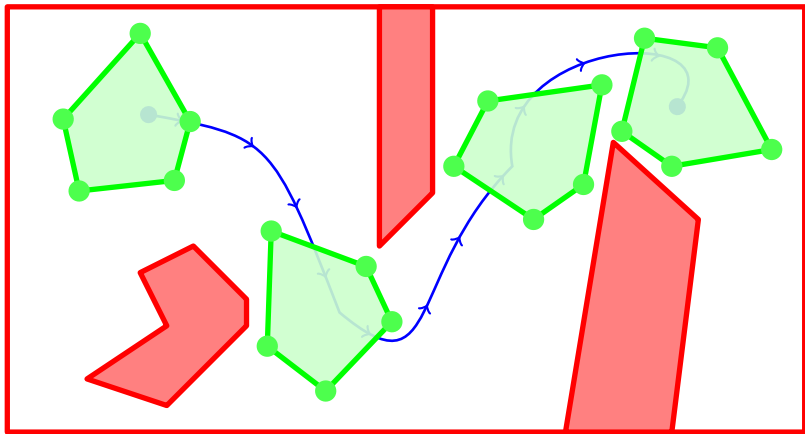


ON SAMPLING BASED ALGORITHMS FOR SOLVING THE MOTION PLANNING PROBLEM

Dror Atariah @ Game Duell

September 18th, 2014

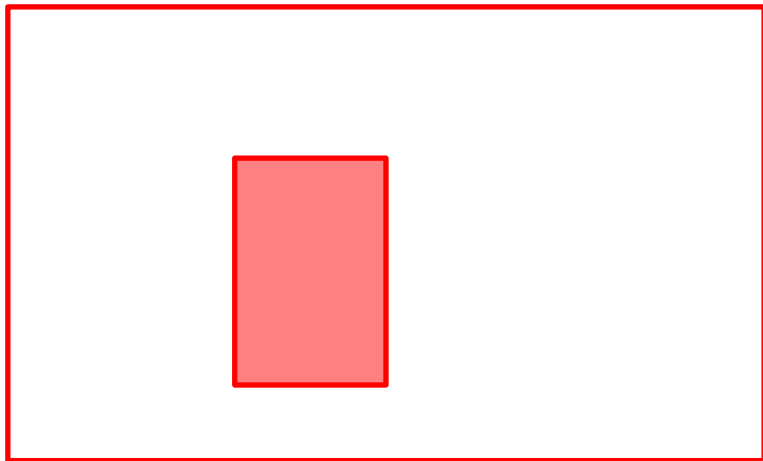
THE MOTION PLANNING PROBLEM



FROM A WORKSPACE TO A CONFIGURATION SPACE

SIMPLE CASE

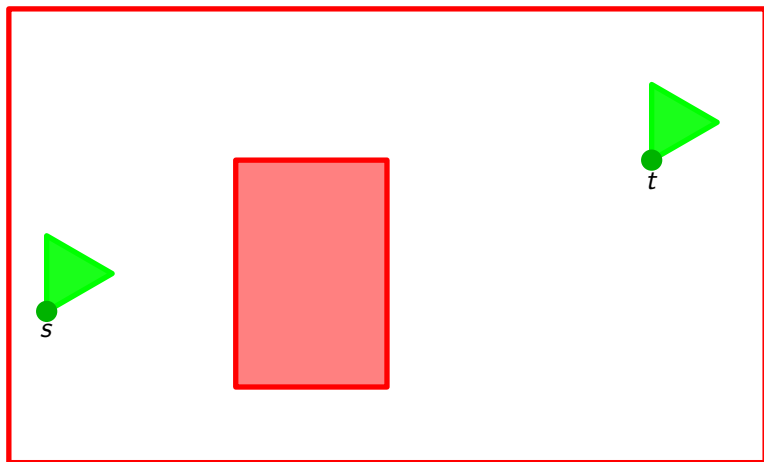
The robot can only translate (2 degrees of freedom)



FROM A WORKSPACE TO A CONFIGURATION SPACE

SIMPLE CASE

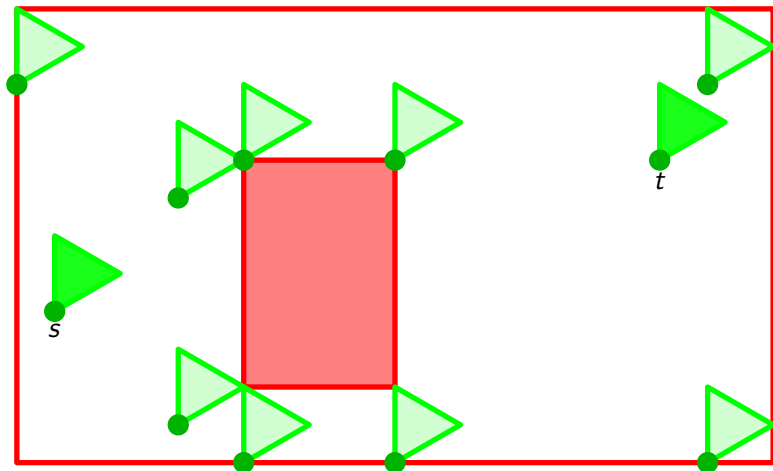
The robot can only translate (2 degrees of freedom)



