Data Structures and Algorithms
(資料結構與演算法)

Lecture 9: Binary Tree

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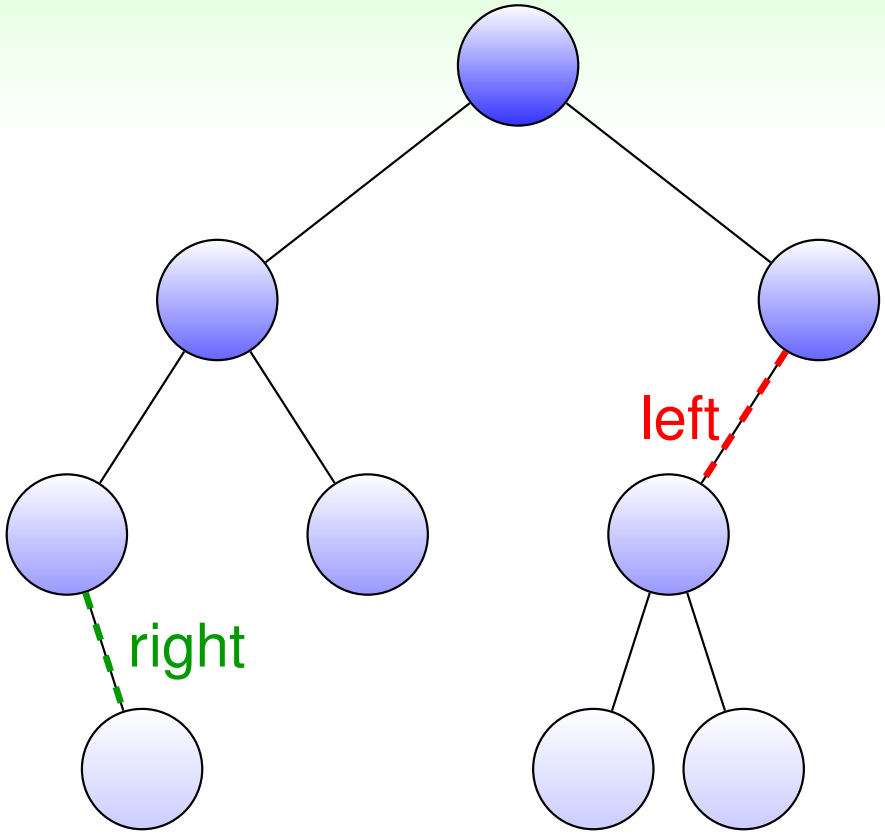
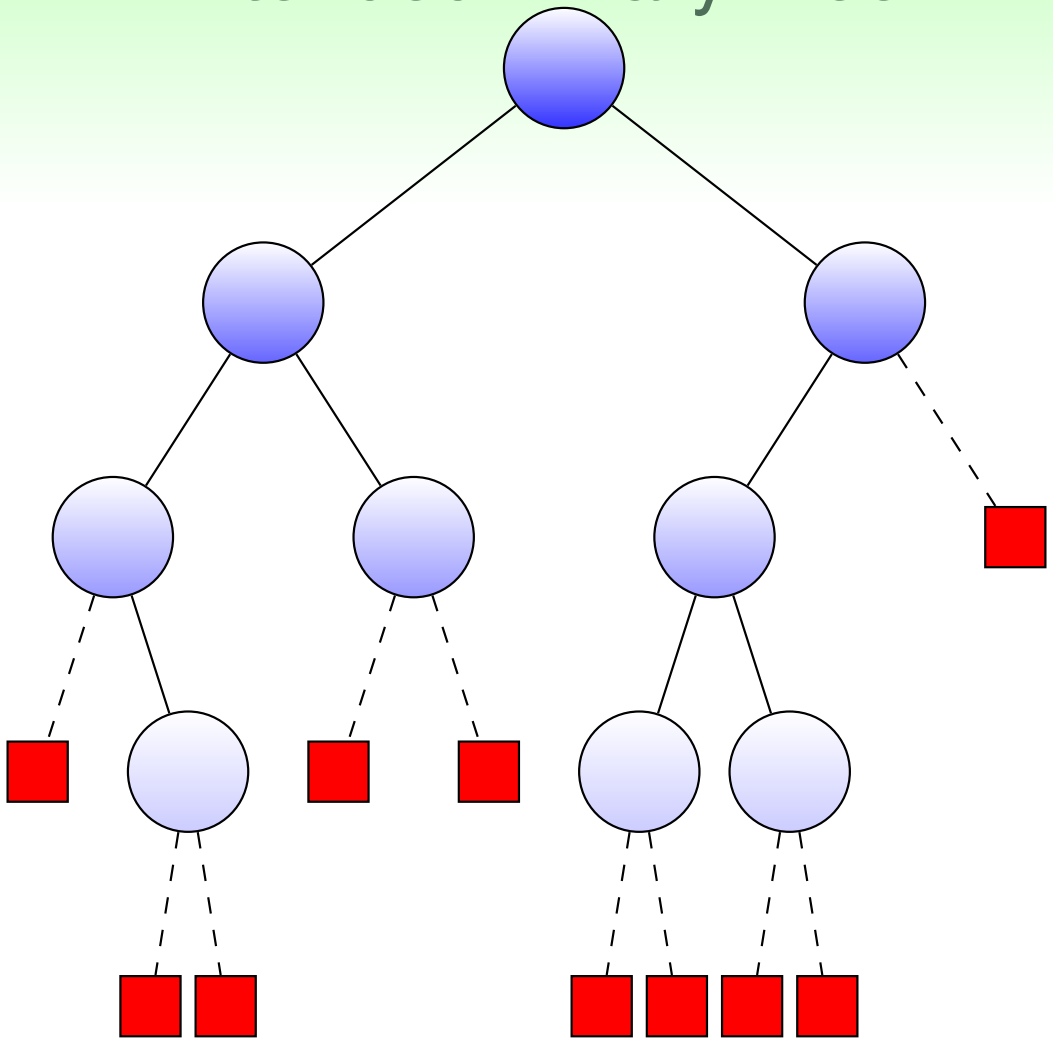
terminologies

