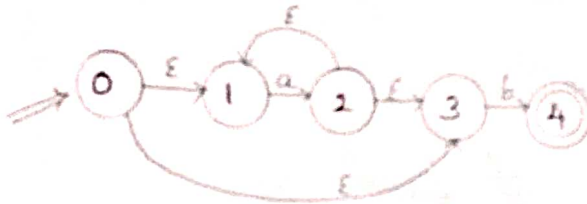


# THEORY OF COMPUTATION

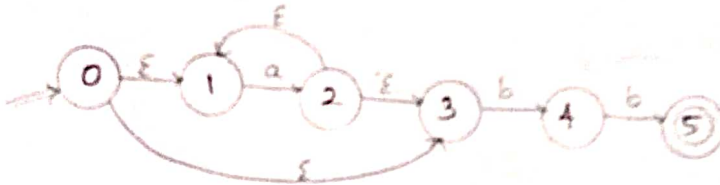
## 1-NFA

Construct NFA for the following regular expressions

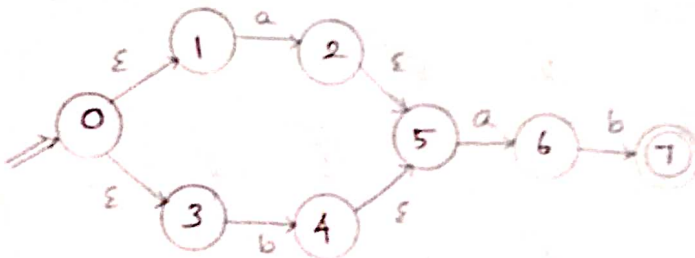
(1)  $a^*b$



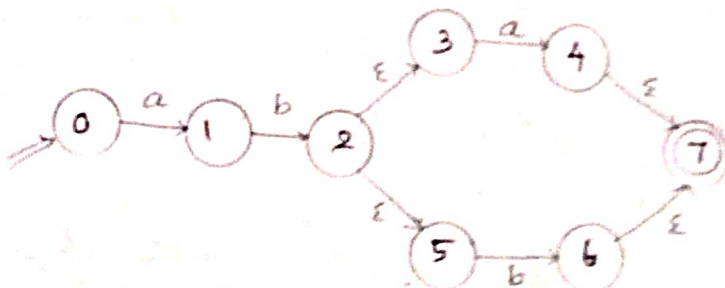
(2)  $a^*bb$



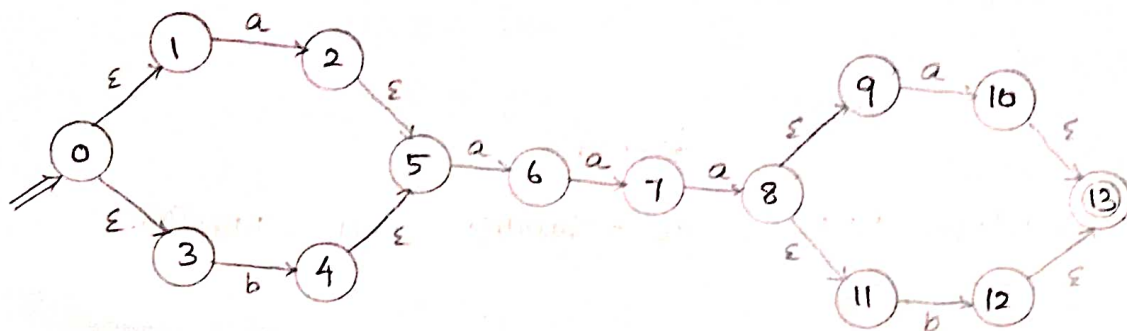
(3)  $(a|b)ab$



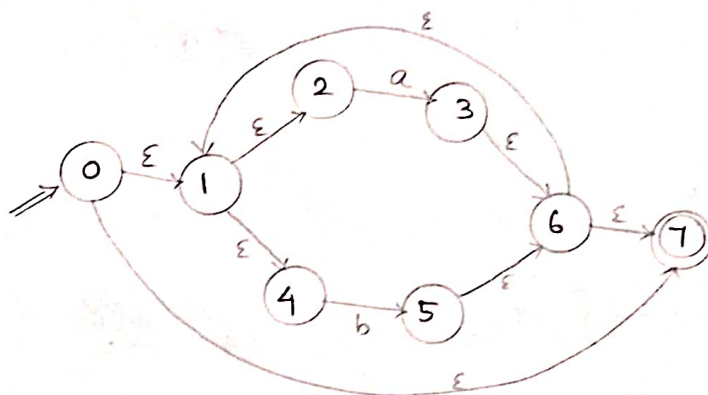
(4)  $ab(a|b)$



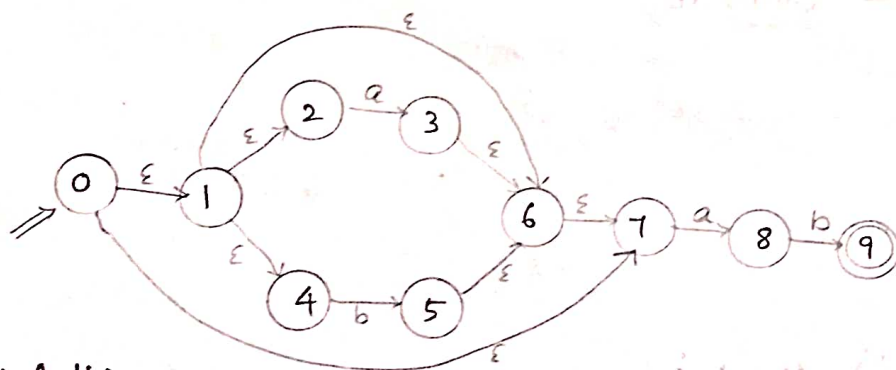
(5)  $(a|b)aaa(a|b)$



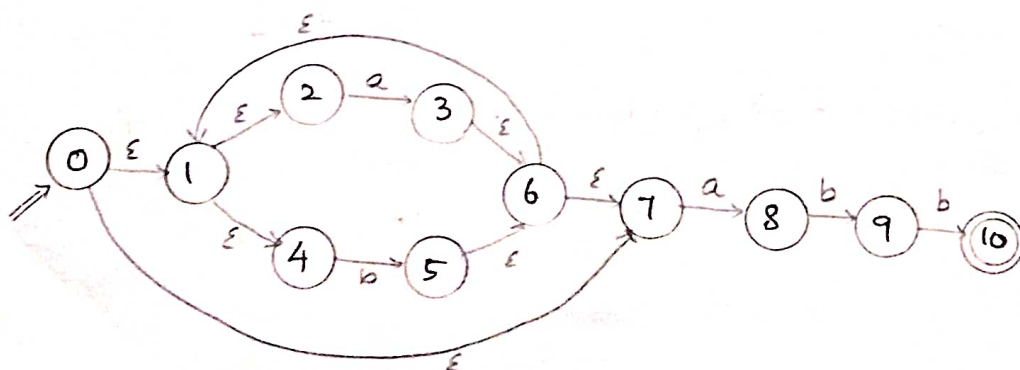
