#### **Pr. 1.** (sol/hsj62)

If A or B is 0 then the problem is trivial, so assume they are both nonzero.

Denote the compact SVDs of A and B by  $A = U\Sigma V'$  and  $B = X\Omega Y'$ , where we omit the usual compact SVD subscripts for simplicity. Here  $\Sigma$  and  $\Omega$  are square and symmetric and invertible (but possibly different sizes if A and B have different ranks).

Now AB' = 0 means  $U\Sigma V'Y\Omega X' = 0$ . Multiplying on the left by  $\Sigma^{-1}U'$  and on the right by  $X\Omega^{-1}$  shows that V'Y = 0. So  $\begin{bmatrix} V & Y \end{bmatrix}$  is a matrix with orthonormal columns.

Likewise A'B = 0 leads to U'X = 0, so  $\begin{bmatrix} U & X \end{bmatrix}$  is a matrix with orthonormal columns.

Therefore, the following decomposition is a valid compact SVD of A + B, to within permutations for sorting the singular values:

$$A+B=U\Sigma V'+X\Omega Y'=egin{bmatrix} U & X\end{bmatrix}egin{bmatrix} \Sigma & 0 \ 0 & \Omega\end{bmatrix}egin{bmatrix} V & Y\end{bmatrix}'.$$

(This "SVD of a sum" can be useful itself.) Thus

$$\|A + B\|_* = \|\Sigma\|_* + \|\Omega\|_* = \|A\|_* + \|B\|_*$$
 .

Now there is one more subtle point we must address here. When A and B are both  $M \times N$ , we need to be sure that  $\operatorname{size}(\Sigma) + \operatorname{size}(\Omega) \le \min(M, N)$ . Specifically, if U is  $M \times r_1$  and X is  $M \times r_2$ , then we need  $r_1 + r_2 \le M$  for the above compact SVD if A + B to be valid. This inequality is assured by the condition U'X = 0 because U and X are each orthonormal bases in  $\mathbb{F}^M$ , so if the sum of their dimensions were to exceed M then their spans would have a nontrivial intersection which would contradict U'X = 0. Likewise for V and Y.

## **Pr. 2.** (sol/hsj31)

- (a)  $\{e_1, e_4\}$  is an orthonormal basis for the null space of X.
- (b)  $\{e_2, e_3\}$  is an orthonormal basis for the orthogonal complement of the null space of X.
- (c) The projection of  $\boldsymbol{x}$  onto the orthogonal complement of the null space of  $\boldsymbol{X}$  is  $\boldsymbol{P}_{\mathcal{N}(\boldsymbol{X})}^{\perp}\boldsymbol{x} = \begin{bmatrix} 0 \\ 2 \\ 3 \\ 0 \end{bmatrix}$
- (d) Using: Y = ones(3,3); (U,s,V) = svd(Y); V0=V[:,2:end] an orthonormal basis for the null space of Y is  $\left\{\begin{bmatrix} 0\\-1\\1 \end{bmatrix} \middle/ \sqrt{2}, \begin{bmatrix} 2\\-1\\-1 \end{bmatrix} \middle/ \sqrt{6}\right\}$

An orthonormal basis for the orthogonal complement of the null space of  $\boldsymbol{Y}$  is  $\{\boldsymbol{v}_1\},\,\boldsymbol{v}_1=\begin{bmatrix}1\\1\\1\end{bmatrix}/\sqrt{3}$ 

The projection of w onto the orthogonal complement of the null space of Y is  $P_{\mathcal{N}(Y)}^{\perp}w = v_1(v_1'w) = \begin{bmatrix} 2\\2\\2 \end{bmatrix}$ 

(e) A possible Julia implementation is

```
# Parse inputs
x = vec(x) \# make sure it is a column
(\underline{\ }, s, V) = svd(A)
r = rank(A) # unfortunate redundancy with svdvals, could avoid by using s
# orthonormal basis for orthogonal complement of the null space of A:
Vr = V[:,1:r]
return Vr * (Vr' * x) # output size is n by the number of columns of x
```

The above solution uses a syd call. The SVD itself calls a QR decomposition that is related to Gram-Schmidt orthogonalization. It is likely that an even more efficient solution exists by using the QR decomposition directly to find the basis  $V_r$  for the orthogonal complement of the null space of A. A QR approach is not required for full credit because we have not covered the QR decomposition.

