5\}\}$$

Proposition: $\forall I \in \mathcal{F}_k, |I| \geq \lceil \frac{k}{2} \rceil$

Proposition: If $|S|=0$ or 1 $\sum_{G, \text{val}}(x_S) = 0$

Corollary: $\forall S: |S| \geq 2, |\sum_{G, \text{val}}(x_S)| \leq 2^k \cdot \lambda^{\lceil \frac{k}{2} \rceil} \leq (4\lambda)^{\frac{k}{2}}$

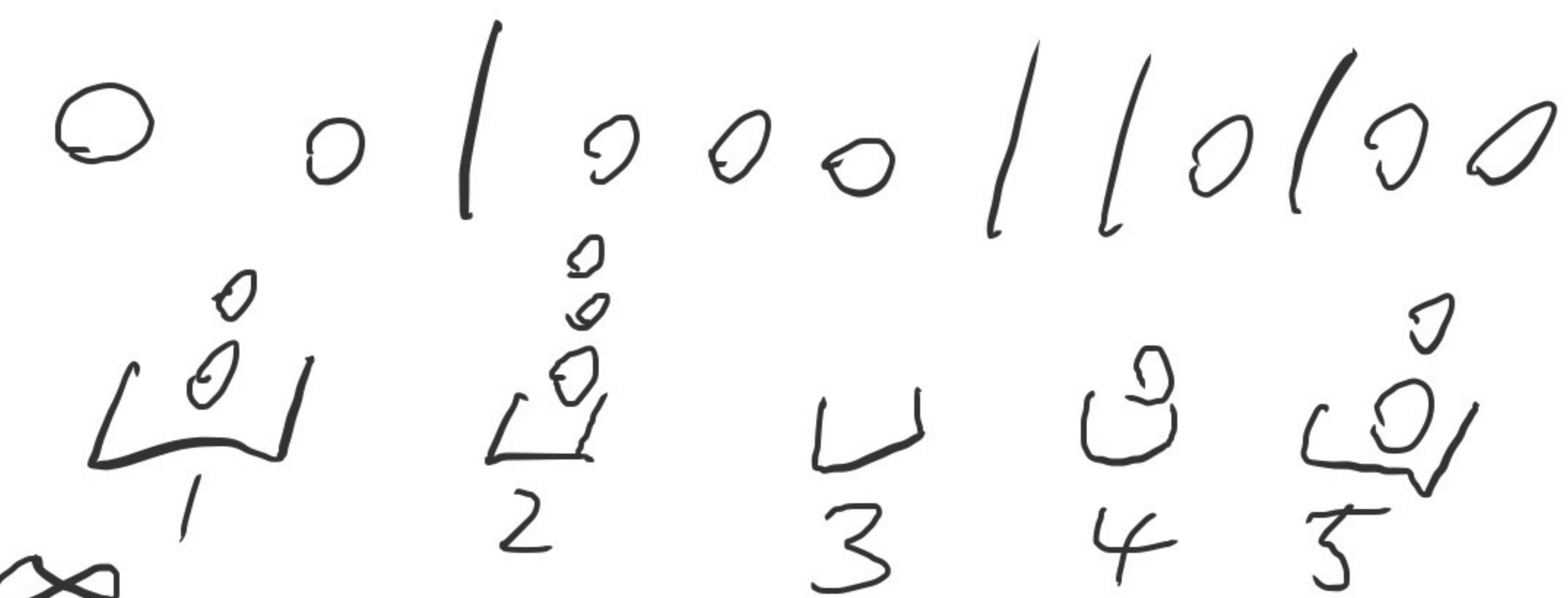
Corollary: If $f \in L_1^+(b)$ then

$$\sum_{f \in \mathcal{L}_1^+(b)} |\sum_{G, \text{val}}(x_S)| \leq \sum_{k=2}^{\infty} (4\lambda)^{\frac{k}{2}} \cdot b^k \text{ which is } O(\lambda b^2).$$

Review: Unlabeled Balls in Labeled Bins

Proposition: There are $\binom{n+k-1}{k-1}$ ways to put n unlabeled balls in k labeled bins.

Trick: Imagine putting $r-1$ dividing lines among the n balls



Claim 3.1: $\sum_{j=r}^{\infty} \binom{j}{r} \lambda^j = \frac{\lambda^r}{(1-\lambda)^{r+1}}$ if $|\lambda| < 1$

Proof: $\frac{1}{1-\lambda} = 1 + \lambda + \lambda^2 + \lambda^3 + \dots$

What is the coefficient of λ^j in $\frac{\lambda^r}{(1-\lambda)^{r+1}}$? Claim: This is like putting $j-r$ balls in $r+1$ labeled bins. There are $\binom{j-r+r+1-1}{r+1-1} = \binom{j}{r}$ ways to do this.

