

# Chapter 2: Operating-System Structures



# Chapter 2:Outline

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- ❑ Operating System Services
- ❑ User and Operating System Interface
- ❑ System Calls
- ❑ System Services
- ❑ Linkers and Loaders
- ❑ Why Applications are Operating System Specific?
- ❑ Operating-System Design and Implementation
- ❑ Operating System Structure
- ❑ Building and Booting an Operating System
- ❑ Operating System Debugging



# Objectives

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- ❑ Identify **services** provided by an operating system
- ❑ Illustrate how **system calls** are used to provide operating system services
- ❑ Compare and contrast **monolithic**, **layered**, **microkernel**, **modular**, and **hybrid** strategies for designing operating systems
- ❑ Illustrate the process for **booting an operating system**
- ❑ Apply tools for **monitoring operating system performance**
- ❑ Design and implement **kernel modules** for interacting with a Linux kernel



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# OPERATING SYSTEM SERVICES



# Operating System Services

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- ❑ Operating systems provide an *environment for execution* of programs and *services* to programs and users
- ❑ A set of operating-system services provides *functions* that are helpful to the user:
  - ❑ *User interface* - Almost all operating systems have a user interface (**UI**).
    - ▶ Varies between *Command-Line Interface (CLI)*, *Graphical User Interface (GUI)*, *Touch-screen*
  - ❑ *Program execution* - The system must be able to *load* a program into memory, to *run* that program, and *end* execution, either normally or abnormally (indicating error)
  - ❑ *I/O operations* - A running program may require I/O, which may involve a file or an I/O device



# Operating System Services (Cont.)

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- **File-system manipulation** - The file system is of particular interest. Programs need to read and write files and directories, create and delete them, search them, list file information, permission management.
- **Communications** – Processes may exchange information, on the same computer or between computers over a network
  - ▶ Communications may be via *shared memory* or through *message passing* (packets moved by the OS)
- **Error detection** – OS needs to be constantly aware of possible errors
  - ▶ May occur in the CPU and memory, hardware, in I/O devices, in user program
  - ▶ For each type of error, OS should take the *appropriate action* to ensure correct and consistent computing
  - ▶ *Debugging facilities* can greatly enhance the user's and programmer's abilities to efficiently use the system



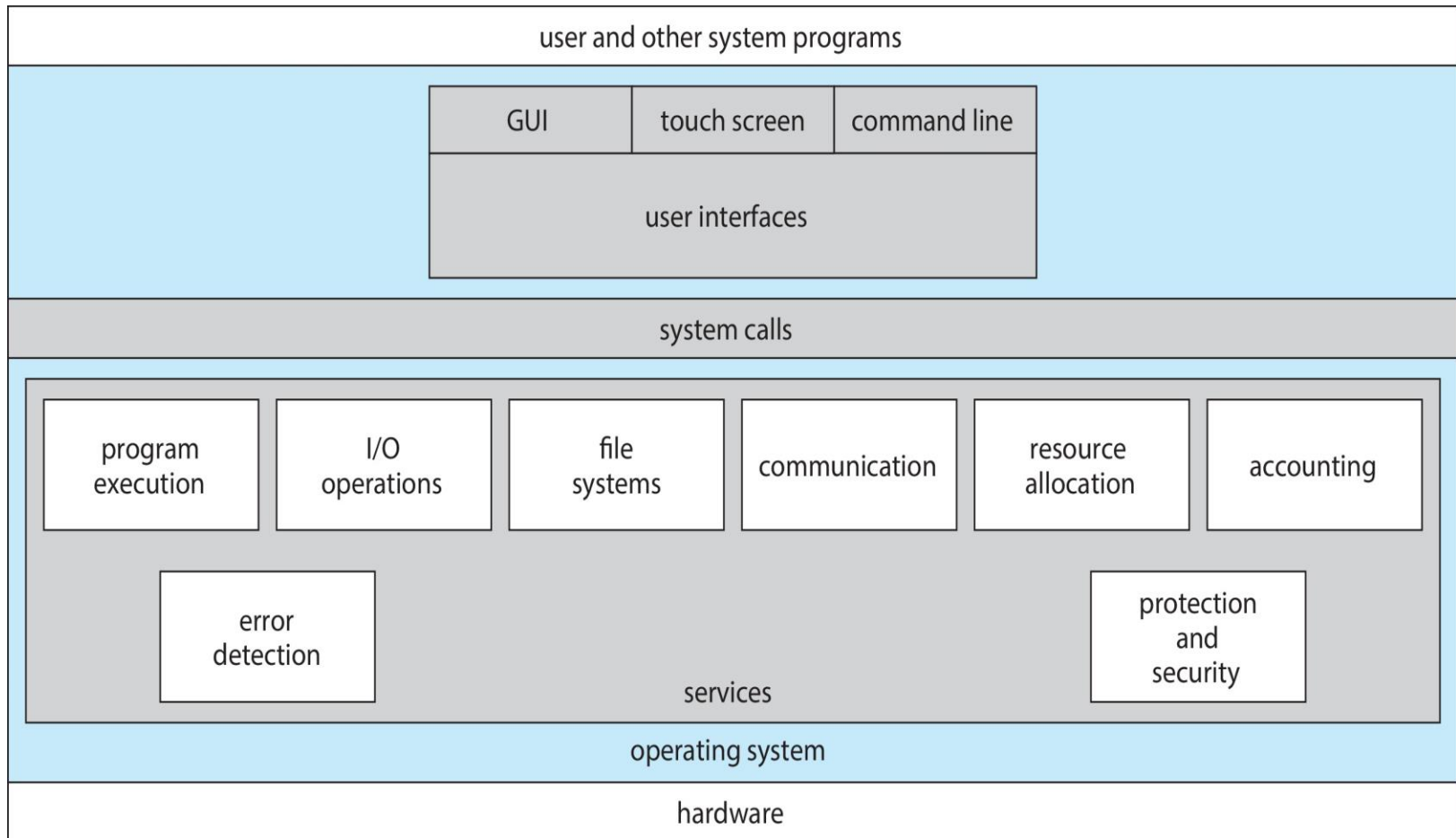
# Operating System Services (Cont.)

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- Another set of *OS functions* exists for ensuring the efficient operation of the system itself via *resource sharing*
  - *Resource allocation* - When multiple users or multiple jobs running concurrently, resources must be allocated to each of them
    - ▶ Many *types of resources* - CPU cycles, main memory, file storage, I/O devices.
  - *Logging* - To *keep track* of which users use how much and what kinds of computer resources
  - *Protection and security* - The owners of information stored in a multiuser or networked computer system may want to control use of that information, concurrent processes should not interfere with each other
    - ▶ *Protection* involves ensuring that all access to system resources is controlled
    - ▶ *Security* of the system from outsiders requires user authentication, extends to defending external I/O devices from invalid access attempts



# A View of Operating System Services





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# USER OPERATING SYSTEM INTERFACE



# User Operating System Interface - CLI

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- CLI or *command interpreter* allows direct command entry
  - Primarily *fetches a command from user* and *executes it*
  - Sometimes implemented in kernel, sometimes by *system programs*
  - Sometimes multiple flavors implemented – *shells*
  - Sometimes commands built-in, sometimes just names of programs
    - ▶ If the latter, adding new features doesn't require shell modification



# Bourne Shell Command Interpreter

```
1. root@r6181-d5-us01:~ (ssh)
X root@r6181-d5-u...  1 X ssh  2 X root@r6181-d5-us01...  3
Last login: Thu Jul 14 08:47:01 on ttys002
iMacPro:~ pbg$ ssh root@r6181-d5-us01
root@r6181-d5-us01's password:
Last login: Thu Jul 14 06:01:11 2016 from 172.16.16.162
[root@r6181-d5-us01 ~]# uptime
 06:57:48 up 16 days, 10:52,  3 users,  load average: 129.52, 80.33, 56.55
[root@r6181-d5-us01 ~]# df -kh
Filesystem      Size  Used Avail Use% Mounted on
/dev/mapper/vg_ks-lv_root
                  50G   19G   28G  41% /
tmpfs            127G  520K  127G   1% /dev/shm
/dev/sda1        477M   71M  381M  16% /boot
/dev/dssd0000    1.0T  480G  545G  47% /dssd_xfs
tcp://192.168.150.1:3334/orangefs
                  12T   5.7T   6.4T  47% /mnt/orangefs
/dev/gpfs-test   23T   1.1T   22T   5% /mnt/gpfs
[root@r6181-d5-us01 ~]#
[root@r6181-d5-us01 ~]# ps aux | sort -nrk 3,3 | head -n 5
root      97653 11.2  6.6 42665344 17520636 ?    S<Ll  Jul13 166:23 /usr/lpp/mmfs/bin/mmfstd
root      69849  6.6  0.0      0      0 ?    S      Jul12 181:54 [vpthread-1-1]
root      69850  6.4  0.0      0      0 ?    S      Jul12 177:42 [vpthread-1-2]
root       3829  3.0  0.0      0      0 ?    S      Jun27 730:04 [rp_thread 7:0]
root       3826  3.0  0.0      0      0 ?    S      Jun27 728:08 [rp_thread 6:0]
[root@r6181-d5-us01 ~]# ls -l /usr/lpp/mmfs/bin/mmfstd
-r-x----- 1 root root 20667161 Jun  3  2015 /usr/lpp/mmfs/bin/mmfstd
[root@r6181-d5-us01 ~]#
```