Graders: accept solutions that use QR instead of SVD, as long as no inv or pinv or other expensive operations are used and as long as the basis matrix  $V_r$  is used efficiently with parentheses like this: Vr \* (Vr) \* xBecause  $\mathcal{N}^{\perp}(A) = \mathcal{R}(A')$ , having a basis  $V_r$  for the range space of A' suffices. The following code also passes.

```
function orthcompnull(A,x)
(Q, \sim) = qr(A') \# QR  approach to getting basis for range (A')
return Q * (Q' * x);
end
```

# **Pr. 3.** (sol/hsj43)

Given training data  $(x_n, y_n)$ , n = 1, ..., N consisting of pairs of features  $x_n \in \mathbb{R}^M$  and responses  $y_n \in \mathbb{R}$ , we train a linear "artificial neuron" to minimize the average MSE loss by solving the following optimization problem:

$$\hat{\boldsymbol{w}} = \operatorname*{arg\,min}_{\boldsymbol{w}} L(\boldsymbol{w}), \quad L(\boldsymbol{w}) = \frac{1}{N} \sum_{n=1}^{N} \left\| y_n - \boldsymbol{w}' \boldsymbol{x}_n \right\|_2^2 = \frac{1}{N} \left\| \boldsymbol{y}' - \boldsymbol{w}' \boldsymbol{X} \right\|_{\mathrm{F}}^2 = \frac{1}{N} \left\| \boldsymbol{y} - \boldsymbol{X}' \boldsymbol{w} \right\|_2^2,$$

where  $\boldsymbol{X} = \begin{bmatrix} \boldsymbol{x}_1 & \dots & \boldsymbol{x}_N \end{bmatrix} \in \mathbb{R}^{M \times N}$  and  $\boldsymbol{y} = (y_1, \dots, y_N)' \in \mathbb{R}^N$ . Assuming X has full rank, the LS solution is simply

$$\hat{\boldsymbol{w}} = (\boldsymbol{X} \boldsymbol{X}')^{-1} \boldsymbol{X} \boldsymbol{y} = \boldsymbol{K}_x^{-1} \boldsymbol{K}'_{yx}$$

where  $K_x = \frac{1}{N} \sum_{n=1}^{N} x_n x_n'$  and  $K_{yx} = \frac{1}{N} \sum_{n=1}^{N} y_n x_n'$ . We need  $N \ge M$  for there to be any chance that  $K_x$  is invertible. Otherwise we need a pseudo-inverse or some form of regularization.

#### **Pr. 4.** (sol/hsj72)

- (a) We already know that  $B'B \succeq 0$  so to show  $B'B \succ 0$  we simply must show that  $x \neq 0 \Rightarrow x'B'Bx \neq 0$ . Suppose there is a nonzero x such that x'B'Bx = 0. Then ||Bx|| = 0 so Bx = 0, but this would contradict the linear independence of the columns of B.
- (b)  $A \succ 0$  means that all eigenvalues are positive, hence nonzero, so the matrix is invertible.
- (c) x'(A+B)x = x'Ax + x'Bx > 0 for all x.
- (d) x'(A+B)x = x'Ax + x'Bx > 0 for all  $x \neq 0$  because A is positive definite.
- (e) Because B has full column rank,  $B'B \succeq 0$  so A'A + B'B is positive definite by the previous property and thus invertible by an earlier subproblem.

