# Semantics and First-Order Predicate Calculus

11-711 Algorithms for NLP

28 November 2017

(With thanks to Noah Smith)

#### **Key Challenge of Meaning**

 We actually say very little - much more is left unsaid, because it's assumed to be widely known.

- Examples:
  - Reading newspaper stories
  - Using restaurant menus
  - Learning to use a new piece of software

#### Meaning Representation Languages

- Symbolic representation that does two jobs:
  - Conveys the meaning of a sentence
  - Represents (some part of) the world
- We're assuming a very literal, context-independent, inference-free version of meaning!
  - Semantics vs. linguists' "pragmatics"
  - "Meaning representation" vs some philosophers' use of the term "semantics".
- Today we'll use first-order logic. Also called First-Order Predicate Calculus. Logical form.

#### A MRL Should Be Able To ...

- Verify a query against a knowledge base: Do CMU students follow politics?
- Eliminate ambiguity: CMU students enjoy visiting Senators.
- Cope with vagueness: Sally heard the news.
- Cope with many ways of expressing the same meaning (canonical forms): The candidate evaded the question vs. The question was evaded by the candidate.
- Draw conclusions based on the knowledge base: Who could become the 45th president?
- Represent all of the meanings we care about

#### **Model-Theoretic Semantics**

- Model: a simplified representation of (some part of) the world: objects, properties, relations (domain).
- Non-logical vocabulary
  - Each element denotes (maps to) a well-defined part of the model
  - Such a mapping is called an interpretation

#### A Model

- Domain: Noah, Karen, Rebecca, Frederick, Green Mango, Casbah, Udipi, Thai, Mediterranean, Indian
- Properties: Green Mango and Udipi are crowded; Casbah is expensive
- Relations: Karen likes Green Mango, Frederick likes Casbah, everyone likes Udipi, Green Mango serves Thai, Casbah serves Mediterranean, and Udipi serves Indian
- n, k, r, f, g, c, u, t, m, i
- Crowded =  $\{g, u\}$
- Expensive = {c}
- Likes = {(k, g), (f, c), (n, u), (k, u), (r, u), (f, u)}
- Serves =  $\{(g, t), (c, m), (u, i)\}$

#### Some English

- Karen likes Green Mango and Frederick likes Casbah.
- Noah and Rebecca like the same restaurants.
- Noah likes expensive restaurants.
- Not everybody likes Green Mango.

- What we want is to be able to represent these statements in a way that lets us compare them to our model.
- Truth-conditional semantics: need operators and their meanings, given a particular model.

#### First-Order Logic

- Terms refer to elements of the domain: constants, functions, and variables
  - Noah, SpouseOf(Karen), X
- Predicates are used to refer to sets and relations;
   predicate applied to a term is a Proposition
  - Expensive(Casbah)
  - Serves(Casbah, Mediterranean)
- Logical connectives (operators):
  - $\land$  (and),  $\lor$  (or),  $\neg$  (not),  $\Rightarrow$  (implies), ...
- Quantifiers ...

#### **Quantifiers in FOL**

- Two ways to use variables:
  - refer to one anonymous object from the domain (existential; ∃; "there exists")
  - refer to all objects in the domain (universal; ∀; "for all")

- A restaurant near CMU serves Indian food
   ∃x Restaurant(x) ∧ Near(x, CMU) ∧ Serves(x, Indian)
- All expensive restaurants are far from campus
   ∀x Restaurant(x) ∧ Expensive(x) ⇒ ¬Near(x, CMU)

#### Inference

- Big idea: extend the knowledge base, or check some proposition against the knowledge base.
- Forward chaining with modus ponens: given  $\alpha$  and  $\alpha \Rightarrow \beta$ , we know  $\beta$ .
- **Backward chaining** takes a query  $\beta$  and looks for propositions  $\alpha$  and  $\alpha \Rightarrow \beta$  that would prove  $\beta$ .
  - Not the same as backward reasoning (abduction).
  - Used by Prolog
- Both are sound, neither is complete by itself.

