Robert Bielas laboratorium z Netezzy - raport

UWAGA: proszę zwracać uwagę na komentarze pod screenami/tekstami planera, gdyż pokazują one moją interpretację wyników.

Część pierwsza, zad. 11

Najlepsza dystrybucja tabel dla tabel w LABDB: w tabelach joinowanych (orders i lineitem, part i lineitem, supplier i part sup itp.) zaproponowałbym dystrybucję wg kolumny, która łączy odpowiednie pary tabel. Jeśli nie zależy nam na informacjach z kilku tabel, ale z jednej, zwróciłbym uwagę na to, którą operację wyszukiwania będziemy najczęściej wykonywać: np. wg. daty. Wtedy zdystrybuowałbym tabelę, posortowaną wg interesującej nas kolumny, i zdystrybuował po kluczu którego zbiór unikalnych wartości ma moc >= 22 (tyle najwyżej data slice'ów jest w serwerze Netezzy wykorzystywanej na laboratorium).

```
Część druga, zad. 1
select count(*) from paymentreport, rezultat:
100000000

select count(*) from paymentreportdetail, rezultat:
600000000

select count(*) from company, rezultat: 22

select count(*) from COMPANYHEADQUARTER, rezultat: 43
select count(*) from Paymentreport_age
100000000

select count(*) from Paymentreport_gender, rezultat:
1000000000

Wniosek: tabele są pokaźnych rozmiarów
zad. 2 rezultat:
```

DATASLICEID	↓Σ⊽⊅	rows	ĮΣ⊽⊅
1		4546633	
2		4546634	
3		4546634	
4		4546634	
5		4546634	
6		4545108	
7		4545108	
8		4545108	
9		4545108	
10		4545108	
11		4545108	
12		4545108	
13		4545108	
14		4545108	
15		4545108	
16		4545108	
17		4545108	
18		4545107	
19		4545107	
20		4545107	
21		4545107	
22		4545107	

zad. 3
Dystrybucja tabeli paymentreport_gender (po kluczu gender):

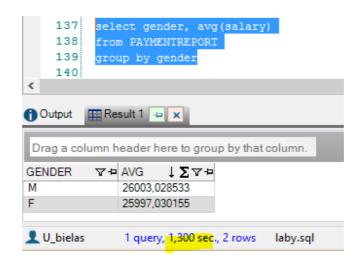
select datasliceid, count(datasliceid) as "rows"
from PAYMENTREPORT_GENDER
group by datasliceid
order by 1

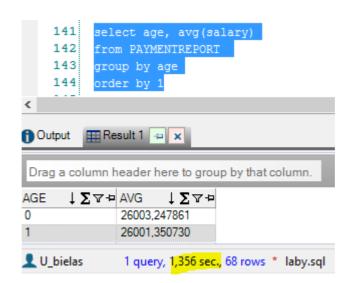
DATASLICEID	↓∑∀⊅ rows	↓Σ⊽≠
10	49999	997
15	50000	003

Dystrybuja tabeli paymentreport_age (po kluczu age) select datasliceid, count(datasliceid) as "rows" from PAYMENTREPORT_AGE group by datasliceid order by 1

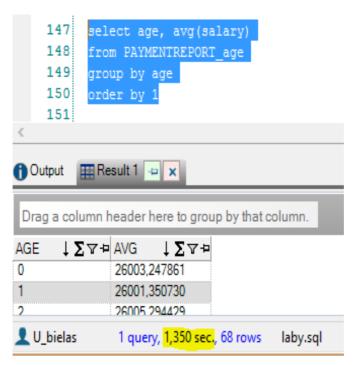
DATASLICEID	↓Σ⊽⊅	rows	↓Σ⊽⊅
1		5000542	
2		5997064	
3		4999551	
4		5003441	
5		4995738	
6		5000843	
7		5000580	
8		5003587	
9		5000312	
10		5002074	
11		4004265	
12		3997798	
13		3999722	
14		4001385	
15		3998297	
16		3995568	
17		4003395	
18		4002063	
19		4999312	
20		3998698	
21		3998174	
22		3997591	

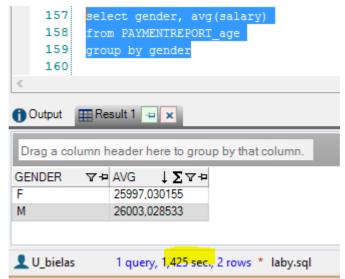
Po utworzeniu tabel w mojej bazie na tych dystrybucjach rozmieszczenie krotek było identyczne.



