Lecture 17: Stochastic Shortest Path Problems

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$$\min \sum_{k=0}^{\infty} E[c(x_k, u_k)|x_0].$$

- ullet Let us assume the states are $\{1,2,\ldots,N\}$
- Also, there is a distinct state called "0" such that
 - (a) if $x_k = 0$, then $x_l = 0 \quad \forall l \ge k$
 - (b) $c(x_k, u_k) = 0$ if $x_k = 0$
- Hence, as the name SSP suggests: find the cheapest path to "0"
- ullet c is assumed to be deterministic, but works for random c as well.

$$\max \sum_{k=0}^{\infty} E[r(x_k, u_k)|x_0],$$

where $r(x_k, u_k) = -c(x_k, u_k)$.

• Bellman Equation:

$$J(i) = \max_{u} r(i, u) + \sum_{j \neq i} P_{ij}(u)J(j)$$

with
$$J(0) = 0$$

• We will prove that T is a contraction mapping in a weighted ∞ -norm. Many of the earlier results for Q-learning continue to hold with the ∞ -norm replaced by the weighted ∞ -norm.

Fixed policy: consider a fixed policy μ . Then,

$$J_{\mu}(i) = r(i, \mu(i)) + \sum P_{ij}(\mu(i))J_{\mu}(j) = r(i) + \sum P_{ij}J_{\mu}(j),$$

we dropped the dependence on u since the policy is fixed. We want to show there exist $w_i>0, \beta\in[0,1)$ such that

$$\max_{i} w_{i} |T_{\mu} Y_{1}(i) - T_{\mu} Y_{2}(i)| \leq \beta \max_{i} w_{i} |Y_{1}(i) - Y_{2}(i)|$$
$$\max_{i} w_{i} |r(i) + \sum_{i} P_{ij} Y_{1}(j) - r(i) - \sum_{i} P_{ij} Y_{2}(j)| \leq \beta \max_{i} w_{i} |Y_{1}(i) - Y_{2}(i)|$$

$$\iff$$

$$\max_{i} w_{i} |\sum_{j} P_{ij}(Y_{1}(j) - Y_{2}(j))| \le \beta \max_{i} w_{i} |Y_{1}(i) - Y_{2}(i)|$$

Start from the LHS:

$$\begin{aligned} \max_{i} w_{i} \left| \sum_{j} \frac{P_{ij}}{w_{j}} w_{j} (Y_{1}(j) - Y_{2}(j)) \right| \\ \leq \max_{i} w_{i} \sum_{j} \frac{P_{ij}}{w_{j}} \underbrace{\max_{j} w_{j} |J_{1}(j) - J_{2}(j)|}_{\|Y_{1} - Y_{2}\|_{\infty, w}} \end{aligned}$$

$$= \underbrace{\max_{i} w_{i} \sum_{j} \frac{P_{ij}}{w_{j}}}_{\|Y_{1} - Y_{2}\|_{\infty, w}} \|Y_{1} - Y_{2}\|_{\infty, w}$$

• We want to show that the quantity in red box is $\leq \beta$ for an appropriate choice of β .

We want

$$\sum_{j} \frac{P_{ij}}{w_j} \le \frac{\beta}{w_i} \quad \forall i$$

 Consider an SSP where all rewards are 1 expect in state 0, where the reward is 0. Then the Bellman equation for this problem is

$$J(i) = 1 + \max_{u} \sum_{j} P_{ij}(u)J(j)$$

$$\geq 1 + \sum_{j} P_{ij}(\mu(i))J(j) \quad \forall \mu$$

$$\forall j \neq 0, J(j) \geq 1$$

$$\sum_{j} P_{ij}(\mu(i))J(j) \le J(i) - 1 \le J(i)\frac{J(i) - 1}{J(i)}$$

Define

$$\beta = \max_{i} \frac{J(i) - 1}{J(i)} < 1$$

Then

$$\sum_{j} P_{ij}(\mu(i))J(j) \le \beta J(i) \implies \sum_{j} \frac{P_{ij}(\mu(i))}{w_{j}} \le \frac{\beta}{w_{i}}$$

ullet Assumption: μ is such that there is a finite time t such that

$$\max_{i} P(x_n \neq 0 | x_0 = i) < 1$$

We used this assumption to ensure that $J_{\mu}(i) < \infty \quad \forall i$. To reason about the optimal policy, we also need

- Assumption: All stationary policies are proper, i.e. satisfying the previous assumption.
- We showed that

$$w_i |T_{\mu}(Y_1)(i) - T_{\mu}(Y_2)(i)| \le \beta \max_j w_j |Y_1(j) - Y_2(j)|$$

$$T(Y_1)(i) \le \max_{\mu} T_{\mu}(Y_1)(i) \le \max_{\mu} T_{\mu}(Y_2)(i) + \beta \frac{1}{w_i} \max_{j} w_j |Y_1(j) - Y_2(j)|$$

Thus,

$$T(Y_1)(i) \le T(Y_2)(i) + \beta \frac{1}{w_i} \max_j w_j |Y_1(j) - Y_2(j)|$$
$$w_i (T(Y_1)(i) - T(Y_2)(i)) \le \beta \max_j w_j |Y_1(j) - Y_2(j)|$$

Switching the roles of Y_1 and Y_2 , we have

$$w_i (T(Y_2)(i) - T(Y_1)(i)) \le \beta \max_j w_j |Y_1(j) - Y_2(j)|$$

$$w_i |T(Y_1)(i) - T(Y_2)(i)| \le \beta \max_j w_j |Y_1(j) - Y_2(j)|$$

$$\max_i w_i |T(Y_1)(i) - T(Y_2)(i)| \le \beta \max_j w_j |Y_1(j) - Y_2(j)|$$

Therefore we conclude that, T is a contraction mapping.

Reference

• This lecture is based on R. Srikant's lecture notes on *Stochastic Shortest Path Problems* available at https://sites.google.com/illinois.edu/mdps-and-rl/lectures?authuser=1

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