Binary Tree



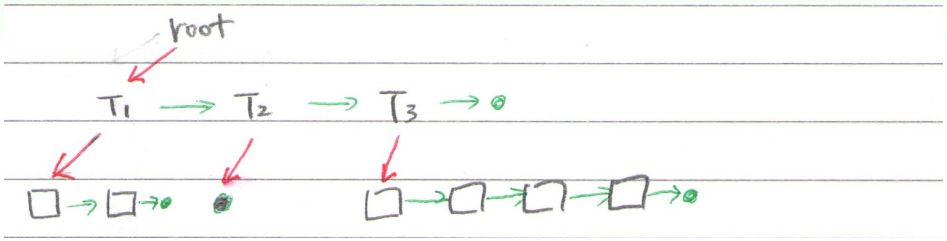
binary tree: rooted tree with node degree ≤ 2
& left/right difference

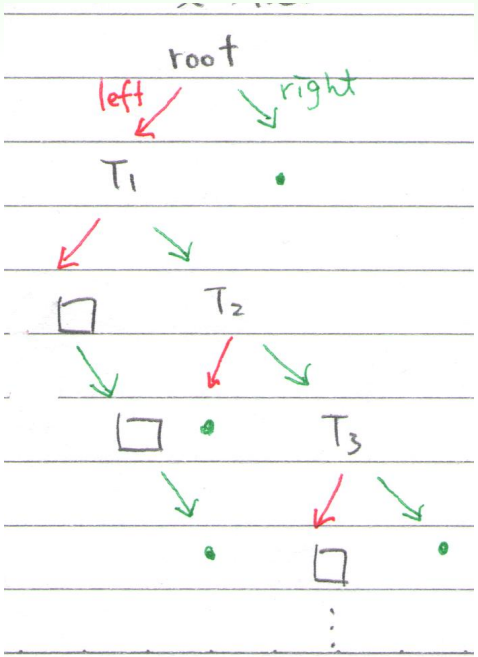
Extended Binary Tree



extended binary tree: added empty external nodes
so all degree = 2 & all original nodes internal