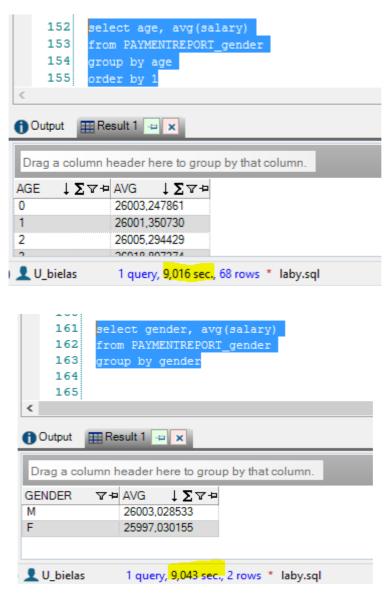


Sprawdzenie czasu wykonania obu agregacji dla dystrybucji po age:





Czas wykonania jest rozsądny, gdyż nie występuje data skew - dane wydystrybuowane po age są równomiernie rozłożone po większej ilości SPU - jak w przypadku dystrybucji po id. Sprawdzenie czasu wykonania obu agregacji dla dystrybucji po gender:



Tutaj nie dziwi nas, że czasy są podobne - w przetwarzanie są zaangażowane tylko dwa S-Blade'y (bo do dwóch wydystrybuowano dane) - reszta śpi.

```
zad. 5
```

```
select count(*) from PAYMENTREPORTDETAIL_ALL_BONUSCODE 1600000000
```

select count(*) from PAYMENTREPORTDETAIL_ALL2 3200000000

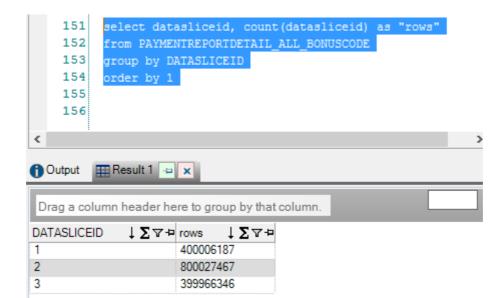
select count(*) from PAYMENTREPORTDETAIL_ALL2_BONUSCODE 3200000000

select count(*) from PAYMENTREPORTDETAIL_ALL
1600000000

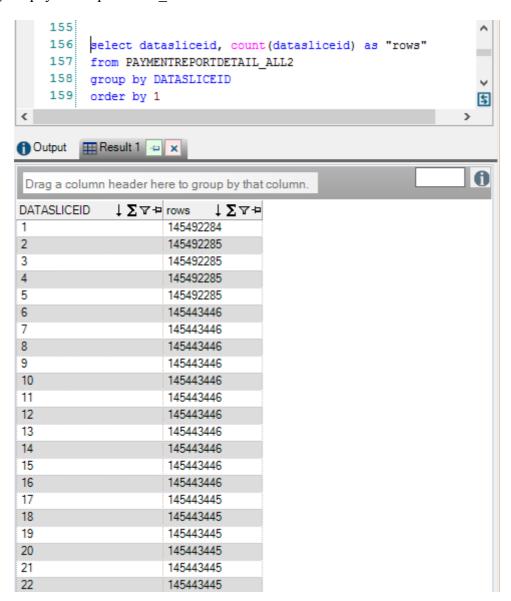
dystrybucja w paymentreportdetail_all

DATASLICEID	↓Σ⊽⊅		↓Σ⊽⊅
1		7274614	
2		72746142	2
3		72746142	2
4		72746142	2
5		72746142	2
6		7272172	3
7		7272172	3
8		7272172	3
9		7272172	3
10		7272172	3
11		72721 <mark>7</mark> 23	3
12		7272172	3
13		7272172	3
14		7272172	3
15		7272172	3
16		7272172	3
17		7272172	3
18		7272172	3
19		7272172	3
20		7272172	3
21		7272172	3
22		7272172	3

dystrybucja w paymentreportdetail_all_bonuscode



dystrybucja w paymentreportdetail all2



dystrybucja w paymentreportdetail_all2_bonuscode:

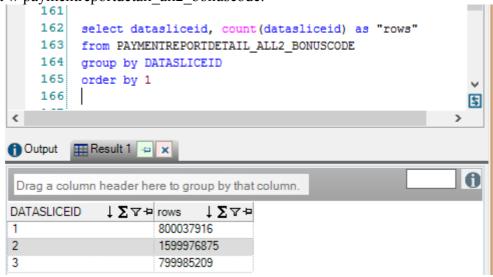


tabela paymentreportdetail_all jest zdystrybuowana na id tabela paymentreportdetail_all2 jest rozdystrybowana na id tabela paymentreportdetail_all2_bonuscode jest rozdystrybowana na bonuscode tabela paymentreportdetail all bonuscode jest rozdystrybowana na bonuscode

Polecenia agregujące:

```
select date, sum(bonus)
from PAYMENTREPORTDETAIL_ALL
group by date

czas: 13.916 sec

select bonuscode, sum(bonus)
from PAYMENTREPORTDETAIL_REPORTID
group by bonuscode

czas: 4.793 sec

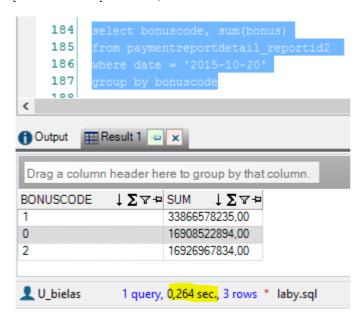
select bonuscode, sum(bonus)
from PAYMENTREPORTDETAIL_ALL_BONUSCODE
group by bonuscode
czas: 85.993 sec (!!!)
```

Ostatnia agregacja wykonywała się najwolniej, ponieważ rozdystrybuowano ją na bonuscode(są 3). Mamy tu do czynienia ze zjawiskiem data_skew: tylko 3 data slice'y były zaangażowane w przetwarzanie (do trzech slice'ów zostały zdystrybuowane dane).

Druga agregacja jest najszybsza, gdyż dane zostały przechowane równomiernie w największej ilości slice'ów. Więcej S-Bladów przetwarzało mniej danych -> wynik marzenie.

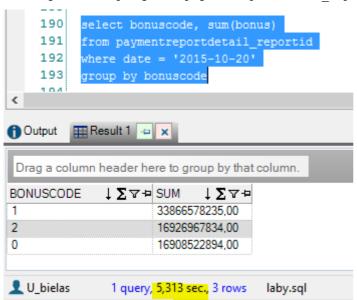
6.

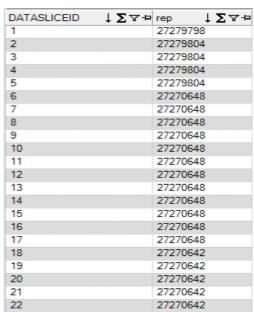
Czas wykonania query dla paymentreportdetail_reportid2 (dystrybucja po payment_report, posortowane po dacie):



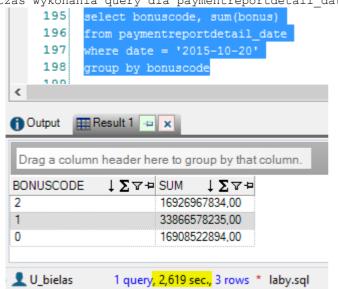
DATASLICEID	↓Σ⊽⊅	rep2	↓Σ⊽⊅
1		2727979	8
2		2727980	4
3		2727980	4
4		2727980	4
5		2727980	4
6		2727064	8
7		2727064	8
8		2727064	8
9		2727064	8
10		2727064	8
11		2727064	8
12		2727064	8
13		2727064	8
14		2727064	8
15		2727064	8
16		2727064	8
17		2727064	8
18		2727064	2
19		2727064	2
20		2727064	2
21		2727064	2
22		2727064	2

Czas wykonania query dla paymentreportdetail reportid





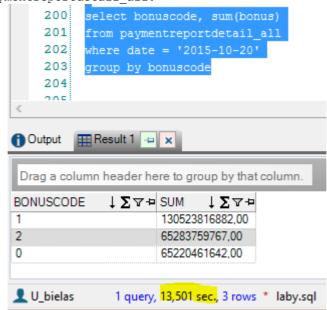
Czas wykonania query dla paymentreportdetail_date:



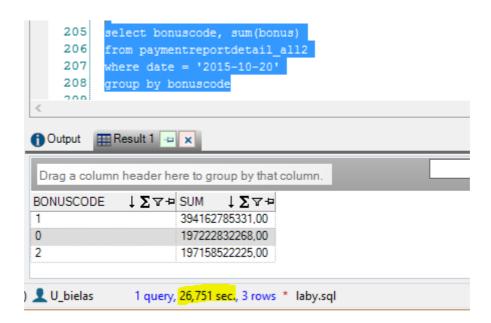
DATASLICEID	Σ⊽⊅	date	↓Σ⊽⊅
1		30000824	4
2		29999864	4
3		30011297	7
4		29998409	9
5		29993986	5
6		29999312	2
7		24997732	2
8		20003703	3
9		20002352	2
10		20000642	2
11		2000583	1
12		20001057	7
13		24997217	7
14		30001279	9
15		30000847	7
16		29997767	7
17		2999990	5
18		30004882	2
19		29990420	5
20		30003865	5
21		2999697	1
22		29991832	2

Najbardziej efektywnie zostało wykonane query dla dystrybucji po payment_report, gdzie dane szostały posortowane po dacie(czyli w tabeli paymentreportdetail_reportid2).

Czsy wykonania dla większych tabel: Czas wykonania dla paymentreportdetail_all:

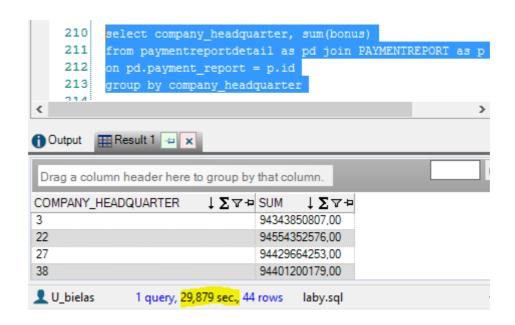


Czas wykonania dla paymentreportdetail all2:



Czas wykonania jest proporcjonalny do ilości krotek w tabeli (w _all jest ich dwa razy mniej niż w _all2 -> w _all czas wykonania był 2 razy mniejszy). Pamiętajmy, że dane w tych tabelach zostały rozdystrybuowane na tym samym kluczu.

7.



QUERY VERBOSE PLAN:

Node 1.

[SPU Sequential Scan table "PAYMENTREPORTDETAIL" as "PD" {(PD."ID")}]

-- Estimated Rows = 599005119, Width = 8, Cost = 0.0 .. 11718.8, Conf = 90.0 (FACT)

Restrictions:

(PD.PAYMENT_REPORT NOTNULL)

Projections:

1:PD.BONUS 2:PD.PAYMENT_REPORT

Cardinality:

PD.PAYMENT REPORT 100.0M (JIT)

[SPU Distribute on {(PD.PAYMENT_REPORT)}] [HashIt for Join]

Node 2.

[SPU Sequential Scan table "PAYMENTREPORT" as "P" {(P."ID")}]

-- Estimated Rows = 100000000, Width = 8, Cost = 0.0 ... 3595.5, Conf = 80.0 (FACT)

Restrictions:

(P."ID" NOTNULL)

Projections:

1:P.COMPANY HEADQUARTER 2:P."ID"

Node 3.

[SPU Hash Join Stream "Node 2" with Temp "Node 1" {(PD.PAYMENT REPORT,P."ID")}]

-- Estimated Rows = 599005119, Width = 8, Cost = 39316.9 .. 84826.1, Conf = 64.0

Restrictions:

(P."ID" = PD.PAYMENT REPORT)

Projections:

1:P.COMPANY_HEADQUARTER 2:PD.BONUS

Cardinality:

PD.PAYMENT REPORT 85.2M (Adjusted)

Node 4.

[SPU Group]

-- Estimated Rows = 44, Width = 8, Cost = 39316.9 .. 98439.9, Conf = 0.0

Projections:

1:P.COMPANY HEADQUARTER 2:PD.BONUS

[SPU Return]

[HOST Merge Group]

Node 5.

