

## Design Analysis and Algorithm – Lab Work

### Week 6

#### Question 1: Write a Program to perform Quick Sort using

##### a. first element

METHOD:

classmate  
Date \_\_\_\_\_  
Page \_\_\_\_\_

Quick sort : first element as pivot

Step 1: 

157	110	147	122	111	149	151	141	123	132	117	133
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

  
P  
Swap 157 with 133

133	110	147	122	111	149	151	141	123	132	117	157
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

All elements are less than pivot so, no swapping in it

Step 2: 

133	110	147	122	111	149	151	141	123	132	117	157
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

  
P      ele > P      swap(147, 117)      ele < P

Step 3: 

133	110	147	122	111	149	151	141	123	132	117	157
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

  
P      ele < P      swap(149, 151)      ele < P      swap(141, 123)      ele < P

Step 4: 

133	110	117	122	111	112	151	141	123	149	147	157
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

  
P      ele < P      swap(151, 123)      ele < P

Step 5: 

133	110	117	122	111	112	123	141	151	149	147	157
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

  
P      ele < P      ele < P

largest  
The element from left has greater index than the smallest element from right so, swap the smallest with pivot.

Step 6: 

123	110	117	122	111	112	133	141	151	149	147	157
-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----

  
P      swap(123, 112)      P      No swap needed



```

        while (arr[j] > pivot)
            j--;

        if (i < j) {
            int temp = arr[i];
            arr[i] = arr[j];
            arr[j] = temp;
        }
    }

    int temp = arr[low];
    arr[low] = arr[j];
    arr[j] = temp;

    return j;
}

void quickSort(int arr[], int low, int high) {
    if (low < high) {
        int p = partition(arr, low, high);
        quickSort(arr, low, p - 1);
        quickSort(arr, p + 1, high);
    }
}

int main() {
    int arr[] = {10, 7, 8, 9, 1, 5};
    int n = sizeof(arr) / sizeof(arr[0]);

    quickSort(arr, 0, n - 1);

    printf("Sorted array:\n");
    for (int i = 0; i < n; i++)
        printf("%d ", arr[i]);

    return 0;
}

```

OUTPUT:

```

C:\Sem-4\DAA\Week-6>gcc QuickSort_1st.c

C:\Sem-4\DAA\Week-6>.\a
Sorted array:
110 111 112 117 122 123 133 141 147 149 151 157

```

## b. middle element

METHOD:

Handwritten notes illustrating the partitioning process for a middle element pivot.

**First Partition:**

Array: 157 110 147 122 111 149 151 141 123 112 117 133

Step 1: Last element as pivot: 133. Array: 157 110 147 122 111 149 151 141 123 112 117 133. Pivot (P) is 133. Elements less than P are 110, 112, 117. Elements greater than P are 157, 147, 122, 111, 149, 151, 141. Swap (147, 117).

Step 2: Array: 157 110 117 122 111 149 151 141 123 112 147 133. Pivot (P) is 133. Elements less than P are 110, 112, 117. Elements greater than P are 157, 149, 151, 141, 123, 147, 122, 111. Swap (147, 117).

**Second Partition:**

Array: 157 110 117 122 111 149 151 141 123 112 117 133

Step 1: Last element as pivot: 133. Array: 157 110 117 122 111 149 151 141 123 112 117 133. Pivot (P) is 133. Elements less than P are 110, 112, 117. Elements greater than P are 157, 149, 151, 141, 123, 147, 122, 111. Swap (147, 117).

