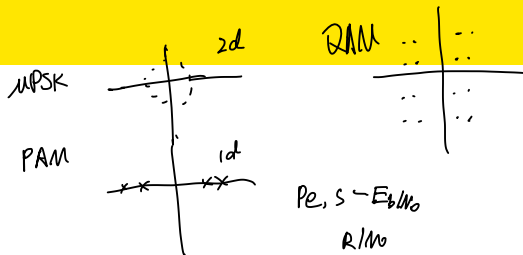


Lecture 12



Goals

- Be able to determine bandwidth efficiency and energy efficiency of orthogonal signals.
- Be able to synthesize different types of orthogonal signals.

Orthogonal Signals

- A set of signals $\{\varphi_i(t) : 0 \leq t \leq T, 0 \leq i \leq M-1\}$ are said to be orthogonal (over the interval $[0, T]$) if

$$\int_0^T \varphi_i(t) \varphi_j(t) dt = 0, \quad i \neq j.$$

- In most cases the signals will have the same energy and it is convenient to normalize the signals to unit energy. A set of signals $\{\varphi_i(t) : 0 \leq t \leq T, 1 \leq i \leq M\}$ are said to be orthonormal (over the interval $[0, T]$) if

$$\int_0^T \varphi_i(t) \varphi_j(t) dt = \begin{cases} 0, & i \neq j \\ 1, & i = j. \end{cases}$$

Orthogonal Signals

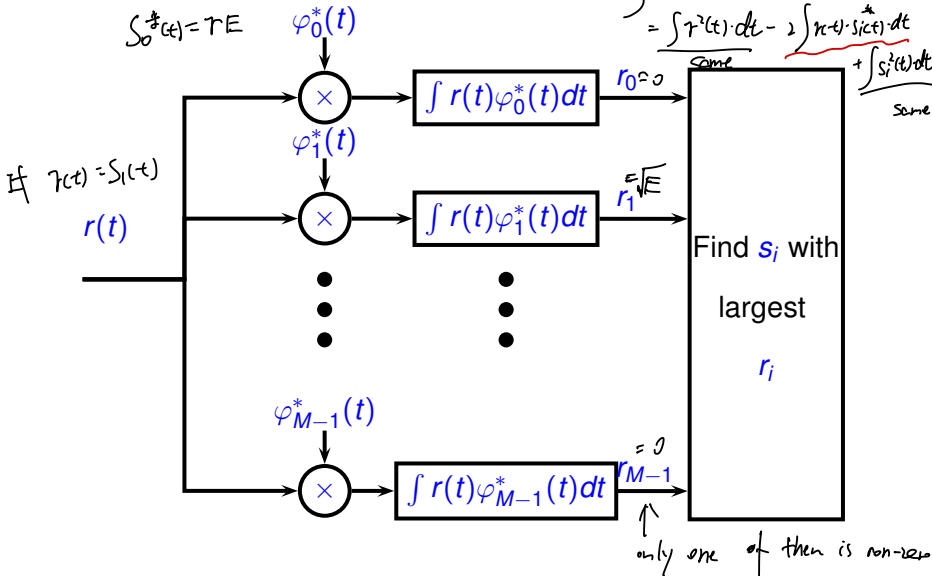
- The set of orthogonal signals can be described by

$$\begin{aligned}
 s_0(t) &= \sqrt{E} \varphi_0(t) & \int s_i^2(t) \cdot dt &= E \\
 s_1(t) &= \sqrt{E} \varphi_1(t) \\
 s_2(t) &= \sqrt{E} \varphi_2(t) \\
 s_{M-1}(t) &= \sqrt{E} \varphi_{M-1}(t)
 \end{aligned}$$

- Clearly we need $N = M$ orthonormal signals to represent M orthogonal signals.
- Orthonormal signals are also equal energy signals (namely energy 1).

Orthogonal Signals Demodulation

$\min_{r(t)} d_E^2(r(t), s_i(t))$
 $= \int (r(t), s_i(t))^2 dt$ *maximum this*
 $= \int r^2(t) dt - 2 \int r(t) \cdot s_i^*(t) dt + \int s_i^2(t) dt$
same *same*



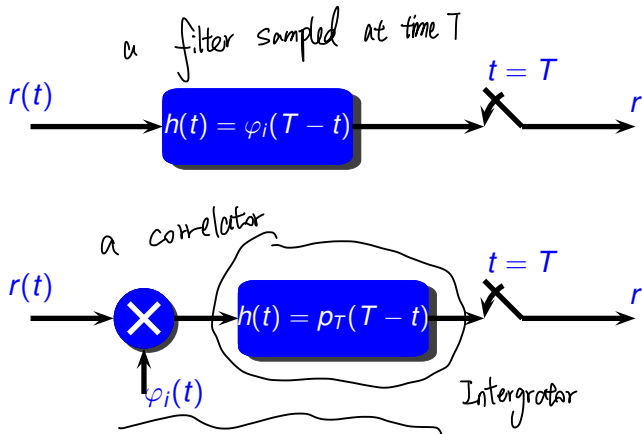
Correlation vs. Filtering

- Assume $\varphi_i^*(t)$ is time limited to the interval $[0, T]$.
- The receiver needs to compute $\int r(t)\varphi_i^*(t)dt$.
- A filter with input $r(t)$ and impulse response $h(t) = \varphi_i^*(T - t)$ sampled at time T has output

$$\begin{aligned}
 r_i &= \int h(T - t)r(t)dt = \int \varphi_i^*(T - (T - t))r(t)dt \\
 &= \int \varphi_i^*(t)r(t)dt
 \end{aligned}$$

- So either a correlator whereby the received signal is correlated with the orthonormal signal can be used to obtain r_i OR a matched filter with impulse response $h(t) = \varphi_i^*(T - t)$ which is sampled at time $t = T$ can be used to obtain r_i .

Correlation vs. Filtering



Orthogonal Modulation Performance

- Suppose signal s_i is transmitted ($0 \leq i \leq M-1$).
- The output of these filters is given by

$$r_m = \begin{cases} \overset{\text{noise along}}{\eta_m}, & m \neq i \\ \sqrt{E} + \eta_m, & m = i \end{cases}$$

signal plus noise

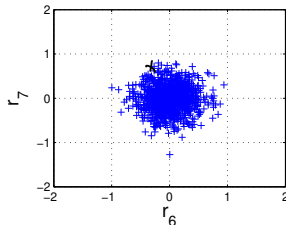
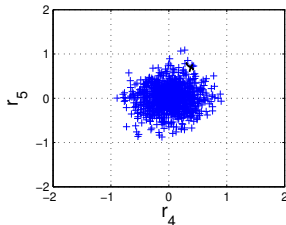
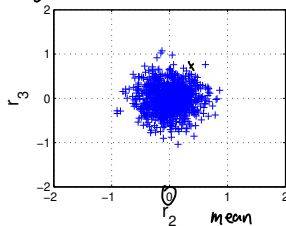
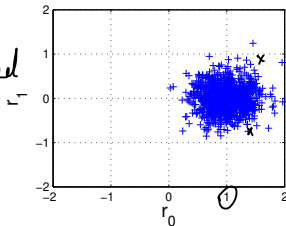
where $\{\eta_m, m = 0, 1, 2, \dots, M-1\}$ is a sequence of independent, identically distributed Gaussian random variables with mean zero and variance $N_0/2$.