FROM A WORKSPACE TO A CONFIGURATION SPACE

SIMPLE CASE

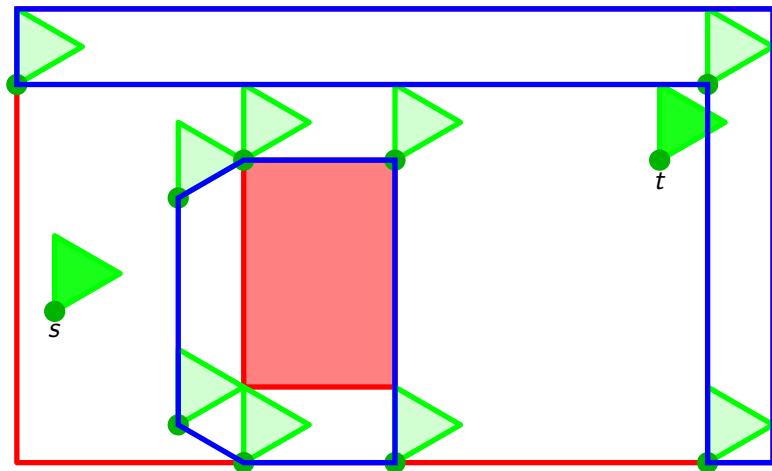
The robot can only translate (2 degrees of freedom)



FROM A WORKSPACE TO A CONFIGURATION SPACE

SIMPLE CASE

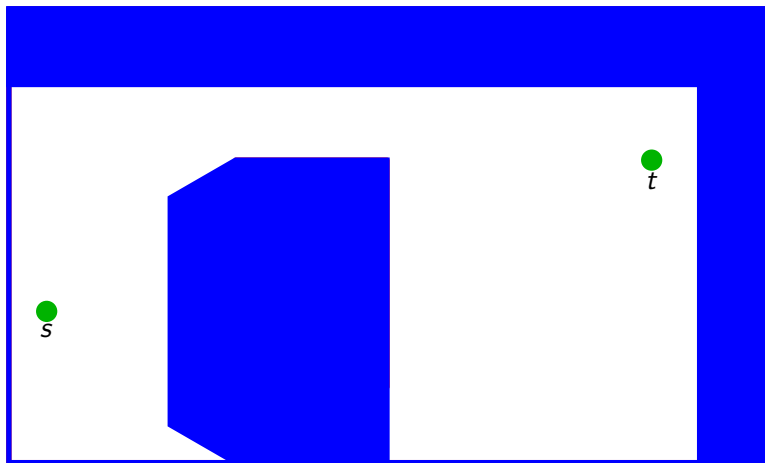
The robot can only translate (2 degrees of freedom)



FROM A WORKSPACE TO A CONFIGURATION SPACE

SIMPLE CASE

The robot can only translate (2 degrees of freedom)

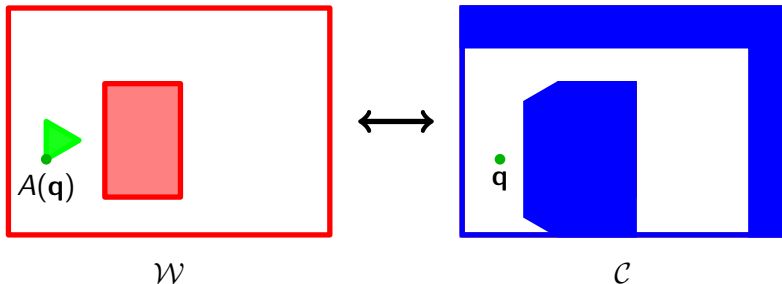


WORKSPACE AND CONFIGURATION SPACE

DEFINITION

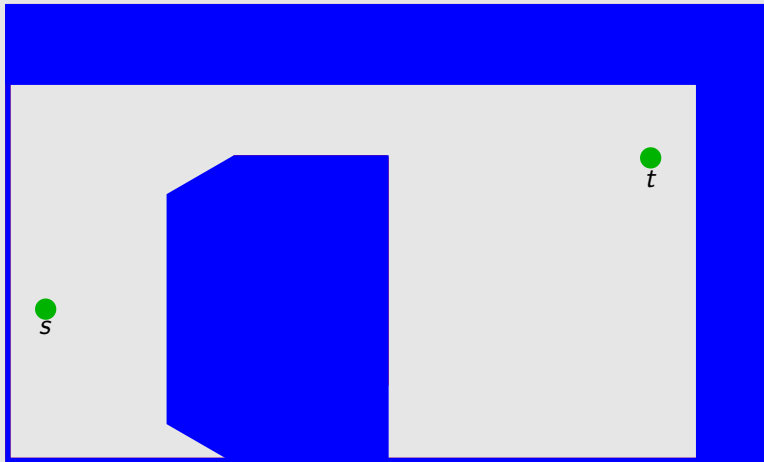
Given a *robot* A , obstacle O and a *workspace* \mathcal{W} , we obtain a *configuration space* denoted by \mathcal{C} .