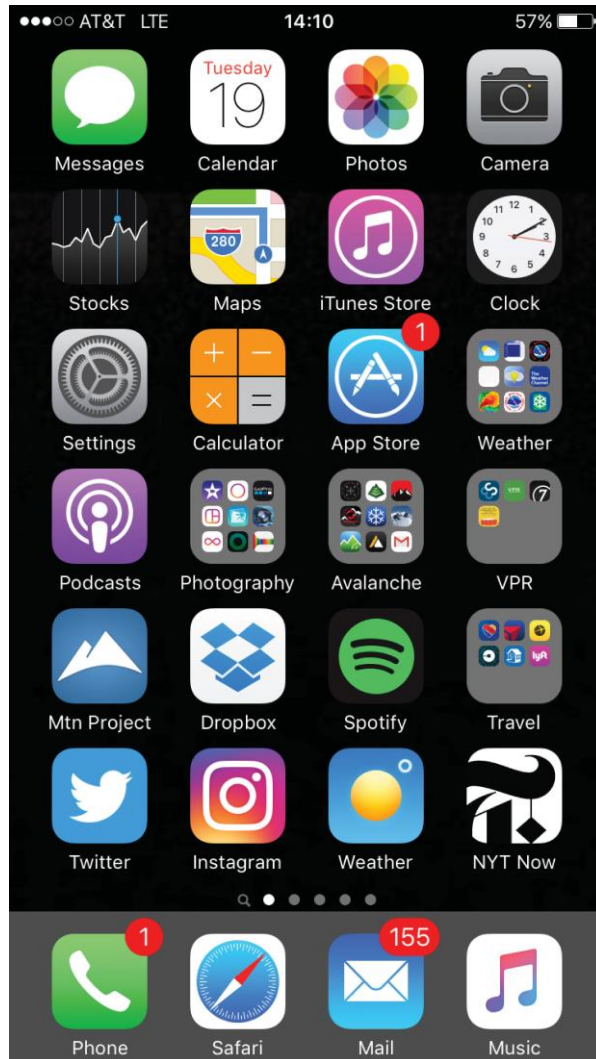
# User Operating System Interface - GUI

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- User-friendly *desktop metaphor interface*
  - Usually mouse, keyboard, and monitor
  - Icons represent files, programs, actions, etc.
  - Various mouse buttons over objects in the interface cause various actions providing information, options, execute function, open directory
  - Invented at Xerox PARC
- Many systems now include *both CLI and GUI interfaces*
  - **Microsoft Windows** is GUI with CLI “command” shell
  - **Apple Mac OS X** is “Aqua” GUI interface with UNIX kernel underneath and shells available
  - **Unix** and **Linux** have CLI with optional GUI interfaces (e.g., CDE, KDE, GNOME)



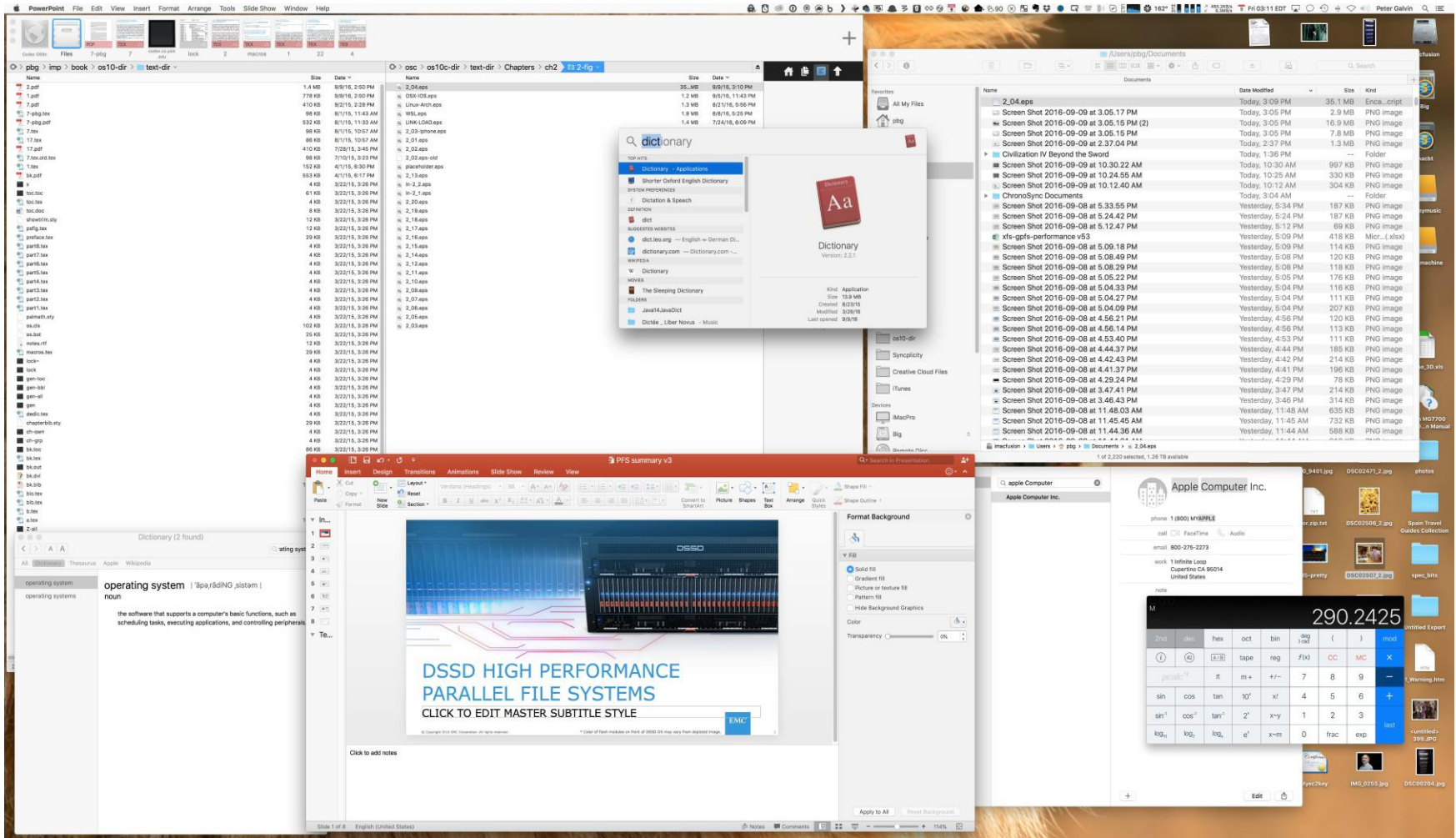
# Touchscreen Interfaces



- Touchscreen devices require new interfaces
  - Mouse not possible or not desired
  - Actions and selection based on *gestures*
  - *Virtual keyboard* for text entry
- Voice commands



# The Mac OS X GUI



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# SYSTEM CALLS





# System Calls

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- *Programming interface* to the *services* provided by the OS
- Typically written in a high-level language (e.g., C or C++)
- Mostly accessed by programs via a high-level **Application Programming Interface (API)** rather than direct use of system call
- Three most common APIs are **Win32 API** for Windows, **POSIX API** for POSIX-based systems (including virtually all versions of UNIX, Linux, and Mac OS X), and **Java API** for the Java Virtual Machine (JVM)

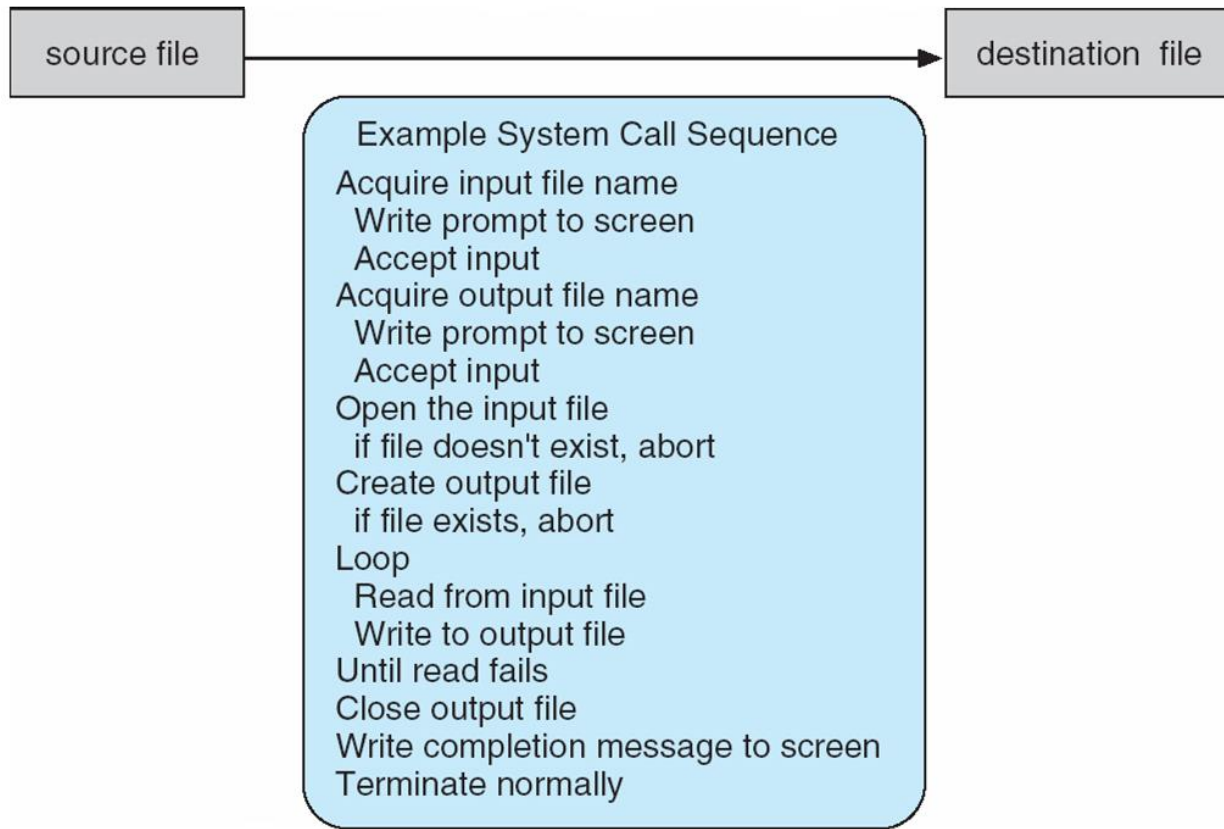
(Note that the system-call names used throughout this text are generic)





# Example of System Calls

- System call sequence to *copy the contents of one file to another file*



# Example of Standard API

## EXAMPLE OF STANDARD API

As an example of a standard API, consider the `read()` function that is available in UNIX and Linux systems. The API for this function is obtained from the man page by invoking the command

```
man read
```

on the command line. A description of this API appears below:

```
#include <unistd.h>
```

```
ssize_t read(int fd, void *buf, size_t count)
```

return  
value

function  
name

parameters

A program that uses the `read()` function must include the `unistd.h` header file, as this file defines the `ssize_t` and `size_t` data types (among other things). The parameters passed to `read()` are as follows:

- `int fd`—the file descriptor to be read
- `void *buf`—a buffer into which the data will be read
- `size_t count`—the maximum number of bytes to be read into the buffer

On a successful read, the number of bytes read is returned. A return value of 0 indicates end of file. If an error occurs, `read()` returns `-1`.



