## Data Structures and Algorithms

(資料結構與演算法)

Lecture 1: Algorithm

Hsuan-Tien Lin (林軒田)

htlin@csie.ntu.edu.tw

Department of Computer Science & Information Engineering

National Taiwan University (國立台灣大學資訊工程系)



# definition of algorithm

# Name Origin of Algorithm

Muhammad ibn Mūsā al-Kwārizmī on a Soviet Union stamp

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### algorithm

- named after al-Kwārizmī (780–850),
   Persian mathematician and father of algebra
- algebra: rules to calculate with symbols
- algorithm: instructions to compute with variables

algorithm: recipe-like instructions for computing

# Recipe for Cooking Dish



a recipe for hamburger on Wikibooks

figure by Gentgeen,

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#### recipe

Wikipedia: a set of instructions that describes how to prepare or make something, especially a dish of prepared food

recipe: instructions to complete a (cooking) task

# Sheet Music for Playing Instrument



first page of the manuscript of Bach's lute suite in G minor

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commons.wikimedia.org/wiki/File:Bachlut1.png

#### sheet music

Wikipedia: handwritten or printed form of musical notation ... to indicate the pitches, rhythms or chords of a song

sheet music: instructions to play instrument (well)

# Kifu for Playing Go



a Japanese kifu

figure by Velobici,

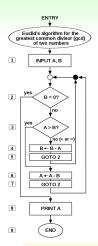
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#### kifu

go game record of steps that describe how the game had been played

kifu: instructions to mimic/learn to play go (professionally)

# Algorithm for Computing



flowchart of Euclid's algorithm for calculating the greatest common divisor (g.c.d.) of two numbers

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#### algorithm

Wikipedia: algorithm is a finite sequence of well-defined, computer-implementable instructions, typically to solve a class of problems or to perform a computation

algorithm ~ computing recipe: (computable) instructions to solve a computing task efficiently/correctly

# pseudo code of algorithm

#### Pseudo Code for GETMININDEX

#### C Version

```
/* return index to min. element
    in arr[0] ... arr[len-1] */
int getMinIndex
        (int arr[], int len) {
    int i;
    int m=0;
    for(i=0;i<len;i++) {
        if (arr[m] > arr[i]) {
            m = i;
        }
    }
    return m;
}
```

#### Pseudo Code Version

```
GET-MIN-INDEX(A)

1 m = 1

2 for i = 2 to A. length

3  // update if i-th element smaller

4  if A[m] > A[i]

5  m = i

6 return m
```

pseudo code: spoken language of programming

#### Bad Pseudo Code: Too Detailed

#### **Unnecessarily Detailed**

```
Get-Min-Index(A)
     m=1
    for i = 2 to A. length
 3
         // update if i-th element smaller
         Am = A[m]
 5
         Ai = A[i]
 6
         if Am > Ai
             m = i
 8
         else
 9
              m = m
10
     return m
```

#### Concise

```
GET-MIN-INDEX(A)

1 m = 1

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

goal of pseudo code: communicate efficiently

# Bad Pseudo Code: Too Mysterious

#### Unnecessarily Mysterious

#### Clear

```
GET-MIN-INDEX(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3  // update if i-th element smaller

4  if A[m] > A[i]

5  m = i
```

goal of pseudo code: communicate correctly

## Bad Pseudo Code: Too Abstract

#### Unnecessarily Abstract

#### Get-Min-Index(A)

- m = 1 // store current min. index
- run a loop through A that updates *m* in every iteration
- return m

#### Concrete

```
Get-Min-Index(A)
```

```
m = 1 // store current min. index
   for i = 2 to A. length
3
        // update if i-th element smaller
        if A[m] > A[i]
             m-i
   return m
```

goal of pseudo code: communicate effectively

#### From Get-Min-Index to Selection-Sort

```
GET-MIN-INDEX(A, \ell, r)

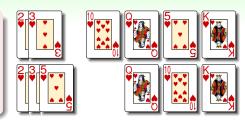
1 m = \ell // store current min. index

2 for i = \ell + 1 to r

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i
```



#### Good Pseudo Code

- modularize, just like coding
- depends on speaker/listener
- usually no formal definition

#### SELECTION-SORT(A)

- 1 for i = 1 to A. length
  - m = Get-Min-Index(A, i, A. length))
- 3 SWAP(A[i], A[m])
- 4 return A // which has been sorted in place

follow any textbook if you really need a definition

# criteria of algorithm

## Criteria of Recipe



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#### Cocktail Recipe: Screwdriver

inputs: 5 cl vodka, 10 cl orange juice

- 1 mix inputs in a highball glass with ice
- 2 garnish with orange slice and serve output: a glass of delicious cocktail

- input: ingredients
- definiteness: clear instructions
- effectiveness:
   feasible instructions
- finiteness: completable instructions
- output: delicious drink

algorithm ~ recipe: same five criteria for algorithm

(Knuth, The Art of Computer Programming)

# Input of Algorithm

... quantities which are given to it initially before the algorithm begins.

These inputs are taken from specified sets of objects. (Knuth, TAOCP)