- If $M = 4$ and $r = (r_0, \dots, r_3) = (0.6, -0.3, \overset{\text{largest}}{1.2}, -0.7)$ then the decision is that signal s_2 was sent because $r_2 = \max(r_i, i = 0, \dots, 3)$

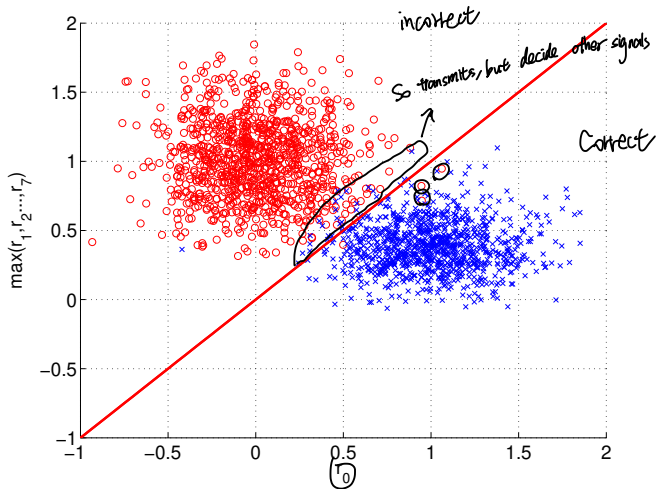
Simulation $M = 8$ (s_0 sent, $E_b/N_0 = 3\text{dB}$)

if s_0 is trans, Receiver will be correct if $r_0 > \max(r_1, r_2, r_3, r_4, r_5, r_6, r_7)$

s_0 is transmitted

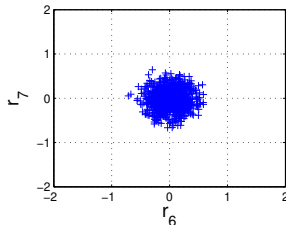
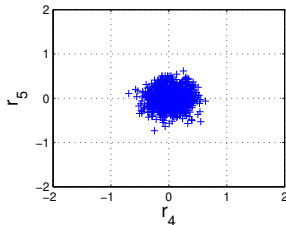
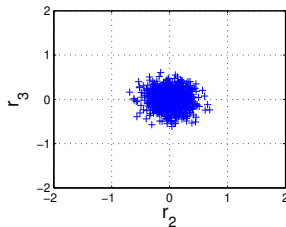
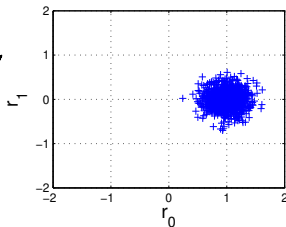


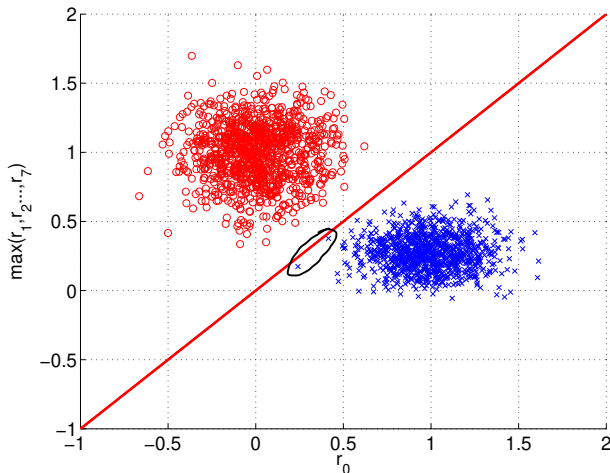
Simulation $M = 8$ (s_0 sent, $E_b/N_0 = 3\text{dB}$)



Simulation $M = 8$ (s_0 sent, $E_b/N_0 = 6\text{dB}$)

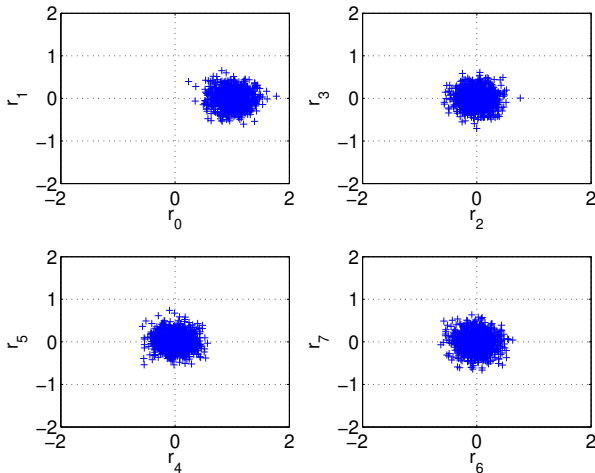
s_0 is transmitted



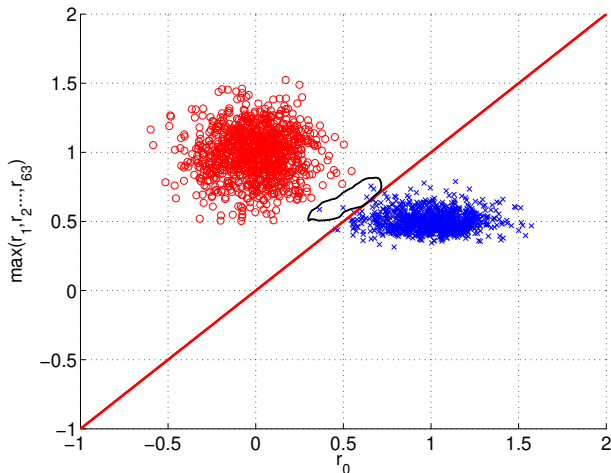
Simulation $M = 8$ (s_0 sent, $E_b/N_0 = 6\text{dB}$)

Simulation $M = 64$ (s_0 sent, $E_b/N_0 = 3\text{dB}$)

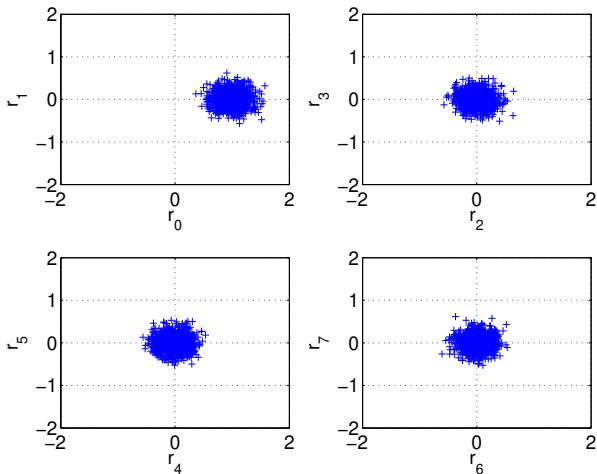
but just show 8 of them



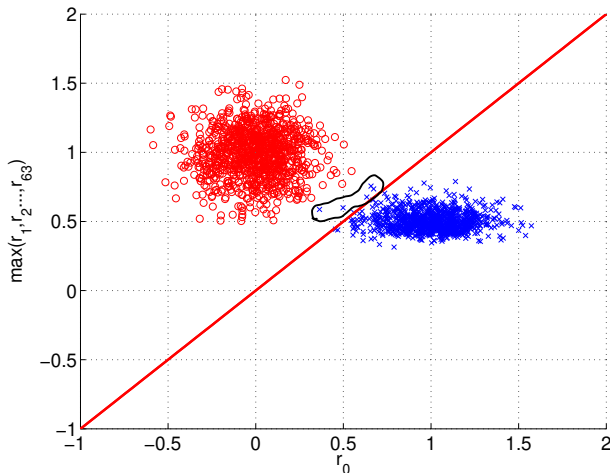
Simulation $M = 64$ (s_0 sent, $E_b/N_0 = 3\text{dB}$)



