

Tutorial 2 Solutions

Wednesday, 12 February 2025

1.

$$A = \{x \in \{0, 1\}^* \mid x = 0^+1\}$$

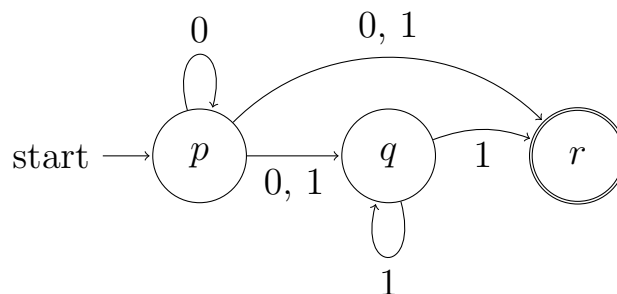
2. The epsilon-closure of a state is the set of states reachable from it using only epsilon transitions.

$$\text{Epsilon-closure}(p) = \{p, q, r\}, \text{Epsilon-closure}(q) = \{q, r\}, \text{Epsilon-closure}(r) = \{r\}$$

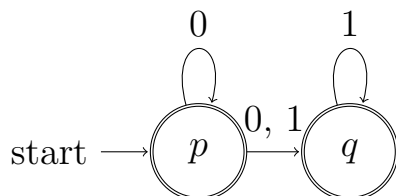
Transition Table for the NFA Without Epsilon Transitions We replace the epsilon transitions with direct transitions based on the epsilon-closure. The updated transitions are:

State	0	1
p	$\{p, q, r\}$	$\{q, r\}$
q	\emptyset	$\{q, r\}$
r	\emptyset	\emptyset

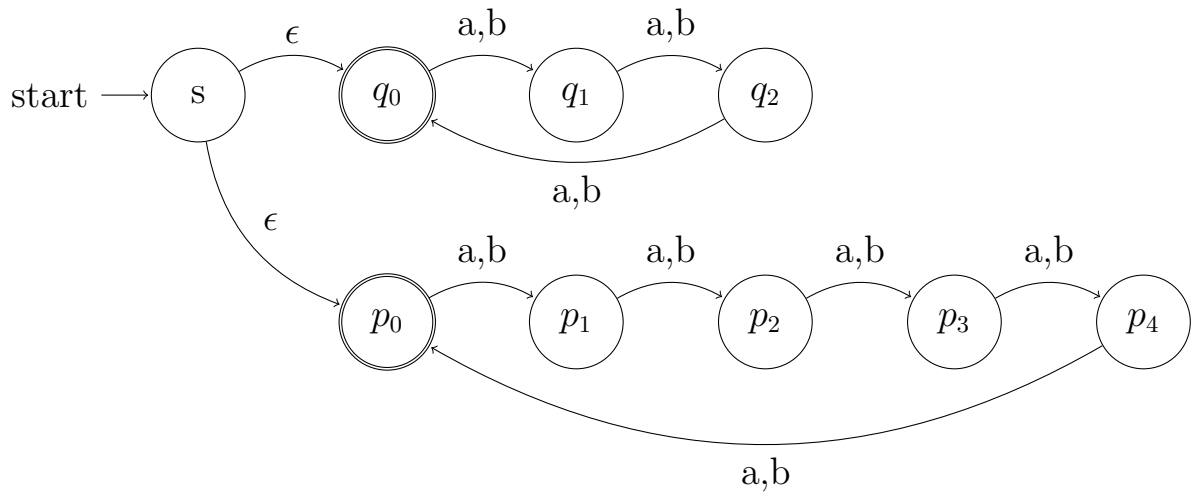
The equivalent NFA without epsilon transitions is constructed as follows: (This doesn't accept the empty string)