- $q \in \mathcal{C} \xLeftrightarrow{\text{bijection}} A(q)$
- $\mathcal{C} = \mathcal{C}_{\text{forb}} \cup \mathcal{C}_{\text{free}}$



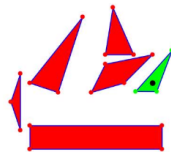
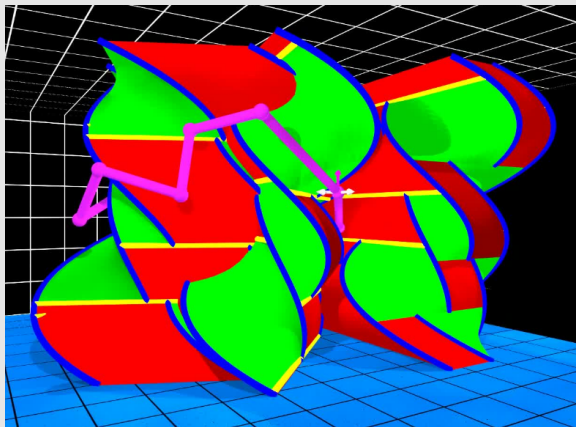
CONFIGURATION SPACE

EXAMPLE



CONFIGURATION SPACE

EXAMPLE



PROBLEM STATEMENT

FORMAL PROBLEM

Navigate a given robot A in a workspace \mathcal{W} that is scattered with obstacles O from a source placement $A(s)$ to a target one $A(t)$.

Equivalently, find a *free path* in \mathcal{C} from s to t .

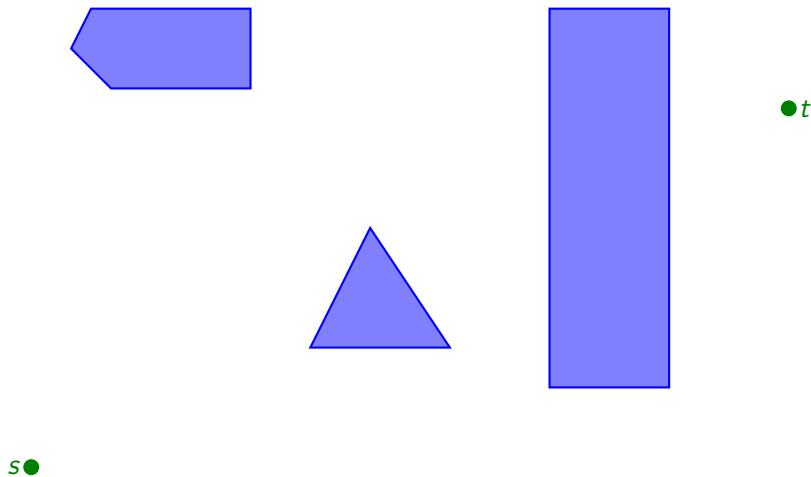
OUTLINE

- 1 PROBABILISTIC ROADMAP (L. KAVRAKI ET AL. 1996)
- 2 MOTION PLANNING VIA MANIFOLD SAMPLES (SALZMAN ET AL. 2013)

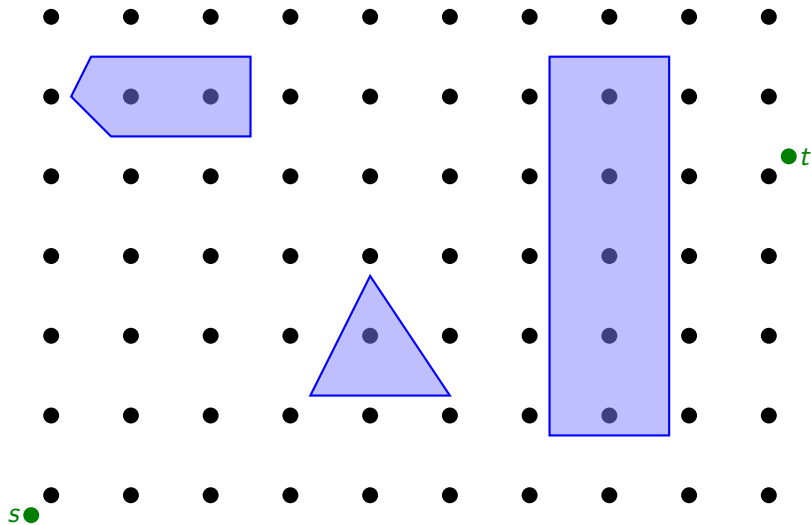
IN A NUT SHELL

- A and \mathcal{W} yield \mathcal{C}
- Sample $\mathcal{C}_{\text{free}}$
- Build a *roadmap graph* $\mathcal{G} = (V, E)$
 - V - sample points in $\mathcal{C}_{\text{free}}$
 - E - free (local) motions
- Connect s and t to \mathcal{G} and find free path

