

Parallel DBSCAN on Spark

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Introduction to DBSCAN

Naive parallel DBSCAN

- Naive partition
- Local DBSCAN
- Merge

Optimization on partition

- Reduce-boundry based partition
- Content based partition

Conclusion

Local DBSCAN algorithm

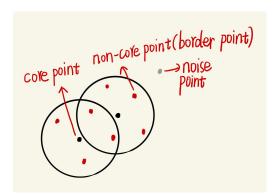


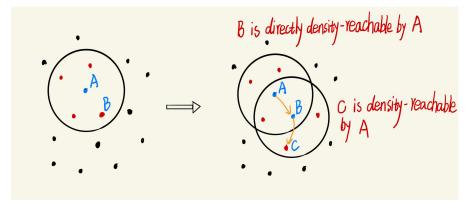


Concepts:

 ϵ (Eps) — the given radius MinPts — the given minimal points if MinPts = 4:









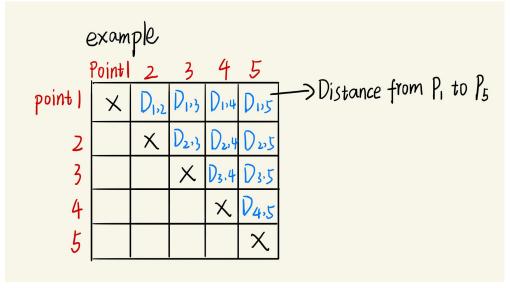
Core idea:

- 1. select a data point p arbitrarily from the data set.
- 2. determine whether p is a core point by Eps and MinPts.If p is a core point, find all the data point which are **density-reachable from p** to form a cluster. If not, select another point to do the same steps as before.

Matrix DBSCAN algorithm



- Introduction: Build a matrix to hold all those point-pair distance value instead of computing every time.
- Pros: In principle, the computational cost could be reduced by half, which is not big algorithm-wise improvement, but a giant increase in engineering, especially it comes to big data problem.



Parallel Matrix DBSCAN algorithm



Two Dimensional Example: (x, y)

Partition:

- define partition number
- 2. define partition block position in the space
 - a. find the minimum x and y and find the maximum x and y
 - b. add eps to upper bound and minus eps to lower bound
 - c. divide the whole area into corresponding partition id (p_id)
- 3. build rdds (pid, dataset_points) by judging whether the points is in the partition area or n

Local DBSCAN:

- 1. each rdd will execute a matrix DBSCAN
- 2. each execution will return ((partition_dataset, core_point list), local_tags)

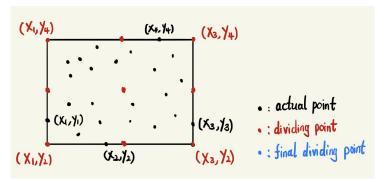
Merge:

- 1. update local tag list into global tag list
- 2. return final global tag list

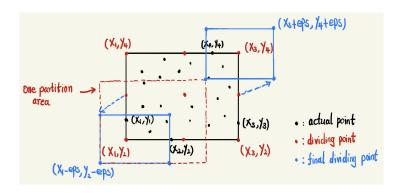
Parallel Matrix DBSCAN algorithm——Partition

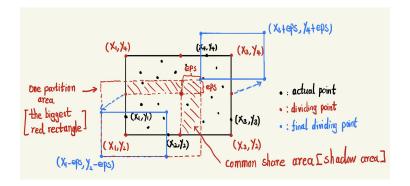


partition number: 4 eps (epsilon): 3 MinPts (minimum points): 2



Define partition block position in the space

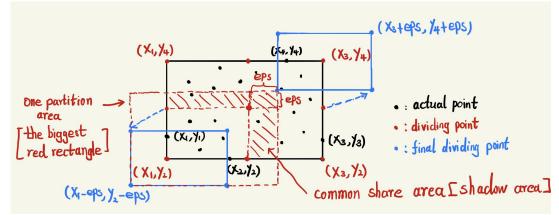




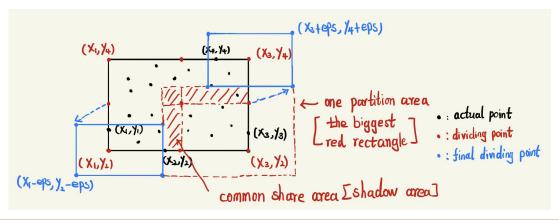
Parallel Matrix DBSCAN algorithm——Local Matrix DBSCAN



partition number: 4 eps (epsilon): 3 MinPts (minimum points): 2



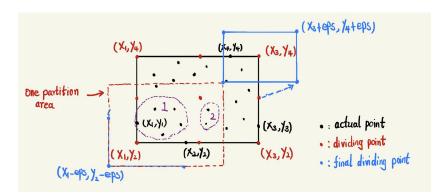
Points in the red shadow area will have multiple labels