#### Inference example

Starting with these facts:

Restaurant(Udipi)

 $\forall$  x Restaurant(x)  $\Rightarrow$  Likes(Noah, x)

We can "turn a crank" and get this new fact:

Likes(Noah, Udipi)

#### **FOL: Meta-theory**

- Well-defined set-theoretic semantics
- Sound: can't prove false things
- Complete: can prove everything that logically follows from a set of axioms (e.g., with "resolution theorem prover")
- Well-behaved, well-understood
- Mission accomplished?

#### FOL: But there are also "Issues"

- "Meanings" of sentences are truth values.
- Only first-order (no quantifying over predicates [which the book does without comment]).
- Not very good for "fluents" (time-varying things, real-valued quantities, etc.)
- Brittle: anything follows from any contradiction(!)
- Goedel incompleteness: "This statement has no proof"!
  - (Finite axiom sets are incomplete w.r.t. the real world.)
- So: Most systems use its descriptive apparatus (with extensions) but not its inference mechanisms.

#### First-Order Worlds, Then and Now

- Interest in this topic (in NLP) waned during the 1990s and 2000s.
- It has come back, with the rise of semi-structured databases like Wikipedia.
  - Lay contributors to these databases may be helping us to solve the knowledge acquisition problem.
- Also, lots of research on using NLP, information extraction, and machine learning to grow and improve knowledge bases from free text data.
  - "Read the Web" project here at CMU.
- And: Semantic embedding/NN/vector approaches.

#### Lots More To Say About MRLs!

- See chapter 17 for more about:
  - Representing events and states in FOL
  - Dealing with optional arguments (e.g., "eat")
  - Representing time
  - Non-FOL approaches to meaning

# **Connecting Syntax and Semantics**

#### **Semantic Analysis**

- Goal: transform a NL statement into MRL (today, FOL).
- Sometimes called "semantic parsing."
- As described earlier, this is the literal, contextindependent, inference-free meaning of the statement

# "Literal, context-independent, inference-free" semantics

- Example: The ball is red
- Assigning a specific, grounded meaning involves deciding which ball is meant
- Would have to resolve indexical terms including pronouns, normal NPs, etc.
- Logical form allows compact representation of such indexical terms (vs. listing all members of the set)
- To retrieve a specific meaning, we combine LF with a particular context or situation (set of objects and relations)
- So LF is a function that maps an initial discourse situation into a new discourse situation (from situation semantics)

## Compositionality

- The meaning of an NL phrase is determined by combining the meaning of its sub-parts.
- There are obvious exceptions ("hot dog," "straw man," "New York," etc.).

 Note: your book uses an event-based FOL representation, but I'm using a simpler one without events.

• Big idea: start with parse tree, build semantics on top using FOL with  $\lambda$ -expressions.

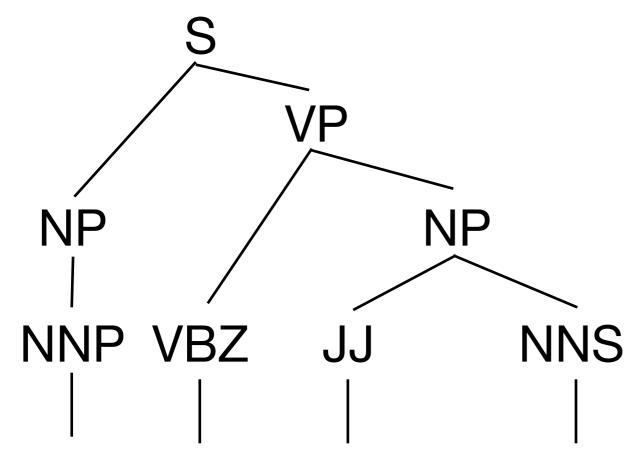
#### **Extension: Lambda Notation**

- A way of making anonymous functions.
- $\lambda x$ . (some expression mentioning x)
  - Example: λx.Near(x, CMU)
  - Trickier example: λx.λy.Serves(y, x)
- Lambda reduction: substitute for the variable.
  - (λx.Near(x, CMU))(LulusNoodles) becomes Near(LulusNoodles, CMU)