[Host Aggregate]

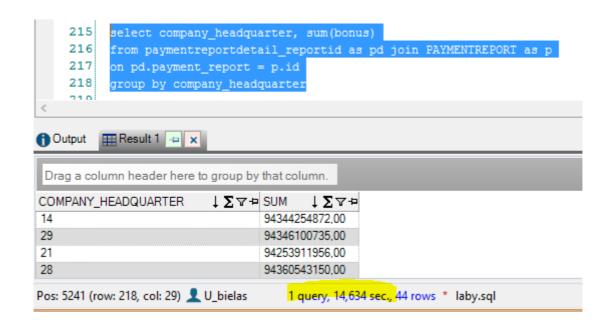
-- Estimated Rows = 44, Width = 20, Cost = 39316.9 ... 98439.9, Conf = 0.0

Projections:

1:P.COMPANY_HEADQUARTER 2:SUM(PD.BONUS)

[Host Return]

Z planu wynika, że jest wykonywana redystrybucja po kluczu paymentreport. Mamy do czynienia z single redistribution. Ponieważ taka sytuacja zachodzi, konieczność owej redystrybucji (wyciąganie odpowiednich danych z data slice'ów i przesłanie ich do odpowiedniego S-Blade'a) wpływa negatywnie na czas wykonania.



OUERY VERBOSE PLAN:

Node 1.

[SPU Sequential Scan table "PAYMENTREPORT" as "P" {(P."ID")}]

-- Estimated Rows = 100000000, Width = 8, Cost = 0.0 ... 3595.5, Conf = 80.0 (FACT)

Restrictions:

(P."ID" NOTNULL)

Projections:

1:P.COMPANY HEADQUARTER 2:P."ID"

[HashIt for Join]

Node 2.

[SPU Sequential Scan table "PAYMENTREPORTDETAIL_REPORTID" as "PD" {(PD.PAYMENT_REPORT)}]

-- Estimated Rows = 599426939, Width = 8, Cost = 0.0 ... 11718.8, Conf = 90.0 (FACT)

Restrictions:

(PD.PAYMENT REPORT NOTNULL)

Projections:

1:PD.BONUS 2:PD.PAYMENT REPORT

Cardinality:

PD.PAYMENT REPORT 100.0M (JIT)

Node 3.

[SPU Hash Join Stream "Node 2" with Temp "Node 1" {(P."ID", PD.PAYMENT REPORT)}]

-- Estimated Rows = 599426939, Width = 8, Cost = 3595.5 .. 72817.7, Conf = 64.0

Restrictions:

(PD.PAYMENT_REPORT = P."ID")

Projections:

1:P.COMPANY HEADQUARTER 2:PD.BONUS

Cardinality:

PD.PAYMENT REPORT 85.2M (Adjusted)

Node 4.

[SPU Group]

-- Estimated Rows = 44, Width = 8, Cost = 3595.5 ... 86441.0, Conf = 0.0

Projections:

1:P.COMPANY HEADQUARTER 2:PD.BONUS

[SPU Return]

[HOST Merge Group]

Node 5.

[Host Aggregate]

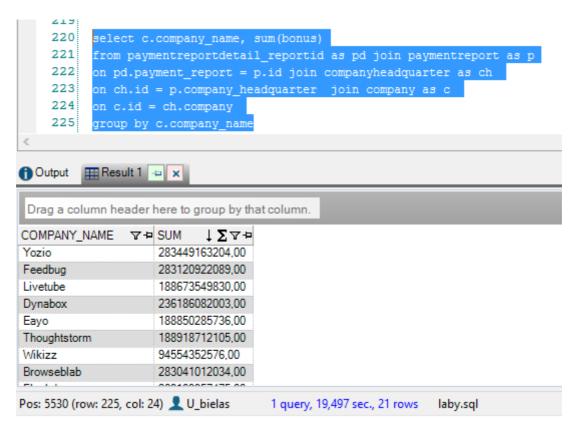
-- Estimated Rows = 44, Width = 20, Cost = 3595.5 .. 86441.0, Conf = 0.0

Projections:

1:P.COMPANY HEADQUARTER 2:SUM(PD.BONUS)

[Host Return]