Simulation $M = 256$ (s_0 sent, $E_b/N_0 = 3\text{dB}$)



Simulation $M = 256$ (s_0 sent, $E_b/N_0 = 3\text{dB}$)



Error Probability

$$\underline{N = M}$$

To determine the probability of error we need to determine the probability that the filter output corresponding to the signal present is smaller than one of the other filter outputs.

$$\begin{aligned}
 P_{e,0} &= 1 - P(r_0 > r_1, r_0 > r_2, \dots, r_0 > r_{M-1} | s_0 \text{ transmitted}) \\
 &= P\{r_0 < \max(r_1, \dots, r_{M-1}) | s_0 \text{ trans}\} \\
 &= 1 - P\{r_1 < r_0, r_2 < r_0, r_3 < r_0, \dots, r_{M-1} < r_0 | s_0 \text{ trans}\} \\
 &= 1 - P\{\eta_1 < \sqrt{E} + \eta_0, \eta_2 < \sqrt{E} + \eta_0, \dots, \eta_{M-1} < \sqrt{E} + \eta_0\} \\
 &\xrightarrow{\text{event}} = 1 - E[P\{\eta_1 < \sqrt{E} + \eta_0, \eta_2 < \sqrt{E} + \eta_0, \dots, \eta_{M-1} < \sqrt{E} + \eta_0 | \eta_0\}] \quad \text{this event is dependent!} \\
 &= 1 - \int P\{\eta_1 < \sqrt{E} + x, \eta_2 < \sqrt{E} + x, \dots, \eta_{M-1} < \sqrt{E} + x | \eta_0 = x\} f_{\eta_0}(x) dx \\
 &= 1 - \int P\{\eta_1 < \sqrt{E} + x\} P\{\eta_2 < \sqrt{E} + x\} \dots P\{\eta_{M-1} < \sqrt{E} + x\} f_{\eta_0}(x) dx
 \end{aligned}$$

$\eta_0, \eta_1, \dots, \eta_{M-1}$ are independent

$$r_0 = \sqrt{E} + \eta_0$$

$$r_1 = \eta_1$$

$$\vdots$$

$$r_{M-1} = \eta_{M-1}$$

Error Probability

$$\sigma = \sqrt{\frac{N_0}{2}}$$

$$P_{e,0} = 1 - \int_{-\infty}^{\infty} \Phi^{M-1}\left(\frac{\sqrt{E} + x}{\sqrt{N_0/2}}\right) f_{\eta_0}(x) dx$$

$$= 1 - \int_{-\infty}^{\infty} \Phi^{M-1}\left(\frac{\sqrt{E} + x}{\sqrt{N_0/2}}\right) \frac{1}{\sqrt{\pi N_0}} \exp\left\{-\frac{x^2}{N_0}\right\} dx$$

$$\frac{y^2}{2} = \frac{x^2}{N_0}$$

$$y = \frac{x}{\sqrt{N_0/2}}$$

integration by parts ↪

$$= 1 - \int_{-\infty}^{\infty} \Phi^{M-1}(\sqrt{2E/N_0} + y) \frac{1}{\sqrt{2\pi}} \exp\left\{-\frac{y^2}{2}\right\} dy$$

$$= (M-1) \int_{-\infty}^{\infty} \Phi^{M-2}(\sqrt{2E/N_0} + y) \frac{1}{\sqrt{2\pi}} e^{-(y + \sqrt{2E/N_0})^2/2} \Phi(y) dy$$

$$= \frac{M-1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \Phi\left(z - \sqrt{\frac{2E}{N_0}}\right) \Phi^{M-2}(z) e^{-z^2/2} dz$$

Error Probability

$\Phi(u)$ is the distribution function of a zero mean, variance 1, Gaussian random variable given by

$$\Phi(u) = \frac{1}{2\pi} \int_{-\infty}^u e^{-x^2/2} dx.$$

The last step in the derivation is obtained by using the integration by parts formula.

$$\int u dv = uv - \int v du$$

where $u = \Phi^{M-1}(\sqrt{2E/N_0} + y)$ and $dv = \frac{1}{\sqrt{2\pi}} \exp\{-y^2/2\} dy$.

Performance of Orthogonal Signals

$$P_{e,s} = \frac{M-1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \Phi(u - \sqrt{\frac{2E}{N_0}}) \phi^{M-2}(u) e^{-u^2/2} du$$

- The formula for $P_{e,s}$ can be interpreted as the expected value of a conditional probability where
 - we condition on one of the incorrect correlator outputs being a value (say u),
 - the correct output being less than u and
 - $M-2$ incorrect correlator outputs also being less than u .
- Since there are $M-1$ possible incorrect outputs we need to multiply this by $M-1$.

Energy per information bit

Normally a communication engineer is more concerned with the energy transmitted per bit rather than the energy transmitted per signal, E . If we let E_b be the energy transmitted per bit then these are related as follows

$$E_b = \frac{E \text{ joules/signal}}{\log_2 M \text{ bit/signal}} \cdot \log_2 M \text{ bit/signal}$$

$M=64 \Rightarrow 6 \text{ bits}$
 $M=32 \Rightarrow 5 \text{ bits}$

Symbol Error Probability

$$P_{e,s} = \frac{M-1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \Phi(u - \sqrt{\frac{2E_b \log_2(M)}{N_0}}) \phi^{M-2}(u) e^{-u^2/2} du$$

use L'Hospital rules

As $M \rightarrow \infty$ the symbol error probability of orthogonal signal sets approaches 1 or 0.

$$\lim_{M \rightarrow \infty} P_{e,s} = \begin{cases} 1 & E_b/N_0 < \ln(2) \\ 0 & E_b/N_0 > \ln(2) \end{cases}$$

\Rightarrow Arbitrarily small error prob ($\rightarrow 0$) if we make m large
as long as $\frac{E_b}{N_0} > \ln(2) = 0.69$ (-1.59 dB)

$$\text{bits/dimension} = \frac{\log_2 M}{M} \rightarrow 0$$

The rate of communications (bits/dimension) as $M \rightarrow \infty$ goes to 0.

$$\lim_{M \rightarrow \infty} \frac{\log_2(M)}{M} = 0$$

That is, orthogonal signals for very large M achieves Shannon's fundamental limit of performance at the point where the rate is zero. That is at a bandwidth efficiency of zero, orthogonal signals have arbitrarily small error probability provided $E_b/N_0 > \ln(2)$.