```
GET-MIN-INDEX(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

one algorithm, many uses (on different legal inputs)

## Definiteness of Algorithm

Each step of an algorithm must be precisely defined; the actions to be carried out must be rigorously & unambiguously specified. (Knuth, TAOCP)

#### Clear

```
GET-MIN-INDEX(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

#### **Ambiguous**

definiteness: clarity of algorithm

# Effectiveness of Algorithm

... all of the operations to be performed in the algorithm must be sufficiently basic that they can in principle be done exactly and in a finite length of time by a man using paper and pencil. (Knuth, TAOCP)

#### Effective

```
GET-MIN-INDEX(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3  // update if i-th element smaller

4  if A[m] > A[i]

5  m = i
```

#### Ineffective

```
GET-SOFT-MIN(A)

1 s = 0 // sum of exponentiated values

2 for i = 1 to A. length

3 s = s + exp(-A[i] \cdot 1126)

4

5

6 return - log(s)/1126
```

floating point errors may make some steps ineffective on some computers

# Finiteness of Algorithm

An algorithm must always terminate after a finite number of steps . . . a very finite number, a reasonable number. (Knuth, TAOCP)

```
GET-MIN-INDEX(A)

1 m=1 // store current min. index

2 for i=2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m=i

6 return m
```

finiteness (& efficiency): often requiring analysis for sophisticated algorithms (to be taught later)

# Output of Algorithm

... quantities which have a specified relation to the inputs (Knuth, TAOCP)

```
GET-MIN-INDEX(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3  // update if i-th element smaller

4  if A[m] > A[i]

5  m = i

6 return m
```

output (correctness): needs proving with respect to requirements

# correctness proof of algorithm

#### Claim



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```
GET-MIN-INDEX(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

#### Correctness of GET-MIN-INDEX

Upon exiting GET-MIN-INDEX(A),

$$A[m] = \min_{1 \le j \le n} A[j]$$

with n = A. length

claim: mathematical statement that declares correctness

#### Invariant

invariants when constructing fractals figures by Johannes Rössel,

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#### Correctness of GET-MIN-INDEX

Upon exiting GET-MIN-INDEX(A),

$$A[m] = \min_{1 \le j \le n} A[j]$$

with n = A. length



#### Get-Min-Index(A)

return m

```
1 m = 1 // store current min. index

2 for i = 2 to A. length

3  // update if i-th element smaller

4  if A[m] > A[i]

5  m = i
```

#### Invariant within GET-MIN-INDEX

Upon finishing the loop with i = k, denote m by  $m_k$ ,

$$A[m_k] \le A[j] \text{ for } j = 1, 2, ..., k$$

(loop) invariant: property that algorithm maintains

## Proof of Loop Invariant

#### Mathematical Induction

#### Base

when i = 2, invariant true because . . .

assume invariant true for i = t - 1; when i = t,

- $m_t = t$  if  $A[m_{t-1}] > A[t]$ •  $A[m_t] = A[t] \le A[t]$ •  $A[m_t] < A[m_{t-1}] \le A[j]$ for other j
- $m_t = m_{t-1}$  if  $A[m_{t-1}] \le A[t]$ •  $A[m_t] = A[m_{t-1}] \le A[t]$ •  $A[m_t] = A[m_{t-1}] \le A[j]$ for other j

—by mathematical induction, invariant true for i = 2, 3, ..., k

#### Get-Min-Index(A)

```
1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

#### Correctness of GET-MIN-INDEX



#### Invariant within GET-MIN-INDEX

Upon finishing the loop with i = k, denote m by  $m_k$ ,

$$A[m_k] < A[j] \text{ for } j = 1, 2, ..., k$$

proof of (loop) invariants ⇒ correctness claim of algorithm

 $\Rightarrow$ 

## Summary

### Lecture 1: Algorithm

- definition of algorithm instructions to complete a task by computer
- pseudo code of algorithm

#### communicate alg. efficiently/correctly/effectively

- criteria of algorithm input, definite, effective, finite, output
- correctness proof of algorithm
   from (loop) invariants to claims