#### Lambda reduction: order matters!

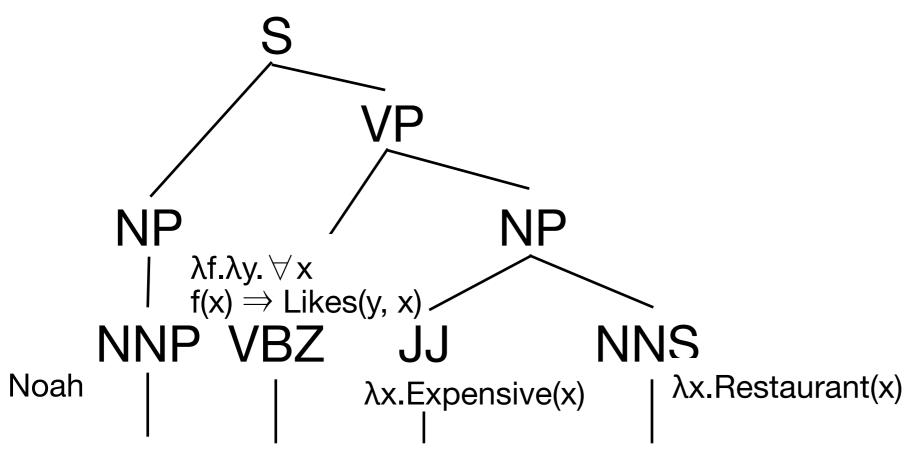
λx.λy.Serves(y, x) (Bill)(Jane) becomes λy.Serves(y, Bill)(Jane)
 Then λy.Serves(y, Bill) (Jane) becomes Serves(Jane, Bill)

λy.λx.Serves(y, x) (Bill)(Jane) becomes λx.Serves(Bill, x)(Jane)
 Then λx.Serves(Bill, x) (Jane) becomes Serves(Bill, Jane)



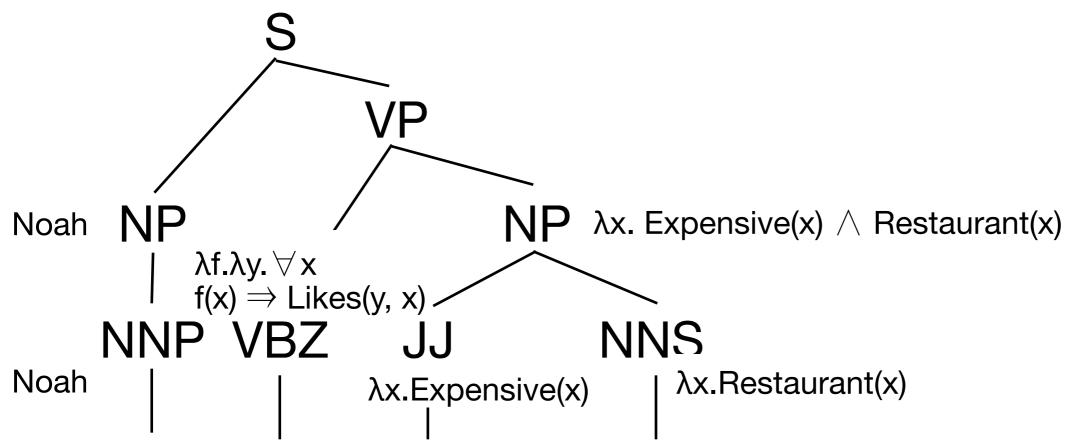
- Noah likes expensive restaurants.
- $\forall$  x Restaurant(x)  $\land$  Expensive(x)  $\Rightarrow$  Likes(Noah, x)