(-1.59 dB)

Bounds on Performance of Orthogonal Signals

The symbol error probability can be upper bounded as

$$P_{e,s} \leq \begin{cases} 1, & \frac{E}{N_0} \leq \ln M \\ \exp \left\{ - \left(\sqrt{\frac{E}{N_0}} - \sqrt{\ln M} \right)^2 \right\}, & \ln M \leq \frac{E}{N_0} \leq 4 \ln M \\ \exp \left\{ - \left(\frac{E}{2N_0} - \ln M \right) \right\}, & \frac{E}{N_0} \geq 4 \ln M. \end{cases}$$

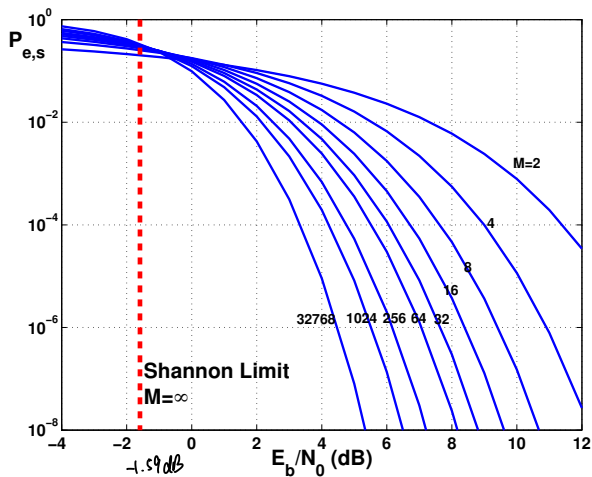
Energy per information bit

The bound on the symbol error probability can be expressed in terms of the energy transmitted per bit as

$$P_{e,s} \leq \begin{cases} 1, & \frac{E_b}{N_0} \leq \ln 2 \\ \exp_2 \left\{ -\log_2 M \left(\sqrt{\frac{E_b}{N_0}} - \sqrt{\ln 2} \right)^2 \right\}, & \ln 2 \leq \frac{E_b}{N_0} \leq 4 \ln 2 \\ \exp_2 \left\{ -\log_2 M \left(\frac{E_b}{2N_0} - \ln 2 \right) \right\}, & \frac{E_b}{N_0} \geq 4 \ln 2 \end{cases}$$

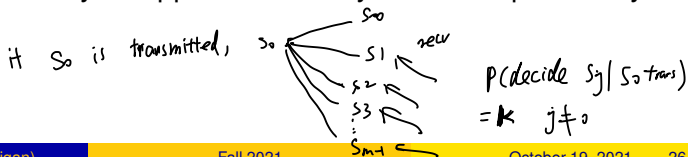
where $\exp_2\{x\}$ denotes 2^x . Note that as $M \rightarrow \infty$, $P_e \rightarrow 0$ if $\frac{E_b}{N_0} > \ln 2 = -1.59\text{dB}$.

Symbol Error Probability for Orthogonal Signals



Bit error probability

- So far we have examined the symbol error probability for orthogonal signals; the probability of deciding the wrong symbol or signal.
- Usually the number of such signals is a power of 2, e.g. 4, 8, 16, 32,
- If so then each transmission of a signal is carrying $k = \log_2 M$ bits of information.
- In this case a communication engineer is usually interested in the bit error probability as opposed to the symbol error probability.

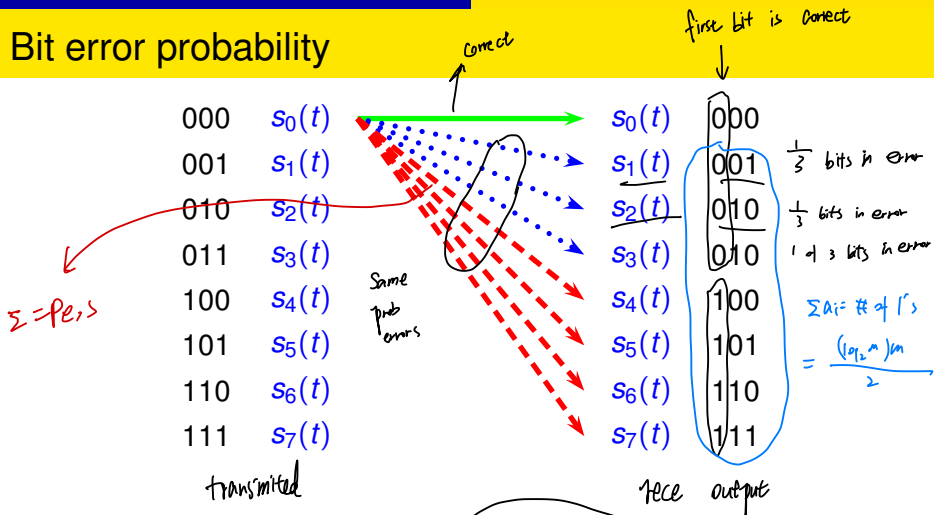


Bit error probability

- Assume signal 0 is transmitted corresponding to the data bits being (000...00).
- If an error occurs and the demodulator chooses one of the incorrect signals than each of the incorrect signals has the same probability.
- Thus the signal corresponding to data bits being (000...01) has the same probability as an error to a signal corresponding to data bits (111...11).
- If signal 0 is transmitted then there will be $M/2$ other signals that will cause a bit error in any particular bit. Thus

$$P_{e,b} = \frac{M}{2(M-1)} P_{e,s} = \frac{2^{k-1}}{2^k - 1} P_{e,s}.$$

Bit error probability



Red line (dashed) corresponds to a symbol error and a bit error for the first bit, green line (solid) corresponds to no error (symbol or bit). Blue line (dotted) corresponds to a symbol error but not a bit error (for the first bit).

Bit error probability

- Consider the case of 8 signals transmitting 3 bits of information.
- Suppose that the bits 000 are mapped to signal s_0 and 111 are mapped to s_7 .
- Then if signal s_0 is transmitted the receiver will either decide correctly or decide incorrectly.
- All the incorrect signals are equally likely to be chosen given an error is made. That is, the probability of the receiver deciding signal s_i given s_0 was transmitted is $P_{e,s}/(M-1)$ for all $i \neq 0$.
- If the receiver decides incorrectly then it has decided one of s_1, \dots, s_7 .
- If it decides s_1 then that will be mapped into bits 001 and one bit error will occur.
- If it decides s_2 then that will be mapped into bits 010 and again one bit error will occur.
- If it decides s_7 then that will be mapped to 111 and three bit errors will occur.

Bit error probability

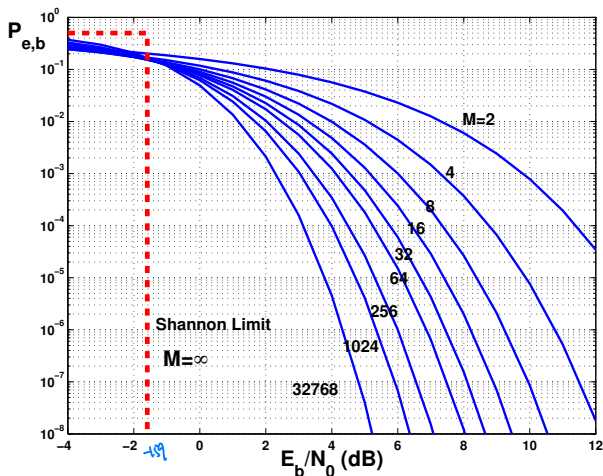
- Let a_i be the number of bit errors that occur when we decide signal s_i .
- Then the bit error probability is

$$\begin{aligned}
 P_{e,b} &= \sum_{i=1}^{M-1} \frac{a_i}{\log_2(M)} \left(\frac{P_{e,s}}{M-1} \right) && \begin{array}{l} \text{Total \# of bits that} \\ \text{are 1} \\ \text{(in error)} \end{array} \\
 &= \frac{P_{e,s}}{M-1} \frac{1}{\log_2(M)} \left(\sum_{i=1}^{M-1} a_i \right) && \text{Total \# of bits that are 1 (in error)} \\
 &= \frac{P_{e,s}}{M-1} \frac{1}{\log_2(M)} \log_2(M) M/2 \\
 &= \frac{MP_{e,s}}{2(M-1)} \approx \frac{1}{2} P_{e,s} \text{ (if } M \text{ is large)}
 \end{aligned}$$

$\rightarrow P(\text{rec dec } s_i)$
 $i=1, 2, \dots, M-1$

$(\log_2 8) \cdot 8 / 2 = 12$

Bit Error Probability for Orthogonal Signals



Asymptotic Error Probability of Orthogonal Signal Sets

- As $M \rightarrow \infty$ the symbol error probability of orthogonal signal sets approaches 1 or 0.

$$\lim_{M \rightarrow \infty} P_{e,s} = \begin{cases} 1 & E_b/N_0 < \ln(2) \\ 0 & E_b/N_0 > \ln(2) \end{cases} \quad \lim_{M \rightarrow \infty} P_{e,b} = \begin{cases} 1/2 & E_b/N_0 < \ln(2) \\ 0 & E_b/N_0 > \ln(2) \end{cases}$$

- The rate of communications (bits/dimension) as $M \rightarrow \infty$ goes to 0.

$$\lim_{M \rightarrow \infty} \frac{\log_2(M)}{M} = 0$$

- That is, orthogonal signals for very large M achieves Shannon's fundamental limit of performance at the point where the rate is zero. That is at a bandwidth efficiency of zero, orthogonal signals have arbitrarily small error probability provided $E_b/N_0 > \ln(2)$.