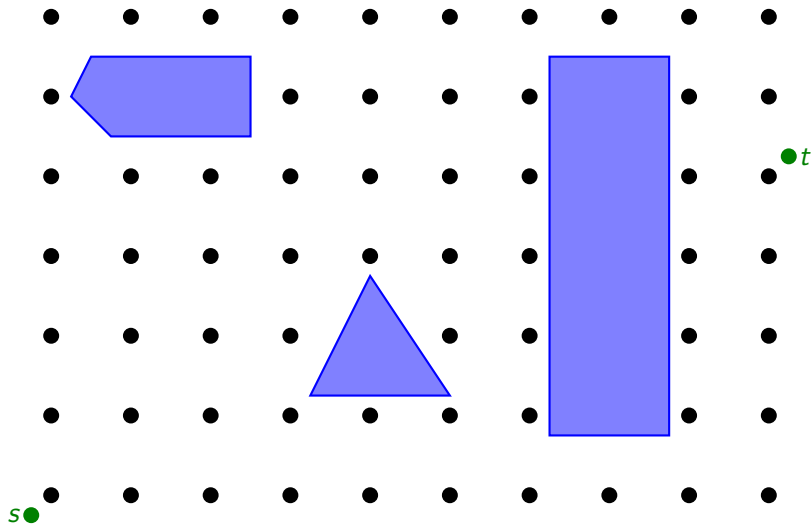
PRM DEMO



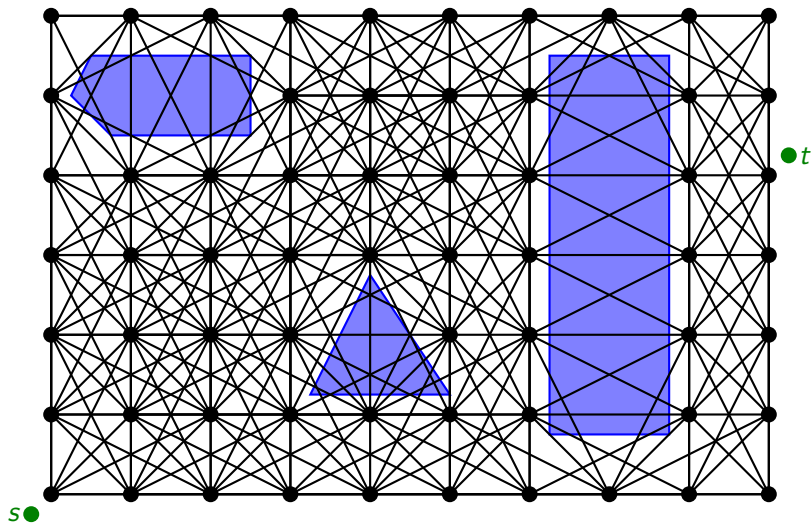
PRM DEMO



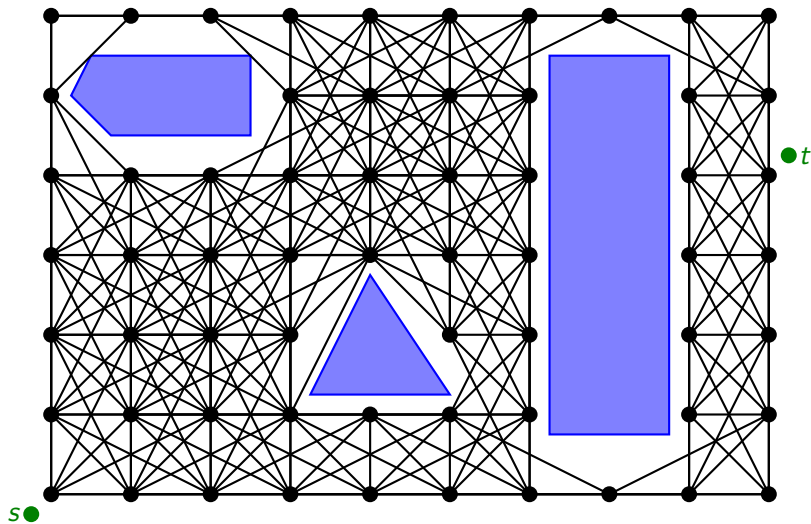
PRM DEMO



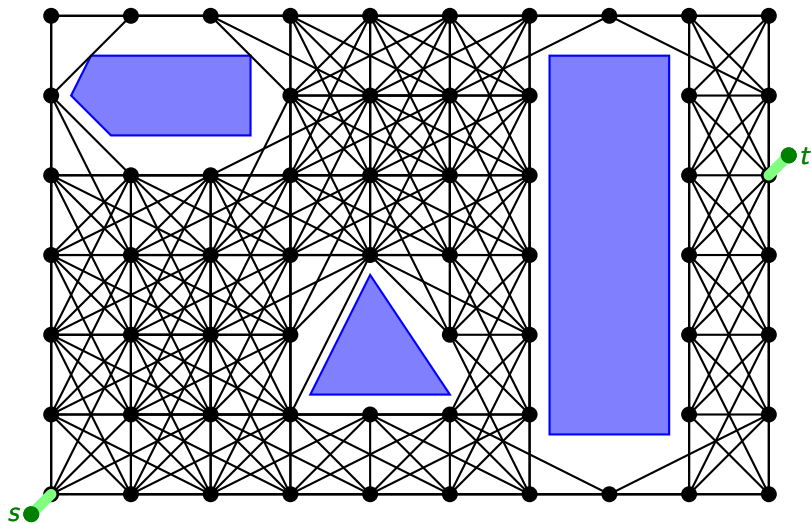
PRM DEMO



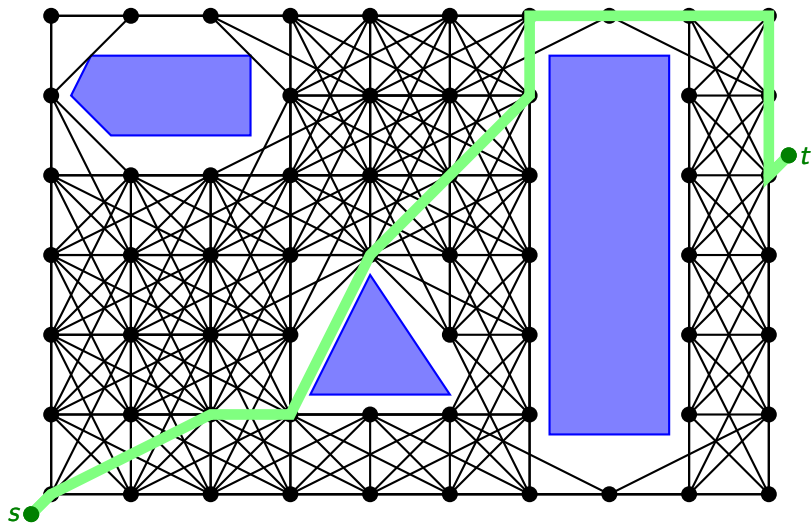
PRM DEMO



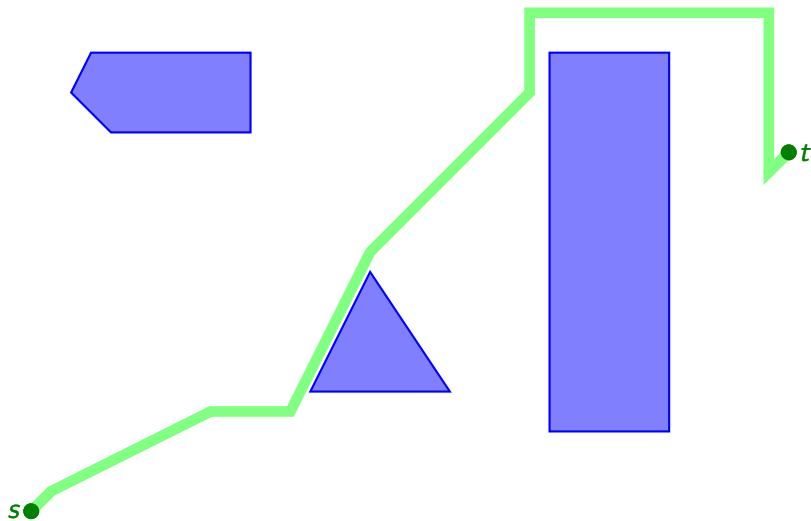
PRM DEMO



PRM DEMO



PRM DEMO



CLASSICAL PRM

L. KAVRAKI ET AL. 1996

Algorithm PRMBuildGraph

```
1:  $V \leftarrow \emptyset$  ;  $E \leftarrow \emptyset$ 
2: loop
3:    $\mathbf{q} \leftarrow$  a point in  $\mathcal{C}_{\text{free}}$ 
4:    $V \leftarrow V \cup \{\mathbf{q}\}$ 
5:    $N_{\mathbf{q}} \leftarrow$  useful neighbors of  $\mathbf{q}$ 
6:   for all  $\mathbf{q}' \in N_{\mathbf{q}}$  do
7:     if  $\mathbf{q}'$  and  $\mathbf{q}$  are not connected in  $\mathcal{G}$  then
8:       if the local planner finds a path between  $\mathbf{q}'$  and  $\mathbf{q}$  then
9:          $E \leftarrow E \cup \overline{\mathbf{q}\mathbf{q}'}$ 
```

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```

PRM IMPLEMENTATION DETAILS

- How to sample?