(6)  $(a|b)^*$



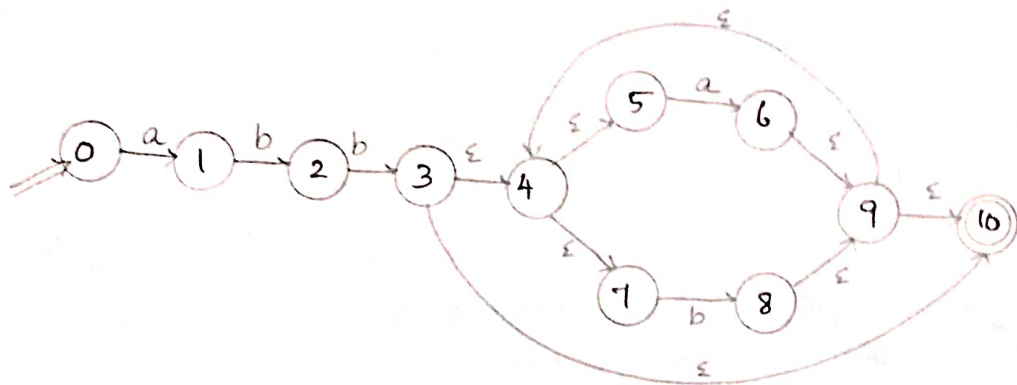
(7)  $(a|b)^*ab$



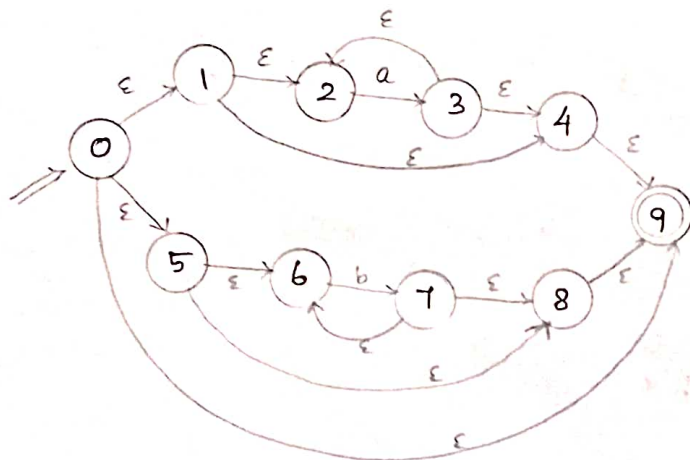
(8)  $(a|b)^*abb$



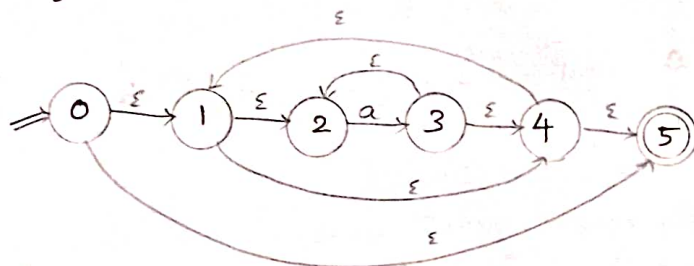
(9)  $abb(a|b)^*$



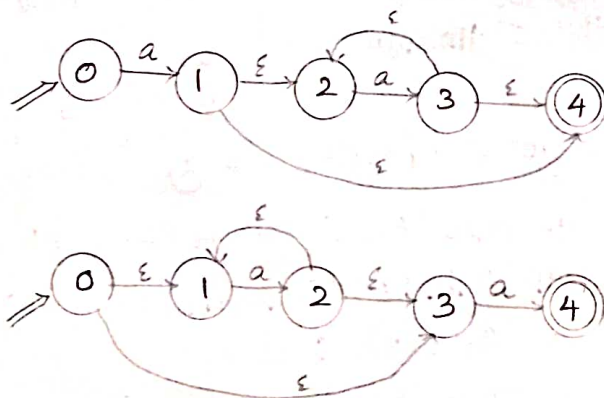
(10)  $a^*|b^*$



(11)  $(a^*)^* = a^*$



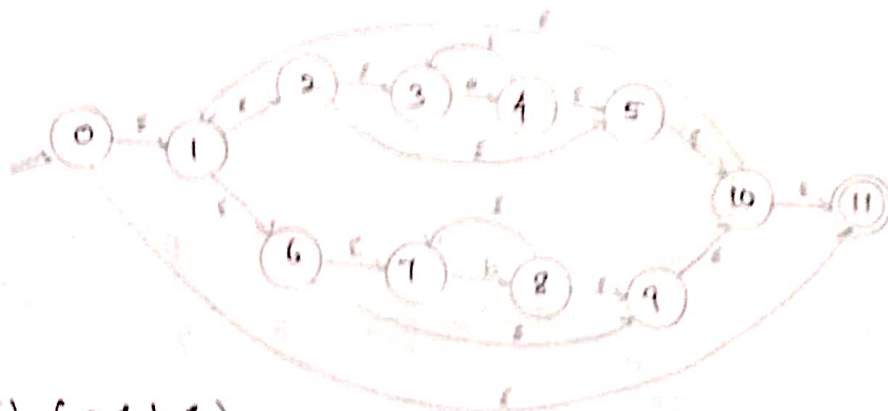
(12)  $a^+ = aa^* = a^*a$



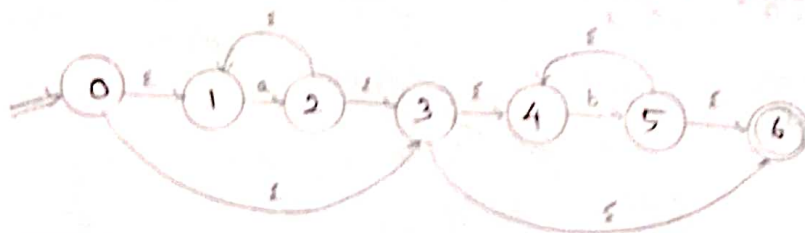
(13)  $(a^* | b^*)abb$



(14)  $(a^* | b^*)^*$



(15)  $(a^* b^*)$



(1)  $a^* b$

2. DFA

Set of states that can be reached from state 0 through  $\epsilon$ -transition.

I/p	starting node	ending node.
a	1	2
b	3	4

$\epsilon$ -closure(0) = {0, 1, 3}  $\rightarrow$  (A)

Apply input 'a' on state (A),  $\epsilon$ -closure(2) = {1, 2, 3}  $\rightarrow$  (E)

Apply input 'b' on state (A),  $\epsilon$ -closure(4) = {4}  $\rightarrow$  (C)

Apply input 'a' on state (E),  $\epsilon$ -closure(2) =  $\rightarrow$  (E)

Apply input 'b' on state (E),  $\epsilon$ -closure(4) =  $\rightarrow$  (C)