Orthogonal Signal Sets

- Below we define several different orthogonal signal sets.
- We will define the bandwidth of a signal set as the minimum difference in carrier frequencies between two such signal sets so that any signal from one set is orthogonal to any signal from the other set.
- We will assume that the double frequency terms are negligible relative to the baseband frequency term.

Review of Integrals of Trig Functions

$$\begin{aligned}
 \rho &= \int_0^T \sin(2\pi f_1 t) \sin(2\pi f_0 t) dt \\
 &= \frac{1}{2} \int_0^T \cos(2\pi(f_1 - f_0)t) - \cos(2\pi(f_0 + f_1)t) dt \\
 &= \frac{T}{2} \left[\frac{\sin(2\pi(f_1 - f_0)T)}{2\pi(f_1 - f_0)T} - \frac{\sin(2\pi(f_0 + f_1)T)}{2\pi(f_0 + f_1)T} \right] \\
 &= \frac{T}{2} \frac{\sin(2\pi(f_1 - f_0)T)}{2\pi(f_1 - f_0)T}
 \end{aligned}$$

This integral will be zero if $2\pi(f_1 - f_0)T = \pi, 2\pi, 3\pi, \dots$. The minimum frequency spacing is $f_1 - f_0 = 1/(2T)$.

Review of Integrals of Trig Functions

$$\begin{aligned}
 \rho &= \int_0^T \sin(2\pi f_1 t + \theta) \sin(2\pi f_0 t) dt \\
 &= \frac{1}{2} \int_0^T \cos(2\pi(f_1 - f_0)t + \theta) - \cos(2\pi(f_0 + f_1)t + \theta) dt \\
 &= \frac{T}{2} \left[\frac{\sin(2\pi(f_1 - f_0)T + \theta)}{2\pi(f_1 - f_0)T} - \frac{\sin(2\pi(f_0 + f_1)T + \theta)}{2\pi(f_0 + f_1)T} \right] \\
 &= \frac{T}{2} \left[\frac{\sin(2\pi(f_1 - f_0)T + \theta)}{2\pi(f_1 - f_0)T} - \frac{\sin(\theta)}{2\pi(f_1 - f_0)T} \right]
 \end{aligned}$$

This integral will be zero if $2\pi(f_1 - f_0)T = 2\pi, 4\pi, 6\pi, \dots$. The minimum frequency spacing is $f_1 - f_0 = 1/(T)$.

Review of Integrals of Trig Functions

$$\begin{aligned}
 \rho &= \int_0^T \cos(2\pi f_1 t) \sin(2\pi f_0 t) dt \\
 &= \frac{1}{2} \int_0^T \sin(2\pi(f_1 + f_0)t) - \sin(2\pi(f_1 - f_0)t) dt \\
 &= \frac{T}{2} \left[\frac{\cos(2\pi(f_1 - f_0)T) - 1}{2\pi(f_1 - f_0)T} - \frac{\cos(2\pi(f_0 + f_1)T) - 1}{2\pi(f_0 + f_1)T} \right] \\
 &= \frac{T}{2} \left[\frac{\cos(2\pi(f_1 - f_0)T) - 1}{2\pi(f_1 - f_0)T} \right]
 \end{aligned}$$

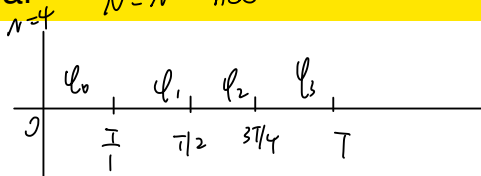
This integral will be zero if $2\pi(f_1 - f_0)T = 2\pi, 4\pi, 6\pi, \dots$. The minimum frequency spacing is $f_1 - f_0 = 1/(T)$.

Review of Integrals of Trig Functions

$$\begin{aligned}
 \rho &= \int_0^T \cos(2\pi f_1 t + \theta) \sin(2\pi f_0 t) dt \\
 &= \frac{1}{2} \int_0^T \sin(2\pi(f_1 + f_0)t + \theta) - \sin(2\pi(f_1 - f_0)t + \theta) dt \\
 &= \frac{T}{2} \left[\frac{\cos(2\pi(f_1 - f_0)T + \theta) - \cos(\theta)}{2\pi(f_1 - f_0)T} - \frac{\cos(2\pi(f_0 + f_1)T + \theta) - \cos(\theta)}{2\pi(f_0 + f_1)T} \right] \\
 &= \frac{T}{2} \left[\frac{\cos(2\pi(f_1 - f_0)T + \theta)}{2\pi(f_1 - f_0)T} - \frac{\cos(\theta)}{2\pi(f_1 - f_0)T} \right]
 \end{aligned}$$

This integral will be zero if $2\pi(f_1 - f_0)T = 2\pi, 4\pi, 6\pi, \dots$. The minimum frequency spacing is $f_1 - f_0 = 1/(T)$.

A. Time-orthogonal

 $N=M$ here

$$\varphi_i(t) = \begin{cases} \sqrt{\frac{2N}{T}} \sin(2\pi f_0 t), & \frac{iT}{N} \leq t < (i+1)T/N \\ 0, & \text{elsewhere} \end{cases}$$

$$i = 0, 1, \dots, N,$$

It is clear that $\varphi_i(t)$ and $\varphi_j(t)$ are orthogonal in time (they do not overlap in time) if

N dimension in bandwidth $N/2T$ time T

A. Time-orthogonal: frequency multiplexing

$\phi_0^{(0)}(t), \phi_1^{(1)}$ are non-zero in $[0, T/N]$

2 signal sets

$$\left\{ \begin{aligned} \phi_i^{(0)}(t) &= \begin{cases} \sqrt{\frac{2N}{T}} \sin(2\pi f_0 t), & \frac{iT}{N} \leq t < (i+1)T/N \\ 0, & \text{elsewhere} \end{cases} \\ \phi_i^{(1)}(t) &= \begin{cases} \sqrt{\frac{2N}{T}} \sin(2\pi f_1 t), & \frac{iT}{N} \leq t < (i+1)T/N \\ 0, & \text{elsewhere} \end{cases} \end{aligned} \right. \quad \text{orthogonal}$$

