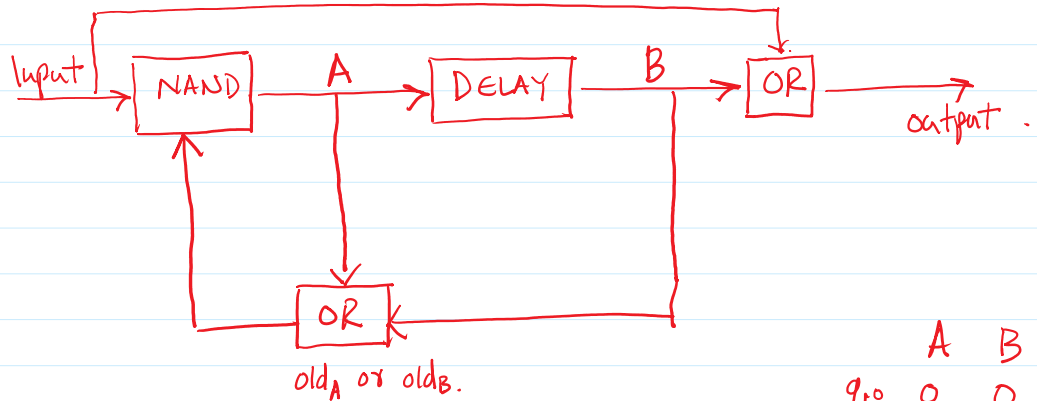


lecture 17:-

Finite Automata In Sequential Circuits. (Mealy Machine).

\Rightarrow - \Rightarrow -
- \Rightarrow -

A	B	A.B	$\neg(A.B)$
0	0	0	1
0	1	0	1
1	0	0	1
1	1	1	0

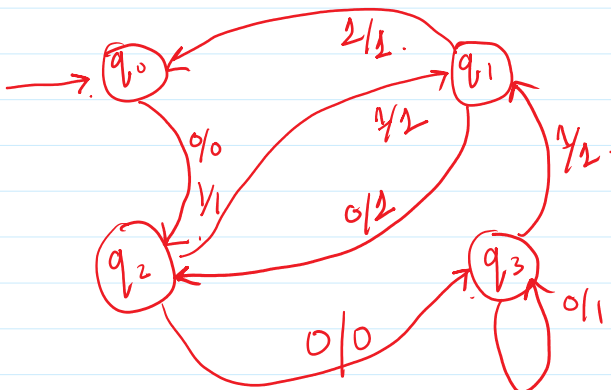


$NewB = OldA$
 $NewA = Input \text{ NAND } (oldA \text{ or } oldB)$
 $Output = OldB \text{ or } Input$

