Assumptions

- 1. There are **p** processors in CPU architecture
- 2. The order in which edges are entering the incidence matrix does not matter
- 3. Storing of the matrix in memory is in row-major format

Observations

- 1. Each cell with positive value will create independent columns in the incidence matrix
 - a. We can parallelize traversal of the adjacency matrix by distributing n/p rows among each processor
 - b. While assigning columns to each edge we can again parallelize it by distributing it into p processors
- 2. Since the memory is stores matrix in the row-major form we can store rows for each edge and then transpose it, this can help us optimize the cache usage.
 - a. If the memory was column-major then we could traverse in adjacency matrix column-wise and distribute columns into processors instead of rows.

Pseudo Code for the Parallel Algorithm

```
// Gloabal Region
N ← #Nodes in graph
P ← #Processors
Local size ← N/P
//Adj_mat and Inc_mat are shared between all the processors
       // Parallel Region
       // Parallelise this loop into processors with N/P rows to each one
       LOOP (0,N):
               LOOP (0,N):
                      // Parallelise this loop
                      While(Adj_mat[i][j]>0):
                              if(i==j)
                                      Adj_mat[i][j] ← Adj_mat[i][j] / 2
                              else
                                     Adj_mat[i][j] \leftarrow Adj_mat[i][j] - 1
                              Temp ← Vector (N,0)
                              Temp[i] \leftarrow Temp[i] + 1
                              Temp[j] \leftarrow Temp[j] + 1
                              Inc_mat.append(Temp)
// Global Region
Inc_mat ← transpose(Inc_mat)
```