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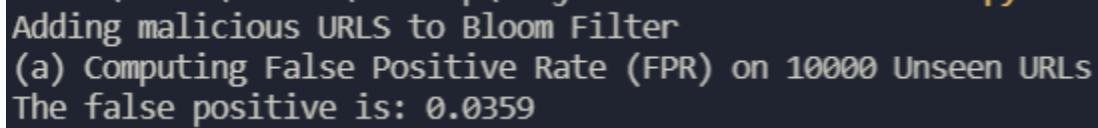
## Project III Report

### Instructions to compile/run program

Run the following in the command line terminal:

```
python bloom_filter.py
```

### Screenshot of program output



```
Adding malicious URLs to Bloom Filter
(a) Computing False Positive Rate (FPR) on 10000 Unseen URLs
The false positive is: 0.0359
```

### Problem b)

Let's compare this approach to using a typical Set data structure. Google wants to store 1 million URLs, with each URL taking (on average) 25 bytes. How much space (in MB, 1 MB = 1 million bytes) is required if we store all the elements in a set? How much space (in MB) is required if we store all the elements in a bloom filter with  $k = 10$  hash functions and  $m = 800,000$  buckets? Recall that 1 byte = 8 bits.

### Answer to b)

Set:

25 bytes \*  $1 \cdot 10^6$  URLs = 25 MB

Bloom Filter:

$k \cdot m / 8 = 10 \text{ hash functions} \cdot 800,000 \text{ buckets} / 8 = 1 \text{ MB}$

### Program Description

This program implements a bloom filter using hash functions for quick lookup time to check if an element is added to the filter or not. The `add()` and `contains()` functions are implemented by me. Since some elements may map to bits that are already set by the hash functions, there can be false positives reported by the `contains()` function.

The `add()` function adds an element to the Bloom filter by setting the bits in the array to true using the hash functions.

The `contains()` function just checks if the bits for an element are set to check if the element might be added or not. If `contains()` returns false, the element is definitely not in the set. If `contains()` return true, the element might be in the set, but there is a chance of a false positive.