$i = 0, 1, \dots, N,$

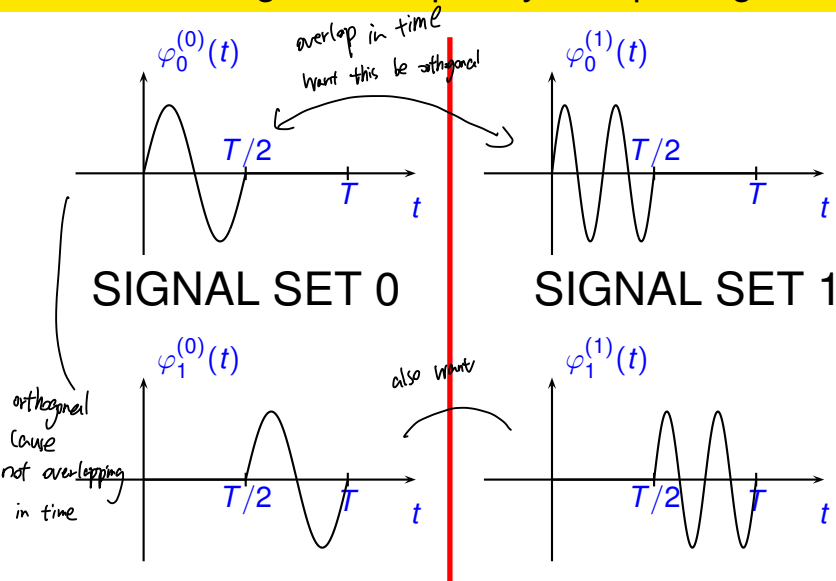
The duration of these signals is T/N . For orthogonality of $\phi_i^{(0)}(t)$ and $\phi_j^{(1)}(t)$ we need

$$f_1 - f_0 = \frac{1}{2(T/N)} = \frac{N}{2T}$$

N dimension in bandwidth
 $\frac{N}{2T}$ in time T

Thus the bandwidth of this signal set is $W = \left[\frac{(n+1)N}{2T} - \frac{nN}{2T} \right] = \frac{N}{2T}$

A. Time-orthogonal: frequency multiplexing



B. Time-orthogonal quadrature-phase

$$\varphi_{2i}(t) = \begin{cases} \sqrt{\frac{2N}{T}} \sin(2\pi f_0 t), & \frac{2iT}{N} \leq t < \frac{2(i+1)T}{N} \\ 0, & \text{elsewhere} \end{cases}$$

$$\varphi_{2i+1}(t) = \begin{cases} \sqrt{\frac{2N}{T}} \cos(2\pi f_0 t), & \frac{2iT}{N} \leq t < \frac{2(i+1)T}{N} \\ 0 & \text{elsewhere} \end{cases}$$

$$i = 0, 1, \dots, \frac{N}{2} - 1, \quad N \text{ even}, \quad f_0 = n \frac{N}{2T},$$

Duration = $T/(2N)$.

$$W = \left[\frac{(n+1)N}{2T} - \frac{nN}{2T} \right] = \frac{N}{2T}$$

C. Frequency-orthogonal

$$\varphi_i(t) = \sqrt{\frac{2E}{T}} \sin[2\pi(f_0 + \frac{i}{2T})t], \quad 0 \leq t \leq T$$

$$i = 0, 1, \dots, N-1, \quad f_0 = \frac{nN}{2T}.$$

$$\begin{aligned} W &= \left[\frac{(n+1)N}{2T} - \frac{nN}{2T} \right] \\ &= \frac{N}{2T} \end{aligned}$$

D. Frequency-orthogonal quadrature-phase

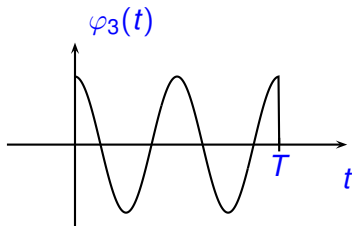
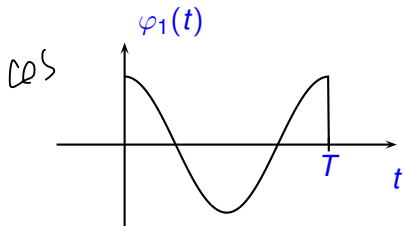
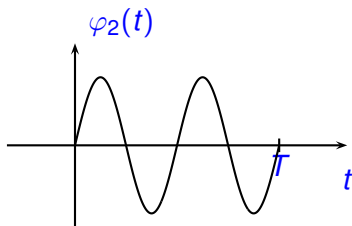
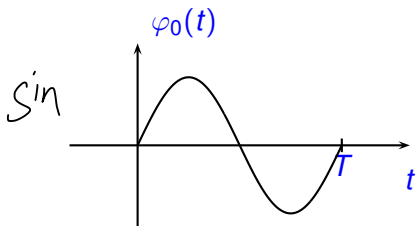
$$\varphi_{2i}(t) = \sqrt{\frac{2E}{T}} \sin[2\pi(f_0 + \frac{i}{T})t] \quad 0 \leq t < T$$

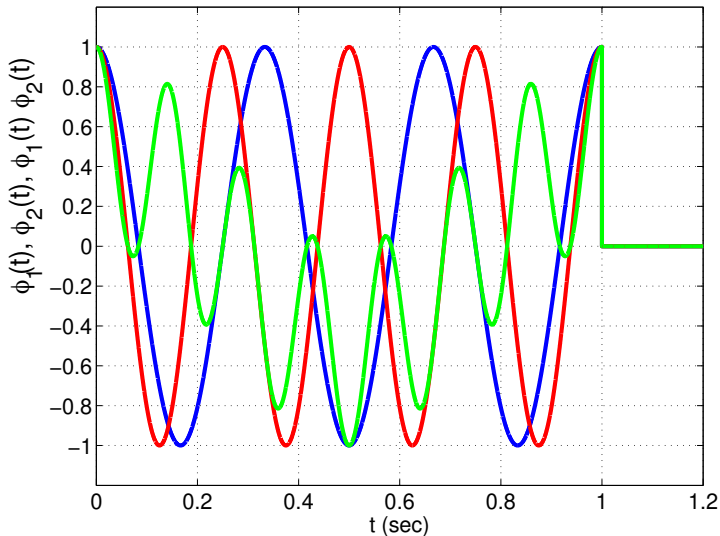
$$\varphi_{2i+1}(t) = \sqrt{\frac{2E}{T}} \cos[2\pi(f_0 + \frac{i}{T})t], \quad 0 \leq t \leq T$$

$$f_0 = \frac{nN}{2T}.$$

$$\begin{aligned} W &= \left[\frac{(n+1)N}{2T} - \frac{nN}{2T} \right] \\ &= \frac{N}{2T} \end{aligned}$$

D. Frequency-orthogonal quadrature-phase



Example ($f_c = 1$, $T = 3$)

E. Hadamard-Walsh Construction

The last construction of orthogonal signals is done via the Hadamard Matrix. The Hadamard matrix is an N by N matrix with components either +1 or -1 such that every pair of distinct rows are orthogonal. We show how to construct a Hadamard when the number of signals is a power of 2 (which is often the case).

Begin with a two by two matrix

$$H_2 = \begin{bmatrix} +1 & +1 \\ +1 & -1 \end{bmatrix} \begin{matrix} \rightarrow \phi_0 \\ \rightarrow \phi_1 \end{matrix} \quad \sum \phi_0(n) \phi_1(n) = 0$$

Then use the recursion

$$H_{2^l} = \begin{bmatrix} +H_{2^{l-1}} & +H_{2^{l-1}} \\ +H_{2^{l-1}} & -H_{2^{l-1}} \end{bmatrix}.$$

E. Hadamard-Walsh Construction

Now it is easy to check that distinct rows in these matrices are orthogonal. The i -th modulated signal is then obtained by using a single (arbitrary) waveform N times in nonoverlapping time intervals and multiplying by the j -th repetition of the waveform by the j th component of the i -th row of the matrix.