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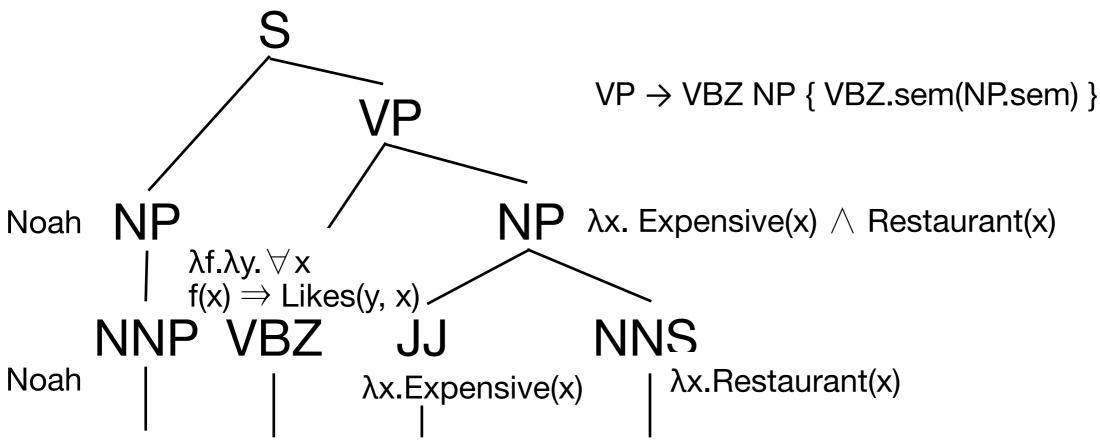