General Tree with Linked Lists ≡ Binary Tree



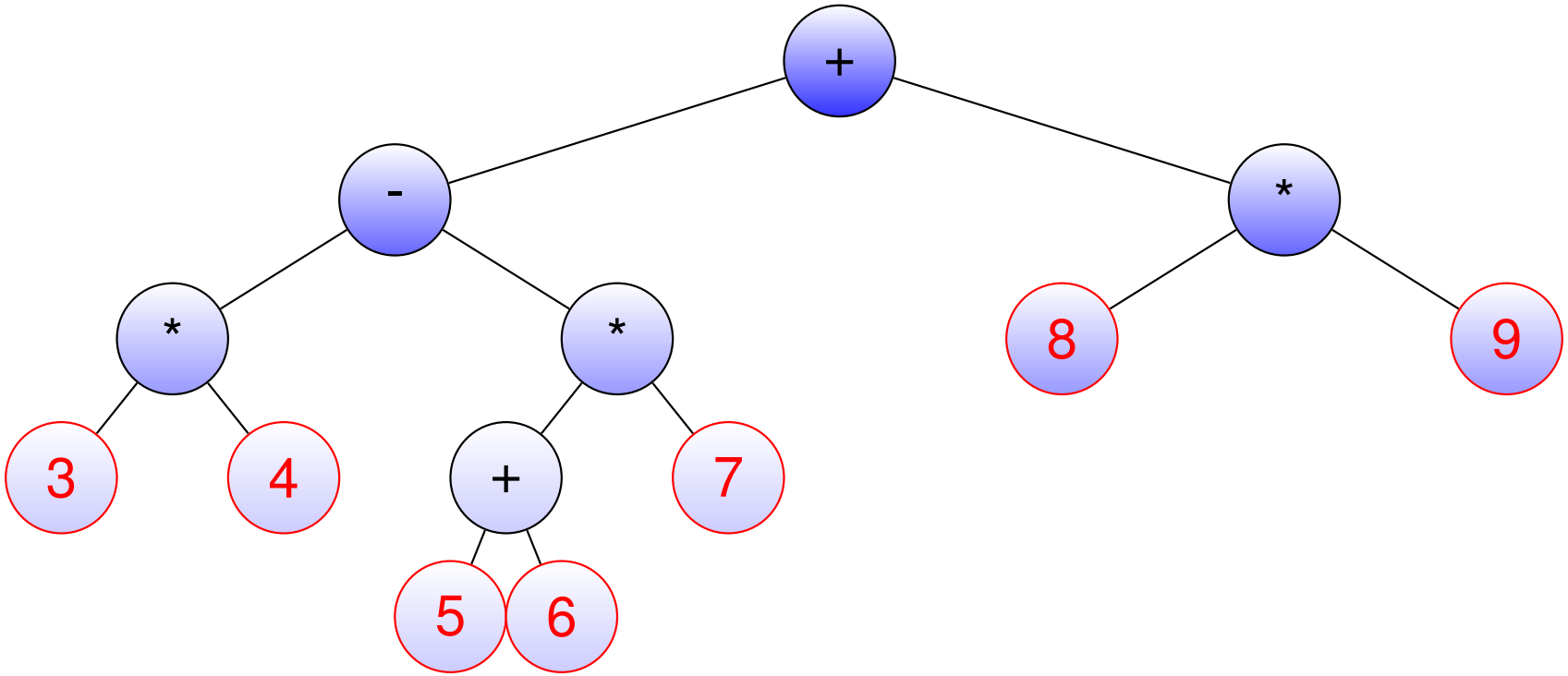


left-child right-sibling implementation of general tree

binary tree traversals

Expression Tree

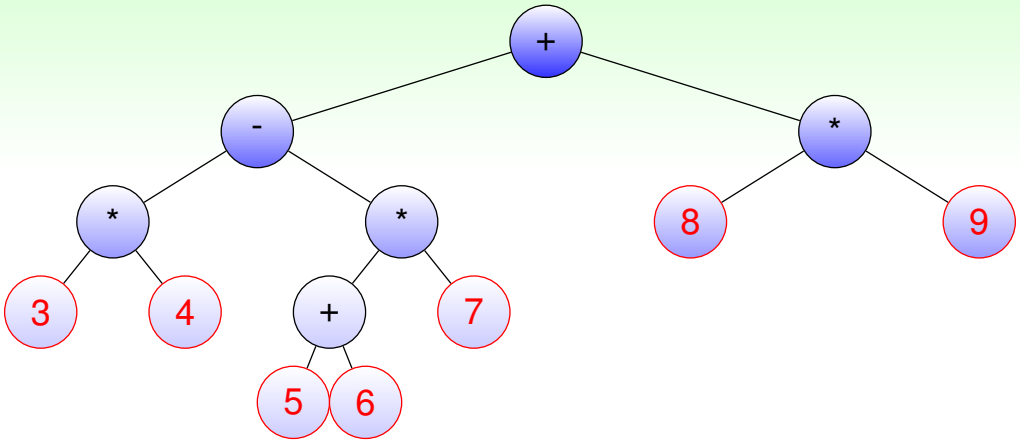
$$3 * 4 - (5 + 6) * 7 + 8 * 9$$



expression tree: (binary) expression
represented by a (binary) tree

Outputting Infix Expression from Expression Tree

want: $3*4-(5+6)*7+8*9$



OUTPUT-INFIX(T)

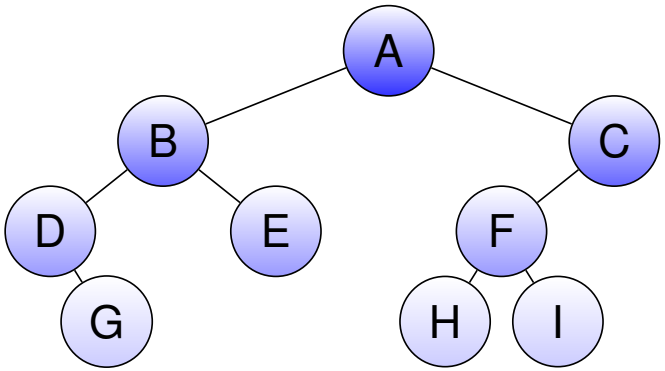
```
 1  if IS-LEAF(T)
 2        print T.data
 3  else
 4        if [T.data higher precedence than T.left.data] print '('
 5        OUTPUT-INFIX(T.left)
 6        if [T.data higher precedence than T.left.data] print ')'
 7        print T.data