SAMPLE THE CONFIGURATION SPACE

OBJECTIVE

Obtain a (uniform) sample of $\mathcal{C}_{\text{free}}$

STEPS

- (Randomly) pick $\mathbf{q} \in \mathcal{C}$
- *Collision detection* for $A(\mathbf{q})$

CONFIGURATION POINTS PICKING

- Random
- Grid
- Cell-Based
- Halton points (pseudo random and deterministic!) [more...](#)
- Random-Halton

CLASSICAL PRM L. KAVRAKI ET AL. 1996

Algorithm PRMBuildGraph

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```

PRM IMPLEMENTATION DETAILS

- How to sample?
- Determine whether a sample point is free?

COLLISION DETECTION

DEFINITION (PREDICATE)

$\phi: \mathcal{C} \rightarrow \{\text{TRUE}, \text{FALSE}\}$, s.t.:

$$\phi(\mathbf{q}) = \begin{cases} \text{TRUE} & \text{if } \mathbf{q} \in \mathcal{C}_{\text{forb}} \\ \text{FALSE} & \text{otherwise} \end{cases}$$

DEFINITION (DISTANCE FUNCTION)

The distance in \mathcal{W} between $A(\mathbf{q})$ and the obstacles is measured by $d: \mathcal{C} \rightarrow [0, \infty)$.

TWO-PHASE COLLISION DETECTION

FACT

Collision detection is an expensive task that is repeatedly invoked!

BREAK-UP THE TASK

- Broad Phase (e.g. using bounding boxes)
- Narrow Phase (e.g. using hierarchical methods; See LaValle 2006, § 5.3.2)