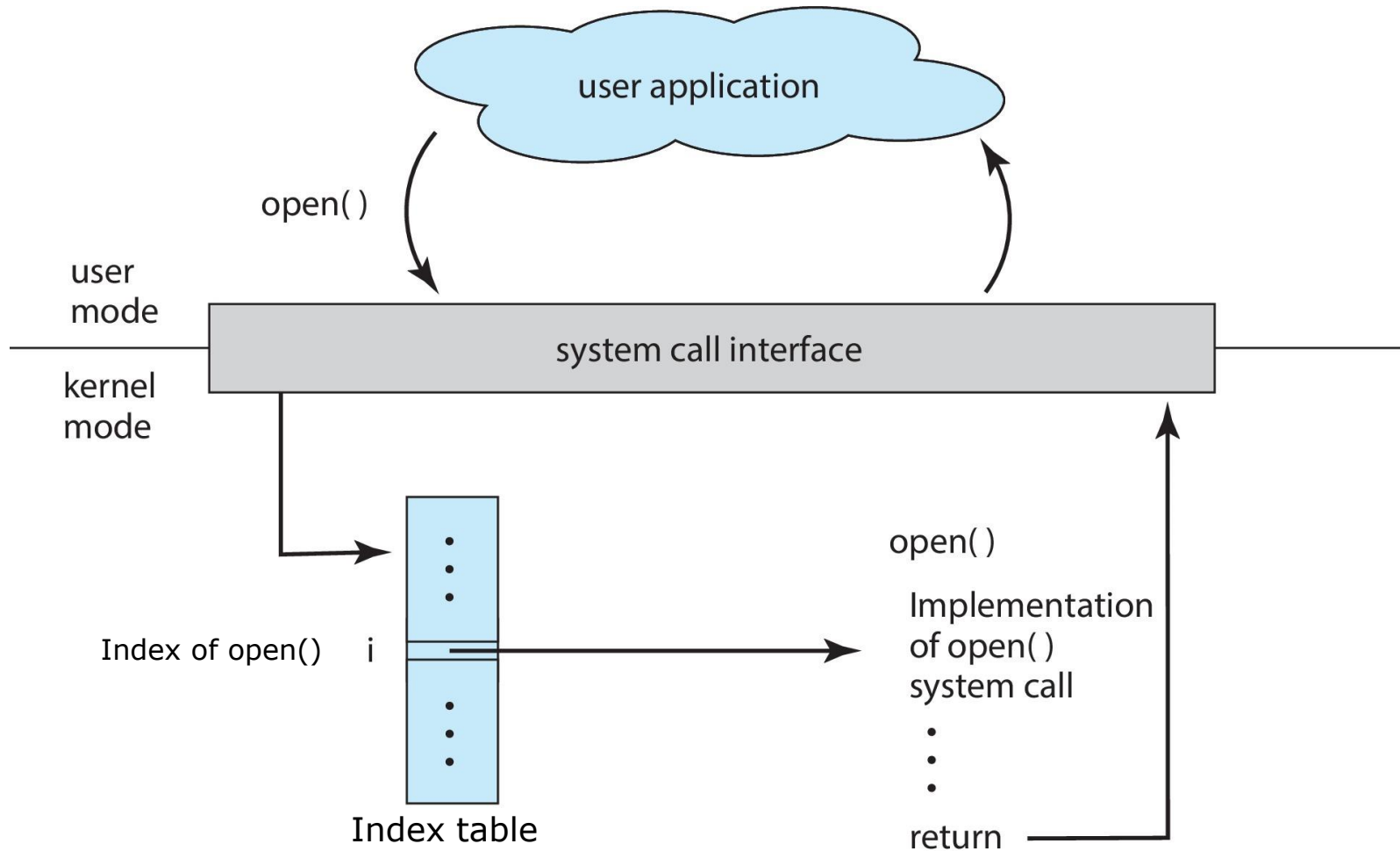
# System Call Implementation

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- Typically, *a number associated with each system call*
  - **System-call interface** maintains a *table* indexed according to these numbers
- The *system call interface invokes* the intended system call in OS kernel and *returns* status of the system call and any return *values*
- *The caller need to know nothing about how the system call is implemented*
  - Just needs to *obey API* and understand what OS will do as a result call
  - Most details of OS interface are *hidden* from programmers by API
    - ▶ Managed by *run-time support library* (set of functions built into libraries included with compiler)



# API – System Call – OS Relationship



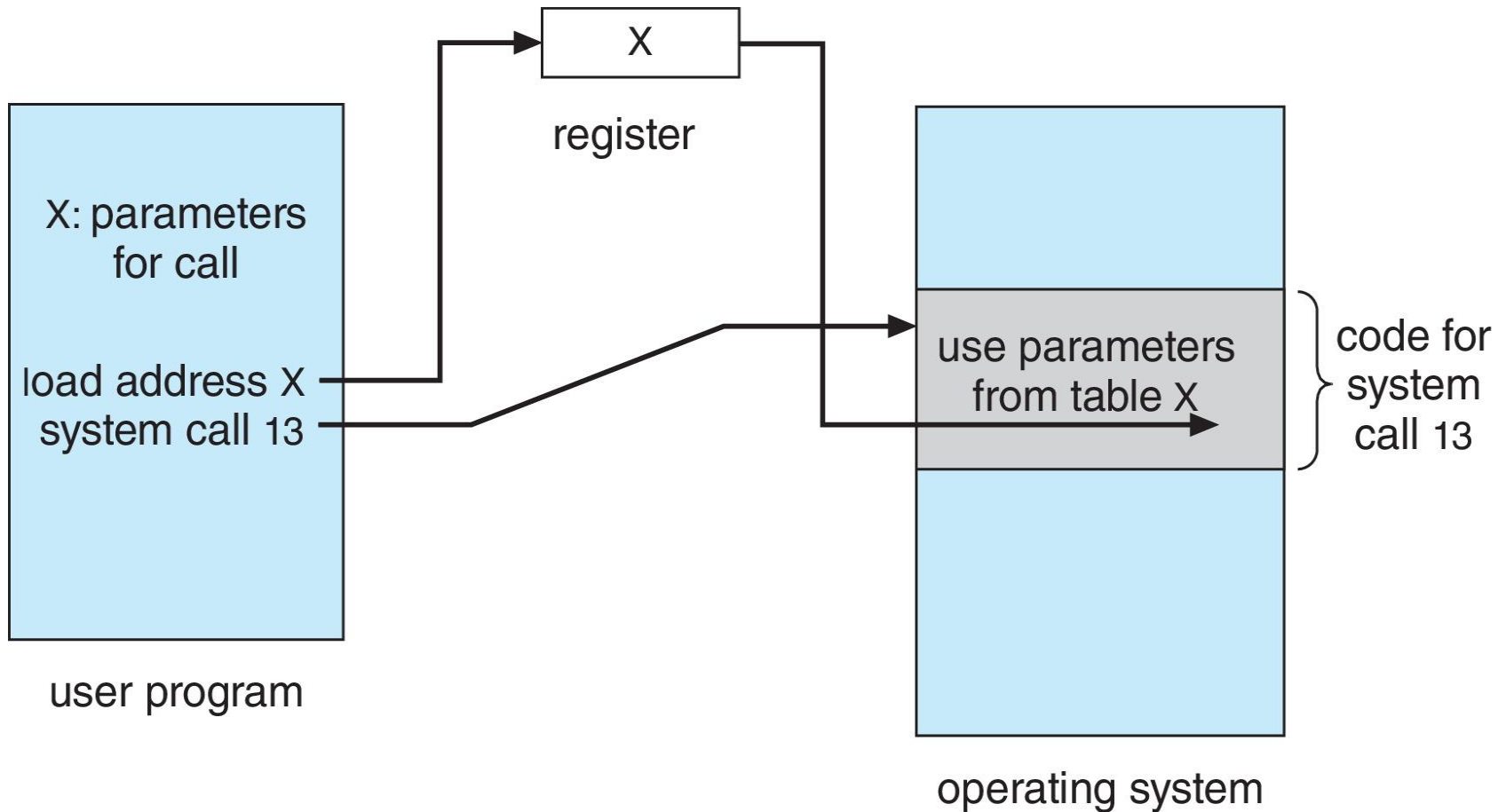
# System Call Parameter Passing

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- Often, *more information* is required than simply identity (i.e., name) of desired system call
  - Exact type and amount of information vary according to OS and call
- Three general methods used to pass *parameters* to the OS
  - **Simplest**: pass the parameters in registers
    - ▶ In some cases, may be more parameters than registers
  - Parameters stored in a *block*, or *table*, in *memory*, and **address of block** passed as a parameter in a register
    - ▶ This approach taken by *Linux* and Solaris
  - Parameters placed, or *pushed*, onto the **stack** (see *Memory structure*) by the *program* and *popped* off the stack by the *operating system*
    - ▶ Block and stack methods *do not limit* the number or length of parameters being passed



# Parameter Passing via Table



# Types of System Calls

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## □ Process control

- create process, terminate process
- end, abort process execution
- load, execute
- get process attributes, set process attributes
- wait for time, wait event, signal event
- allocate and free memory
- dump memory if error
- *debugger* for determining *bugs*, *single step* execution
- *Locks* for managing access to shared data between processes



# Types of System Calls (cont.)

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## □ File management

- create file, delete file
- open, close file
- read, write, reposition
- get and set file attributes

## □ Device management

- request device, release device
- read, write, reposition
- get device attributes, set device attributes
- logically attach or detach devices





# Types of System Calls (Cont.)

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## □ Information maintenance

- get time or date, set time or date
- get system data, set system data
- get and set process, file, or device attributes

## □ Communications

- create, delete communication connection
- *Message passing model*: send, receive messages to host name or process name
- *Shared-memory model*: create and gain access to memory regions
- transfer status information
- attach and detach remote devices



# Types of System Calls (Cont.)

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## □ Protection

- Control access to resources
- Get and set permissions
- Allow and deny user access