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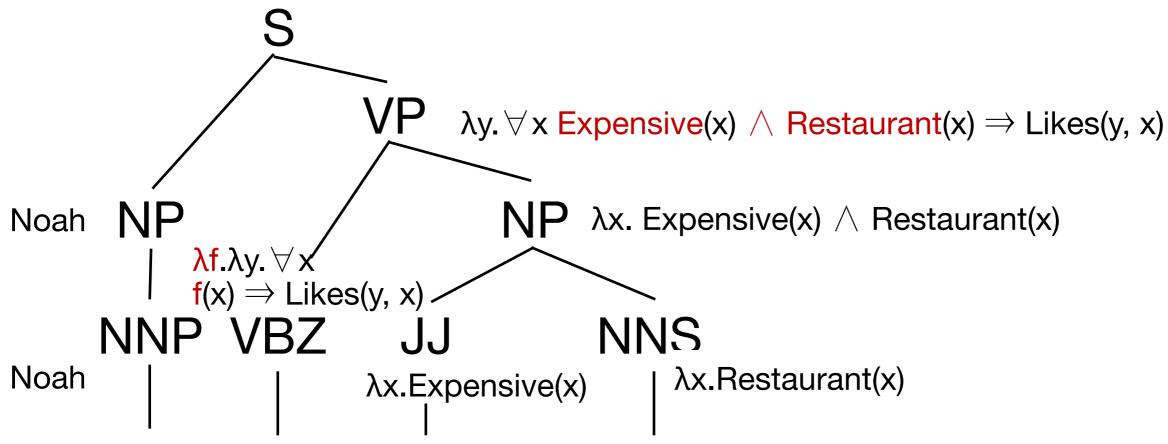
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```
S \Rightarrow \text{NP VP \{ VP.sem(NP.sem) \}} VP \quad \lambda y. \ \forall x \ \text{Expensive}(x) \ \land \ \text{Restaurant}(x) \Rightarrow \text{Likes}(y, x) Noah \quad NP \quad NP \quad \lambda x. \ \text{Expensive}(x) \ \land \ \text{Restaurant}(x) NNP \quad VBZ \quad JJ \quad NNS \quad Noah \quad | \quad \lambda x. \text{Expensive}(x) \quad | \quad \lambda x. \text{Restaurant}(x)
```

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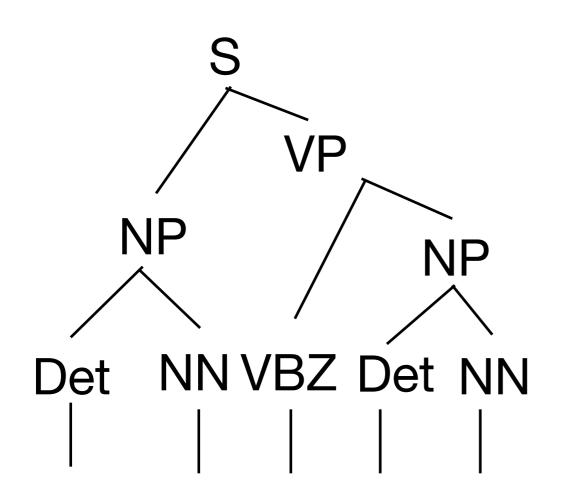
```
S \\ \hline VP \\ \lambda y. \\ \forall x \  \, \text{Expensive}(x) \land \  \, \text{Restaurant}(x) \Rightarrow \text{Likes}(\text{Noah, } x) \\ \hline VP \\ \lambda y. \\ \forall x \  \, \text{Expensive}(x) \land \  \, \text{Restaurant}(x) \Rightarrow \text{Likes}(y, x) \\ \hline NOah \\ \hline NP \\ \lambda x. \  \, \text{Expensive}(x) \land \  \, \text{Restaurant}(x) \\ \hline NNP \\ VBZ \\ JJ \\ NNS \\ Noah \\ \hline \\ \lambda x. \  \, \text{Expensive}(x) \\ \hline \lambda x. \  \, \text{Expensive}(x) \\ \hline \lambda x. \  \, \text{Restaurant}(x) \\ \hline \\ \lambda x. \  \, \text{Expensive}(x) \\ \hline \\ \lambda x. \  \, \text{Restaurant}(x) \\ \hline \\ \lambda x. \  \, \text{Expensive}(x) \\ \hline \\ \lambda x. \  \, \text{Restaurant}(x) \\ \hline \\ \lambda x. \  \, \text{Expensive}(x) \\ \hline \\ \lambda x. \  \, \text{Expensive}(x) \\ \hline \\ \lambda x. \  \, \text{Restaurant}(x) \\ \hline \\ \lambda x. \  \, \text{Expensive}(x) \\ \hline \\ \lambda x. \  \, \text{Restaurant}(x) \\ \hline \\ \lambda x. \  \, \text{Expensive}(x) \\ \hline \\ \lambda x.
```

- Noah likes expensive restaurants.
- $\forall$  x Restaurant(x)  $\land$  Expensive(x)  $\Rightarrow$  Likes(Noah, x)

# Alternative (Following SLP)

- · Noah likes expensive restaurants.
- $\forall$  x Restaurant(x)  $\land$  Expensive(x)  $\Rightarrow$  Likes(Noah, x)

## **Quantifier Scope Ambiguity**



