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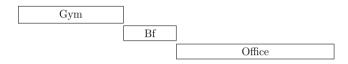
Biological Vision and Applications Module 07-05: Qualitative spatial and temporal relations

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## Allen's Interval Algebra

- Consider three 1D interval events: A, B, C
  - 1. A  $r_{AB}$  B
  - 2. B *r<sub>BC</sub>* C
- Can we infer the relation between A and C?
  - ightharpoonup A  $r_{AC}$  C
  - ► Given  $r_{AB}$ ,  $r_{BC}$ , can we find  $r_{AC}$ ?

### An Intuitive Introduction



- Consider the statements:
  - 1. I went to gym just before having my breakfast: Gym m Bf
  - 2. I went to office immediately after the breakfast: Bf m Office
- We can conclude
  - ► Temporal relation between Gym and Office: Gym *b* Office
- Transition rule:  $(m, m) \rightarrow b$
- An example of temporal sequencing

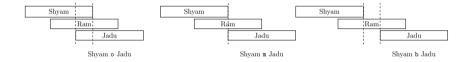
# An Intuitive Introduction (contd)



- Consider the statements:
  - 1. I attended office some time during yesterday: Office d Yday
  - 2. I ate my lunch while at office: Lunch d Office
- We can conclude
  - Temporal relation between Lunch and Yesterday: Lunch d Yday
- Transition rule:  $(d, d) \rightarrow d$
- An example of hierarchical decomposition

# An Intuitive Introduction (contd)

- 1. Ram came in the room while Shyam was there and continued after Shyam left:
  - Shyam o Ram
- 2. Jadu came into the room when Ram was there and continued after Ram left:
  - ▶ Ram *o* Jadu



- The temporal relation between Shyam and Jadu cannot be uniquely resolved
  - ► Shyam b, m, o Jadu
- Transition rule:  $(o, o) \rightarrow \{b, m, o\}$

## Allen's temporal algebra

- Given that A r<sub>AB</sub> B and B r<sub>BC</sub> C
  - 1. where  $r_{XY}$  is one of the Allen's relation
- Temporal constraint between A and C: A  $R_{AC}$  C
  - $\triangleright$  where  $R_{AC}$  is a subset of Allen's relation
- Mapping  $r \times r \xrightarrow{T} R$  is defined over a transitivity table
  - $ightharpoonup R \leftarrow T(r_1, r_2)$
  - ▶ 13 × 13 entries in transitivity table

Allen's transitivity table

### Allen's Interval Algebra

#### Generalizing ...

- $R_{ij}$ ,  $R_{jk}$ ,  $R_{ik}$ : Temporal constraints between event-pairs
- $(E_i, E_j)$ ,  $(E_j, E_k)$  and  $(E_i, E_k)$ 
  - In general, each is a subset of Allen's relations
- The algorithm for computing  $Constraint(R_{ij}, R_{jk}) \neq R_{ik}$ :

### Algorithm 18: Computing relational constraint

```
procedure Constraint(R_{ij}, R_{jk})
C = \emptyset;
for each p \in R_{ij} do
\begin{array}{c|c} & \text{for each } p \in R_{ij} \text{ do} \\ & \hline C \leftarrow C \cup T(p,q); \\ & \text{end} \\ & \text{end} \\ & \text{return } C; \\ \end{array}
```

# Uncertainty with the endpoints



- Where do the shore ends and the sea starts
- A man is walking He falls
- When does he start falling and when does he end falling?

Measurement error

 $\mathbf{B}$ 

A before B?
A meets B?
A overlaps B?

## Approximate qualitative relations

#### Conceptual neighborhood

- Relations organized in 2D defines conceptual neighbors
- Ambiguity in boundary / Measurement error may lead to a relation to be misclassified in it's conceptual neighborhood
- A set of relations in a conceptual neighborhood defines an "approximate relation"
- For fuzzy representation, see book

	$e_A < e_B$			$e_A = e_B$		$e_A > e_B$		
$e_A = s_B$ $e_A < s_B$	ь							
$e_A = s_B$		m						$s_A < s_B$
$e_A > s_B$			0	fi	di			
			s	eq	si			$s_A = s_B$
			d	f	oi			
						mi		$s_A > s_B$
							bi	
	$s_A < e_B$					$s_A = e_B$	$s_A > e_B$	

Semantics of Allen's relations



No quiz for module 07-05

End of Module 07-05