 8        if [T.data higher precedence than T.right.data] print '('
 9        OUTPUT-INFIX(T.right)
10        if [T.data higher precedence than T.right.data] print ')'
```

OUTPUT-INFIX: variant of INORDER-TRAVERSAL

From OUTPUT-INFIX to INORDER-TRAVERSAL

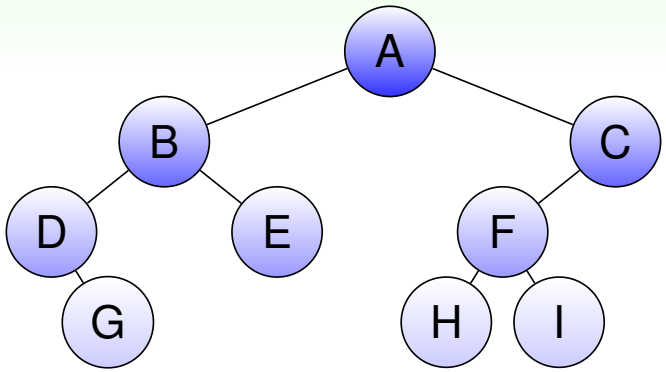
```
OUTPUT-INFIX(T)
 1  if IS-LEAF(T)
 2        print T.data
 3  else
 4        if ...
 5        OUTPUT-INFIX(T.left)
 6        if ...
 7        print T.data
 8        if ...
 9        OUTPUT-INFIX(T.right)
10        if ...
```

```
INORDER-TRAVERSAL(T)
 1  if T is NIL
 2
 3  else
 4
 5        INORDER-TRAVERSAL(T.left)
 6
 7        action with T.data
 8
 9        INORDER-TRAVERSAL(T.right)
10
```



