

#### Ranking the importance of webpages

- Idea 1: View in-links as votes
- Idea 2: Link from an important page counts more
- Idea 3: Each page's vote is evenly split among its out-links.
- Idea 4: Each link's vote is proportional to the importance of the page.



#### Ranking the importance of webpages

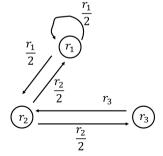
- Idea 1: View in-links as votes
- Idea 2: Link from an important page counts more
- Idea 3: Each page's vote is evenly split among its out-links.
- Idea 4: Each link's vote is proportional to the importance of the page.

Consider a network with three webpages, where  $r_i$  is the importance score of webpage i and the direct edges are hyperlinks.

$$r_1 = \frac{r_1}{2} + \frac{r_2}{2}$$

$$r_2 = \frac{r_1}{2} + r_3$$

$$r_3 = \frac{r_2}{2}$$



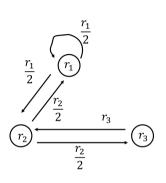


Consider a network with three webpages, where  $r_i$  is the importance score of webpage i and the direct edges are hyperlinks.

$$r_1 = \frac{r_1}{2} + \frac{r_2}{2}$$

$$r_2 = \frac{r_1}{2} + r_3$$

$$r_3 = \frac{r_2}{2}$$

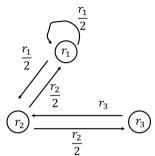


- There is no unique solution.
- Add an additional constraint  $r_1 + r_2 + r_3 = 1$ .
- Unique solution:  $r_1 = \frac{2}{5}$ ,  $r_2 = \frac{2}{5}$ , and  $r_3 = \frac{1}{5}$ .



#### Markov chain perspective

- Imagine a random Web surfer.
- At any time t, surfer is on web i. At time t+1 the surfer picks an out-link uniformly at random and goes to the next web.
- $p_i(t)$ : probability that the surfer is at web i at time t.



$$p_i(t+1) = \sum_{j:j o i} rac{1}{d_j} p_j(t) \Rightarrow p(t+1) = p(t) egin{pmatrix} rac{1}{2} & rac{1}{2} & 0 \ rac{1}{2} & 0 & rac{1}{2} \ 0 & 1 & 0 \end{pmatrix}$$

 $\pi = (\frac{2}{5}, \frac{2}{5}, \frac{1}{5})$  is the stationary distribution of this Markov chain.



#### Consider a random process

$$X_n = X_0 + W_1 + \cdots + W_n,$$

where  $W_1, W_2, \cdots$  are independent random variables with  $\mathbb{P}(W_i = 1) = p$  and  $\mathbb{P}(W_i = -1) = 1 - p$ .

#### Gambler's ruin problem

Assume  $X_0 = k$ . The random process terminates when  $X_n = 0$  (the gambler is ruined) or  $X_n = b$  (the gambler is successful). Define  $S_b$  to be the event that the gambler is successful without being ruined first. What is  $\mathbb{P}(S_b)$ ?

Define 
$$s_k = \mathbb{P}(S_b|X_0 = k)$$
. Clearly  $s_0 = 0$  and  $s_b = 1$ .



It is a discrete-time Markov chain with two absorbing states (i.e. not irreducible).



It is a discrete-time Markov chain with two absorbing states (i.e. not irreducible).

#### Key idea

Conditioning on the value of the first step  $W_1$ , the conditional probability of success, i.e.  $\mathbb{P}(S_b|X_0=k,W_1=w)$  is the same as the unconditional probability of success from initial condition k+w, i.e.,  $\mathbb{P}(S_b|W_0=k+w)$ .

$$s_k = p s_{k+1} + (1-p) s_{k-1}$$

Boundary conditions:  $s_0 = 0$  and  $s_b = 1$ .



$$s_k = ps_{k+1} + (1-p)s_{k-1}$$
 with  $s_0 = 0, b = 1$ .

**Case 1:** 
$$p = \frac{1}{2}$$

$$s_k = \frac{1}{2} \left( s_{k+1} + s_{k-1} \right) \Rightarrow s_k = \frac{k}{b}.$$



$$s_k = ps_{k+1} + (1-p)s_{k-1}$$
 with  $s_0 = 0,_b = 1$ .

**Case 2:**  $p \neq \frac{1}{2}$ 

$$s_k = \frac{1 - \left(\frac{1-p}{p}\right)^k}{1 - \left(\frac{1-p}{p}\right)^b}.$$

For  $p > \frac{1}{2}$ ,

$$\lim_{b \to \infty} s_k = 1 - \left(\frac{1-p}{p}\right)^k$$
.

The probability of ruin decreases geometrically with the initial wealth k.



Assume the gambler wins with probability 0.6 and loses with probability 0.4. We assume the gambler can bit any nonzero value (i.e. any positive real number).

How to bet to maximize the wealth?

Question: Bet a fixed amount or a fixed fraction?



Assume the gambler wins with probability 0.6 and loses with probability 0.4. We assume the gambler can bit any nonzero value (i.e. any positive real number).

The gambler bets  $\alpha$  fraction of the wealth each time. Define  $Z_n$  such that

$$\mathbb{P}(Z_n = 1 + \alpha) = 0.6 \quad \text{and} \quad \mathbb{P}(Z_n = 1 - \alpha) = 0.4.$$

Then

$$W_T = W_0 \prod_n Z_n$$
 or  $\log W_T = \log W_0 + \sum_n \log Z_n$ .

From LLN, we use the following approximation:

$$\frac{\log W_T}{T} \rightarrow 0.6 \log (1 + \alpha) + 0.4 \log (1 - \alpha) \quad (a.s.)$$



$$\max_{\alpha} 0.6 \log \left(1 + \alpha\right) + 0.4 \log \left(1 - \alpha\right)$$



$$\max_{\alpha} 0.6 \log (1+\alpha) + 0.4 \log (1-\alpha)$$

$$\frac{0.6}{1+\alpha} = \frac{0.4}{1-\alpha} \Rightarrow \alpha = 0.2.$$



When betting x dollars, the gambler wins with probability p and gets Ax dollars and loses with probability 1-p and gets 0 dollars.

$$\max_{\alpha} p \log (1 - \alpha + A\alpha) + (1 - p) \log (1 - \alpha)$$



When betting x dollars, the gambler wins with probability p and gets Ax dollars and loses with probability 1-p and gets 0 dollars.

$$\max_{\alpha} p \log (1 - \alpha + A\alpha) + (1 - p) \log (1 - \alpha)$$

$$\frac{d}{d\alpha}p\log(1-\alpha+A\alpha)+(1-p)\log(1-\alpha)=0$$



When betting x dollars, the gambler wins with probability p and gets Ax dollars and loses with probability 1-p and gets 0 dollars.

$$\max_{\alpha} p \log (1 - \alpha + A\alpha) + (1 - p) \log (1 - \alpha)$$

$$\frac{d}{d\alpha}p\log(1-\alpha+A\alpha)+(1-p)\log(1-\alpha)=0$$

$$\frac{p(A-1)}{1-\alpha+A\alpha}-\frac{1-p}{1-\alpha}=0\Rightarrow\frac{p(A-1)-(1-p)}{A-1}$$



Fraction 
$$=\frac{\mathsf{Edge}}{\mathsf{Odd}}$$
.

- Edge: the fraction of money you win on average when betting a unit amount of money.
- Odd: when you win, the profit you make.