CLASSICAL PRM L. KAVRAKI ET AL. 1996

Algorithm PRMBuildGraph

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9:          $E \leftarrow E \cup \overline{\mathbf{q}\mathbf{q}'}$ 
```

PRM IMPLEMENTATION DETAILS

- How to sample?
- Determine whether a sample point is free?
- How to determine whether an *edge* is free or not?

LOCAL PLANNER

OBJECTIVE

Given $\mathbf{q}, \mathbf{q}' \in \mathcal{C}_{\text{free}}$ find a path in $\mathcal{C}_{\text{free}}$ connecting them.

TRADE-OFF

- Accurate and slow
- Inaccurate and quick

LOCAL PLANNER

OBJECTIVE

Given $\mathbf{q}, \mathbf{q}' \in \mathcal{C}_{\text{free}}$ find a path in $\mathcal{C}_{\text{free}}$ connecting them.

TRADE-OFF

- Accurate and slow
- Inaccurate and quick

GENERAL LOCAL PLANNER

L. KAVRAKI ET AL. 1996

- Given $\mathbf{q}, \mathbf{q}' \in \mathcal{C}_{\text{free}}$
- Let $\{\mathbf{c}_i\}$ be points on the *line* connecting \mathbf{q} and \mathbf{q}' , s.t.
 $d(A(\mathbf{c}_i), A(\mathbf{c}_{i+1})) \leq \epsilon$
- $\forall i$: collision detection for $A_\epsilon(\mathbf{c}_i)$

BINARY LOCAL PLANNER

GERAERTS AND OVERMARS 2004

- Given $\mathbf{q}, \mathbf{q}' \in \mathcal{C}_{\text{free}}$
- Let \mathbf{c} be the midpoint of the *line* connecting \mathbf{q} and \mathbf{q}'
 - Collision detection for $A(\mathbf{c})$
 - Repeat on two sides of \mathbf{c}

COMPLETENESS OF THE METHOD

COMPLETENESS OF THE METHOD

DEFINITION

A method is *probabilistically complete* if the probability of answering a query *incorrectly* after building a roadmap tends to zero as the number of samples goes to *infinity*.

INFLUENCING PARAMETERS

- N - number of samples
- $\gamma: [0, L] \rightarrow \mathcal{C}_{\text{free}}$ - free path from $\gamma(0) = \mathbf{s}$ to $\gamma(L) = \mathbf{t}$.
- $R = \inf_{0 \leq t \leq L} \inf_{y \in \mathcal{O}} |\gamma(t) - y|$

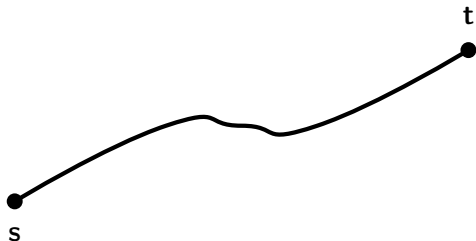
THEOREM (L. E. KAVRAKI ET AL. 1998)

$$P(\text{FAILURE}) \leq \frac{2L}{R} \left(1 - \frac{\pi R^2}{4|\mathcal{C}_{\text{free}}|} \right)^N$$

PRM IS PROBABILISTIC COMPLETE

THEOREM (L. E. KAVRAKI ET AL. 1998)

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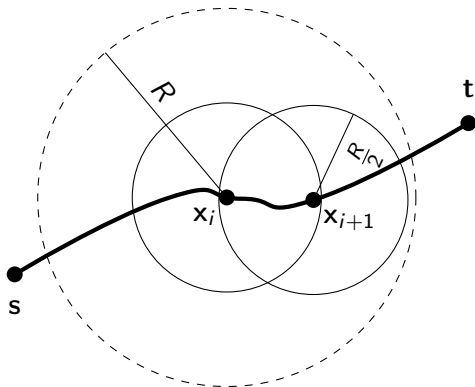


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- $n = \lceil 2L/R \rceil$ points on γ
- $d(\mathbf{x}_i, \mathbf{x}_{i+1}) \leq R/2$ and $B_{R/2}(\mathbf{x}_{i+1}) \subset B_R(\mathbf{x}_i)$

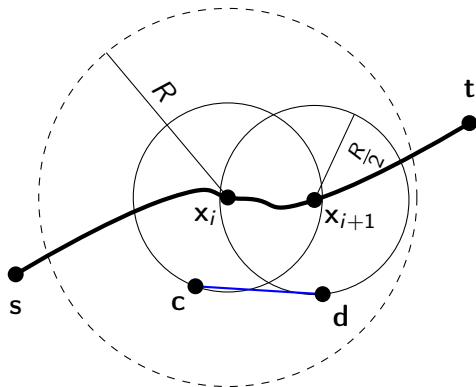


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- $\Rightarrow \overline{\mathbf{cd}}$ is free

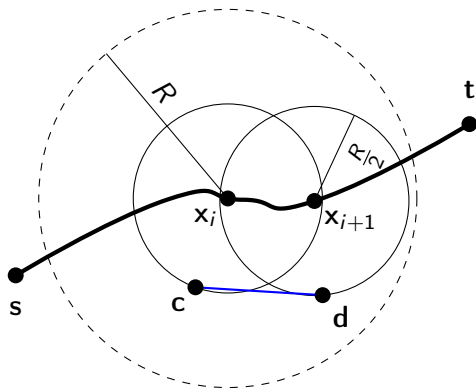


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- $n = \lceil 2L/R \rceil$ points on γ
- $d(\mathbf{x}_i, \mathbf{x}_{i+1}) \leq R/2$ and $B_{R/2}(\mathbf{x}_{i+1}) \subset B_R(\mathbf{x}_i)$
- $\Rightarrow \overline{\mathbf{cd}}$ is free
- PRM succeeds if each $B_{R/2}(\mathbf{x}_j)$ contains a sample
- $P(B_{R/2}(\mathbf{x}_j) \text{ is empty}) = \left(1 - \frac{|B_{R/2}|}{|C_{\text{free}}|} \right)^N$