INORDER-TRAVERSAL visit sequence: DGBEAHFIC

POSTORDER-TRAVERSAL & PREORDER-TRAVERSAL



GDEBHIFCA

```
POSTORDER-TRAVERSAL(T)
1  if T is not NIL
2        POSTORDER-TRAVERSAL(T.left)
3        POSTORDER-TRAVERSAL(T.right)
4        action with T.data
   // appl: evaluate expression tree
```

DGBEAHFIC

```
INORDER-TRAVERSAL(T)
1  if T is not NIL
2        INORDER-TRAVERSAL(T.left)
3        action with T.data
4        INORDER-TRAVERSAL(T.right)
   // appl.: output infix expression
```

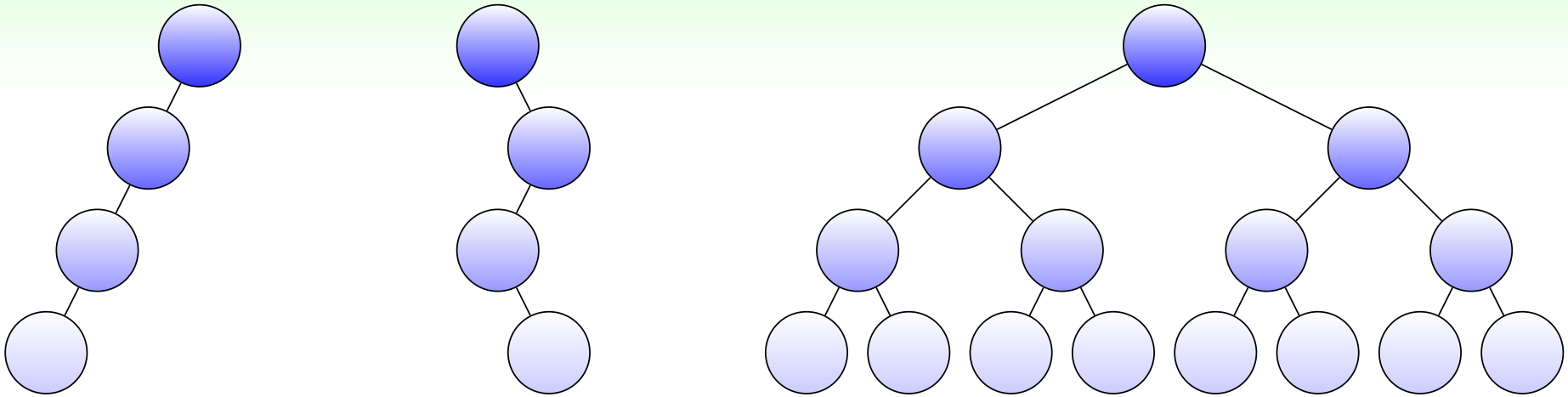
ABDGECFHI

```
PREORDER-TRAVERSAL(T)
1  if T is not NIL
2        action with T.data
3        PREORDER-TRAVERSAL(T.left)
4        PREORDER-TRAVERSAL(T.right)
   // appl.: compare two trees
```

