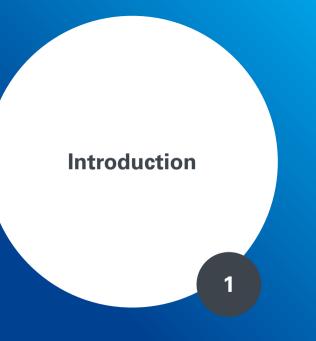


Fachpraktikum Algorithmik für OSM-Daten

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- only using upwards reduces the searchspace

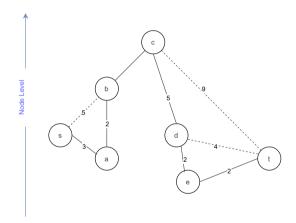


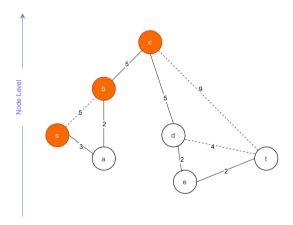
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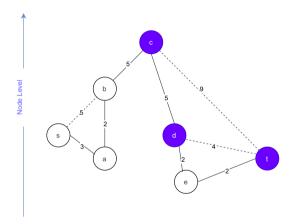
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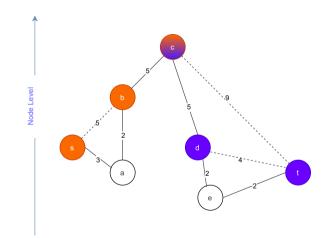
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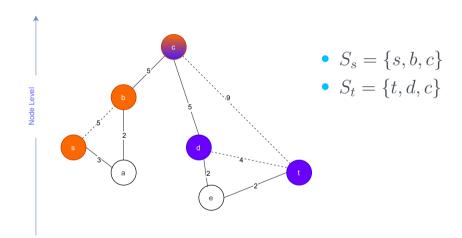
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- find node $n \in S_s \cap \in S_t$ where d(s,n) + d(n,t) is minimal

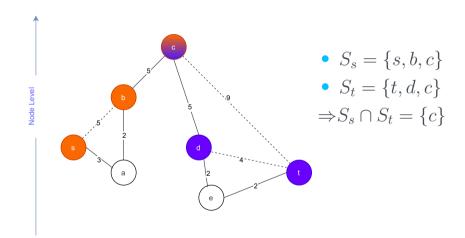


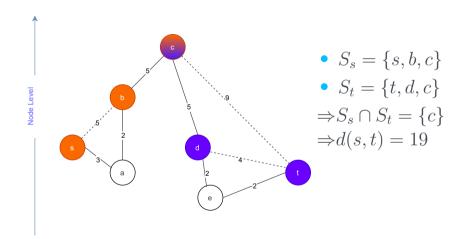


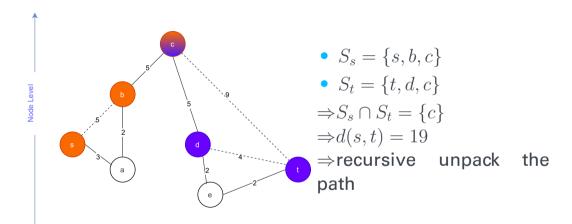












Optimizations

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But...

This optimization cannot be observed well when benchmarking random s-t queries.

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Idea: Nodes $x, y \in V$ can have the same level if $xy = e \notin E$

• Select $U \subseteq V$ such that $\forall x, y \in U : xy \notin E$

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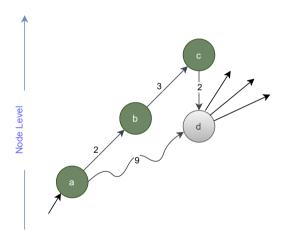
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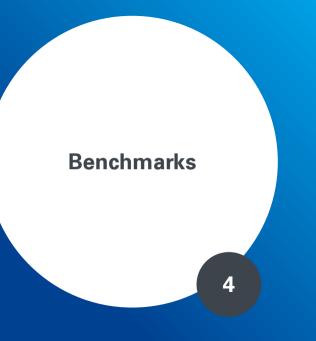
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- add contractions where the $edgedifference \leq avgedgedifference$
- ⇒Reduces the preprocessing time

Stall On Demand

Some Nodes are settled with wrong distances and can never be part of a shortest path.

Stall On Demand

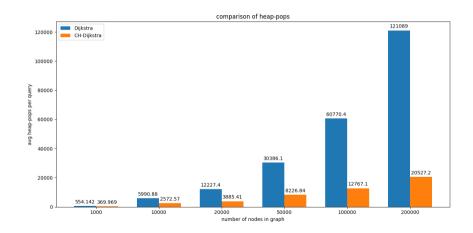




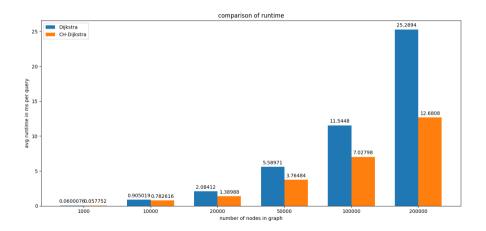
Benchmarks

The benchmarks compare a normal Dijkstra with an optimized CH-Dijkstra

Heap-Pops



Querytime



THE END

Thanks for listening!