# Examples of Windows and Unix System Calls

## EXAMPLES OF WINDOWS AND UNIX SYSTEM CALLS

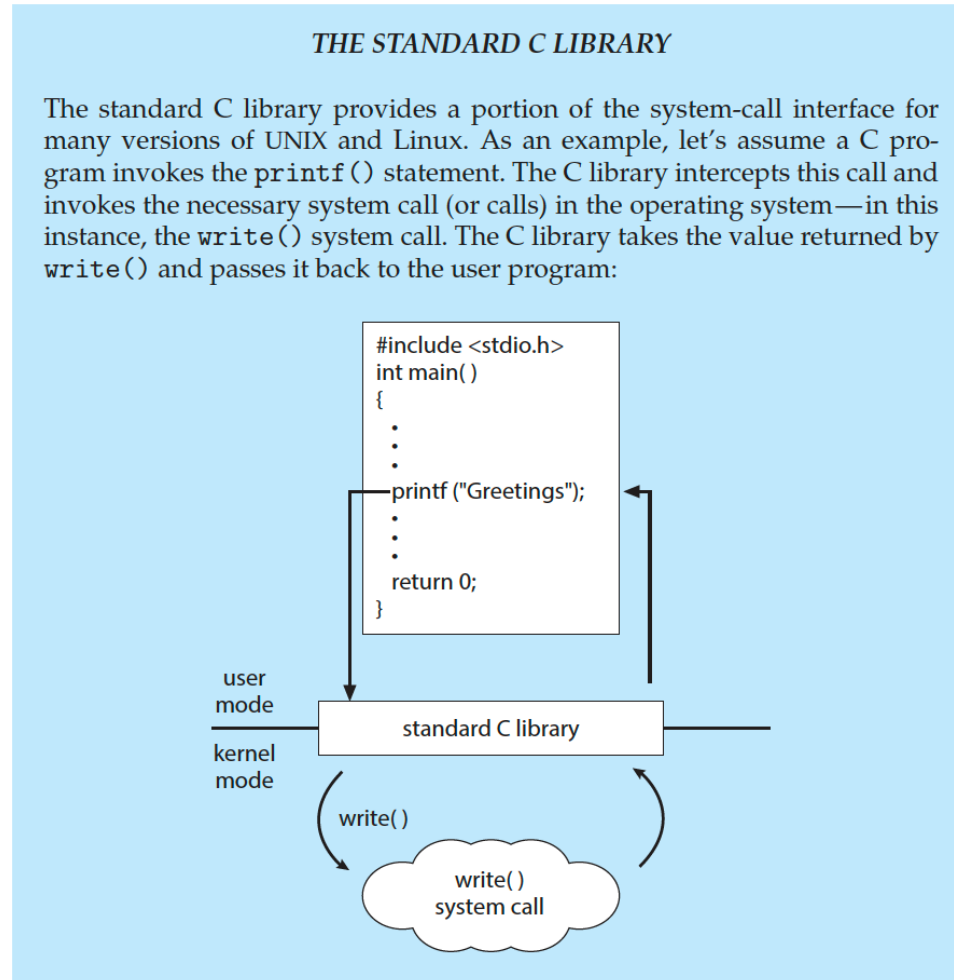
The following illustrates various equivalent system calls for Windows and UNIX operating systems.

	Windows	Unix
<b>Process control</b>	CreateProcess() ExitProcess() WaitForSingleObject()	fork() exit() wait()
<b>File management</b>	CreateFile() ReadFile() WriteFile() CloseHandle()	open() read() write() close()
<b>Device management</b>	SetConsoleMode() ReadConsole() WriteConsole()	ioctl() read() write()
<b>Information maintenance</b>	GetCurrentProcessID() SetTimer() Sleep()	getpid() alarm() sleep()
<b>Communications</b>	CreatePipe() CreateFileMapping() MapViewOfFile()	pipe() shm_open() mmap()
<b>Protection</b>	SetFileSecurity() InitializeSecurityDescriptor() SetSecurityDescriptorGroup()	chmod() umask() chown()



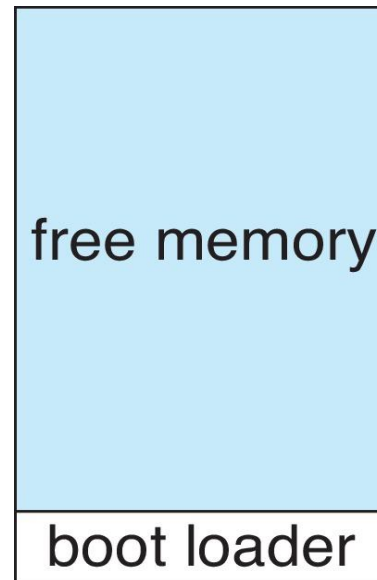
# Standard C Library Example

- C program invoking `printf()` library function, which calls `write()` system call



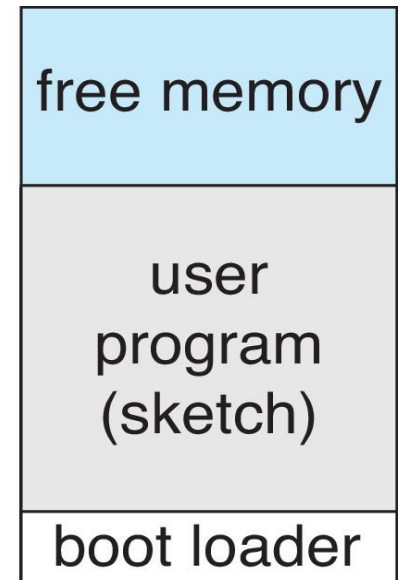
# Example: Arduino

- ❑ *Single-tasking*
- ❑ *No operating system*
- ❑ *Single memory space*
- ❑ Programs (sketch) loaded via USB into flash memory
- ❑ Boot loader loads program
- ❑ Program exit → shell reloaded



(a)

At system startup

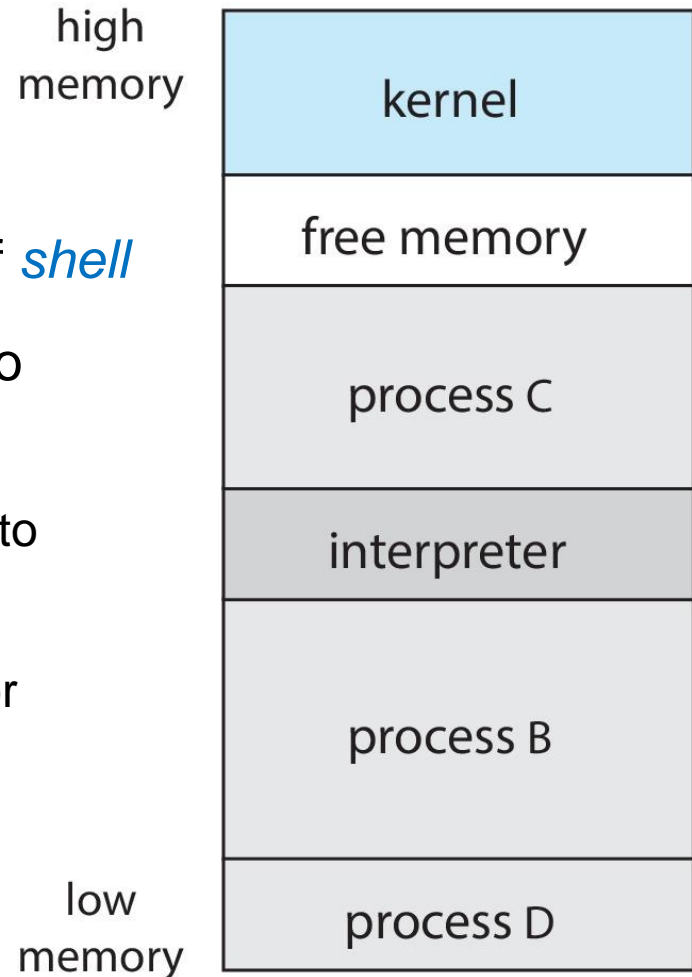


(b)

running a program

# Example: FreeBSD

- Unix variant
- *Multitasking*
- User login -> invoke user's choice of *shell*
- Shell executes `fork()` system call to create *process*
  - Executes `exec()` to load program into process
  - Shell waits for process to terminate or continues with user commands
- Process exits with:
  - *code = 0 – no error*
  - *code > 0 – error code*



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# SYSTEM SERVICES



# System Services

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- ❑ System programs provide a convenient environment for *program development* and *execution*. They can be divided into:
  - ❑ File manipulation
  - ❑ Status information sometimes stored in a file
  - ❑ Programming language support
  - ❑ Program loading and execution
  - ❑ Communications
  - ❑ Background services
  - ❑ Application programs
- ❑ Most *users' view* of the operation system is defined by *system programs*, not the actual *system calls*





# System Services (Cont.)

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- Some of *system services* are simply *user interfaces to system calls*; others are considerably more complex
- **File management** - Create, delete, copy, rename, print, dump, list, and generally manipulate files and directories
- **Status information**
  - Some ask the system for *information* (e.g., date, time, amount of available memory, disk space, number of users)
  - Others provide detailed *performance*, *logging*, and *debugging* information
  - Typically, these programs *format* and *print* the output to the terminal or other output devices
  - Some systems implement a *registry* - used to store and retrieve *configuration information*



# System Services (Cont.)

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## □ File modification

- Text *editors* to create and modify files
- Special *commands* to search contents of files or perform transformations of the text

## □ Programming-language support - *Compilers, assemblers, debuggers* and *interpreters* are sometimes provided

## □ Program loading and execution - Absolute *loaders, relocatable loaders, linkage editors*, and *overlay-loaders, debugging systems* for higher-level and machine language

## □ Communications - Provide the mechanism for creating *virtual connections* among processes, users, and computer systems

- Allow users to send messages to one another's screens, browse web pages, send electronic-mail messages, log in remotely, transfer files from one machine to another



# System Services (Cont.)

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## □ Background Services

- Launch *at boot time*
  - ▶ Some for the system startup, then terminate
  - ▶ Some from system boot to shutdown
- Provide *facilities* like disk checking, process scheduling, error logging, printing
- Run in *user context* not kernel context
- Known as *services*, *subsystems*, *daemons*

## □ Application programs

- Don't pertain to system, *run by users*
- *Not* typically considered *part of OS*
- Launched by command line, mouse click, finger poke