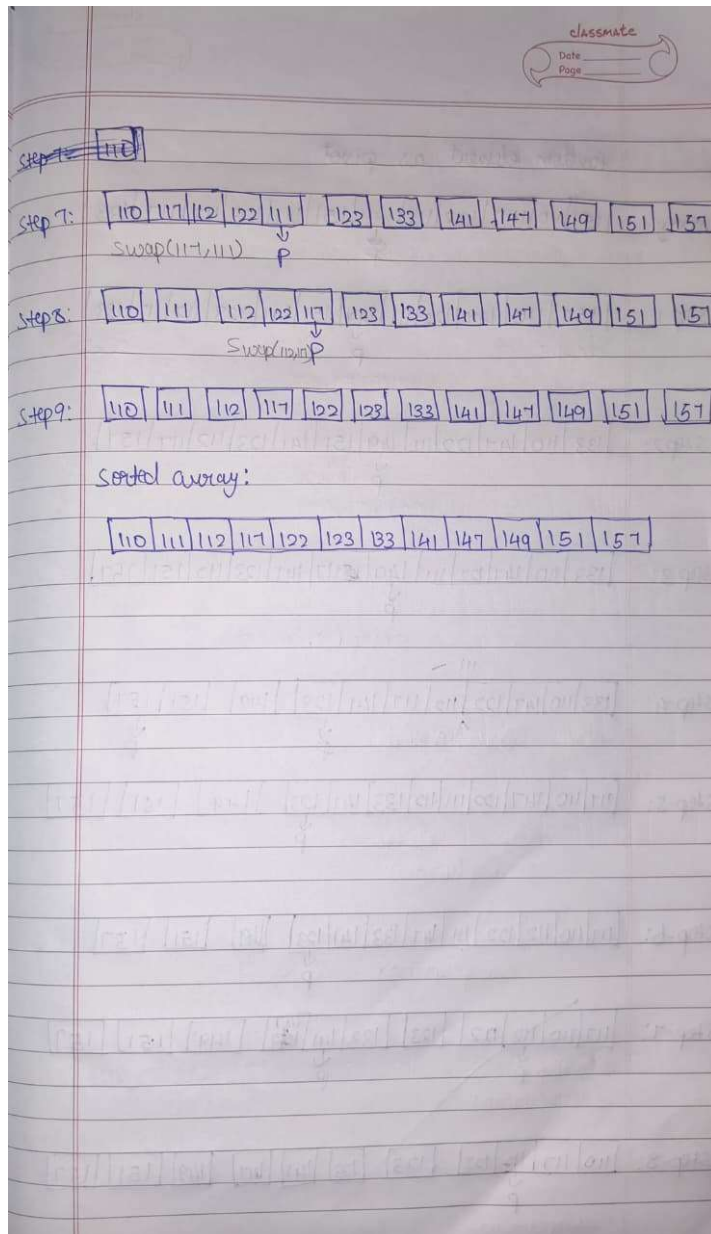
Step 2: Array: 117 110 147 122 111 149 151 141 123 112 157 133. Pivot (P) is 133. Elements less than P are 110, 112, 117. Elements greater than P are 157, 149, 151, 141, 123, 147, 122, 111. Swap (147, 112).

Step 3: Array: 117 110 112 122 111 149 151 141 123 147 157 133. Pivot (P) is 133. Elements less than P are 110, 112, 117. Elements greater than P are 157, 149, 151, 141, 123, 147, 122, 111. Swap (149, 123).

Step 4: Array: 117 110 112 122 111 123 151 141 149 147 157 133. Pivot (P) is 133. Elements less than P are 110, 112, 117. Elements greater than P are 157, 149, 151, 141, 123, 147, 122, 111. Swap (151, 133).

Step 5: Array: 117 110 112 122 111 123 133 141 149 147 157 151. Pivot (P) is 133. Elements less than P are 110, 112, 117. Elements greater than P are 157, 149, 151, 141, 123, 147, 122, 111. Swap (157, 151).

Step 6: Array: 117 110 112 122 111 123 133 141 149 147 151 157. Pivot (P) is 133. Elements less than P are 110, 112, 117. Elements greater than P are 157, 149, 151, 141, 123, 147, 122, 111. Swap (147, 110).