traversal: template for designing tree algorithms

full and complete binary trees

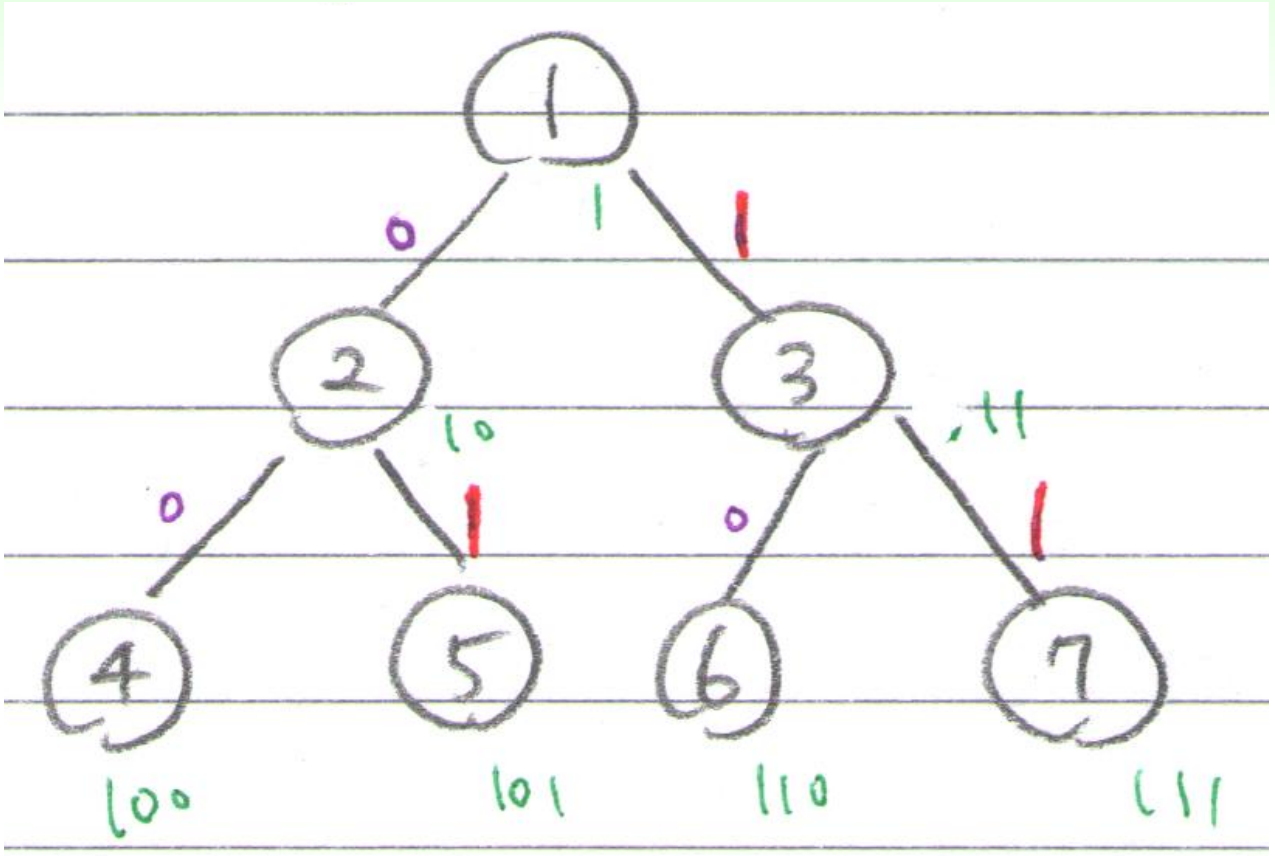
Nodes in Binary Trees



height	min	max
1	1	1
2	2	3
3		
...		
h	h (skewed)	$2^h - 1$ full

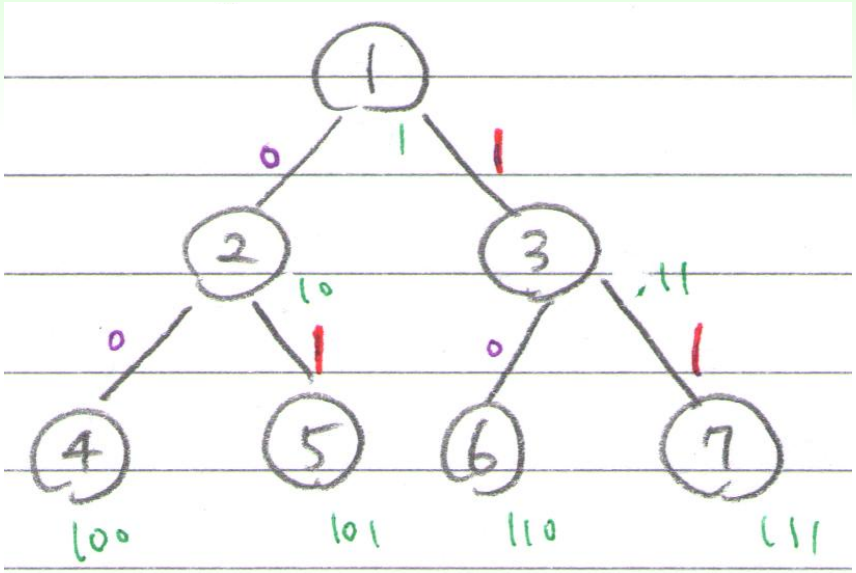
$$h \leq n \leq 2^h - 1 \Leftrightarrow \lg(n+1) \leq h \leq n$$