```
S \rightarrow NP \ VP \ \{ \ NP.sem(VP.sem) \}
NP \rightarrow Det \ NN \ \{ \ Det.sem(NN.sem) \}
VP \rightarrow VBZ \ NP \ \{ \ VBZ.sem(NP.sem) \}
Det \rightarrow every \ \{ \ \lambda f. \lambda g. \ \forall \ u \ f(u) \Rightarrow g(u) \}
Det \rightarrow a \ \{ \ \lambda m. \lambda n. \ \exists \ x \ m(x) \ \land \ n(x) \}
NN \rightarrow man \ \{ \ \lambda v.Man(v) \}
NN \rightarrow woman \ \{ \ \lambda y.Woman(y) \}
VBZ \rightarrow loves \ \{ \ \lambda h. \lambda k.h(\lambda w. \ Loves(k, w)) \}
```

Every man loves a woman.

•  $\forall$  u Man(u)  $\Rightarrow \exists$  x Woman(x)  $\land$  Loves(u, x)

## This Isn't Quite Right!

- "Every man loves a woman" really is ambiguous.
  - $\forall$  u Man(u)  $\Rightarrow \exists$  x Woman(x)  $\land$  Loves(u, x)
  - $\exists x \text{ Woman}(x) \land \forall u \text{ Man}(u) \Rightarrow \text{Loves}(u, x)$

- This gives only one of the two meanings.
  - Extra ambiguity on top of syntactic ambiguity
- One approach is to delay the quantifier processing until the end, then permit any ordering.

#### **Quantifier Scope**

- A seat was available for every customer.
- A toll-free number was available for every customer.

- A secretary called each director.
- A letter was sent to each customer.

- Every man loves a woman who works at the candy store.
- Every 5 minutes a man gets knocked down and he's not too happy about it.

#### What Else?

- Chapter 18 discusses how you can get this to work for other parts of English (e.g., prepositional phrases).
- Remember attribute-value structures for parsing with more complex things than simple symbols?
  - You can extend those with semantics as well.
- No time for ...
  - Statistical models for semantics
  - Parsing algorithms augmented with semantics
  - Handling idioms

#### **Generalized Quantifiers**

- In FOL, we only have universal and existential quantifiers
- One formal extension is type-restriction of the quantified variable: Everyone likes Udipi:

```
\forall x Person(x) \Rightarrow Likes(x, Udipi) becomes
```

 $\forall$  x | Person(x).Likes(x, Udipi)

- English and other languages have a much larger set of quantifiers: all, some, most, many, a few, the, ...
- These have the same form as the original FOL quantifiers with type restrictions:

<quant><var>|<restriction>.<body>

#### Generalized Quantifier examples

- Most dogs bark
  - Most  $x \mid Dog(x)$ . Barks(x)
- Most barking things are dogs
  - Most  $x \mid Barks(x) \cdot Dog(x)$
- The dog barks
  - The  $x \mid Dog(x)$ . Barks(x)
- The happy dog barks
  - The x | (Happy(x)  $\land$  Dog(x)) . Barks(x)
- Interpretation and inference using these are harder...

#### **Speech Acts**

- Mood of a sentence indicates relation between speaker and the concept (proposition) defined by the LF
- There can be operators that represent these relations:
  - ASSERT: the proposition is proposed as a fact
  - YN-QUERY: the truth of the proposition is queried
  - COMMAND: the proposition describes a requested action
  - WH-QUERY: the proposition describes an object to be identified

#### **ASSERT (Declarative mood)**

The man eats a peach

ASSERT(The x | Man(x) . (A y | Peach(y) . Eat(x,y)))

#### YN-QUERY (Interrogative mood)

Does the man eat a peach?

YN-QUERY(The x | Man(x) . (A y | Peach(y) . Eat(x,y)))

# **COMMAND** (Imperative mood)

Eat a peach, (man).

COMMAND(A y | Peach(y) . Eat(\*HEARER\*,y))

#### **WH-QUERY**

- What did the man eat?
  - WH-QUERY(The x | Man(x) . (WH y | Thing(y) . Eat(x,y)))
- One of a whole set of new quantifiers for wh-questions:
  - What: WH x | Thing(x)
  - Which dog: WH x | Dog(x)
  - Who: WH x | Person(x)
  - How many men: HOW-MANY x | Man(x)

#### Other complications

- Relative clauses are propositions embedded in an NP
  - Restrictive versus non-restrictive: the dog that barked all night vs. the dog, which barked all night
- Modal verbs: non-transparency for truth of subordinate clause: Sue thinks that John loves Sandy
- Tense/Aspect
- Plurality
- Etc.