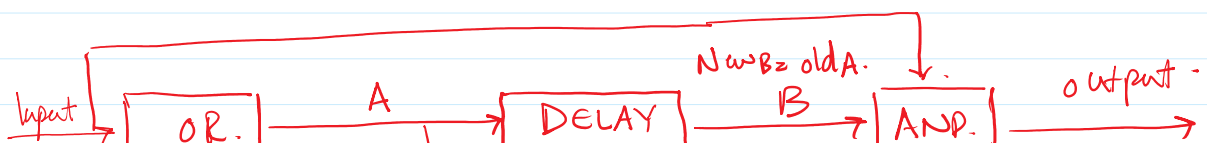
	A	B
q_0	0	0
q_1	0	1
q_2	1	0
q_3	1	1

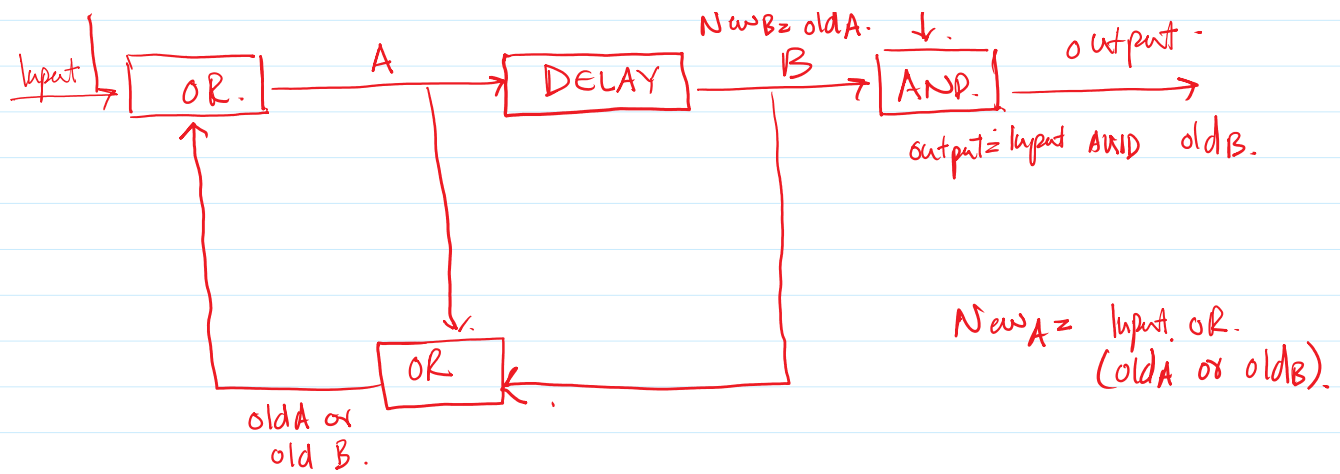
	old oldA	State oldB	Input 0 New State NewA	Input 0 New State NewB	output
$\rightarrow q_0$	0	0	(1	0)	0
q_1	0	1	(1	0)	1
$\rightarrow q_2$	1	0	(1	1)	0
q_3	1	1	(1	1)	1

	Input 1 New State NewA	Input 1 New State NewB	output
$(1$	0)	q_2	1
(0	0)	q_0	1
(0	1)	q_1	1
(0	1)	q_1	1



A	B	A and B	A NAND B
0	0	0	1
0	1	0	1
1	0	0	1
1	1	1	0





New A = input OR (old A or old B).

New B = old A

output = input AND old B.

A B.

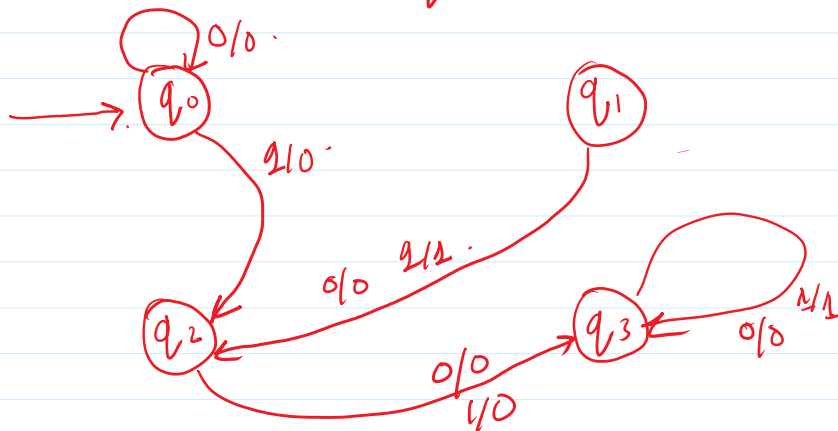
q₀ 0 0

q₁ 0 1

q₂ 1 0

q₃ 1 1.

		INPUT 0		Output	INPUT 1.		Output
old	State	New	State		New	State	
oldA	oldB	NewA	NewB		NewA	NewB	
q ₀ 0	0	(0 0)	q ₀	0	(1 0)	q ₂	0
q ₁ 0	1	(1 0)	q ₂	0	(1 0)	q ₂	1.
q ₂ 1	0	(1 1)	q ₃	0	(1 1)	q ₃	0
q ₃ 1	1	(1 1)	q ₃	1	(1 1)	q ₃	1.