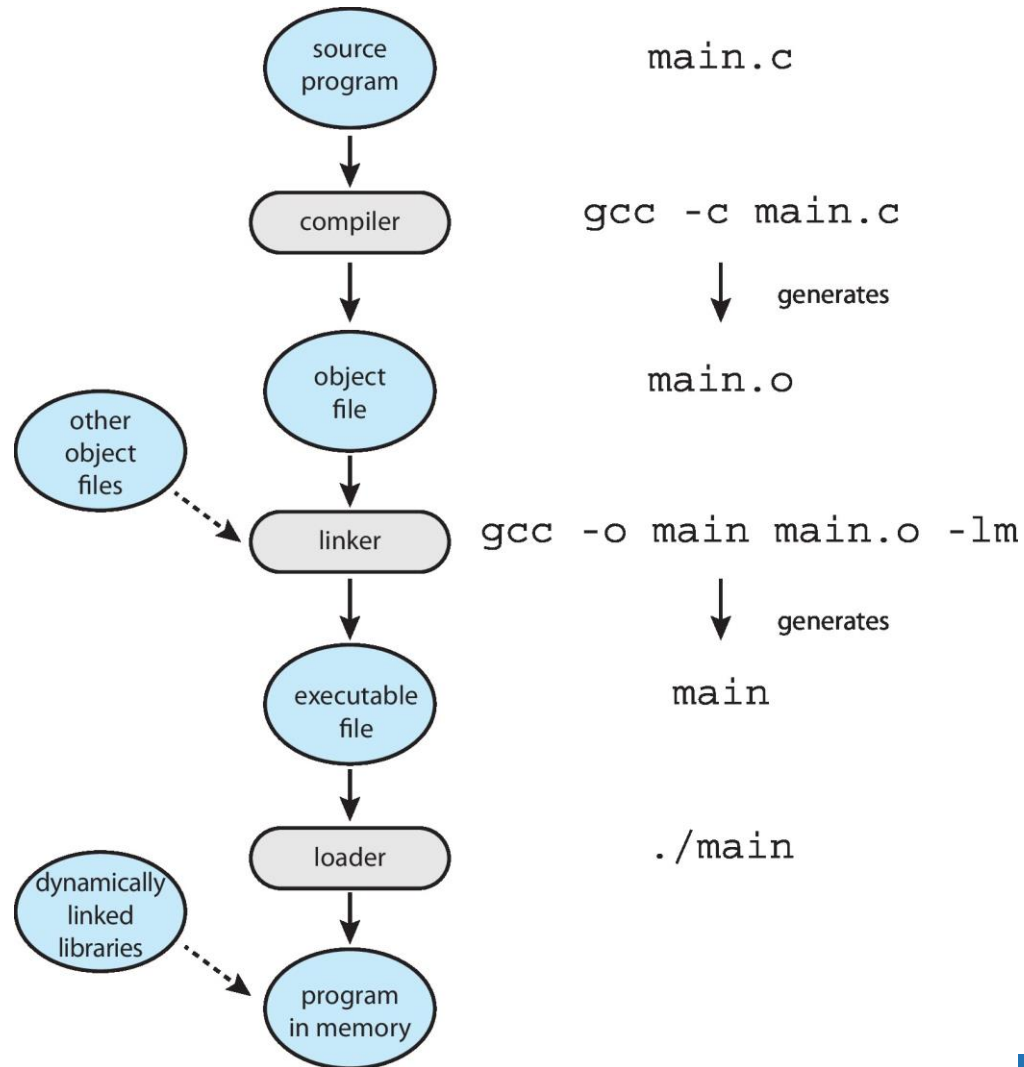
# Linkers and Loaders

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- ❑ *Source code* compiled into *object files* (designed to be loaded into any physical memory location) – *relocatable object file*
- ❑ **Linker** combines these files (also, brings in libraries) into single *binary executable file*
- ❑ *Program* resides on secondary storage as *binary executable* and must be brought into memory by **loader** to be executed
  - ❑ *Relocation* assigns final addresses to program parts and adjusts code and data in program to match those addresses
- ❑ *Modern general-purpose systems don't link libraries into executables*
  - ❑ Rather, *dynamically linked libraries* (in Windows, *DLLs*) are *loaded* as needed, *shared* by all that use the same version of that same library
- ❑ Object, executable files have *standard formats*, so operating system knows how to *load* and *start* them



# The Role of the Linker and Loader



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# OPERATING SYSTEM DESIGN AND IMPLEMENTATION



# Why Applications are Operating System Specific

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- ❑ *Apps compiled on one system usually not executable on other operating systems*
- ❑ Each operating system provides *its own unique system calls, own file formats*, etc.
- ❑ Apps can be *multi-operating system*
  - ❑ Written in *interpreted* language like Python, Ruby, and *interpreter* available on multiple operating systems
  - ❑ App written in language that includes a VM containing the running app
  - ❑ Use standard language (like C), *compile separately* on each operating system to run
- ❑ **Application Binary Interface (ABI)** is an architecture equivalent of API, defines how different *components of binary code* can interface for a given operating system on a given architecture, CPU, etc.



# Operating System Design and Implementation

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- Design and implementation of OS *not “solvable”*, but some approaches have proven successful
- Internal structure of different Operating Systems can *vary widely*
- Affected by *choice of hardware*, *type of system*
- Start the design by defining *goals* and *specifications*
  - **User goals** – operating system should be *convenient to use*, *easy to learn*, *reliable*, *safe*, and *fast*
  - **System goals** – operating system should be *easy to design*, *implement*, and *maintain*, as well as *flexible*, *reliable*, *error-free*, and *efficient*





# Operating System Design and Implementation (Cont.)

- Important principle is to separate

**Policy:** *What* will be done?

**Mechanism:** *How* to do it?

- *Mechanisms* determine how to do something; *policies* decide what will be done
- The *separation of policy from mechanism* is a very important principle, it allows *maximum flexibility* if policy decisions are to be changed later (e.g., timer)
- Specifying and designing an OS is highly creative task of *software engineering*



# OS Implementation

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- ❑ Much variation
  - ❑ Early, OSes in *assembly language*
  - ❑ Then, *system programming languages* like Algol, PL/1, and now C, C++
- ❑ Usually, a mix of languages
  - ❑ *Lowest levels* in assembly
  - ❑ *Main body* in C
  - ❑ *Systems programs* in C, C++, scripting languages like Perl, Python, shell scripts
- ❑ More high-level language easier to *port* to other hardware, but slower
- ❑ *Emulation* can allow an OS to run on *non-native hardware*



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# OPERATING SYSTEM STRUCTURE



# Operating System Structure

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- *General-purpose OS* is a very large program
- Various ways to structure ones
  - Simple structure – **MS-DOS**
  - More complex – **UNIX**
  - Layered – an abstraction
  - Microkernel – **Mach**



# Monolithic Structure – Original UNIX

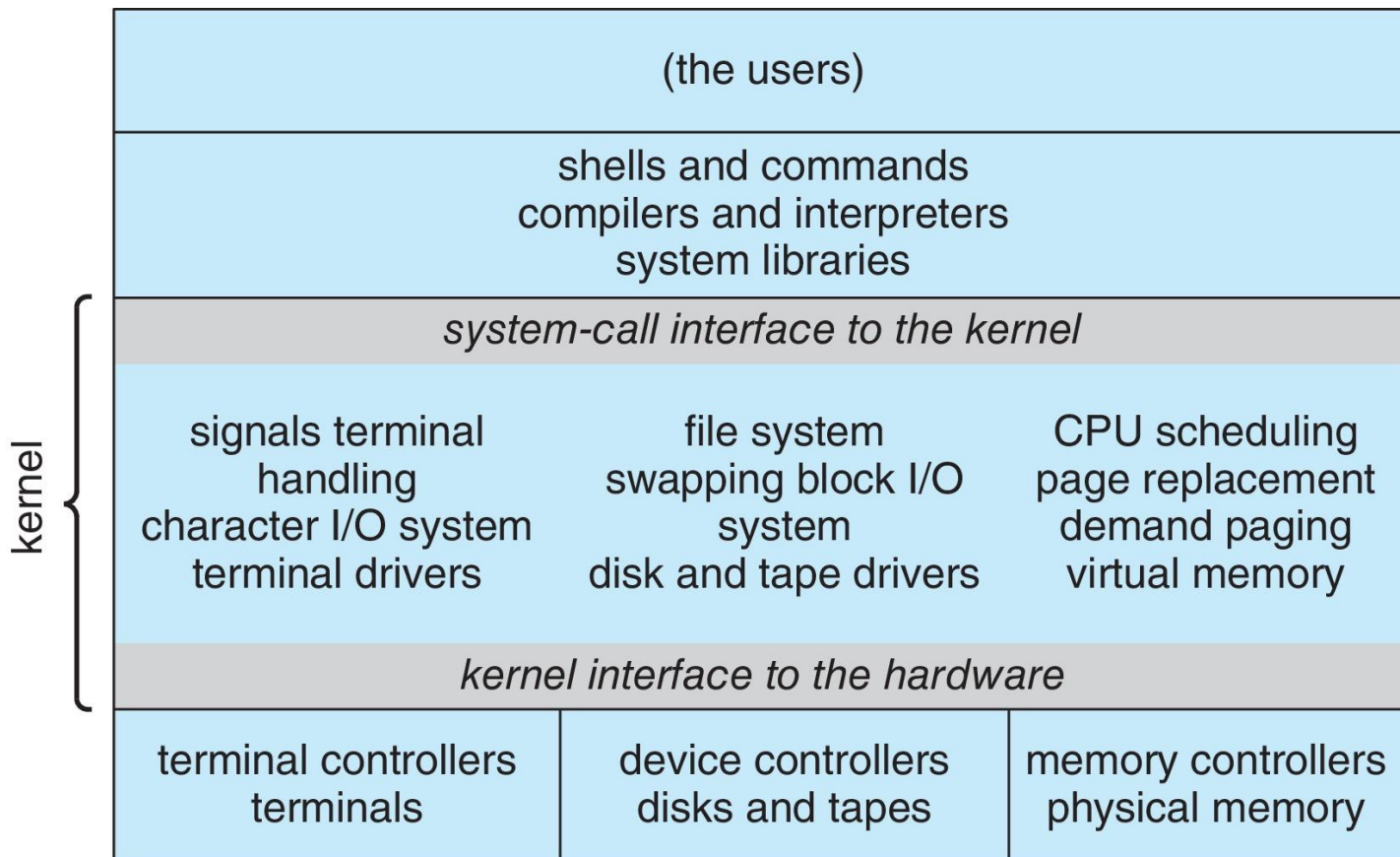
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- UNIX – limited by hardware functionality, the original UNIX operating system had limited structuring. The UNIX OS consists of two separable parts
  - **Systems programs**
  - **Kernel**
    - ▶ Consists of everything below the *system-call interface* and above the *physical hardware*
    - ▶ Provides the *file system*, *CPU scheduling*, *memory management*, and *other operating-system functions*; a large number of functions *for one level*



