Operating system structures

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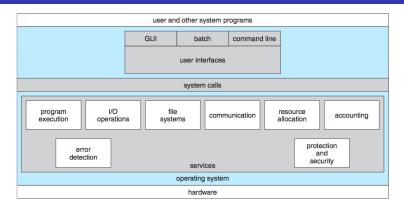
Outline

- 1 Operating system services
- 2 System calls and programs
- 3 Operating system structure
- 4 Advanced issues

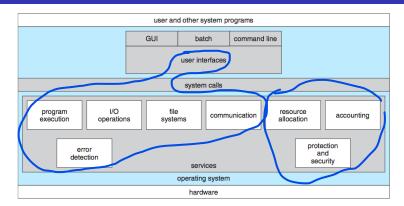
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Operating system services



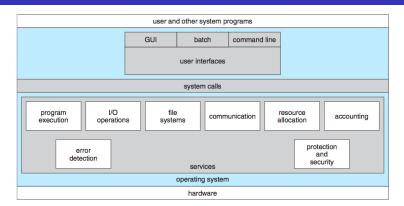
Operating system services



Functional services

Non-functional services

Operating system services



- Functional services
 - User interface (Graphical User Interface, Batch Interface, Command Line Interface), Program execution, I/O operations, File-system manipulation, Communications. Error detection

Non-functional services

Resource allocation, Accounting, Protection and Security

Program execution

What is a program?

A computer program is a collection of instructions that performs a specific task when executed by a computer.

Program execution

What is a program?

A computer program is a collection of instructions that performs a specific task when executed by a computer.

- OS is able to load a program into memory and to run that program
- The program is able to end its execution (either normally or abnormally)

I/O operations & File-system manipulation

I/O operations

- A running program may access I/O devices, e.g., recording DVD)
- For efficiency and protection, OS must provide a means (for programs) to do I/O

I/O operations & File-system manipulation

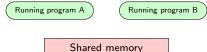
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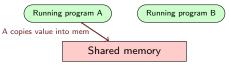
File-system manipulation

■ Programs need operations on files/directories: list, create, delete, read, write, permission management

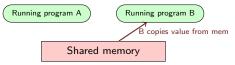
- There are many communications
 - between running programs on a machine
 - between running program on different machines
- There are two main types of communications
 - shared memory: read/write on a shared of memory
 - message passing: packets with predefined formats are exchanged between running programs



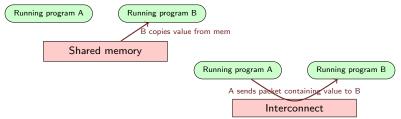
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Error detection

- Errors can occur everywhere
 - Hardware: CPU, memory, I/O devices,...
 - User program: arithmetic overflow, illegal access of memory, division by zero,...
- OS should take appropriate actions to detect and correct errors constantly.

Non-functional services

Resource allocation

- OS manages multiple resources (hardware, software) and allocates them to multiple users and multiple running programs
- Special codes are needed to make allocation efficiently, e.g., CPU scheduling, printers allocation

Accounting

- Recording which users to use how much and what kind of resources
- Usage statistics is useful to reconfigure for improvement of computing services

Protection and security

• Information of multiple users on a networked computer should be controlled by its owner.

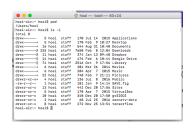
User interface

- Command interpreter a special program to allows users to directly enter commands (or programs) (by text) to be performed by the operating system
 - Multiple command interpreters (also shell) in a modern operating system Example: Bourne shell, Bash shell, C shell, Bourne-Again shell, Korn shell

```
↑ hoai — -bash — 80×24
hoai-air:~ hoai$ owd
/Users/hoai
hoai-air:~ hoai$ ls -1
            5 hoai staff 170 Jul 14 2015 Applications
            5 hoai staff
                           178 Feb 9 18:27 Desktop
drwx----+ 16 hoai staff
                           544 Aug 31 18:40 Documents
drex-----@ 225 hoai staff 7650 Feb 8 12:04 Downloads
drwx----- 9 11 hoai staff
                           374 Jan 13 89:48 Dropbox
           14 hoai staff
                           476 Feb 6 10:44 Google Drive
drwx-----9 74 hoai staff 2516 Oct 9 17:56 Library
            6 hoai staff
                           284 Mar 26 2014 Movies
                           284 Apr 7 2015 Music
            6 hoai staff
drwx----+ 22 hoai staff
                           748 Feb 7 21:11 Pictures
drwxr-xr-x+ 4 hoai staff
                           136 Jul 8 2016 Public
            1 hoai staff
                           181 Jan 9 14:14 SAVE.fig
drwxr-xr-x 13 hosi staff
                           442 Dec 28 17:56 Sites
drwxr-xr-x 5 hoai staff
                           170 Apr 7 2015 VirtualBox
drwxr-xr-x 10 hoai staff
                           348 Dec 28 17:58 gvSIGCE
drwxr-xr-x 2 hoai staff
                           68 Jul 25 2016 socnety-data
drwxr-xr-x 8 hosi staff
                           272 Nov 25 13:51 tensorflow
hoai-air:~ hoai$ Ⅱ
```