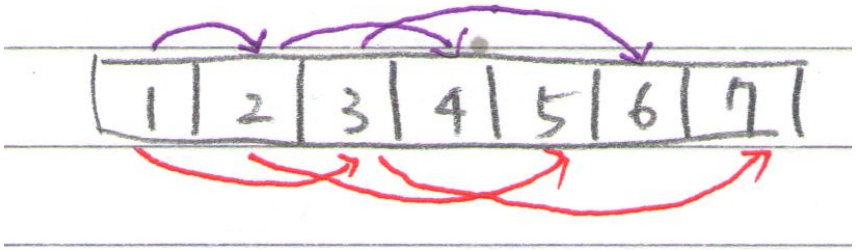
Node Index in Full Binary Tree



node index $= (1\,path\ code)_2$

Representing/Packing Full Binary Tree in Array

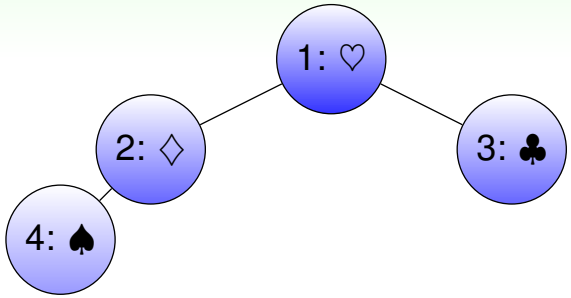




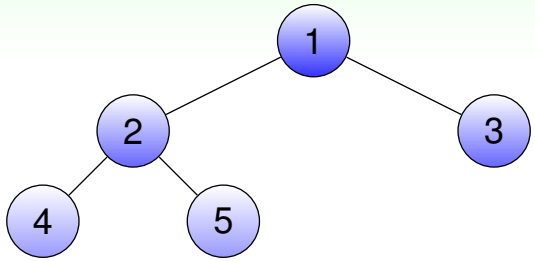
- implicit links: no need for explicit pointers
- can similarly pack any binary tree if NIL can represent NO-DATA (with **space wasting** in data field)

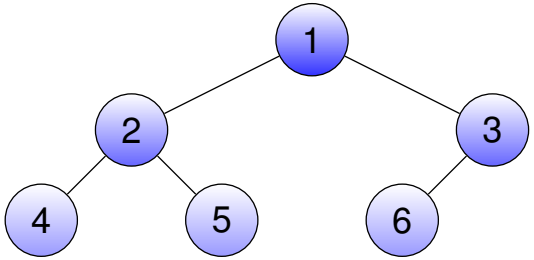
complete binary tree: full binary tree with first n nodes (**no waste** with array representation)

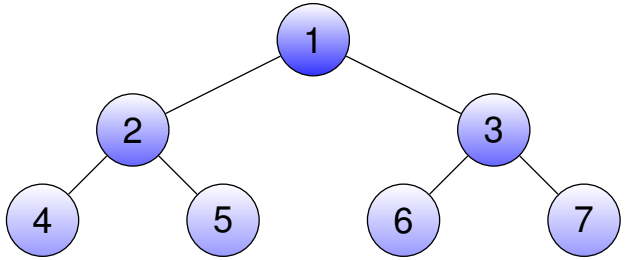
Complete Binary Trees



1	2	3	4
♡	♢	♣	♠







will use this **tree-in-array** property next time

Summary

Lecture 9: Binary Tree

- terminologies

  binary tree $\equiv$ tree with left/right sub-trees
- binary tree traversals

  templates for (recursive) tree algorithm design
- full and complete binary trees

can be implemented in array without wasting space