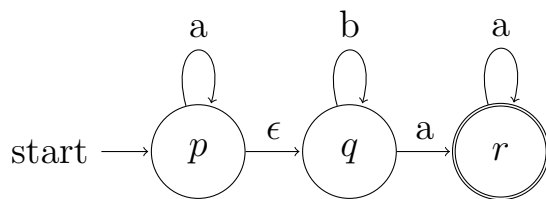
An equivalent NFA that also accepts the empty string:



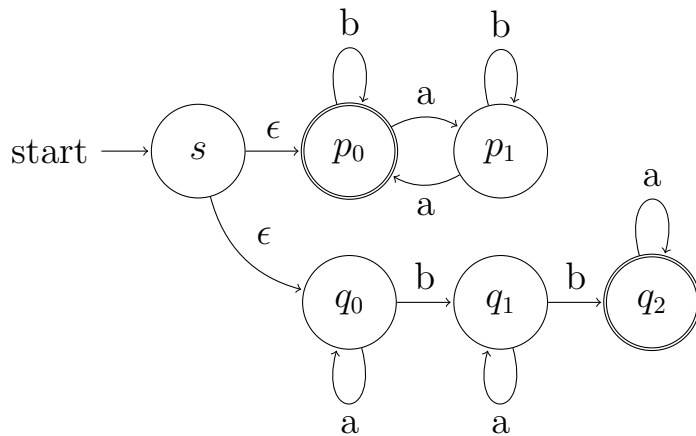
3. (a) $S = \{s \mid \#s \text{ is divisible by either 3 or 5}\}$



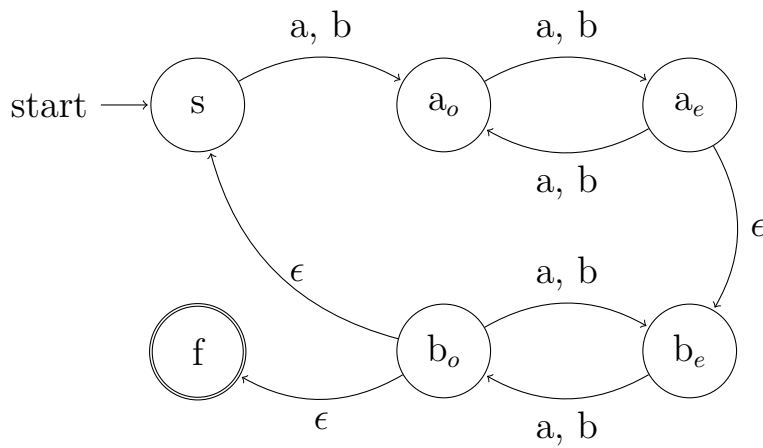
(b) Language $a^*b^*a^+$.



(c) The language comprising all strings that contain an even number of a 's or contain exactly two b 's.



4. L : Shuffle of languages A and B.



5.

$A = \{x \in \{0, 1\}^* \mid \text{the second symbol from the right is } 0\}$.

	0	1
\emptyset	\emptyset	\emptyset
$\rightarrow \{p\}$	$\{p, q\}$	$\{p\}$
$\{q\}$	$\{r\}$	$\{r\}$
$\{r\}F$	\emptyset	\emptyset
$\{p, q\}$	$\{p, q, r\}$	$\{p, r\}$
$\{p, r\}F$	$\{p, q\}$	$\{p\}$
$\{q, r\}F$	$\{r\}$	$\{r\}$
$\{p, q, r\}F$	$\{p, q, r\}$	$\{p, r\}$

Starting from state $\{p\}$, we observe that the states $\{q, r\}$, $\{q\}$, $\{r\}$, and \emptyset can never be reached. Therefore, these states can be ignored, and the transition table simplifies as follows:

	0	1
$\rightarrow \{p\}$	$\{p, q\}$	$\{p\}$
$\{p, q\}$	$\{p, q, r\}$	$\{p, r\}$
$\{p, r\}F$	$\{p, q\}$	$\{p\}$
$\{p, q, r\}F$	$\{p, q, r\}$	$\{p, r\}$

