Informatics II, Spring 2024, Exercise 3

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Learning Goal

- Understand the efficiency of algorithms like binary search.
- Learn how to analyze algorithmic complexity and asymptotic complexity of algorithms.
- Learn how to analyze special case and correctness of algorithms.

Task 1: Linear Search and Binary Search [Easy]

Consider an array A with n distinct integers that are sorted in an ascending order and an integer t.

- a) The C funtion linear_search traverses the integers in A, one after another, from the beginning. If t is found in A linear_search returns 1, otherwise 0. Complete the C funtion linear_search(int A[], int n, int t) in task1.c file.
- b) The C funtion binary_search(int A[], int n, int t) that employs binary search to find integer t in A. Reference the following pseudocode for binary search to implement the binary_search function. If t is found in A binary_search returns 1, otherwise 0. Complete the C funtion binary_search(int A[], int n, int t) in task1.c file.

```
Algo: BinSearch1(A,v)

Input: sequence A[1..n] of length n, value v

Output: either index i such that v == A[i] or NIL

l = 1; r = n;

m = \lfloor (l+r)/2 \rfloor;

while l \le r \land v \ne A[m] do

if v < A[m] then r = m-1 else l = m+1;

m = \lfloor (l+r)/2 \rfloor;

if l \le r then return m else return NIL;
```

c) What are the asymptotic complexity for the C funtions linear_search and binary_search.

- d) Compile task1.c file. Run your codes with the following parameters for n and t:
 - n = 1000000, t = 1000000
 - n = 10000000, t = 10000000
 - n = 1000000000, t = 1000000000.

Report the run time growth for linear_search and binary_search, respectively.

Task 2: Algorithmic Complexity [Easy]

Below is a pseudocode of a function named whatDoesItDo, which takes an array A[1..n] of n integers and an integer k as inputs.

Note: In the above pseudocode, for j = i to n by k means we do not increase j by 1, but each time, we increase it by k, i.e., j=j+k.

- a) Perform exact analysis of the running time of the algorithm.
- b) Determine the asymptotic complexity of the algorithm?

Task 3: Asymptotic Complexity [Easy]

a) Calculate the asymptotic tight bound for the following functions and rank them by their order of growth (lowest first). Clearly work out the calculation step by step in your solution.

$$f_1(n) = (2n+3)!$$

$$f_2(n) = 2\log(6^{\log n^2}) + \log(\pi n^2) + n^3$$

$$f_3(n) = 4^{\log_2 n}$$

$$f_4(n) = 12\sqrt{n} + 10^{223} + \log 5^n$$

$$f_5(n) = 10^{\log 20} n^4 + 8^{229} n^3 + 20^{231} n^2 + 128n \log n$$

$$f_6(n) = \log n^{2n+1}$$

$$f_7(n) = \log^2(n) + 50\sqrt{n} + \log(n)$$

$$f_8(n) = 14400$$

b) Assume $f_1(n) = O(1)$, $f_2(n) = O(N^2)$, and $f_3(n) = O(N \log N)$. From these complexities it follows that $f_1(n) + f_2(n) + f_3(n) = O(N \log N)$.

Answer:	☐ True	\Box False

Task 4: Special Case and Correctness Analysis [Medium]

Consider the algorithm algo1. The input parameters are an array A[1..n] with n distinct integers and $k \leq n$.

```
Algo: algo1(A, n, k)

sum = 0;

for i = 1 to k do

maxi = i;

for j = i to n do

if A[j] > A[maxi] \text{ then}

maxi = j;

sum = sum + A[maxi];

swp = A[i];

A[i] = A[maxi];

A[maxi] = swp;

return sum
```

- a) Specify the pre/post conditions of the algo1 algorithm.
- b) For the two for loops in the algorithm:
 - i. Determine if the loop is up loop or down loop.
 - ii. Determine the invariants of these two loops and verify whether they are hold in three stages: initialization, maintenance and termination.
- c) Identify some edges cases of the algorithm and verify if the algorithm has the correct output.
- d) Conduct an exact analysis of the running time of algorithm algo1.
- e) Determine the best and the worst case of the algorithm. What is the running time and asymptotic complexity in each case?