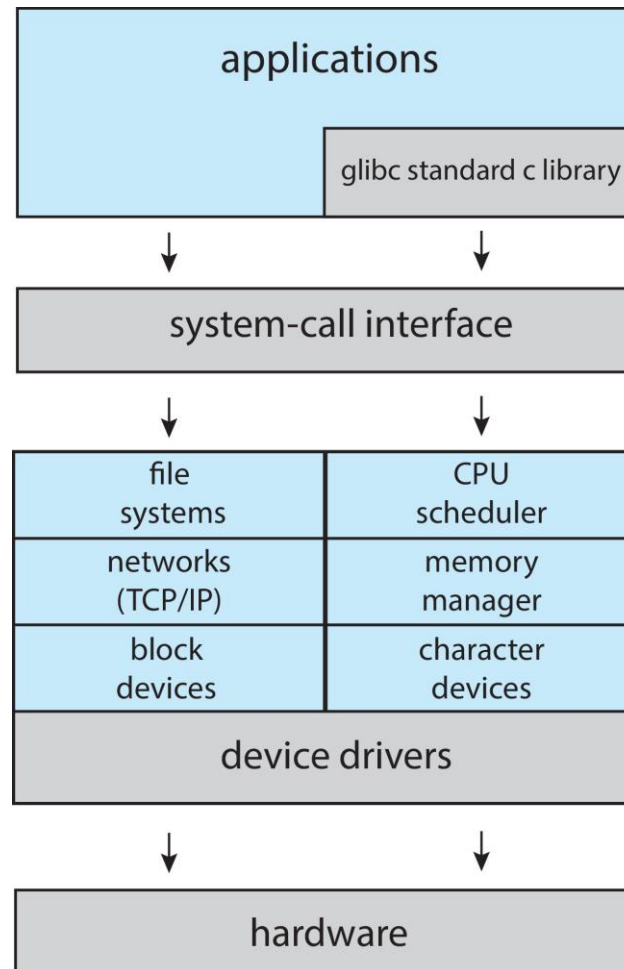
# Traditional UNIX System Structure

- Beyond simple but *not fully layered*



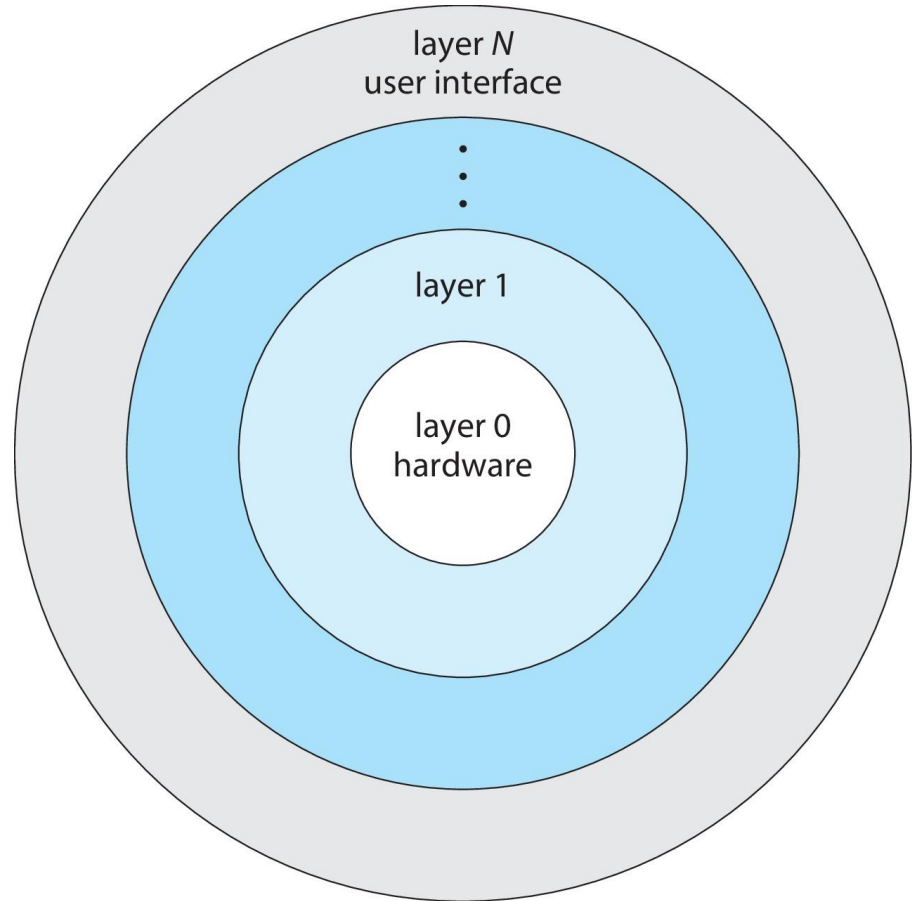
# Linux System Structure

## □ *Monolithic* plus *modular design*



# Layered Approach

- The operating system is divided into a *number of layers* (levels), each built on top of lower layers. The bottom layer (layer 0), is the *hardware*; the highest (layer N) is the *user interface*.
- With *modularity*, layers are selected such that *each uses functions (operations) and services of only lower-level layers*





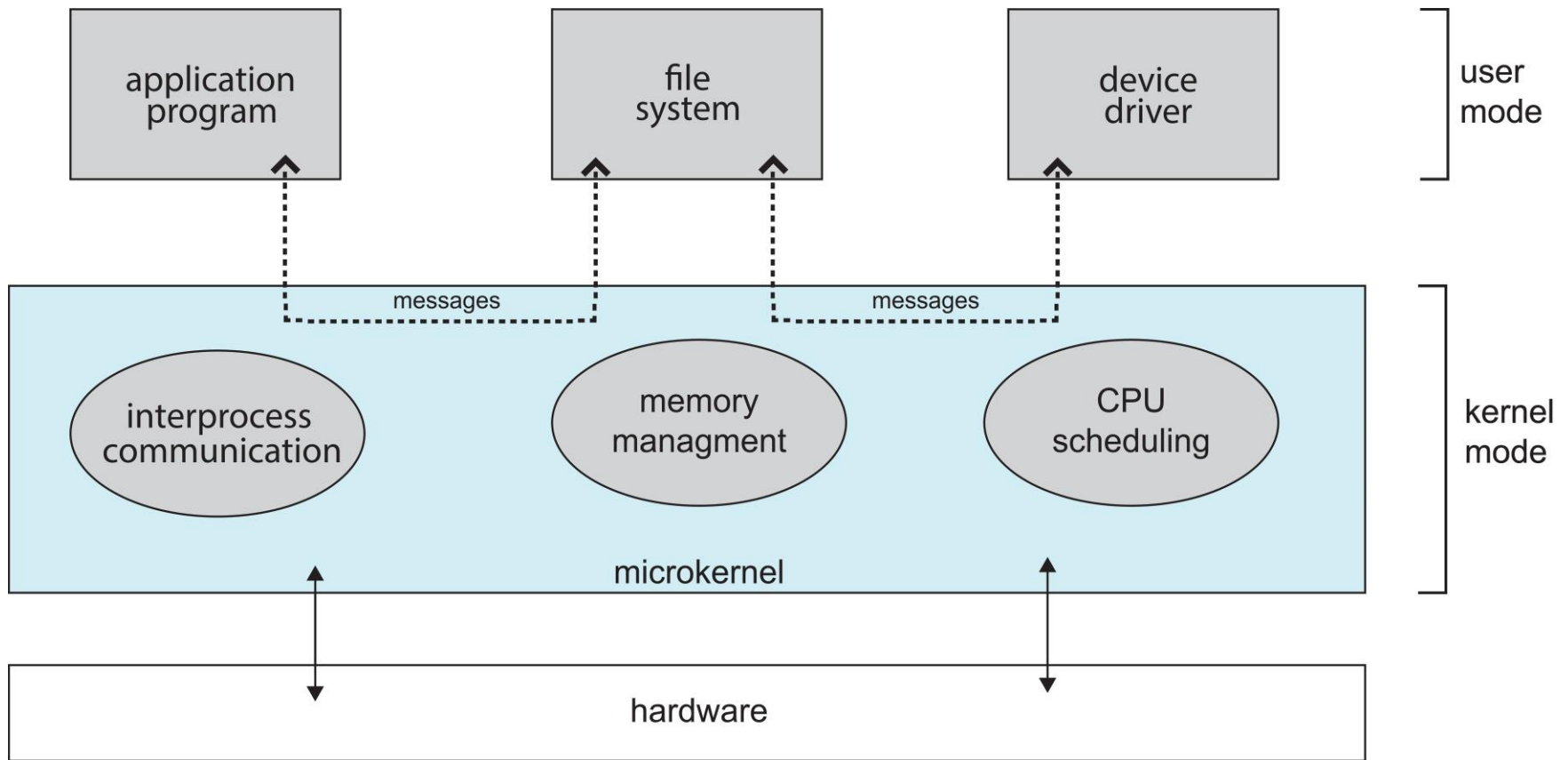
# Microkernels

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- Moves as much from the *kernel* into *user space*
- **Mach** is an example of *microkernel*
  - Mac OS X kernel (i.e., **Darwin**) partly based on Mach
- Communication takes place between user modules using *message passing model*
- Benefits
  - Easier to *extend* a microkernel
  - Easier to *port* the operating system to new architectures
  - More *reliable* (less code is running in kernel mode), more *secure*
- Detriments: *performance overhead* of user space to kernel space communication



# Microkernel System Structure



# Modules

---

- ❑ Many modern operating systems implement **Loadable Kernel Modules (LKMs)**
  - ❑ Uses *object-oriented approach*
  - ❑ Each *core component* is separate
  - ❑ Each *talks* to the others *over known interfaces*
  - ❑ Each is *loadable* as needed within the kernel
- ❑ Overall, *similar to layers* but with *more flexible*
  - ❑ E.g., Linux, Solaris, etc.