E. Hadamard-Walsh Construction

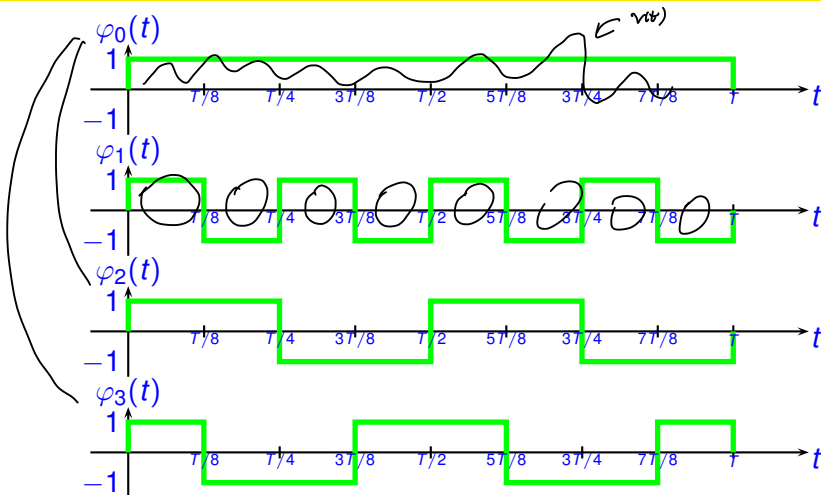
Example ($N = 4$):

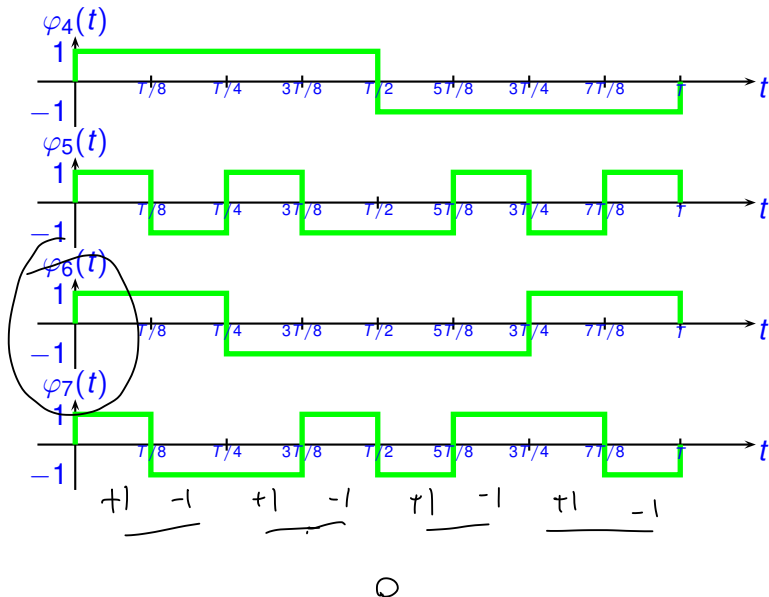
$$\begin{aligned}
 H_4 &= \begin{bmatrix} \underline{H_2} & \underline{H_2} \\ \underline{H_2} & \underline{-H_2} \end{bmatrix} \\
 &= \begin{bmatrix} +1 & +1 & +1 & +1 \\ +1 & -1 & +1 & -1 \\ +1 & +1 & -1 & -1 \\ +1 & -1 & -1 & +1 \end{bmatrix}.
 \end{aligned}$$

E. Hadamard-Walsh Construction, $M = 8$

$$\begin{aligned}
 H_8 &= \begin{bmatrix} H_4 & H_4 \\ H_4 & -H_4 \end{bmatrix} \\
 &= \begin{bmatrix} H_2 & H_2 & H_2 & H_2 \\ H_2 & -H_2 & H_2 & -H_2 \\ H_2 & H_2 & -H_2 & -H_2 \\ H_2 & -H_2 & -H_2 & H_2 \end{bmatrix} \\
 &= \begin{bmatrix} +1 & +1 & +1 & +1 & +1 & +1 & +1 & +1 \\ +1 & -1 & +1 & -1 & +1 & -1 & +1 & -1 \\ +1 & +1 & -1 & -1 & +1 & +1 & -1 & -1 \\ +1 & -1 & -1 & +1 & +1 & -1 & -1 & +1 \\ +1 & +1 & +1 & +1 & -1 & -1 & -1 & -1 \\ +1 & -1 & +1 & -1 & -1 & +1 & -1 & +1 \\ +1 & +1 & -1 & -1 & -1 & -1 & +1 & +1 \\ +1 & -1 & -1 & +1 & -1 & +1 & +1 & -1 \end{bmatrix}.
 \end{aligned}$$

E. Hadamard-Walsh Construction, $M = 8$





Processing of Hadamard Generated Orthogonal Signals

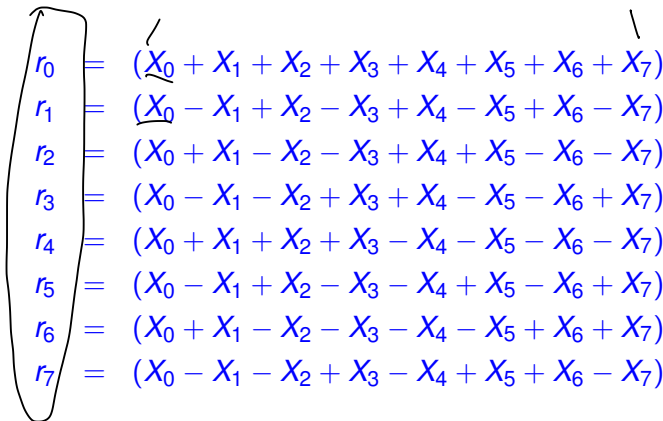
$$\int r(t) \cdot \varphi_i(t) \cdot dt \quad i = 0, 1, \dots, N-1$$

Let

$$X_i = \int_{iT/N}^{(i+1)T/N} r(t) dt, \quad i = 0, 1, \dots, N-1.$$

Then the correlations of $r(t)$ and $\varphi_i(t)$ for $i = 0, 1, \dots, 7$ can be calculated from X_0, X_1, \dots, X_7 .

Processing of Hadamard Generated Orthogonal Signals

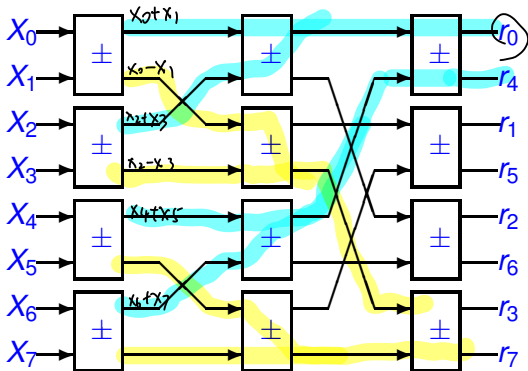


$$\begin{aligned}
 r_0 &= (\underline{X_0} + X_1 + X_2 + X_3 + X_4 + X_5 + X_6 + X_7) \\
 r_1 &= (\underline{X_0} - X_1 + X_2 - X_3 + X_4 - X_5 + X_6 - X_7) \\
 r_2 &= (X_0 + X_1 - X_2 - X_3 + X_4 + X_5 - X_6 - X_7) \\
 r_3 &= (X_0 - X_1 - X_2 + X_3 + X_4 - X_5 - X_6 + X_7) \\
 r_4 &= (X_0 + X_1 + X_2 + X_3 - X_4 - X_5 - X_6 - X_7) \\
 r_5 &= (X_0 - X_1 + X_2 - X_3 - X_4 + X_5 - X_6 + X_7) \\
 r_6 &= (X_0 + X_1 - X_2 - X_3 - X_4 - X_5 + X_6 + X_7) \\
 r_7 &= (X_0 - X_1 - X_2 + X_3 - X_4 + X_5 + X_6 - X_7)
 \end{aligned}$$