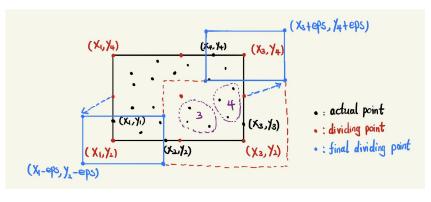


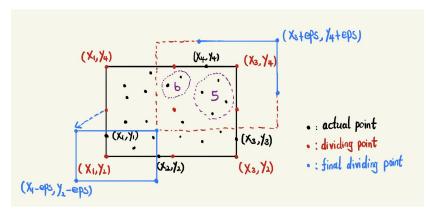
Parallel Matrix DBSCAN algorithm——Local Matrix DBSCAN

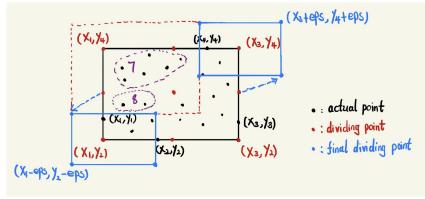


partition number: 4 eps (epsilon): 3 MinPts (minimum points): 2





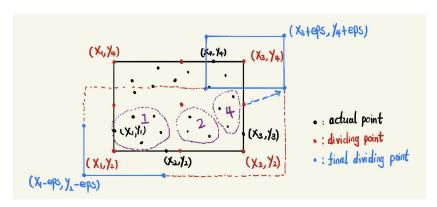


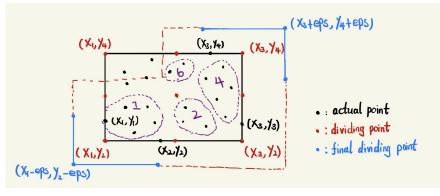


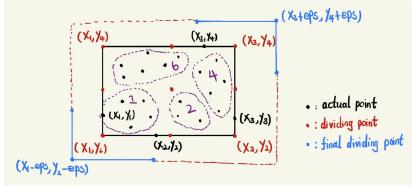
Parallel Matrix DBSCAN algorithm——Merge



partition number: 4 eps (epsilon): 3 MinPts (minimum points): 2







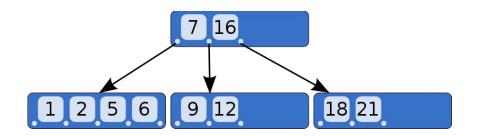
Using local labels to update global labels for all clusters

Optimization on partition——R tree



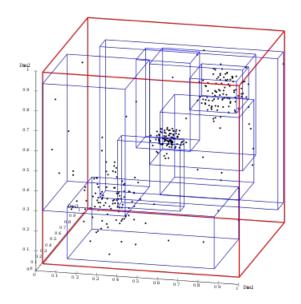
B-tree:

A self-balancing tree data structure



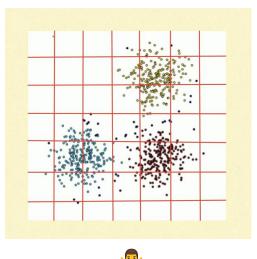
R-tree:

A tree data structures used for spatial access methods

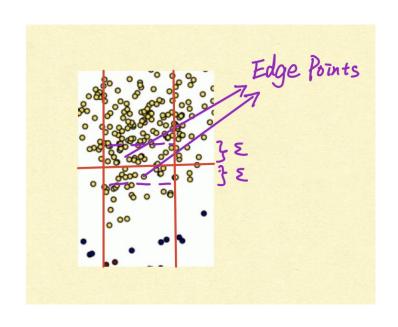


Optimization on partition——R tree (Boundary-based Strategy)