User interface

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■ GUI - a user friendly graphical interface, input/output is performed in a more interactive way

Choice of interface

Interface choice

Kind of interface (CLI or GUI) is mostly one of personal preference

- System administratiors or power users often prefer CLI
 - With deep system knowledge, using CLI is more efficient and secure
 - Multiple shell commands can be combined into a program, called shell script to perform more complex tasks.
- Normal users often choose GUI

Choice of interface

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UI can vary from system to system. Therefore, it is **not** a **direct function** of operating system.

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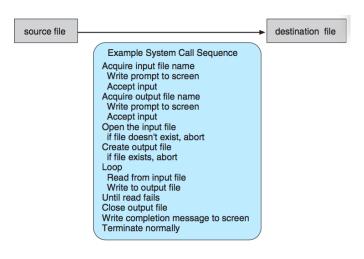
Example on system call demand

File-to-file copy

source file destination file Example System Call Sequence Acquire input file name Write prompt to screen Accept input Acquire output file name Write prompt to screen Accept input Open the input file if file doesn't exist, abort Create output file if file exists, abort Loop Read from input file Write to output file Until read fails Close output file Write completion message to screen Terminate normally

Example on system call demand

File-to-file copy



Even a simple program executes many system calls

API vs. system call interface

Application Programming Interface (API)

Set of functions available to application programmers

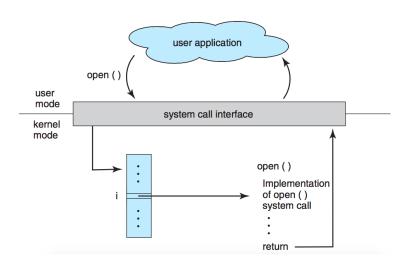
- Using API to increase program portability: ability to compile and run on systems with the same API without code modification
- Working with API is easier than with actual system calls
- 3 most common APIs: Windows API, POSIX API, Java API

System-call interface

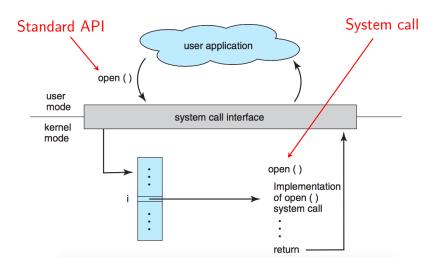
A run-time support system (e.g., libc) serves as a link between API and system calls in operating system

■ Each system call is identified by an index number

A call to system call open()



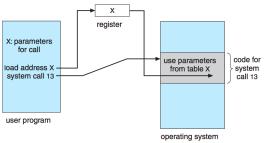
A call to system call open()



Parameters passing to system call

There are 3 ways.

- By registers: not enough for large parameters
- Stored in block (or table), passing address of the block in parameter



Pushed onto program's stack, and popped off by OS

Types of system calls

- Process control
- File management
- Device management
- Information maintenance
- Communications
- Protection

Process control

What is a program ?

A computer program is a collection of instructions that performs a specific task when executed by a computer.

What is a process?

