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Section: BSCS-6B

Assignment: Amdahl's law for stand alone system

Problem 1 → A system has a processor which takes 40% of the total execution time & can be improved by 50%. What is the maximum theoretical speedup we can achieve on this system?

Solution: According to Amdahl's law the maximum theoretical speedup is given by:

$$\text{Max speedup} = 1 / [1 - (1 + (1 - P) / S)]$$

where P is the fraction of the program that cannot be parallelized & S is the speedup factor.

$$P = 0.4 (40\%), \quad S = 1 + 0.5 = 1.5$$

$$\text{max speedup} = 1 / [1 - (0.4 + (1 - 0.4) / 1.5)]$$

$$= 1 / [1 - 0.8] = 5$$

∴ Therefore the max theoretical speedup can be achieved is $\boxed{5}$

Problem 2 A computer program has a serial portion that takes up 60% of the total execution time. If we can parallelize the remaining 40% of the program & achieve a speedup factor of 2. What is the maximum theoretical speedup we can achieve on this system?

Solution

2

According to Amdahl's Law

$$\text{Max Speedup} = \frac{1}{(1-P) + \frac{P}{S}}$$

where P is the fraction of program that can't be parallelized & S is speedup factor of parallelized part

$$P = 0.6 \text{ (60\%)} \quad S = 2$$

$$\begin{aligned} \text{Max speedup} &= \frac{1}{(1-0.6) + \frac{0.6}{2}} \\ &= \frac{1}{(1-0.6)} \\ &= 5 \end{aligned}$$

Therefore the maximum theoretical speedup we can achieve on this system is $\boxed{5}$

Amdahl's law for Standalone system:-

It is a formula that calculate the maximum speedup that can be achieved when performing ~~the~~ task on system with multiprocessor & cores for Standalone systems. It states that the max speedup is limited by the serial portion of the program.

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Amdahl's law for ³ distributed system

Explanation: In a distributed system, Amdahl's law takes into account both communication overhead & the serial portion of the program. The formula is similar to the standalone but with an additional term that accounts for the time spent communicating b/w the processor.

Problem #1 A distributed system has a task that requires two sub-tasks to be completed. Sub-task A cannot be parallelized & takes 60% of the total time, while sub-task B can be parallelized & takes 40% of the total time. If we add 4 more nodes to the system & sub-task B can be perfectly parallelized, what is the max theoretical speedup we can achieve on this system.

Solution

In a distributed system, Amdahl's law can be applied to each sub-task individually. & then combined to calculate the overall max speedup.

For sub-task A, since it can't be parallelized,

$$\text{max speedup} = \frac{1}{1 - 0.6} = 2.5$$

For sub-task B with 4 nodes, the speedup factor is equal to the number of nodes added

$$S = 4 + 1 = 5$$

$$\boxed{S = 5}$$

Therefore max speedup we can achieve is

$$= \frac{1}{\left(\frac{0.4}{5} + 0.6\right)} = \boxed{2.5}$$

$$= \frac{1}{P + \left(\frac{1-P}{S-A} + \frac{1-P}{S-B} \right)}$$

So $P = 0.6$, $S-A = 2.5$, $\&$ $S-B = 2.5$

now speedup = $\frac{1}{0.6 + \left(\frac{0.4}{2.5} + \frac{0.4}{2.5} \right)}$

$$= \frac{1}{0.6 + 0.16 + 0.16}$$

$$= \frac{1}{0.92}$$

$$= \boxed{1.09}$$

Therefore, the maximum theoretical speedup we can achieve with 4 additional nodes is $\boxed{1.09}$.

Problem #2: A distributed system has a task that can be divided into 4 sub-tasks each taking an equal amount of time. The system has 8 nodes, but due to communication overhead only 95% of the sub-task can be parallelized. What is the max theoretical speedup we can achieve on this system?

Solution:

In this case, we use Amdahl's law to calculate the maximum theoretical speedup for each sub-task.

$$\text{max speedup} = \frac{1}{P + \left(\frac{1-P}{S} \right)}$$

where P is the fraction of the sub-task that can't be parallelized & S is the speedup factor.

Example 2
 $p = 0.05$ (45% can be parallelized)

$N = 8 + 1 = 9$ (8 nodes + 1 for serial portion)

max speedup for each sub-task = $\frac{1}{0.05 + \left(\frac{1-0.05}{4}\right)}$

$= 1.82$

the overall maximum theoretical speedup is given by the combined effect of all 4 sub-tasks

max speedup = $\frac{1}{\dots}$

$$p + \left(\frac{1-p}{S_1}\right) + \left(\frac{1-p}{S_2}\right) + \left(\frac{1-p}{S_3}\right) + \left(\frac{1-p}{S_4}\right)$$

where p is fraction of the program that can't be parallelized & S_1, S_2, S_3, S_4 are speedup factor for each of the 4 sub-tasks.

Now $p = 0.05$

$$S_1 + S_2 + S_3 + S_4 = 4 \times 1.82$$

Now speedup = $\frac{1}{\dots}$

$$0.05 + \left(\frac{1-0.05}{1.82}\right) + \left(\frac{1-0.05}{1.82}\right) + \left(\frac{1-0.05}{1.82}\right) + \left(\frac{1-0.05}{1.82}\right)$$
$$= 3.48$$

Therefore the maximum theoretical speedup we can achieve on this system is

3.48

$P = 0.05$ (95% can be parallelized)

$S = 8 + 1 = 9$ (8 nodes + 1 for serial portion)

$$\text{max speedup for each sub-task} = \frac{1}{0.05 + \left(\frac{1 - 0.05}{9} \right)} = 1.82$$

The overall maximum theoretical speedup is given by the combined effect of all 4 sub-tasks

$$\text{max speedup} = \frac{1}{P + \left(\frac{1-P}{S_1} \right) + \left(\frac{1-P}{S_2} \right) + \left(\frac{1-P}{S_3} \right) + \left(\frac{1-P}{S_4} \right)}$$

where P is fraction of the program that can't be parallelized & S_1, S_2, S_3, S_4 are speedup factor for each of the 4 sub-tasks.

Here $P = 0.05$

$$S_1 = S_2 = S_3 = S_4 = 1.82$$

$$\begin{aligned} \text{Max speedup} &= \frac{1}{0.05 + \left(\frac{1 - 0.05}{1.82} \right) + \left(\frac{1 - 0.05}{1.82} \right) + \left(\frac{1 - 0.05}{1.82} \right) + \left(\frac{1 - 0.05}{1.82} \right)} \\ &= 3.48 \end{aligned}$$

Therefore the maximum theoretical speedup we can achieve on this system is

$$\boxed{3.48}$$