Reduce the amount of edge points



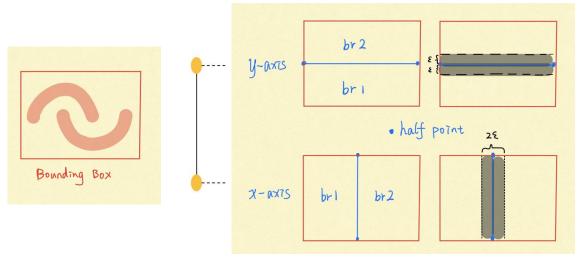
Make sure workload of each worker as balanced as possible

Optimization on partition——R tree (Boundary-based Strategy)

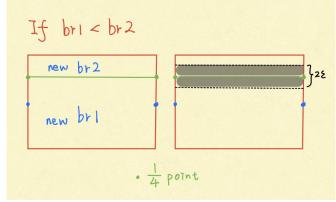




Set max points number in partition - maxpoints



As for y-axis:



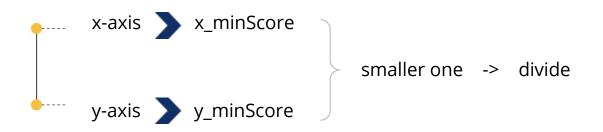
minimize Score = |# in br1 - # in br2| * (# in Black part)

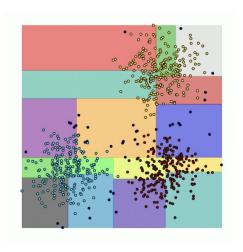
Optimization on partition——R tree (Boundary-based Strategy)





Max points number in partition - maxpoints





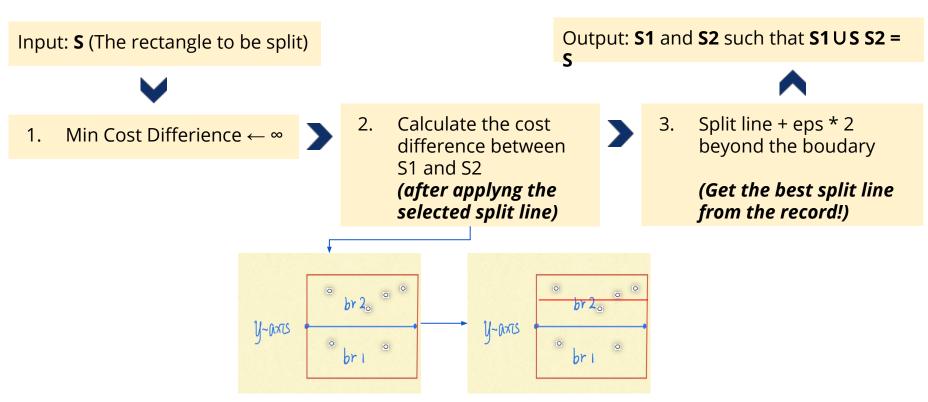
Put final br1 & br2 to queue (right hand)

Take out the leftmost mbr in the queue in turn to do the above calculation

Optimization on partition——R tree (Cost-based Strategy)

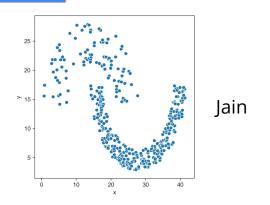


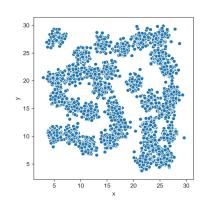
Goal: Make the workload (# of points) of partitioned region as balanced as possible.



Conclusion







D31

Partition Number	4	8	10	Partition Number	4	8	10
Naive	1179ms	797ms	808ms	Naive	26410ms	19415ms	18483ms
Parallel Matrix DBSCAN	422ms	385ms	380ms	Parallel Matrix DBSCAN	6730ms	5590ms	4564ms
Rtree - rbp	360ms	330ms	289ms	Rtree - rbp	7993ms	5722ms	5328ms
Rtree - cbp	308ms	284ms	271ms	Rtree - cbp	6977ms	5652ms	5164ms

Conclusion



- Performance
- Sclability
- Bottle neck (Master)
- Know your data

Future Work:

Now: using number of points to measure workload when doing partition

Future: using the real cost (maybe considering density) of DBSCAN to measure workload

More expreiments on large volume datasets with large partitions

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Rtree - cbp	308ms	284ms	271ms	Rtree - cbp	6977ms	5652ms	5264ms

Reference



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