(f) It suffices to show that A'A + B'B > 0 when  $\mathcal{N}(A) \cap \mathcal{N}(B) = \{0\}$ . We already have shown that  $A'A + B'B \geq 0$  so we just need to verify that  $x'(A'A + B'B)x \neq 0$  for any  $x \neq 0$ .

Suppose the contrary, *i.e.*, that x'(A'A + B'B)x = 0 for some  $x \neq 0$ . then it follows that ||Ax|| = 0 and ||Bx|| = 0 so Ax = 0 and Bx = 0. But this means that x is in the null space of both A and B and that contradicts the assumption that  $\mathcal{N}(A) \cap \mathcal{N}(B) = \{0\}$ .

(g)  $B \succ 0$  implies that B is invertible, so if x is any nonzero vector, then y = Bx is nonzero, because otherwise B would be singular.

For any nonzero vector y,  $A > 0 \Rightarrow y'Ay > 0$  by definition of a positive definite matrix.

Thus  $A \succ 0$ ,  $B \succ 0 \Rightarrow x'B'ABx > 0$ ,  $\forall x \neq 0 \Rightarrow BAB \succ 0$ .

## **Pr. 5.** (sol/hsj61)

As described in the course notes, because  $\|\cdot\|_2$  and  $\|\cdot\|_*$  are untarily invariant:

$$\hat{\boldsymbol{X}} = \boldsymbol{U}_r \hat{\boldsymbol{\Sigma}}_r \boldsymbol{V}_r', \quad \hat{\boldsymbol{\Sigma}}_r = \operatorname*{arg\,min}_{\boldsymbol{S} \succeq \boldsymbol{0}} \frac{1}{2} \left\| \boldsymbol{\Sigma}_r - \boldsymbol{S} \right\|_2^2 + \beta \left\| \boldsymbol{S} \right\|_*, \quad \boldsymbol{S} = \operatorname{diag}(s_1, \dots, s_r)$$

So here we must solve

$$\underset{s_{1},\dots,s_{r}\geq0}{\arg\min}\left\{\left(\frac{1}{2}\max_{k}\left|\sigma_{k}-s_{k}\right|^{2}\right)+\beta\sum_{k=1}^{r}s_{k}\right\}.$$

Having  $s_k > \sigma_k$  would only increase the cost, so we must solve

$$\underset{s_1,\ldots,s_r}{\arg\min} \left\{ \left( \frac{1}{2} \max_k \left( \sigma_k - s_k \right)^2 \right) + \beta \sum_{k=1}^r s_k \right\}, \text{ s.t. } 0 \le s_k \le \sigma_k, \forall k.$$

Consider first the case where r=1, then we have simply

$$\hat{\sigma}_1 = \underset{0 \le s_1 \le \sigma_1}{\arg \min} f(s_1, \sigma_1, \beta) = \max(\sigma_1 - \beta, 0), \qquad f(s_1, \sigma_1, \beta) \triangleq \frac{1}{2} (\sigma_1 - s_1)^2 + \beta s_1$$

by differentiating and setting to zero and minding the constraints. So the rank-1 minimizer here is

$$\hat{\boldsymbol{X}} = \max(\sigma_1 - \beta, 0)\boldsymbol{u}_1\boldsymbol{v}_1',$$

which is the same solution as when using the Frobenius norm.

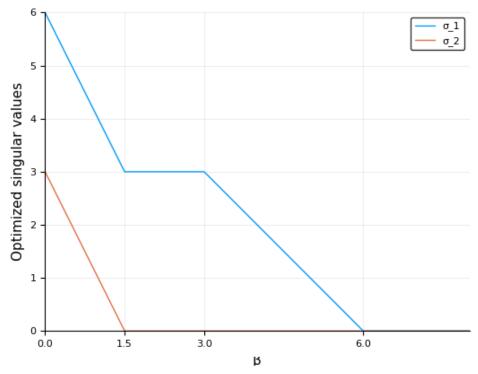
Grader: give full credit for correct solutions to the rank-1 case.