Processing of Hadamard Generated Orthogonal Signals

$M=8$

Using in
2G and
3G system



$$r_7 = X_0 - X_1 - (X_2 - X_3) - [(X_4 - X_5) - (X_6 - X_7)]$$

Efficient method of correlating for H-W orthogonal signals

Bandwidth of Orthogonal Signals

If we define bandwidth of M signals as minimum frequency separation between two such signal sets such that any signal from one signal set is orthogonal to every signal from a frequency adjacent signal set then for all of these examples of M orthogonal signals the bandwidth is

$$W = \frac{M}{2T}.$$

$$\frac{N}{WT} = 2 \quad \frac{\text{dimension}}{\text{Hz}}$$

Orthogonal Signals in Bandwidth W , Time T .

There are $N = 2WT$ orthogonal signals in bandwidth W and time duration T .

Equivalently, the bandwidth efficiency of a modulation technique (in bits/second/Hz) can be computed from the rate (in bits/dimension) by multiplying by 2.

$$2 \frac{\log_2 M}{M}$$

$$R/W(\text{bits/second/Hz}) = \mathcal{R}(\text{bits/dimension}) \times (2 \text{ dimensions/second/Hz})$$

Capacity

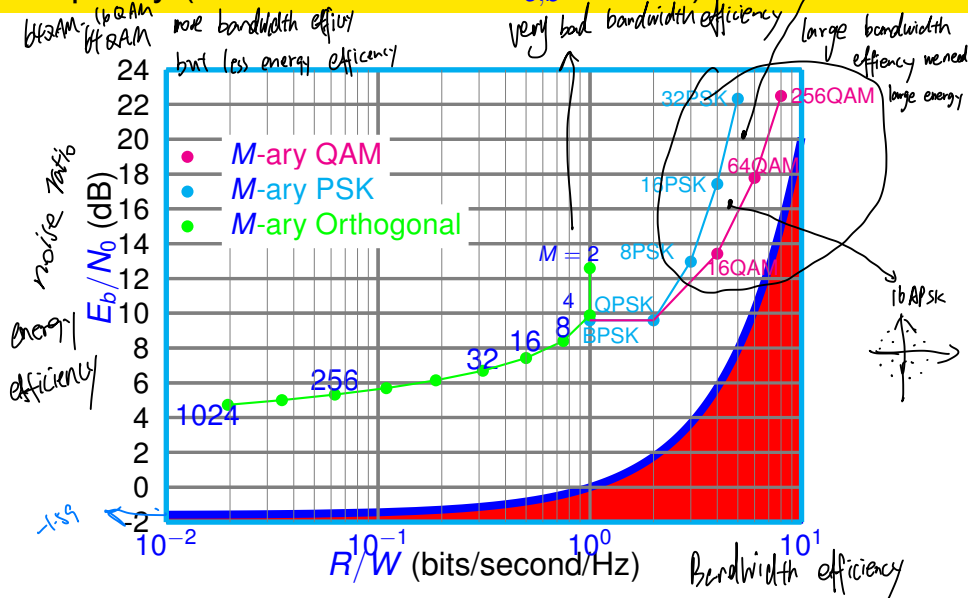
$$W = \frac{M}{2T}$$

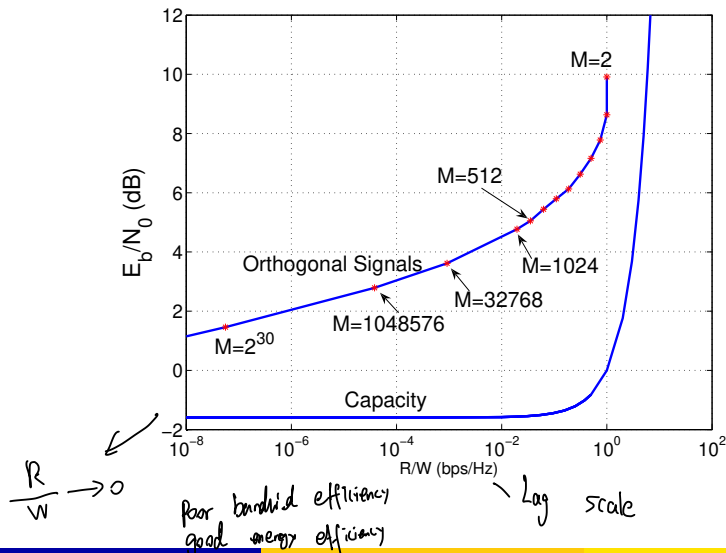
The bandwidth of a set of M orthogonal signals is $W = M/(2T)$. If we transmit $\log_2(M)$ bits in T seconds then the data rate is $R = \log_2(M)/T$. The bandwidth efficiency is

$$R = \frac{\log_2 M}{T}$$

$$R/W = \frac{\log_2(M)/T}{M/2T} = \frac{2 \log_2(M)}{M}$$

M	R/W
2	1.00000000
4	1.00000000
8	0.75000000
16	0.50000000
32	0.31250000
64	0.18750000
128	0.10937500
256	0.06250000
512	0.03515625
1024	0.01953125

Capacity (vs. Modulation at $P_{e,b} = 10^{-5}$)

Capacity vs. Orthogonal Signals @ $P_{e,b} = 0.001$ 

• opt rec: Find signal closest to rec signal
 $\min_i d \in (r, s_i)$

• opt rec if all signal have the same energy is
 find signal with max correlation $\max (r(t), s_i(t))$

• mpsk, QAM, orthogonal there is a known formula $P_{e,s}$ $P_{e,b}$

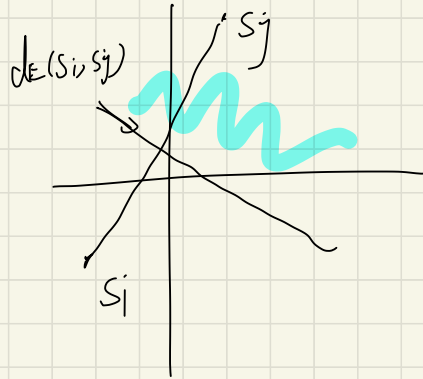
• For other modulation scheme $P_{e,s}$, $P_{e,b}$ is
 hard to calculate

Union bound:

$$P_{e,s} \leq \frac{1}{M} \sum_{i=0}^{N-1} \sum_{j \neq i} P_e(s_i \rightarrow s_j)$$

Pairwise
error

$$\Rightarrow P_e(s_i \rightarrow s_j) = Q\left(\frac{d_i(s_i, s_j)}{\sigma}\right)$$



• M-ary orthogonal signal have

$P_{e,s}, P_{e,b} \rightarrow 0$ as $M \rightarrow \infty$ if $\frac{E_b}{N_0} = 1/n \text{ dB}$
 $= -1.59 \text{ dB}$

but rate (bit/second) $\rightarrow 0$ as $M \rightarrow \infty$

$$\bullet \frac{R}{W} \text{ (bits/second)} = \underset{\substack{\downarrow \\ \text{bits/dimension}}}{\geq r}$$