Czas wykonywania jest krótszy, gdyż nie trzeba dokonywać redystrybucji (jak wynika z planera – nie ma wzmianki distribute on(...)). 8



```
QUERY VERBOSE PLAN:
Node 1.
[SPU Sequential Scan table "PAYMENTREPORTDETAIL REPORTID" as "PD"
{(PD.PAYMENT REPORT)}]
-- Estimated Rows = 599426939, Width = 8, Cost = 0.0 .. 11718.8, Conf = 90.0 (FACT)
Restrictions:
(PD.PAYMENT REPORT NOTNULL)
Projections:
1:PD.BONUS 2:PD.PAYMENT REPORT
Cardinality:
PD.PAYMENT REPORT 100.0M (JIT)
[HashIt for Join]
Node 2.
[SPU Sequential Scan table "COMPANYHEADQUARTER" as "CH" {(CH."ID")}]
-- Estimated Rows = 43, Width = 8, Cost = 0.0 ... 0.0, Conf = 100.0
Restrictions:
((CH."ID" NOTNULL) AND (CH.COMPANY NOTNULL))
Projections:
1:CH."ID" 2:CH.COMPANY
[HashIt for Join]
Node 3.
[SPU Sequential Scan table "COMPANY" as "C" {(C."ID")}]
-- Estimated Rows = 22, Width = 54, Cost = 0.0 ... 0.0, Conf = 80.0
Restrictions:
(C."ID" NOTNULL)
Projections:
1:C.COMPANY NAME 2:C."ID"
Node 4.
[SPU Hash Join Stream "Node 3" with Temp "Node 2" {(C."ID")}]
-- Estimated Rows = 22, Width = 54, Cost = 0.0 ... 0.0, Conf = 64.0
Restrictions:
(C."ID" = CH.COMPANY)
Projections:
1:CH."ID" 2:C.COMPANY NAME
Cardinality:
CH."ID" 22 (Adjusted)
CH.COMPANY 22 (Adjusted)
[SPU Broadcast]
[HashIt for Join]
```

-- Estimated Rows = 99935252, Width = 8, Cost = 0.0 .. 1139.2, Conf = 90.0 [BT: MaxPages=890 TotalPages=19580] (JIT-Stats)

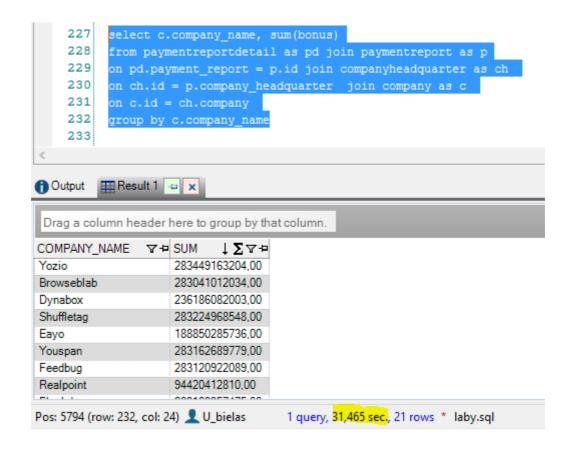
[SPU Sequential Scan table "PAYMENTREPORT" as "P" {(P."ID")}]

Node 5.

```
Restrictions:
((P."ID" NOTNULL) AND (P.COMPANY HEADQUARTER NOTNULL))
Projections:
1:P."ID" 2:P.COMPANY HEADQUARTER
Cardinality:
P.COMPANY HEADQUARTER 43 (Adjusted)
Node 6.
[SPU Hash Join Stream "Node 5" with Temp "Node 4" {(P."ID")}]
-- Estimated Rows = 51129664, Width = 54, Cost = 0.0 ... 10451.4, Conf = 51.2
Restrictions:
(P.COMPANY HEADQUARTER = CH."ID")
Projections:
1:C.COMPANY NAME 2:P."ID"
Cardinality:
CH."ID" 22 (Adjusted)
CH.COMPANY 22 (Adjusted)
P.COMPANY HEADQUARTER 22 (Adjusted)
Node 7.
[SPU Hash Join Stream "Node 6" with Temp "Node 1" {(PD.PAYMENT REPORT,P."ID")}]
-- Estimated Rows = 302930461, Width = 54, Cost = 36343.7 .. 76782.2, Conf = 46.1
Restrictions:
(P."ID" = PD.PAYMENT REPORT)
Projections:
1:C.COMPANY NAME 2:PD.BONUS
Cardinality:
CH.COMPANY 22 (Adjusted)
Node 8.
[SPU Group]
-- Estimated Rows = 22, Width = 54, Cost = 36343.7 ... 83667.0, Conf = 0.0
Projections:
1:C.COMPANY NAME 2:PD.BONUS
[SPU Return]
[HOST Merge Group]
Node 9.
[Host Aggregate]
-- Estimated Rows = 22, Width = 66, Cost = 36343.7 .. 83667.0, Conf = 0.0
Projections:
```

