

1. 6a

$$1.0\text{E}6 = 10^6$$

$$1\text{ GHz} = 1 \times 10^9\text{Hz}$$

		10%	20%	50%	20%
P1	2.5GHz	1	2	3	3
P2	3GHz	2	2	2	2

$$T_{P1} = \frac{(0.1 \times 10^6 + 0.2 \times 10^6 \times 2 + 0.5 \times 10^6 \times 3 + 0.2 \times 10^6 \times 3)}{2.5 \times 10^9} = 1.04 \times 10^{-3}\text{s}$$

$$T_{P2} = \frac{(0.1 \times 10^6 \times 2 + 0.2 \times 10^6 \times 2 + 0.5 \times 10^6 \times 2 + 0.2 \times 10^6 \times 2)}{3 \times 10^9} = 0.67 \times 10^{-3}\text{s}$$

$$\therefore T_{P1} > T_{P2}$$

\therefore P2 is faster than P1

$$\text{CPI}(P1) = \frac{1.04 \times 10^{-3} \times 2.5 \times 10^9}{10^6} = 2.6$$

$$\text{CPI}(P2) = \frac{0.67 \times 10^{-3} \times 3 \times 10^9}{10^6} = 2.0$$

\therefore P1' s global CPI is 2.0

P2' s global CPI is 2.6

1. 6b

$$\text{Clock Cycle}_{P1} = 0.1 \times 10^6 + 0.2 \times 10^6 \times 2 + 0.5 \times 10^6 \times 3 + 0.2 \times 10^6 \times 3 \\ = 2.6 \times 10^6$$

$$\text{Clock Cycle}_{P2} = 0.1 \times 10^6 \times 2 + 0.2 \times 10^6 \times 2 + 0.5 \times 10^6 \times 2 + 0.2 \times 10^6 \times 2 \\ = 2.0 \times 10^6$$

\therefore P1 's Clock Cycle is 2.6×10^6 s and P2' s Clock Cycle is 2.0×10^6

1.9.1

1 Million = 10^6

$$T_{P1} = \frac{2.56 \times 10^9 + 1.28 \times 10^9 \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 9.60s$$

$$T_{P2} = \frac{\frac{2.56 \times 10^9}{0.7 \times 2} + \frac{1.28 \times 10^9}{0.7 \times 2} \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 7.02s$$

$$R_{P2} = \frac{9.6}{7.02} = 1.37$$

$$T_{P4} = \frac{\frac{2.56 \times 10^9}{0.7 \times 4} + \frac{1.28 \times 10^9}{0.7 \times 4} \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 3.86s$$

$$R_{P4} = \frac{9.6}{3.86} = 2.49$$

$$T_{P8} = \frac{\frac{2.56 \times 10^9}{0.7 \times 8} + \frac{1.28 \times 10^9}{0.7 \times 8} \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 2.25s$$

$$R_{P8} = \frac{9.6}{2.25} = 4.27$$

The total execution time for this program on 1, 2, 4, and 8 processors is 9.6s, 7.02s, 3.86s and 2.25s. The relative speedup of the 2, 4, and 8 processor is 1.37, 2.49 and 4.27.

1.9.2

$$T_{P1b} = \frac{2.56 \times 10^9 \times 2 + 1.28 \times 10^9 \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 10.88s$$

$$T_{P2b} = \frac{\frac{2.56 \times 10^9}{0.7 \times 2} \times 2 + \frac{1.28 \times 10^9}{0.7 \times 2} \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 7.95s$$

$$T_{P4b} = \frac{\frac{2.56 \times 10^9}{0.7 \times 4} \times 2 + \frac{1.28 \times 10^9}{0.7 \times 4} \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 4.29s$$

$$T_{P8b} = \frac{\frac{2.56 \times 10^9}{0.7 \times 8} \times 2 + \frac{1.28 \times 10^9}{0.7 \times 8} \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 2.47s$$

$$\therefore T_{P1b} > T_{P1} \quad T_{P2b} > T_{P2} \quad T_{P4b} > T_{P4} \quad T_{P8b} > T_{P8}$$

∴ The execution time of the program on 1, 2, 4, or 8 processors will increase.

1. 9. 3

$$\therefore T_{P4} = \frac{\frac{2.56 \times 10^9}{0.7 \times 4} + \frac{1.28 \times 10^9}{0.7 \times 4} \times 12 + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 3.86s$$

Let the new CPI of load/store instructions be x.

Then,

$$\frac{2.56 \times 10^9 + 1.28 \times 10^9 x + 2.56 \times 10^8 \times 5}{2 \times 10^9} = 3.86$$

$$0.64x + 1.92 = 3.86$$

$$x = 3.03$$

$$\therefore CPI_{\text{reduced}} = \frac{3.03}{12} = 25\%$$

The CPI of load/store instructions should be reduced to 25% in order for a single processor to match the performance of four processors using the original CPI values.

1.12.1

$$5.0\text{E}9 = 5 \times 10^9$$

$$T_{P1} = \frac{5 \times 10^9}{4 \times 10^9} = 1.25\text{s}$$

$$T_{P2} = \frac{0.75 \times 10^9}{3 \times 10^9} = 0.25\text{s}$$

$$\therefore T_{P2} < T_{P1}$$

\therefore P1 has smaller performance than P2.

\therefore It's false that the computer with the largest clock rate as having the largest performance.

1.12.2

$$T_{P1} = \frac{0.9 \times 10^9}{4 \times 10^9} = 0.225\text{s}$$

Let instructions to execute in the same time as P1 be x .

Then

$$\frac{0.75x}{3 \times 10^9} = 0.225$$

$$x = 0.9 \times 10^9$$

$$\therefore x < 0.9 \times 10^9$$

The number of instructions that P2 can execute in the same time that P1 needs to execute $1.0\text{E}9$ instructions is 0.9×10^9 .

1. 14. 1

$$T = \frac{50 \times 10^6 + 110 \times 10^6 + 80 \times 10^6 \times 4 + 16 \times 2 \times 10^6}{2 \times 10^9} = 0.256s$$

If the program changed to be run 1 times faster than before:

Let us Consider the total number of clocks for FP instructions is x

$$0.178 = \frac{50 \times 10^6 x + 110 \times 10^6 + 80 \times 10^6 \times 4 + 16 \times 2 \times 10^6}{2 \times 10^9}$$

$$x = -4.12$$

$$\therefore x < 0$$

\therefore The CPI of FP instructions can not be improve.

1. 14. 2

$$T = \frac{50 \times 10^6 + 110 \times 10^6 + 80 \times 10^6 \times 4 + 16 \times 2 \times 10^6}{2 \times 10^9} = 0.256s$$

If the program changed to be run 1 times faster than before:

Let us Consider the total number of clocks for FP instructions is x

$$0.178 = \frac{50 \times 10^6 + 110 \times 10^6 + 80 \times 10^6 \times x + 16 \times 2 \times 10^6}{2 \times 10^9}$$

$$x = 0.8$$

$$\therefore x > 0$$

\therefore The CPI of L/S instructions can be improved by 0.8 times if the programs to be run two times faster.

1. 14. 3

$$T_2 = \frac{0.6 \times 50 \times 10^6 + 0.6 \times 110 \times 10^6 + 0.7 \times 80 \times 10^6 \times 4 + 0.7 \times 16 \times 2 \times 10^6}{2 \times 10^9} = 0.17s$$

$$\frac{0.256}{0.17} = 1.505$$

\therefore the improvement of the execution time of the program is 1.505 times.