REGULAR LANGUAGES:-

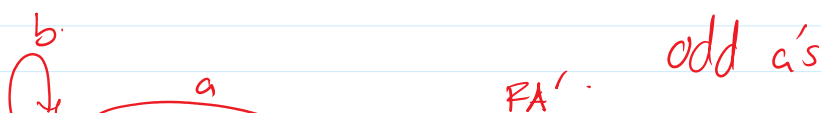
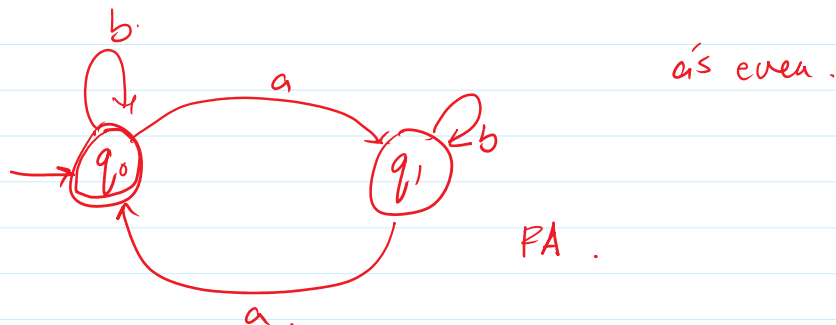
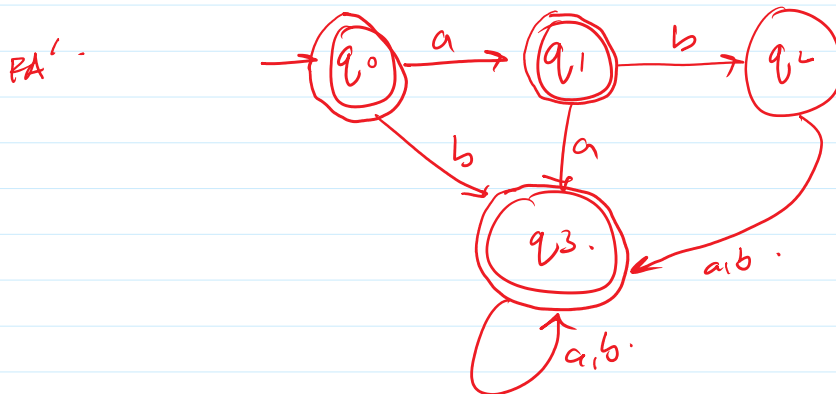
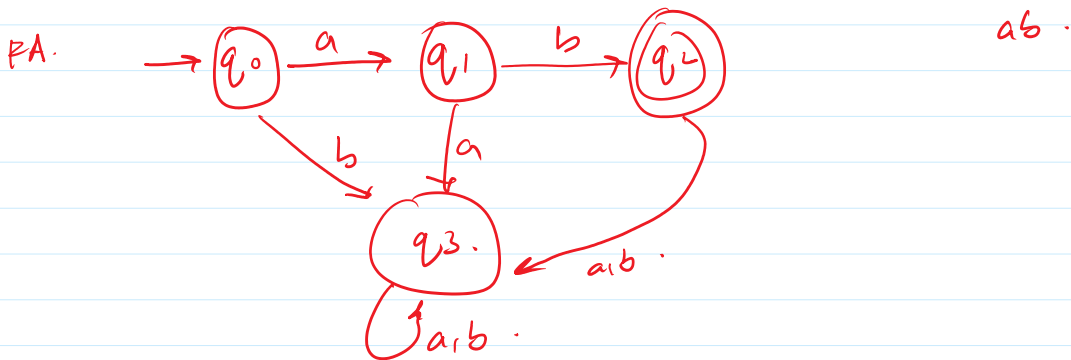
A language is called Regular if \exists a Regular Expression defining the language.

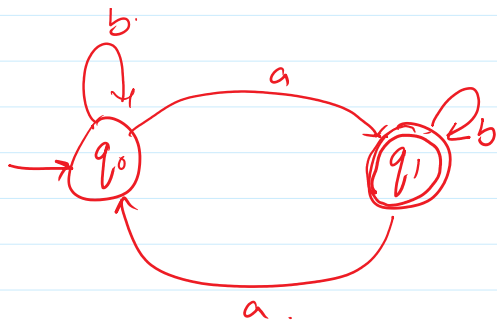
palindrome.

Properties:- 1- closure.
2- Complement & Intersection.