# Hybrid Systems

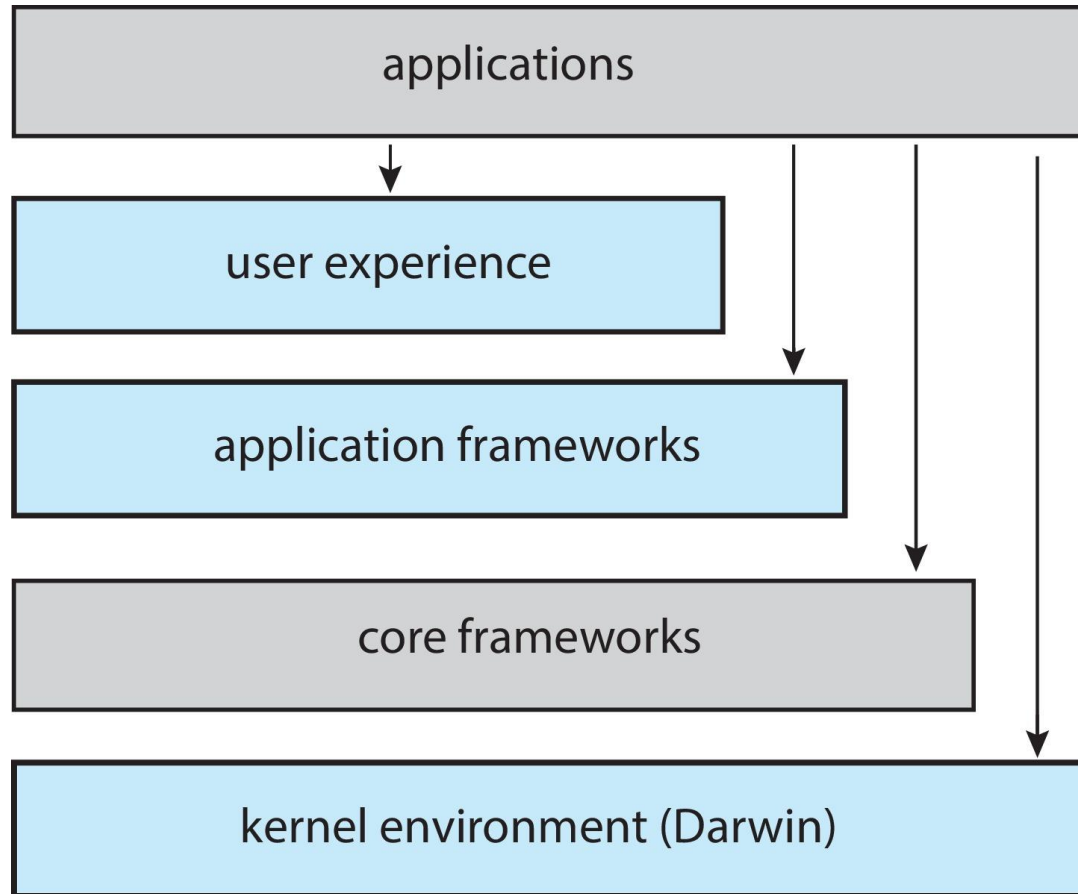
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- Most modern operating systems are actually *not one pure model*
- **Hybrid system** combines multiple approaches to address *performance, security, usability needs*
  - **Linux** and **Solaris** kernels in kernel address space, so *monolithic*, plus *modular* for *dynamic loading of functionality*
  - **Windows** mostly *monolithic*, plus *microkernel* for different *subsystem personalities*
  - **Apple Mac OS X** *hybrid, layered, Aqua UI* plus *Cocoa programming environment*
    - ▶ Below is kernel consisting of *Mach microkernel* and BSD Unix parts, plus *I/O kit* and *dynamically loadable modules* (called *kernel extensions*)

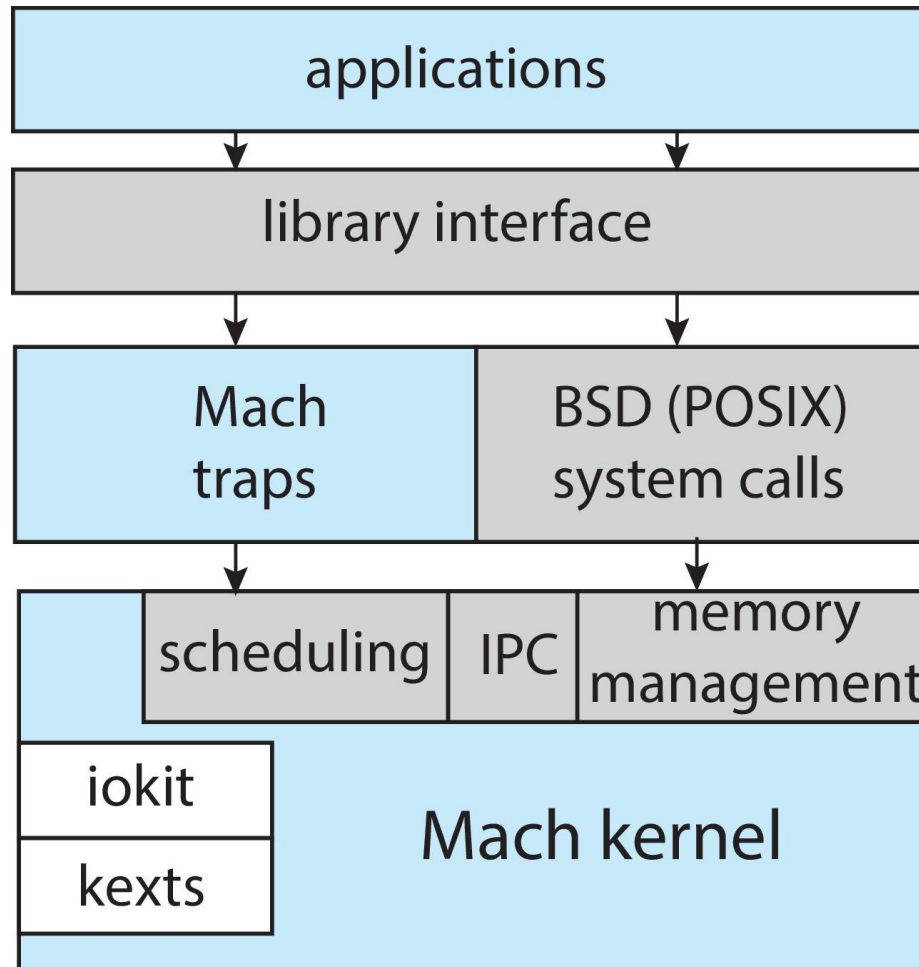


# macOS and iOS Structure

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# Darwin



# iOS

- Apple mobile OS for *iPhone, iPad*
  - Structured on Mac OS X, added functionality
  - Does not run OS X applications natively
    - ▶ Also runs on different CPU architecture (ARM vs. Intel)
  - **Cocoa Touch Objective-C API** for developing apps
  - **Media services** layer for graphics, audio, video
  - **Core services** provides cloud computing, databases
  - **Core operating system**, based on Mac OS X kernel

Cocoa Touch

Media Services

Core Services

Core OS



# Android

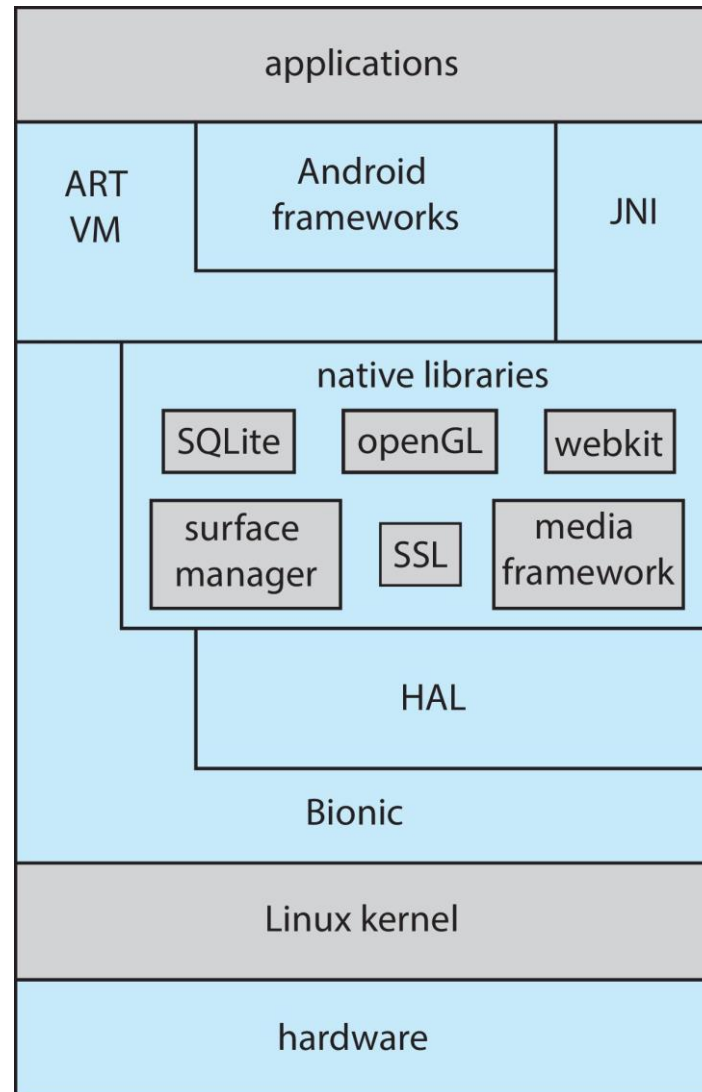
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- ❑ Developed by **Open Handset Alliance** (mostly Google), *open source*
- ❑ Similar stack to iOS
- ❑ Based on **Linux kernel** but modified
  - ❑ Provides process, memory, device-driver management
  - ❑ Adds *power management*
- ❑ **Runtime environment** includes *core set of libraries* and *Dalvik virtual machine* (VM), apps developed in **Java** plus **Android API**
  - ▶ *Java class files* compiled to *Java bytecode* then translated to *executable* than runs in *Dalvik VM*
- ❑ **Libraries** include *frameworks for web browser* (*webkit*), *database* (*SQLite*), *multimedia*, smaller *libc*, *etc.*





# Android Architecture



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# BUILDING AND BOOTING AN OPERATING SYSTEM



# Building and Booting an Operating System

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- Operating systems generally *designed* to run on a *class of systems* with *variety of peripherals*
- Commonly, operating system already *installed* on purchased computer
  - But can build and install some other operating systems
  - If generating an operating system *from scratch*
    - ▶ **Write** the operating system source code
    - ▶ **Configure** the operating system for the system on which it will run
    - ▶ **Compile** the operating system
    - ▶ **Install** the operating system
    - ▶ **Boot** the computer and its new operating system



# Building and Booting Linux

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- *Download* Linux source code (from <http://www.kernel.org>)
- *Configure* kernel via “**make menuconfig**”
- *Compile* the kernel using “**make**”
  - Produces **vmlinuz**, the kernel image
- *Compile* kernel modules via “**make modules**”
- *Install* kernel modules into **vmlinuz** via “**make modules\_install**”
- *Install* new kernel on the system via “**make install**”



# System Boot

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- ❑ When power on system, execution *starts at a fixed memory location*
- ❑ Operating system must be made available to hardware so hardware can start it
  - ▶ Small piece of code – **bootstrap loader**, **BIOS**, stored in **ROM** or **EEPROM** *locates the kernel, loads it into memory, and starts it*
  - ▶ Sometimes two-step process where **boot block** at fixed location loaded by ROM code, which loads bootstrap loader from disk
  - ▶ Modern systems replace BIOS with **Unified Extensible Firmware Interface (UEFI)**
- ❑ Common bootstrap loader, **GRUB**, allows selection of kernel from multiple disks, versions, kernel options
- ❑ Kernel loads and system is then *running*
- ❑ Boot loaders frequently allow *various boot states*