1:C.COMPANY NAME 2:SUM(PD.BONUS)

[Host Return]



QUERY VERBOSE PLAN:

Node 1.

[SPU Sequential Scan table "PAYMENTREPORTDETAIL" as "PD" {(PD."ID")}]

-- Estimated Rows = 599005119, Width = 8, Cost = 0.0 .. 11718.8, Conf = 90.0 (FACT)

Restrictions:

(PD.PAYMENT_REPORT NOTNULL)

Projections:

1:PD.BONUS 2:PD.PAYMENT REPORT

Cardinality:

PD.PAYMENT REPORT 100.0M (JIT)

[SPU Distribute on {(PD.PAYMENT_REPORT)}]

[HashIt for Join]

Node 2.

[SPU Sequential Scan table "COMPANYHEADQUARTER" as "CH" {(CH."ID")}]

-- Estimated Rows = 43, Width = 8, Cost = 0.0 .. 0.0, Conf = 100.0

Restrictions:

((CH."ID" NOTNULL) AND (CH.COMPANY NOTNULL))

Projections:

1:CH."ID" 2:CH.COMPANY

[HashIt for Join]

Node 3.

[SPU Sequential Scan table "COMPANY" as "C" {(C."ID")}]

-- Estimated Rows = 22, Width = 54, Cost = 0.0 ... 0.0, Conf = 80.0

```
Restrictions:
(C."ID" NOTNULL)
Projections:
1:C.COMPANY NAME 2:C."ID"
Node 4.
[SPU Hash Join Stream "Node 3" with Temp "Node 2" {(C."ID")}]
-- Estimated Rows = 22, Width = 54, Cost = 0.0 ... 0.0, Conf = 64.0
Restrictions:
(C."ID" = CH.COMPANY)
Projections:
1:CH."ID" 2:C.COMPANY NAME
Cardinality:
CH."ID" 22 (Adjusted)
CH.COMPANY 22 (Adjusted)
[SPU Broadcast]
[HashIt for Join]
Node 5.
[SPU Sequential Scan table "PAYMENTREPORT" as "P" {(P."ID")}]
-- Estimated Rows = 99935252, Width = 8, Cost = 0.0 .. 1139.2, Conf = 90.0 [BT: MaxPages=890
TotalPages=19580] (JIT-Stats)
Restrictions:
((P."ID" NOTNULL) AND (P.COMPANY HEADQUARTER NOTNULL))
Projections:
1:P."ID" 2:P.COMPANY HEADQUARTER
Cardinality:
P.COMPANY HEADQUARTER 43 (Adjusted)
Node 6.
[SPU Hash Join Stream "Node 5" with Temp "Node 4" {(P."ID")}]
-- Estimated Rows = 51129664, Width = 54, Cost = 0.0 ... 10451.4, Conf = 51.2
Restrictions:
(P.COMPANY HEADQUARTER = CH."ID")
Projections:
1:C.COMPANY NAME 2:P."ID"
Cardinality:
CH."ID" 22 (Adjusted)
CH.COMPANY 22 (Adjusted)
P.COMPANY HEADQUARTER 22 (Adjusted)
Node 7
[SPU Hash Join Stream "Node 6" with Temp "Node 1" {(PD.PAYMENT REPORT,P."ID")}]
-- Estimated Rows = 302717287, Width = 54, Cost = 36327.2 .. 86000.2, Conf = 46.1
Restrictions:
(P."ID" = PD.PAYMENT REPORT)
Projections:
1:C.COMPANY NAME 2:PD.BONUS
```

Cardinality:

CH.COMPANY 22 (Adjusted)

Node 8.