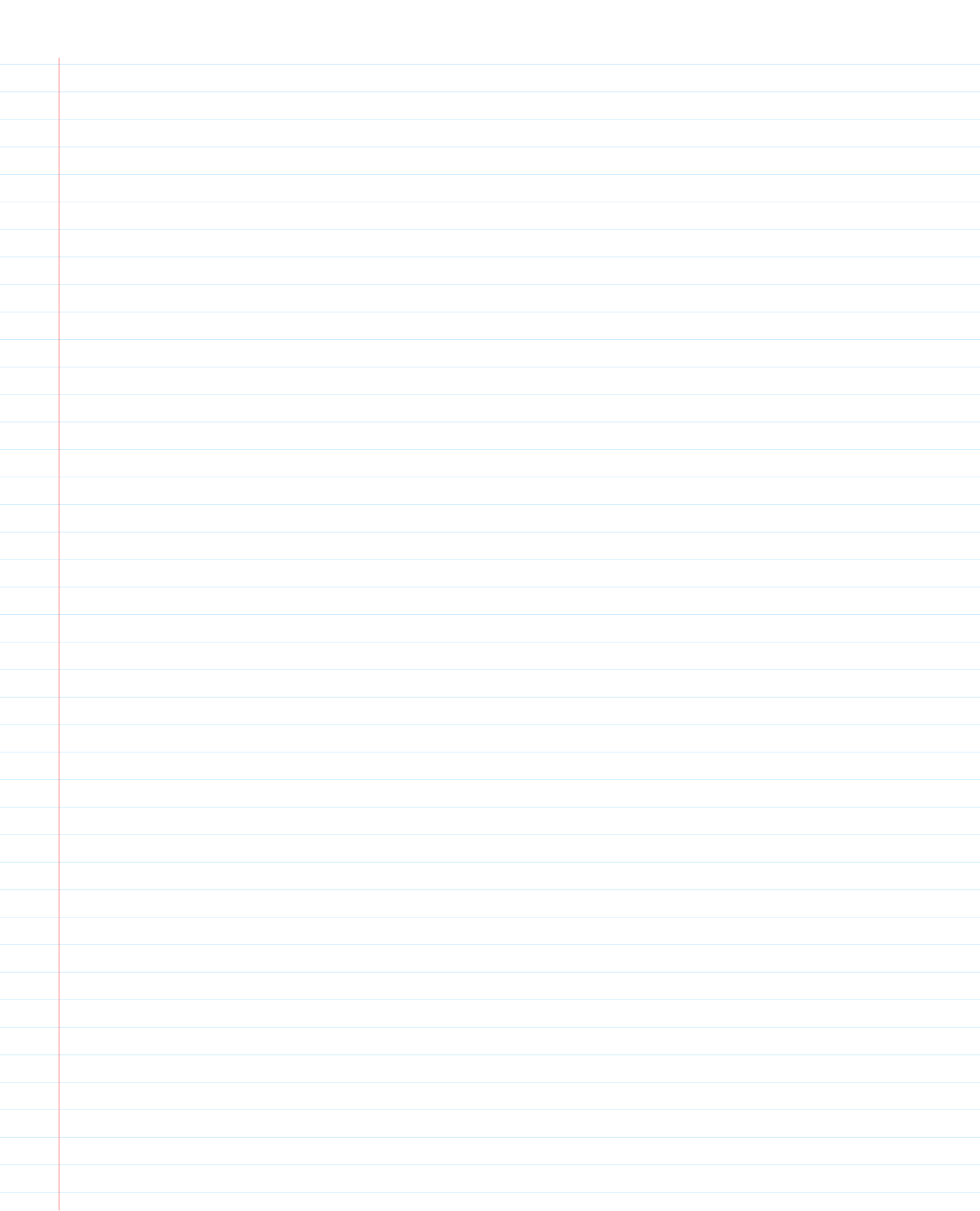
PA, let PA' be the Complement.
 $PA' =$ All the words that are not accepted by PA.

Procedure for determining a Complement.
 H.W. { 1- change all final states to Non final.
 2- change all non final to final.





RA' . odd a's



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