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Global Cardinality Constraint

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Solving the Model

- ⌘ Solving ... solving ... solving ...
- ⌘ Nothing came out after 5 mins
- ⌘ A solution finally came out in 10.5 minutes
- ⌘ Changing the number of days to 6, the model can find an optimal solution only after 7 seconds

What went wrong?



Shift Requirement Constraints

- There are o soldiers on night shift each day

```
forall(d in DAY)
  (sum(s in SOLDIER)
    (roster[s,d] = NIGHT) = o);
```

- There are between l and u soldiers on each evening shift

```
forall(d in DAY)
  (sum(s in SOLDIER)
    (roster[s,d] = EVE) >= l);
forall(d in DAY)
  (sum(s in SOLDIER)
    (roster[s,d] = EVE) <= u);
```

3

Shift Requirement Constraints

- There are o soldiers on day shift each night

```
forall(d in DAY)
  (sum(s in SOLDIER)
    (roster[s,d] = NIGHT) = o);
```

- There are between l and u soldiers on each evening shift

```
forall(d in DAY)
  (sum(s in SOLDIER)
    (roster[s,d] = EVE) >= l);
forall(d in DAY)
  (sum(s in SOLDIER)
    (roster[s,d] = EVE) <= u);
```

Common Subexpressions

4



Intermediate Variables

Intermediate variables

- Store values of expressions that are reused
- Are dependent on decisions
- (note: intermediate parameters too!)

```
array[DAY] of var 0..card(SOLDIER): onEve;  
onEve = [sum(s in SOLDIER) (roster[s,d] = EVE)  
        | d in DAY];  
forall(d in DAY)  
    (onEve[d] >= 1 /\ onEve[d] <= u);
```

Choose bounds for intermediates well

Or simply

```
array[DAY] of var 1..u: onEve;  
onEve = [sum(s in SOLDIER) (roster[s,d] = EVE)  
        | d in DAY];
```

5

Another Common Subexpression

Objective

```
var int: tOnEve = sum(d in DAY)  
    (sum(s in SOLDIER) (roster[s,d] = EVE));  
solve maximize (tOnEve);
```

6



The Patrol Model (patrolV2.mzn)

```
enum SOLDIER;
enum SHIFT;
int: nDays;
set of int: DAY = 1..nDays;
int: o;
int: l;
int: u;

array[SOLDIER, DAY] of var SHIFT: roster;

constraint forall(d in 1..(nDays-2),
  s in SOLDIER) ((roster[s,d] = NIGHT) /\
    (roster[s,d+1] = NIGHT)
  -> (roster[s,d+2] != NIGHT));
constraint forall(d in 1..(nDays-1),
  s in SOLDIER) ((roster[s,d] = EVE) ->
    (roster[s,d+1] != NIGHT));
```

7

The Patrol Model (patrolV2.mzn)

```
constraint forall(d in DAY) (sum (s in SOLDIER)
  ((roster[s,d] = NIGHT)) = o);
array[DAY] of var l..u: onEve;
constraint onEve = [sum (s in SOLDIER)
  (roster[s,d]=EVE) | d in DAY];
solve maximize sum(onEve);
```

8



Solving the New Model

- ⌘ The new version of the model after removing the common subexpression is solved in 4.5 seconds

9

Partitioning Problems

- ⌘ Many times when we are partitioning a set we have to partition it with bounds on the size of the partitions
 - e.g. $\#NIGHT = o, 1 \leq \#EVE \leq u$
- ⌘ We have special constraints for partitioning with size bounds
 - `global_cardinality(x, v, c)`
 - Sizes of `v` and `c` are the same
 - Constrains $c_i = \sum_{j \in 1..n} (x_j = v_i)$
 - Collects the counts, and bounds the count of the number of occurrences of v_i
 - e.g. `global_cardinality(x, [1,2], [2,1])`
`x = [1,1,2,3]` ✓, `[1,2,3,4]` ✗

10



Global Cardinality

Replace

```
forall(d in DAY) (sum(s in SOLDIER)
  (roster[s,d] = NIGHT) = o);
forall(d in DAY) (sum(s in SOLDIER)
  (roster[s,d] = EVE) >= 1);
forall(d in DAY) (sum(s in SOLDIER)
  (roster[s,d] = EVE) <= u);
```

By

```
array[DAY] of var l..u: onEve;
constraint forall(d in DAY)
  (global_cardinality(
    [roster[s,d] | s in SOLDIER],
    [NIGHT, EVE],
    [o, onEve[d]]));
```

11

Global Cardinality Variants

- There are a number of variants of global cardinality

- Collecting counts, and requiring every value is counted

- `global_cardinality_closed(x, v, c)`
- Values for x_i must be from v , i.e. $\forall i. \exists j. x_i = v_j$

- Bounding the count of the number of occurrences

- `global_cardinality_low_up(x, v, lo, hi)`
- Constrains $lo_i \leq (\sum_{j \text{ in } 1..n} (x_j = v_i)) \leq hi_i$

- Bounding the count, and requiring that every value is counted

```
global_cardinality_low_up_closed(x,v,l,u)
```

12



The Patrol Model (patrolV3.mzn)

```
enum SOLDIER;  
enum SHIFT;  
int: nDays;  
set of int: DAY = 1..nDays;  
int: o;  
int: l;  
int: u;  
  
array[SOLDIER, DAY] of var SHIFT: roster;  
  
constraint forall(d in 1..(nDays-2),  
  s in SOLDIER) ((roster[s,d] = NIGHT) /\  
    (roster[s,d+1] = NIGHT)  
  -> (roster[s,d+2] != NIGHT));  
constraint forall(d in 1..(nDays-1),  
  s in SOLDIER) ((roster[s,d] = EVE) ->  
    (roster[s,d+1] != NIGHT));
```

13

The Patrol Model (patrolV3.mzn)

```
array[DAY] of var l..u: onEve;  
constraint forall(d in DAY) (global_cardinality(  
  [roster[s,d] | s in SOLDIER],  
  [NIGHT, EVE], [o, onEve[d]]));  
  
solve maximize sum(onEve);
```

14



Solving the Global Constraint Model

- ⌘ The global constraint version of the model is solved in 1.5 seconds
- ⌘ In subsequent courses, we will explain how modeling can affect solving efficiency

15

Summary

- ⌘ Many discrete optimization problems involve
 - deciding a function $f: \text{DOM} \rightarrow \text{COD}$
 - this can be seen as partitioning DOM
- ⌘ The partition view is useful when we reason about the sets $F(c)$ (pseudo-inverse of f)
- ⌘ Partitioning with cardinality constraints is a substructure captured by the
 - `global_cardinality` family
- ⌘ Common subexpressions
- ⌘ Logical connectives

16



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