# Operating-System Debugging

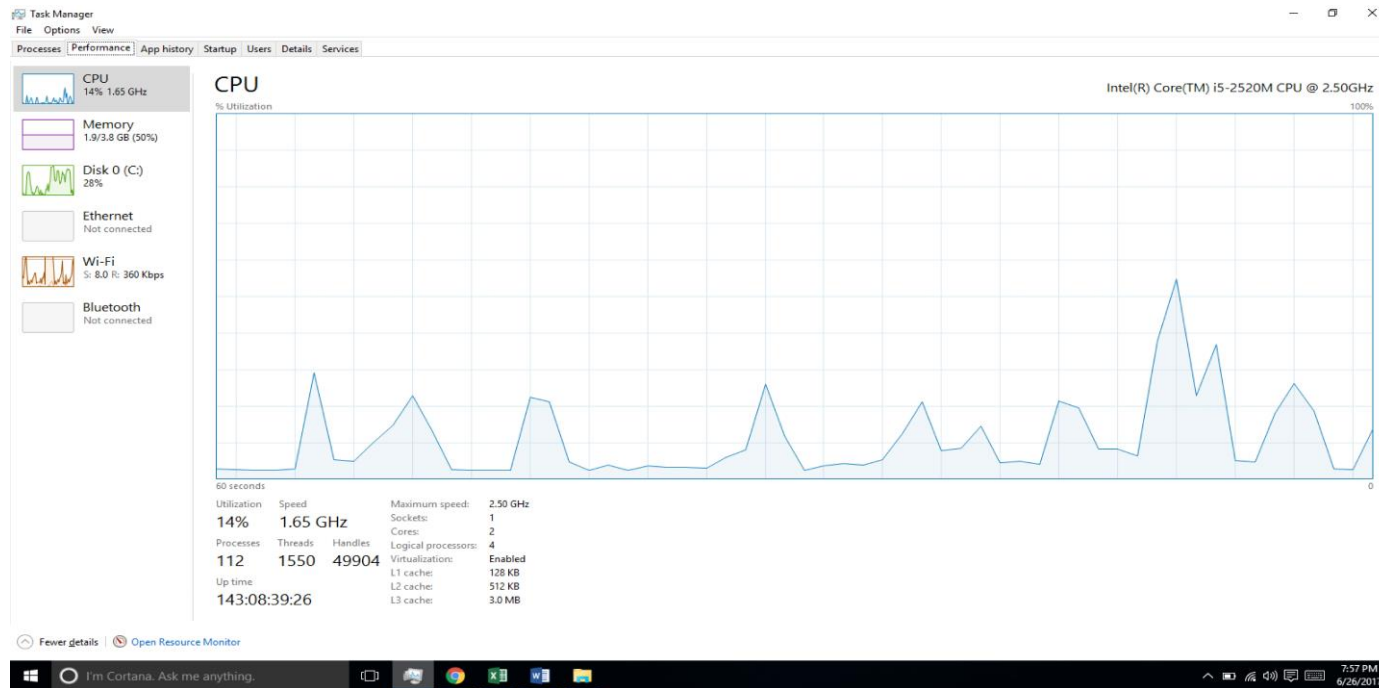
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- *Debugging* is finding and fixing errors, or *bugs*
  - Also, *performance tuning*
- OS generate *log files* containing error information
  - Application failure can generate *core dump file* capturing memory of the process
  - Operating system failure can generate *crash dump file* containing kernel memory
- Beyond crashes, *performance tuning* can optimize system performance
  - Sometimes using *trace listings* of activities, recorded for analysis
  - *Profiling* is *periodic sampling of instruction pointer* to look for statistical trends

Kernighan's Law: “*Debugging is twice as hard as writing the code in the first place. Therefore, if you write the code as cleverly as possible, you are, by definition, not smart enough to debug it.*”



# Performance Tuning



- Improve performance by *removing bottlenecks*
- OS must provide means of *computing and displaying measures* of system behavior
  - E.g., “top” Linux program or **Windows Task Manager**



# Tracing

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- ❑ *Collects data for a specific event*
  - ❑ E.g., steps involved in a system call invocation
- ❑ *Tools* include
  - ❑ **strace** – trace system calls invoked by a process
  - ❑ **gdb** – source-level debugger
  - ❑ **perf** – collection of Linux performance tools
  - ❑ **tcpdump** – collects network packets





# BCC

- *Debugging interactions* between user-level and kernel code nearly impossible without toolset that understands both and instrument their actions
- **BPF Compiler Collection (BCC)** is a rich *toolkit* providing tracing features for Linux
  - See also the original **DTrace**
  - For example, `disksnoop.py` traces disk I/O activity

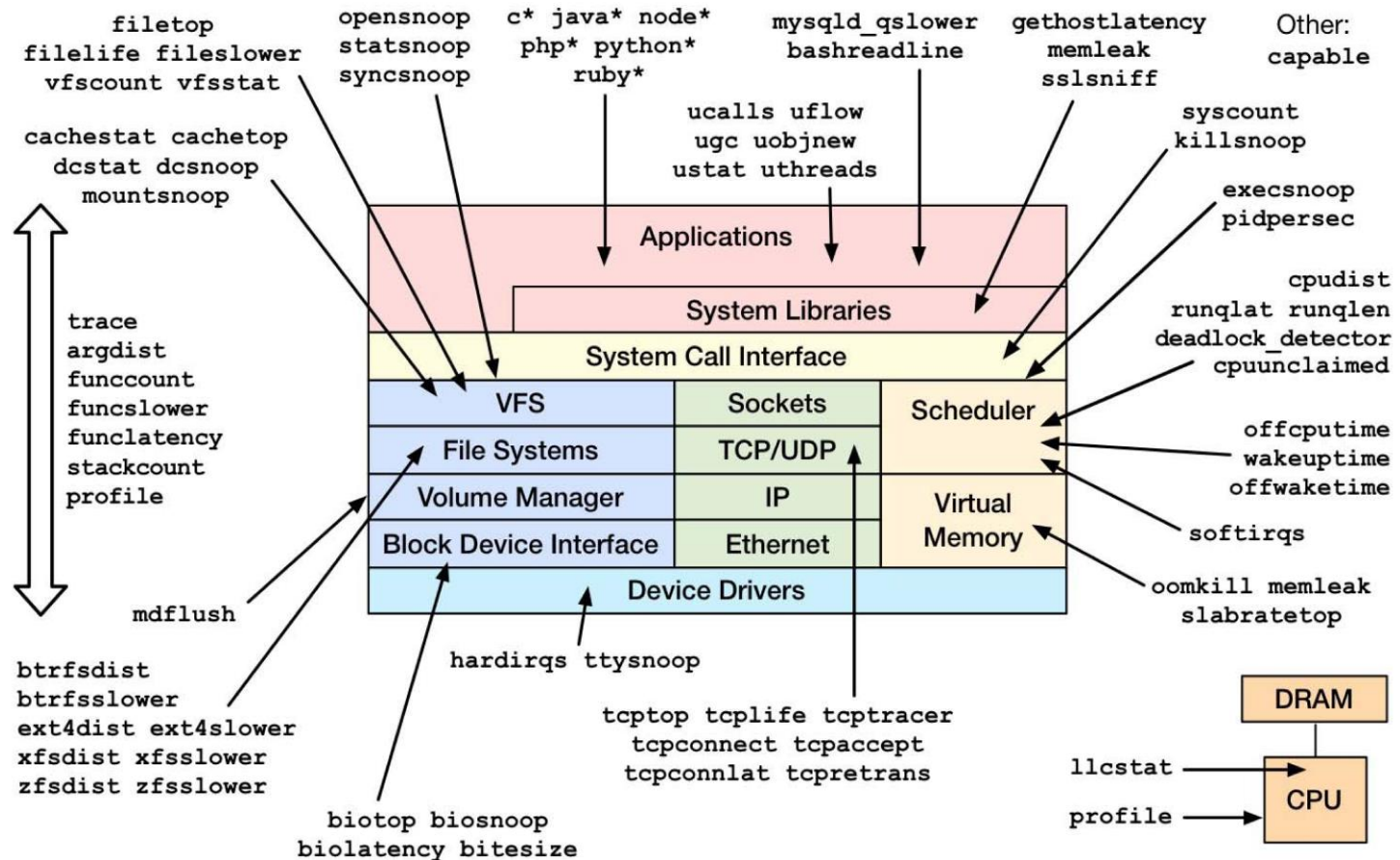
TIME(s)	T	BYTES	LAT(ms)
1946.29186700	R	8	0.27
1946.33965000	R	8	0.26
1948.34585000	W	8192	0.96
1950.43251000	R	4096	0.56
1951.74121000	R	4096	0.35

- Many other tools (next slide)



# Linux BCC/BPF Tracing Tools

## Linux bcc/BPF Tracing Tools



<https://github.com/iovisor/bcc#tools> 2017



# Summary

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- ❑ An **operating system** provides an environment for the execution of programs by providing services to users and programs.
- ❑ The **three primary approaches for interacting** with an operating system are (1) command interpreters, (2) graphical user interfaces, and (3) touch-screen interfaces.
- ❑ **System calls** provide an interface to the services made available by an operating system. Programmers use a system call's application programming interface (API) for accessing system-call services.
- ❑ System calls can be divided into **six major categories**: (1) process control, (2) file management, (3) device management, (4) information maintenance, (5) communications, and (6) protection.
- ❑ The **standard C library** provides the system-call interface for UNIX and Linux systems.



# Summary (Cont.)

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- ❑ Operating systems also include a **collection of system programs** that provide utilities to users.
- ❑ A **linker** combines several relocatable object modules into a single binary executable file. A **loader** loads the executable file into memory, where it becomes eligible to run on an available CPU.
- ❑ There are several reasons why **applications are operating-system specific**. These include different binary formats for program executables, different instruction sets for different CPUs, and system calls that vary from one operating system to another.
- ❑ An operating system is **designed** with specific goals in mind. These goals ultimately determine the **operating system's policies**. An operating system implements these policies through specific mechanisms.



# Summary (Cont.)

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- A **monolithic operating system** has no structure; all functionality is provided in a single, static binary file that runs in a single address space. Although such systems are difficult to modify, their primary benefit is efficiency.
- A **layered operating system** is divided into a number of discrete layers, where the bottom layer is the hardware interface, and the highest layer is the user interface.
  - Although layered software systems have had some success, this approach is generally not ideal for designing operating systems due to performance problems.
- The **microkernel** approach for designing operating systems uses a minimal kernel; most services run as user-level applications. Communication takes place via message passing.



# Summary (Cont.)

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- A **modular** approach for designing operating systems provides operating-system services through modules that can be loaded and removed during run time. Many contemporary operating systems are constructed as **hybrid systems** using a combination of a monolithic kernel and modules.
- A **boot loader** loads an operating system into memory, performs initialization, and begins system execution.
- The **performance** of an operating system can be monitored using either counters or tracing. **Counters** are a collection of system-wide or per-process statistics, while **tracing** follows the execution of a program through the operating system.



# End of Chapter 2