A program loaded into memory and executing is called a process

- end(), abort()
- load(), execute()
- create_process(),
 terminate_process()
- get_process_attributes(), set_process_attributes()

- wait_for_time()
- wait_event(), signal_event()
- allocate and free memory

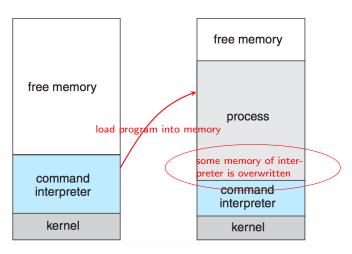
Execution on single-task OS MS-DOS

free memory free memory process command command interpreter interpreter kernel kernel

At system startup

Running a program

Execution on single-task OS MS-DOS

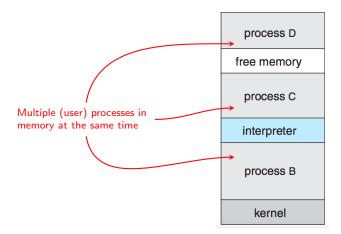


At system startup

Running a program

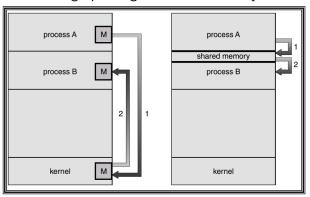
Execution on multiple-task OS

FreeBSD



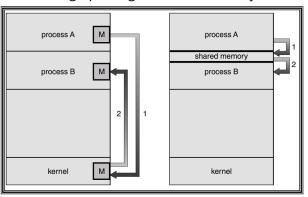
Interprocess communication

 Interprocess communication may take place using either message passing or shared memory



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Message passing

Shared memory

System programs

$System\ programs = system\ utilities$

Providing a convenient environment for program development and execution

There are categories as follows.

- File manipulation
- Status information
- File modification
- Programming-language support
- Program loading and execution
- Communications
- Background services

System programs

System programs = system utilities

Providing a convenient environment for program development and execution

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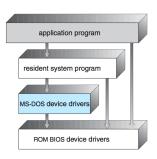
The view of most users on the operating system is defined by system programs, not actual system calls.

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Simple structure MS-DOS

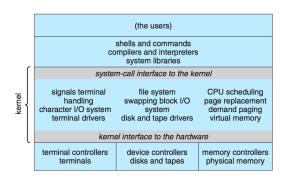
- MS-DOS: written to provide the most functionality in the least space
 - not divided into modules
 - interfaces and levels of functionality are not well separated



Simple structure

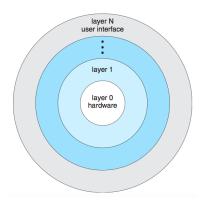
Traditional UNIX

- Limited structure with limited functionality
- Consists of 2 parts: kernel and system programs



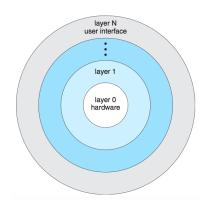
- Considered to be layered to some extent: system-call inteface, kernel and hardware interface
- lacktriangle Enormous amount of functionality combined in one level ightarrow monolithic structure

Layered approach



- OS is broken into several layers
- Layer M provides data structures
 & routines for upper layer, and
 can invoke operations of lower
 layer

Layered approach



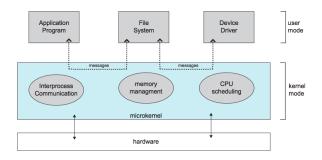
Advantages (due to modularity): easy to debug; simple in design and implementation; secure

- OS is broken into several layers
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 layer

Disadvantages: Not easy to define layers (which layer is above/below which layer); inefficiency

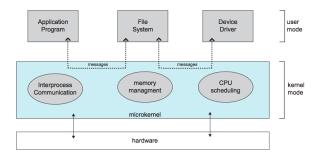
Microkernel structure

- Moves as much from the kernel into "user" space
- Communication takes place between user modules using message passing



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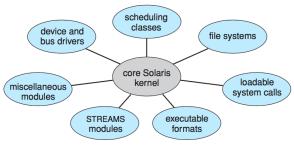


Advantages (due to modularity): easy to extend kernel; easier to port OS to new architectures; more reliable (less code running in kernel); more secure

Disadvantages: system-function overhead

Modules-based structure

- Best current methodology is loadable kernel modules
 - Only need core services
 - Additional services can be loaded as modules in boot time and run time
 - Additional services have to be recompiled to add new features
- Idea of modules-based kernel more flexible than layered approach and also similar to microkernel structure



Hybrid structure

In practice, very few OS adopt a single, stricted defined structure

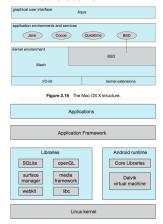


Figure 2.18 Architecture of Google's Android.



Figure 2.17 Architecture of Apple's iOS.

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