

2.1

a)

Code:

```
#!/usr/bin/env stack
-- stack --install-ghc runghc
f :: [[Int]] -> [[Int]]
f x = map(filter(even)) x
main = do
    print(f [[3,2,4],[1]])
    print(f [[0,0,0],[3,9,10,5]])
    let a = f [[0,0,0],[3,9,10,5]]
    print(a)
```

Output:

```
[[2,4],[1]]
[[0,0,0],[10]]
[[0,0,0],[10]]
```

b)

```
#!/usr/bin/env stack
-- stack --install-ghc runghc
my_append :: [a] -> [a] -> [a]
my_append xs ys = foldr (:) ys xs

main = do
    print(my_append [1,5,7] [-1,-3,-5])
```

Output:

```
laallen@LAPTOP-3ERSVLR2:/mnt/c/Users/Mordi/work/haskell/assign2$ ./2_1_b.hs
[1,5,7,-1,-3,-5]
```

2.2)

a) Defines a tree, which is either a leaf or a Node and two trees, which represent the connection between the node and its two children. A leaf or a node might be empty, but if it's not, they're all the same type. But that type can be anything.

b)

```
#!/usr/bin/env stack
```

```
-- stack --install-ghc runghc
```

```
data MTree a = MLeaf (Maybe a)
             | MNode (Maybe a) (MTree a) (MTree a)
             deriving (Show)
```

```
mTreeMap :: (a -> b) -> MTree a -> MTree b
```

```
mTreeMap f (MLeaf Nothing)    = MLeaf Nothing
```

```
mTreeMap f (MLeaf (Just a))   = MLeaf (Just(f(a)))
```

```
mTreeMap f (MNode Nothing t1 t2) = MNode Nothing (mTreeMap f t1) (mTreeMap f t2)
```

```
mTreeMap f (MNode (Just n) t1 t2) = MNode (Just(f(n))) (mTreeMap f t1) (mTreeMap f t2)
```

```
mTreeFilter :: (a -> Bool) -> MTree a -> MTree a
```

```
mTreeFilter f (MLeaf Nothing) = MLeaf Nothing
```

```
mTreeFilter f (MLeaf (Just a)) = if f(a)
                                then MLeaf(Just(a))
                                else
```

```
                                MLeaf(Nothing)
```

```
                                MLeaf(Nothing)
```

```
mTreeFilter f (MNode Nothing t1 t2) = MNode Nothing (mTreeFilter f t1) (mTreeFilter f t2)
```

```
mTreeFilter f (MNode (Just a) t1 t2) = if f(a)
                                then (MNode (Just(a)) (mTreeFilter f t1) (mTreeFilter f t2))
                                else (MNode Nothing (mTreeFilter f t1) (mTreeFilter f t2))
```

```
main =
```

```
do
```

```
  let a = MNode Nothing (MNode (Just 3) (MLeaf (Just 2)) (MLeaf (Just 1))) (MNode (Just 2) (MLeaf (Just 3)) (MLeaf (Just 4)))
```

```
  print "a:"
```

```
  print a
```

```
  let b = mTreeMap (* 2) a
```

```
  print "b:"
```

```
  print (b)
```

```
  print "c:"
```

```
  let c = mTreeFilter(==2) a
```

```
  print (c)
```

Output:

```
laallen@LAPTOP-3ERSVLR2:/mnt/c/Users/Mordi/work/haskell/assign2$ ./2_2.hs
"a:"
MNode Nothing (MNode (Just 3) (MLeaf (Just 2)) (MLeaf (Just 1))) (MNode (Just 2) (MLeaf (Just 3)) (MLeaf (Just 4)))
"b:"
MNode Nothing (MNode (Just 6) (MLeaf (Just 4)) (MLeaf (Just 2))) (MNode (Just 4) (MLeaf (Just 6)) (MLeaf (Just 8)))
"c:"
MNode Nothing (MNode Nothing (MLeaf (Just 2)) (MLeaf Nothing)) (MNode (Just 2) (MLeaf Nothing) (MLeaf Nothing))
```

2.3)

Prove: $\text{map } f (\text{map } g \text{ l}) = \text{map } g (\text{map } f \text{ l})$

Know: $f(g(x)) = g(f(x))$ for all numbers

Base Case: $x = []$, x is an empty list

By how map is defined, we know that:

$\text{map } g \text{ x} = []$

and

$\text{map } f \text{ x} = []$

so:

$\text{map } f (\text{map } g \text{ x}) = \text{map } f ([]) = []$

and

$\text{map } g (\text{map } f \text{ x}) = \text{map } g ([]) = []$

Thus:

$\text{map } f (\text{map } g \text{ x}) = \text{map } g (\text{map } f \text{ x})$

Inductive hypothesis:

$\text{map } f (\text{map } g \text{ n}) = \text{map } g (\text{map } f \text{ n})$, where n is a list of arbitrary size

Prove:

$\text{map } f (\text{map } g \text{ m}) = \text{map } g (\text{map } f \text{ m})$, where m is a list of size $(n) + 1$

Proof:

$\text{map } f (\text{map } g \text{ m}) = \text{map } f (\text{map } g \text{ m:n})$, where $m:n$ is the first element of m , followed by the other n elements

$\text{map } f (\text{map } g \text{ m:n}) = f(g(m)) ++ \text{map } f (\text{map } g \text{ n})$, by definition of map

$f(g(m)) ++ \text{map } f (\text{map } g \text{ n}) = g(f(m)) ++ \text{map } f (\text{map } g \text{ n})$, by the Know condition

$g(f(m)) ++ \text{map } f (\text{map } g \text{ n}) = g(f(m)) ++ \text{map } g (\text{map } f \text{ n})$, by Inductive hypothesis

$g(f(m)) ++ \text{map } g (\text{map } f \text{ n}) = \text{map } g (\text{map } f \text{ m:n})$, by definition of concatenation, and map

Thus:

$\text{map } f (\text{map } g \text{ m}) = \text{map } g (\text{map } f \text{ m})$