However, for r=2, the solution (found numerically) is different because of the max<sub>k</sub> term:

$$\hat{\sigma}_{2} = [\sigma_{2} - 2\beta]_{+}$$

$$\hat{\sigma}_{1} = \begin{cases} [\sigma_{1} - 2\beta]_{+}, & 0 \leq \beta \leq \sigma_{2}/2\\ \sigma_{1} - \sigma_{2}, & \sigma_{2}/2 \leq \beta \leq \sigma_{2}\\ [\sigma_{1} - \beta]_{+}, & \sigma_{2} \leq \beta \leq \sigma_{1}\\ 0, & \sigma_{1} \leq \beta \end{cases}$$

The following figure illustrates.



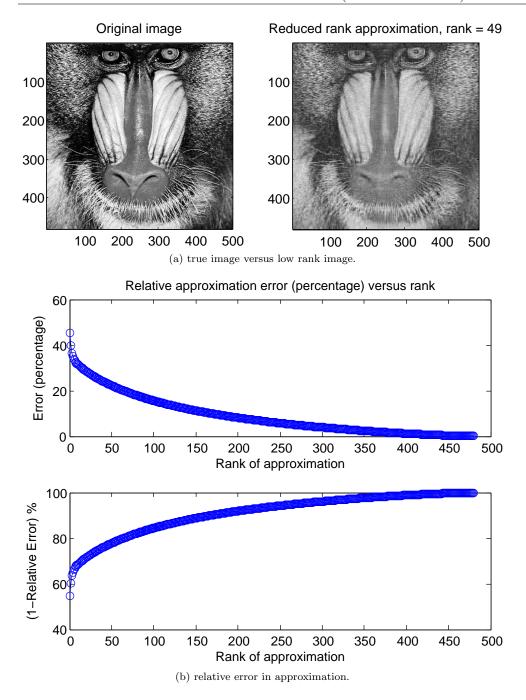
I am unsure how to solve it for a general rank. If you think you have a solution, first have a friend in the class check it, and if you both agree it looks correct, then please tell me about it!

## **Pr. 6.** (sol/hs042)

(a) The idea is to view the image as a matrix  $\mathbf{A} = \sum_{i=1}^n \sigma_i \mathbf{u}_i \mathbf{v}_i'$  where, for the  $500 \times 480$  mandrill image shown below, n = 480. We know that the optimal (with respect to any unitarily invariant norm) rank-k approximation to  $\mathbf{A}$  is  $\mathbf{A}_k = \sum_{i=1}^k \sigma_i \mathbf{u}_i \mathbf{v}_i'$ . This approximation requires 1+480+500 = 981 real numbers for each additional term. The original matrix stores  $480 \times 500 = 240000$ . One-fifth of this total is 48000. Thus the number of terms that gives 48000 real numbers is  $48.93 \approx 49$ .

Note that if you just took one fifth of the number of terms in the full image, that would give you 96 terms, but that would entail a roughly 40% reduction as opposed to the 20%.

In general, for an  $m \times n$  matrix, the rank r approximation requires  $r \times (1 + m + n)$  real numbers. Hence, for a given compression fraction p, we want to choose the maximum value of r such that  $\frac{r(1+m+n)}{mn} \leq p$ . I.e., r can be chosen as the *floor* of  $\frac{p \times mn}{1+m+n}$ .