[SPU Group]

-- Estimated Rows = 22, Width = 54, Cost = 36327.2 ... 92880.1, Conf = 0.0

Projections:

1:C.COMPANY NAME 2:PD.BONUS

[SPU Return]

[HOST Merge Group]

Node 9.

[Host Aggregate]

-- Estimated Rows = 22, Width = 66, Cost = 36327.2 ... 92880.1, Conf = 0.0

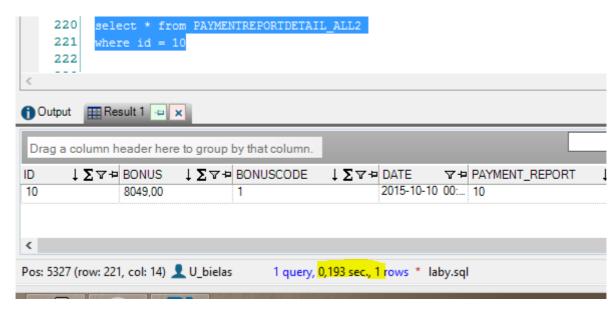
Projections:

1:C.COMPANY_NAME 2:SUM(PD.BONUS)

[Host Return]

W obu powyższych sytuacjach pojawia się wzorzec, o którym wspomniano w zad.7: tam gdzie jest konieczność zastosowania redystrybucji, negatywnie (czasowo) odbija się to na wydajności.

9. ciekawostka:



QUERY VERBOSE PLAN:

Node 1.

[SPU Sequential Scan table "PAYMENTREPORTDETAIL_ALL2" {(PAYMENTREPORTDETAIL ALL2."ID")}]

-- Estimated Rows = 32000000, Width = 24, Cost = 0.0 .. 68181.8, Conf = 80.0

Restrictions:

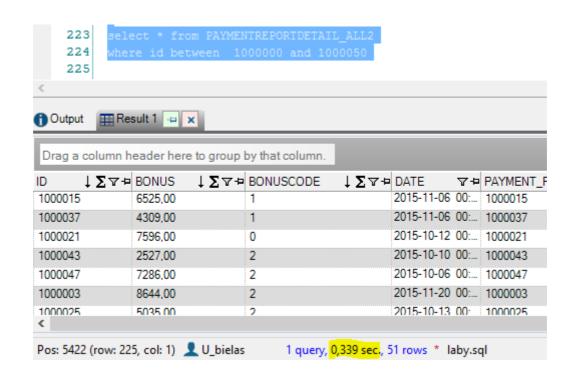
(PAYMENTREPORTDETAIL ALL2."ID" = 10)

Projections:

- 1:PAYMENTREPORTDETAIL ALL2."ID" 2:PAYMENTREPORTDETAIL ALL2.BONUS
- 3:PAYMENTREPORTDETAIL ALL2.BONUSCODE
- 4:PAYMENTREPORTDETAIL ALL2.DATE
- 5:PAYMENTREPORTDETAIL ALL2.PAYMENT REPORT

[SPU Return]

[Host Return]



QUERY VERBOSE PLAN:

Node 1.

[SPU Sequential Scan table "PAYMENTREPORTDETAIL_ALL2" {(PAYMENTREPORTDETAIL ALL2."ID")}]

-- Estimated Rows = 51, Width = 24, Cost = 0.0 ... 68181.8, Conf = 64.0

Restrictions:

((PAYMENTREPORTDETAIL_ALL2."ID" <= 1000050) AND (PAYMENTREPORTDETAIL_ALL2."ID" >= 1000000))

Projections:

- 1:PAYMENTREPORTDETAIL ALL2."ID" 2:PAYMENTREPORTDETAIL ALL2.BONUS
- 3:PAYMENTREPORTDETAIL ALL2.BONUSCODE
- 4:PAYMENTREPORTDETAIL ALL2.DATE
- 5:PAYMENTREPORTDETAIL ALL2.PAYMENT REPORT

[SPU Return]

[Host Return]

Czasy wyszukiwania jednego rekordu są porównywalne z czasem wyszukiwania 50 rekordów. Jest to prawdopodobnie zasługa zastosowanych zone map.