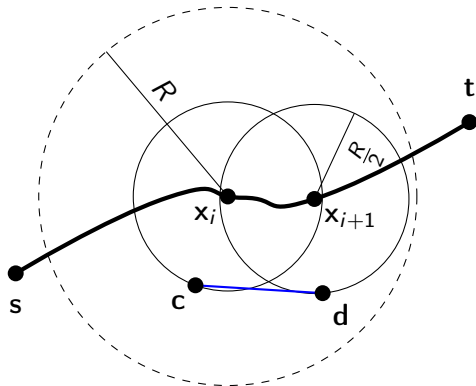


PRM IS PROBABILISTIC COMPLETE

THEOREM (L. E. KAVRAKI ET AL. 1998)

$$P(\text{FAILURE}) \leq \frac{2L}{R} \left(1 - \frac{\pi R^2}{4|C_{\text{free}}|} \right)^N$$

$$\begin{aligned} P(\text{FAILURE}) &\leq P(\text{some ball is empty}) \\ &\leq \sum_{j=1}^{n-1} P(j\text{-th ball is empty}) \\ &= \left(\left\lceil \frac{2L}{R} \right\rceil - 1 \right) \left(1 - \frac{|B_{R/2}|}{|C_{\text{free}}|} \right)^N \end{aligned}$$



IN A NUT SHELL

- A and \mathcal{W} yield \mathcal{C}
- Sample $\mathcal{C}_{\text{free}}$
- Build a *roadmap graph* $\mathcal{G} = (V, E)$
 - V - sample points in $\mathcal{C}_{\text{free}}$
 - E - free (local) motions
- Connect s and t to \mathcal{G} and find free path

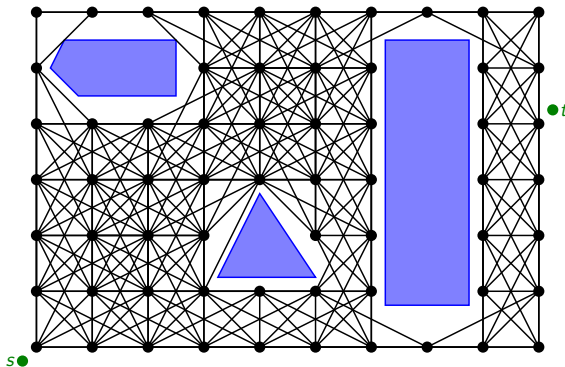
OUTLINE

- 1 PROBABILISTIC ROADMAP (L. KAVRAKI ET AL. 1996)
- 2 MOTION PLANNING VIA MANIFOLD SAMPLES (SALZMAN ET AL. 2013)

MOTION PLANNING VIA MANIFOLD SAMPLES SALZMAN ET AL.

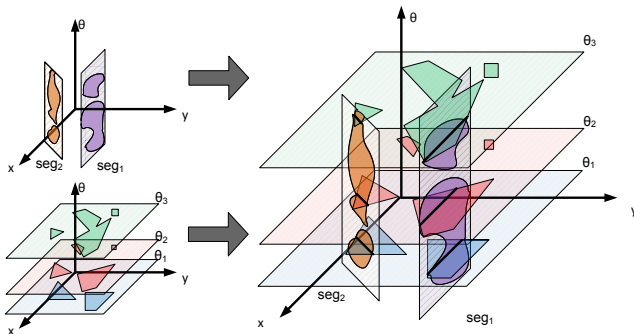
2013

- In PRM connectivity of \mathcal{C} is captured via point samples
- Try to obtain “bigger” samples and improve the capturing



OUTLINE OF MMS

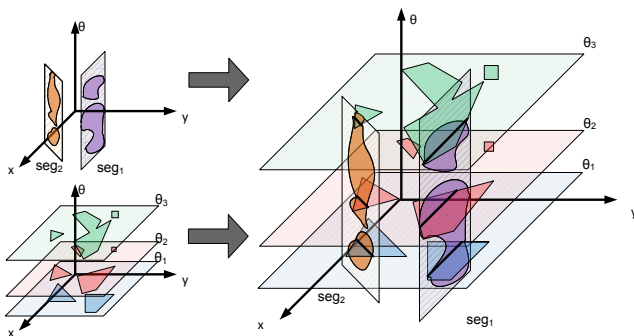
- A and \mathcal{W} yield \mathcal{C}
- Sample \mathcal{C}
- Build a *connectivity graph* $\mathcal{G} = (V, E)$
 - V - Free Space Cells (FSC)
 - E - Between intersecting FSC's
- Connect s and t to \mathcal{G} and find free path



Courtesy of Oren Salzman

PREPROCESSING

- Family of constraints $\Psi \Rightarrow$ manifolds in \mathcal{C}
 Example: Translating and rotating planar robot
 - Horizontal planes
 - Vertical slabs
- Decompose the manifolds into *Free Space Cells*
- Connect, in \mathcal{G} , intersecting FSC's



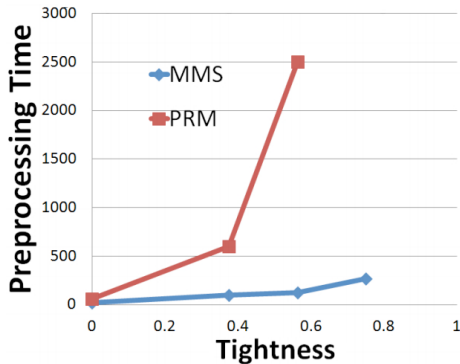
Courtesy of Oren Salzman

DISCUSSION

- MMS is also probabilistic complete

DISCUSSION

- MMS is also probabilistic complete
- No significant improvements in simple cases
- 20-fold speedup in a coordination tight setting
- Significant improvement in tight cases



Courtesy of Oren Salzman

SUMMARY

- Introduced the motion planning problem, and
- Sample based methods to solve it:
 - Probabilistic Roadmap Method (PRM)
 - Motion Planning via Manifold Samples (MMS)
 - Rapidly-exploring Random Trees (RRT)
due to Lavalle and Kuffner 2000
 - PRM*, RRT* and RRG (=Rapidly-exploring Random Graph)
due to Karaman and Frazzoli 2011

THANK YOU FOR YOUR ATTENTION!

- Twitter: @drorata
- LinkedIn: www.linkedin.com/in/atariah

HALTON POINTS

- $k \in \mathbb{Z}$ and a prime $p \Rightarrow k = \sum_{i=0}^r a_i p^i$ s.t. $0 \leq a_i < p$
- Let $\Phi_p(k) = \sum_{i=0}^r \frac{a_i}{p^{i+1}}$
- For primes $p_1 < p_2 < \dots < p_{d-1}$, then the k -th d -dimensional *Halton point* is

$$\left(\frac{k}{n}, \Phi_{p_1}(k), \Phi_{p_2}(k), \dots, \Phi_{p_{d-1}}(k) \right) \in [0, 1]^d,$$

where $k = 0, 1, \dots, n-1$

- For further details see Wong et al. 1997; Chazelle 1998 and LaValle 2006, § 5.2

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