# A possible Julia implementation is

```
# Parse inputs
m, n = size(A)

# Compute compression factor
r = Int64(floor(p * m * n / (m + n + 1)))
r = minimum([r, m, n]) # just to be safe

# Compute compressed image
U, s, V = svd(A)
Ac = U[:, 1:r] * Diagonal(s[1:r]) * V[:, 1:r]'
return Ac, r
end
```

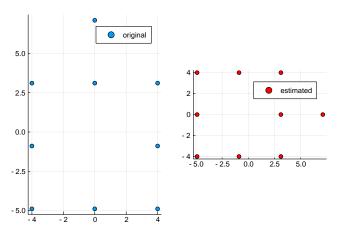
(b) The MIT logo without the lettering has rank 4 if the background white space is represented as a zero in the matrix; otherwise it is rank 5. Note that the singular values  $\sigma_k$  for  $k \geq 6$  are not numerically 0, but that they are very, very small. Either way it is nearly perfectly compressible with a low-rank approximation. In contrast, the image of the logo with the lettering is full rank, so the approximation error decreases as the approximation rank increases. Being able to spot low-rank patterns is a valuable skill!

### **Pr. 7.** (sol/hs052)

(a) A possible Julia implementation is

```
using LinearAlgebra: I, svd, Diagonal
function dist2locs(D. d)
# Syntax:
               Xr = dist2locs(D, d)
                D is an n x n matrix such that D[i, j] is the distance from
# Inputs:
                object i to object j
                d is the desired embedding dimension.
                Xr is an n x d matrix whose rows contains the relative
                coordinates of the n objects
# Note:
                MDS is only unique up to rotation and translation,
                so we enforce the following conventions on Xr in this order:
        [ORDER] Xr[:,i] corresponds to ith largest eigenpair of C * C'
       [CENTER] The centroid of the coordinates is zero
         [SIGN] The largest magnitude element of Xr[:, i] is positive
   # Parse inputs
   n = size(D, 1)
   # Compute correlation matrix - ensures [CENTER]
   S = D .* D
   S = 0.5 * (S + S') # force symmetry (in case of noise)
   P = I - ones(n, n) / n
   CCt = -0.5 * (P * S * P)
   # Compute relative coordinates
    _, s, V = svd(CCt) # using SVD ensures [ORDER]
    Xr = V[:, 1:d] * Diagonal(sqrt.(s[1:d]))
    # Apply [SIGN]
   Xr . \star = sign.(Xr[findmax(abs.(Xr), dims=1)[2]])
    return Xr
end
```

(b) Applying the MDS method to the given distance matrix produces the right figure below. The left figure is the original coordinates.



# Optional problem(s) below

# **Pr. 8.** (sol/hs034)

Here  $\mathbf{A} \in \mathbb{R}_8^{19 \times 48}$  has rank r=8. We are looking for the number of linearly independent solutions to the homogeneous linear system  $\mathbf{A}\mathbf{x} = \mathbf{0}$ . All vectors  $\mathbf{x}$  such that  $\mathbf{A}\mathbf{x} = \mathbf{0}$  must belong to the nullspace of  $\mathbf{A}$ . Thus, the number of linearly independent solutions to  $\mathbf{A}\mathbf{x} = \mathbf{0}$  is precisely the dimension of the nullspace of  $\mathbf{A}$ . From the rank-plus-nullity theorem (Corollary 3.18), we have that  $n = \dim \mathcal{N}(A) + \dim \mathcal{R}(A)$ . Here n = 48,  $r = \operatorname{rank}(\mathbf{A}) = \dim \mathcal{R}(A) = 8$  so that we must have  $\dim \mathcal{N}(A) = 48 - 8 = 40$ . Thus there are 40 linearly independent solutions of  $\mathbf{A}\mathbf{x} = \mathbf{0}$ .