CODE:

```
#include <stdio.h>

int partition(int arr[], int low, int high) {
    int pivot = arr[high];
    int i = low - 1;

    for (int j = low; j < high; j++) {
        if (arr[j] <= pivot) {
            i++;
            int temp = arr[i];
            arr[i] = arr[j];
        }
    }
}
```

```

        arr[j] = temp;
    }
}

int temp = arr[i + 1];
arr[i + 1] = arr[high];
arr[high] = temp;

return i + 1;
}

void quickSort(int arr[], int low, int high) {
    if (low < high) {
        int p = partition(arr, low, high);
        quickSort(arr, low, p - 1);
        quickSort(arr, p + 1, high);
    }
}

int main() {
    int arr[] = {10, 7, 8, 9, 1, 5};
    int n = sizeof(arr) / sizeof(arr[0]);

    quickSort(arr, 0, n - 1);

    printf("Sorted array:\n");
    for (int i = 0; i < n; i++)
        printf("%d ", arr[i]);

    return 0;
}

```

OUTPUT:

```
C:\Sem-4\DAA\Week-6>gcc QuickSort_last.c
```

```
C:\Sem-4\DAA\Week-6>.\a
```

```
Sorted array:
```

```
110 111 112 117 122 123 133 141 147 149 151 157
```

c. random element.

METHOD:

random element as pivot

array: 157 110 147 122 111 149 151 141 123 112 117 133

step 1: 157 110 147 122 111 149 151 141 123 112 117 133  
 ↓ P  
 swap(157, 133)

step 2: 133 110 147 122 111 149 151 141 123 112 117 157  
 ↓ P  
 swap(151, 117)

step 3: 133 110 147 122 111 149 157 141 123 112 151 157  
 ↓ P  
 swap(149, 112)

step 4: 133 110 147 122 112 117 141 123 149 151 157  
 ↓ P  
 swap(149, 151)

step 5: 117 110 147 122 111 112 133 141 123 149 151 157  
 ↓ P  
 swap(147, 151)

step 6: 117 110 112 122 111 147 133 141 123 149 151 157  
 ↓ P  
 swap(147, 123)

step 7: 117 110 112 122 123 133 141 147 149 151 157  
 ↓ P  
 swap(117, 110)

step 8: 110 117 112 122 123 133 141 147 149 151 157  
 ↓ P  
 swap(117, 112)



```

        int temp = arr[i];
        arr[i] = arr[j];
        arr[j] = temp;
    }
}

    int temp = arr[i + 1];
    arr[i + 1] = arr[high];
    arr[high] = temp;

    return i + 1;
}

int randomPartition(int arr[], int low, int high) {
    int randomIndex = low + rand() % (high - low + 1);

    int temp = arr[randomIndex];
    arr[randomIndex] = arr[high];
    arr[high] = temp;

    return partition(arr, low, high);
}

void quickSort(int arr[], int low, int high) {
    if (low < high) {
        int p = randomPartition(arr, low, high);
        quickSort(arr, low, p - 1);
        quickSort(arr, p + 1, high);
    }
}

int main() {
    srand(time(NULL));

    int arr[] = {10, 7, 8, 9, 1, 5};
    int n = sizeof(arr) / sizeof(arr[0]);

    quickSort(arr, 0, n - 1);

    printf("Sorted array:\n");
    for (int i = 0; i < n; i++)
        printf("%d ", arr[i]);

    return 0;
}

```

OUTPUT:

```
C:\Sem-4\DAA\Week-6>gcc QuickSort_random.c
C:\Sem-4\DAA\Week-6>.\a
Sorted array:
110 111 112 117 122 123 133 141 147 149 151 157
```

### Time Complexity:

Quick Sort has a best-case and average-case time complexity of  $O(n \log n)$ , which occurs when the pivot element divides the array into two nearly equal subarrays at each recursive step. However, in the worst case, when the pivot selection results in highly unbalanced partitions (such as when the smallest or largest element is consistently chosen as the pivot), the time complexity degrades to  $O(n^2)$ . Although different pivot selection strategies can reduce the likelihood of the worst case, they do not change the theoretical time complexity bounds of the algorithm.

### Space Complexity:

Quick Sort is an in-place sorting algorithm and therefore requires  $O(1)$  auxiliary space for data storage. However, due to its recursive nature, additional space is required for the recursion stack. In the best and average cases, the depth of recursion is  $O(\log n)$ , resulting in a space complexity of  $O(\log n)$ , whereas in the worst case of highly unbalanced partitions, the recursion depth can grow to  $O(n)$ , leading to a space complexity of  $O(n)$ .