Definition: Define $\beta_k = \sum_{\substack{S \subseteq [+] \\ |S|=k}} \sum_{\text{6, val}} (\chi_S)$

Lemma 4.4: $|\beta_k| \leq 2^k \binom{k+1}{\lfloor \frac{k}{2} \rfloor} \left(\frac{\lambda}{1-\lambda}\right)^{\lceil \frac{k}{2} \rceil}$

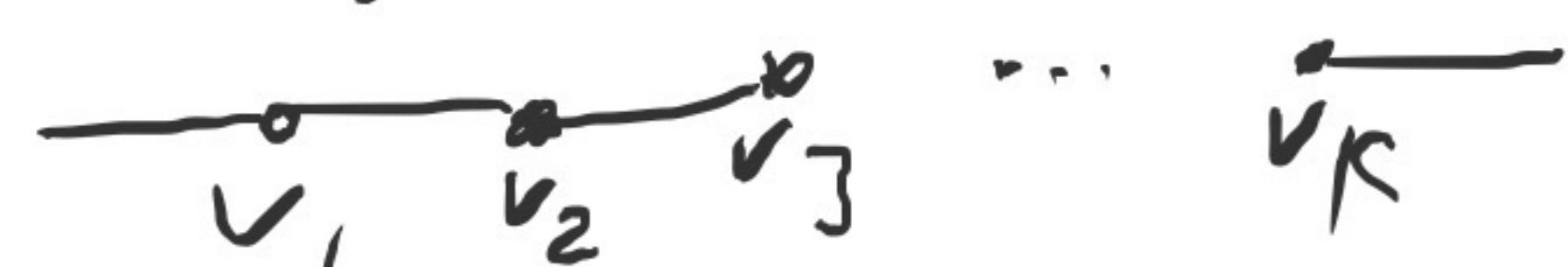
Claim 4.5: $|\beta_k| \leq 2^k \sum_{m=\lceil \frac{k}{2} \rceil}^{k+1} \binom{m+1}{\lceil \frac{k}{2} \rceil - 1} \binom{k-m}{k-\lceil \frac{k}{2} \rceil} \lambda^m$

By symmetry,
this only depends
on $|I|$ and
is decreasing
in I .

Proof: $|\beta_k| \leq \sum_{\substack{S \subseteq [+] \\ |S|=k}} \sum_{I \in F_k} \lambda^{\sum_{j \in I} \Delta_j(S)} = \sum_{I \in F_k} \sum_{\substack{S \subseteq [+] \\ |S|=k}} \lambda^{\sum_{j \in I} \Delta_j(S)}$
 Choosing a particular $I^* \in F_k$ where $|I^*| = \lceil \frac{k}{2} \rceil$, $|\beta_k| \leq 2^k \sum_{\substack{S \subseteq [+] \\ |S|=k}} \lambda^{\sum_{j \in I^*} \Delta_j(S)}$

Q: What is the coefficient of λ^m in this expression?

We need to split m factors of λ among (I^*) Δ_j where each Δ_j is at least 1. This is = putting $m - |I^*|$ balls in $|I^*|$ labeled bins.
 We also need to divide the remaining $+ - 1 - m$ steps of the path among $k+1 - |I^*|$ other Δ . The middle Δ must be at least 1.



This is = putting $+ - 1 - m - (k+1 - |I^*| - 2)$ unlabeled balls

This gives a coefficient of $\binom{m-1}{|I^*|-1} \binom{t-m}{k-|I^*|}$ for λ^m

$$\text{Claim 4.5: } |\beta_k| \leq 2^k \sum_{n_0=\lceil \frac{k}{2} \rceil}^{+\lfloor \frac{k}{2} \rfloor} \binom{m-1}{\lceil \frac{k}{2} \rceil - 1} \binom{t-m}{k - \lceil \frac{k}{2} \rceil} \lambda^m$$

$$\leq 2^k \binom{+1}{k - \lceil \frac{k}{2} \rceil} \lambda \sum_{n=\lceil \frac{k}{2} \rceil}^{\infty} \binom{m-1}{\lceil \frac{k}{2} \rceil - 1} \lambda^{m-1}$$

$$\text{Claim 3.1: } \sum_{j=r}^{\infty} \binom{j}{r} \lambda^j = \frac{\lambda^r}{(1-\lambda)^{r+1}}$$

$\therefore \frac{\lambda^{\lceil \frac{k}{2} \rceil - 1}}{(1-\lambda)^{\lceil \frac{k}{2} \rceil}}$

$$r = \lceil \frac{k}{2} \rceil - 1$$

Lemma 4.4

$$|\beta_k| \leq 2^k \binom{+1}{k - \lceil \frac{k}{2} \rceil} \cdot \frac{\lambda^{\lceil \frac{k}{2} \rceil}}{(1-\lambda)^{\lceil \frac{k}{2} \rceil}}$$

$$\text{Lemma 4.4: } |\beta_k| \leq 2^k \binom{+1}{\lfloor \frac{k}{2} \rfloor} \left(\frac{\lambda}{1-\lambda} \right)^{\lceil \frac{k}{2} \rceil}$$

For a symmetric function f , defining \hat{f}_k to be \hat{f}_S for some S of size k

$$\mathbb{E}_\lambda(f) \leq \sum_{k=2}^{+} |\beta_k| \cdot \hat{f}_k$$

How large can \hat{f}_k be?

There are $\binom{+}{k}$ subsets S of size k .

Since $\sum_{S: |S|=k} \hat{f}_S^2 = \binom{+}{k} \hat{f}_k^2 \leq 1$ $\hat{f}_k \leq \frac{1}{\sqrt{\binom{+}{k}}}$

$$\mathbb{E}_\lambda(f) \leq \sum_{k=2}^{+} 2^k \frac{\binom{+1}{\lfloor \frac{k}{2} \rfloor}}{\sqrt{\binom{+}{k}}} \left(\frac{\lambda}{1-\lambda} \right)^{\lceil \frac{k}{2} \rceil}$$

$O(\frac{\lambda^{\lceil \frac{k}{2} \rceil}}{\sqrt{+}})$ k is odd

$O(\lambda^{\frac{k}{2}})$ k is even