# **Pr. 9.** (sol/hs038)

If  $\mathbf{A} \in \mathbb{R}^{M \times N}$  and  $\mathbf{B} \in \mathbb{R}^{N \times M}$ , then

$$\operatorname{vec}(\boldsymbol{A}^T) = \begin{bmatrix} A(1,:)^T \\ A(2,:)^T \\ \vdots \\ A(M,:)^T \end{bmatrix} \quad \text{and} \quad \operatorname{vec}(\boldsymbol{B}) = \begin{bmatrix} B(:,1) \\ B(:,2) \\ \vdots \\ B(:,M) \end{bmatrix} \Rightarrow \operatorname{vec}(\boldsymbol{A}^T)^T \operatorname{vec}(\boldsymbol{B}) = \sum_{m=1}^M A(m,:)B(:,m).$$

Furthermore,

$$\begin{aligned} & \operatorname{trace}(\boldsymbol{A}\boldsymbol{B}) = \operatorname{trace} \left( \begin{bmatrix} A(1,:) \\ A(2,:) \\ \vdots \\ A(M,:) \end{bmatrix} \begin{bmatrix} B(:,1) & B(:,2) & \dots & B(:,M) \end{bmatrix} \right) \\ & = \operatorname{trace} \left( \begin{bmatrix} A(1,:)B(:,1) & A(1,:)B(:,2) & \dots & A(1,:)B(:,M) \\ A(2,:)B(:,1) & A(2,:)B(:,2) & \dots & A(2,:)B(:,M) \\ & & \ddots & \\ A(M,:)B(:,1) & A(M,:)B(:,2) & \dots & A(M,:)B(:,M) \end{bmatrix} \right) \\ & = \sum_{m=1}^{M} A(m,:)B(:,m) = \operatorname{vec}(\boldsymbol{A}^T)^T \operatorname{vec}(\boldsymbol{B}), \end{aligned}$$

using the preceding equality. This is the general property.

When  $\boldsymbol{A}$  is symmetric (and  $\boldsymbol{B}$  is square with the same size), then  $\boldsymbol{A} = \boldsymbol{A}^T$ , so that  $\operatorname{trace}(\boldsymbol{A}\boldsymbol{B}) = \operatorname{vec}(\boldsymbol{A})^T \operatorname{vec}(\boldsymbol{B})$ .

#### **Pr. 10.** (sol/hsj51)

Show that the weighted Euclidean norm  $\|x\|_{W}$  is a valid norm iff W is a positive definite matrix. First we show sufficiency: if W is positive definite, then we show that  $\|x\|_{W} = \sqrt{x'Wx}$  is a valid norm.

• Clearly  $\boldsymbol{x} = \boldsymbol{0} \Rightarrow \|\boldsymbol{x}\|_{\boldsymbol{W}} = 0$ 

- • By definition of a positive definite matrix,  $\| \boldsymbol{x} \|_{\boldsymbol{W}} = 0 = \sqrt{\boldsymbol{x}' \boldsymbol{W} \boldsymbol{x}} \Rightarrow \boldsymbol{x} = \boldsymbol{0}$
- $\bullet \ \left\|\alpha \boldsymbol{x}\right\|_{\boldsymbol{W}} = \sqrt{(\alpha^*)\boldsymbol{x}'\boldsymbol{W}(\alpha\boldsymbol{x})} = \left|\alpha\right|\sqrt{\boldsymbol{x}'\boldsymbol{W}\boldsymbol{x}} = \left|\alpha\right|\left\|\boldsymbol{x}\right\|_{\boldsymbol{W}}$
- Because W is (Hermitian) positive definite, it has a unitary eigendecomposition of the form  $W = V\Lambda V'$  where the eigenvalues in  $\Lambda$  are all positive. Let  $S = V\Lambda^{1/2}V'$  so that W = SS = S'S. Then  $\|x\|_W = \sqrt{x'Wx} = \sqrt{x'S'Sx} = \|Sx\|$  so  $\|x + y\|_W = \|S(x + y)\| \le \|Sx\| + \|Sy\| = \|x\|_W + \|y\|_W$

Now we show necessity. If  $\|x\|_{W}$  is a valid norm, then for all  $x \ 0 \le \|x\|_{W}^{2} = x'Wx$  and for  $x \ne 0$ :  $0 < \|x\|_{W}^{2} = x'Wx$ . These are the two conditions for a (Hermitian) symmetric matrix to be positive definite.