The Past, Present and Future of TIPC Protocol

Ying Xue October 18, 2015

















Agenda

- TIPC Introduction
- TIPC Evolution History
- Demo



- TIPC = Transparent Inter Process Communication
- A communication protocol designed for **Clustered** computers

Cluster == Cloud



TIPC == Simple

- Service can run anywhere without changing its address
 - Service address is always valid
 - Service address is decoupled from physical location
 - Service address space can be limited
- Services can start to talk anytime whether its peer is launched or not



TIPC == Powerful

- Service can run as usual even if one of two interfaces it depends on is down
 - Redundant dual links
 - Link load sharing vs active/standby
- Service can be timely aware of any change of network topology or services if it wants
 - Service and topology tracking function
 - No more heat-beating



TIPC == Reliable

- Reliable datagram unicast and multicast
 - No real flow control, message may be still rejected.
 - Rejected message may be dropped or returned
 - Multicast cannot be made returnable
- Connections with stream or message transport
 - Traffic control algorithm



Where is TIPC?

Main TIPC website: http://tipc.sourceforge.net/

- Linux
 - TIPC 1.6.x in Linux kernel as of 2.6.16
 - TIPC available as add-on for 2.6.16-2.6.34
 - TIPC 2.0 available as add-on for 2.6.35+
- VxWorks
 - TIPC 1.7.5 in VxWorks 6.7 (earlier versions from VxWorks 6.1)
 - TIPC 1.7.7 available as add-on for all 6.x



When is TIPC released?

- 1990's
 - Single cluster proprietary versions created at Ericsson
- 2000
 - TIPC 1.3/1.4: Initial Open Source Linux version
- 2004
 - Wind River TIPC 1.0: TIPC subset ported to VxWorks
- 2006
 - TIPC 1.5/1.6 integrated into Linux & VxWorks
- 2007
 - TIPC 1.7: multi-cluster capability



When is TIPC released(Cont.)?

- 2012
 - TIPC 2.x: Linux kernel



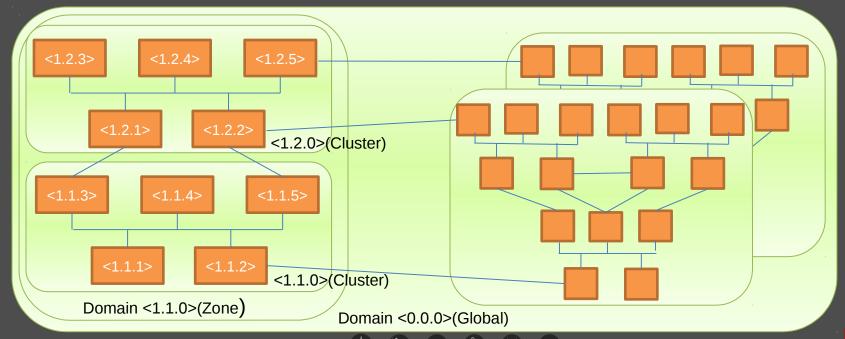
Who is doing TIPC?

- Developer
 - Jon Maloy at Ericsson
 - Ying Xue at Wind River
 - Erik Hugne at Ericsson
 - Richard Alpe at Ericsson
- Users
 - Telecom equipment manufactures
 - HA (High Available), such as OpenSAF



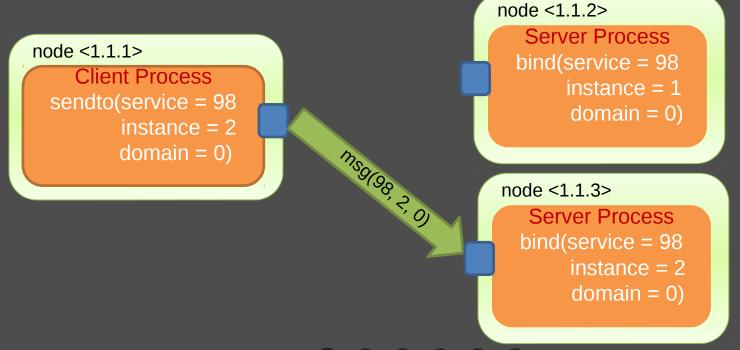
TIPC Network Topology

32bit domain identifier can be assigned to node, cluster and zone with <Z.C.N> structure



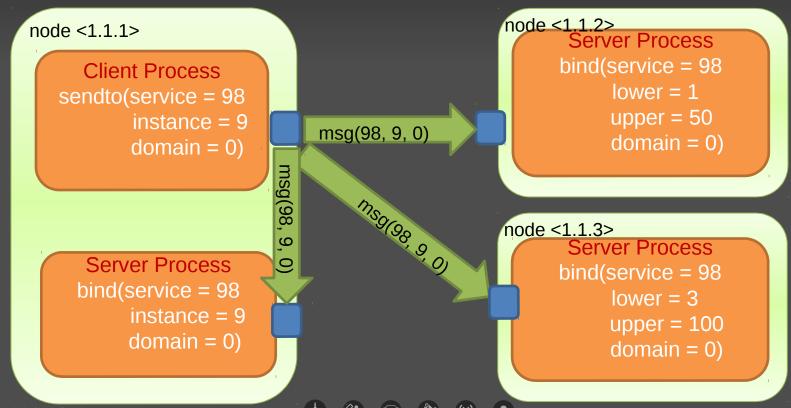
Location Transparency

Physical location of server is not known by client





Reliable Multicast



Service Subscription

node <1.1.1>
Client Process
subscribe(service = 98
lower = 0
upper =
100)

event(98, 1, <1.1.27 up)

event(98, 2, \(\sigma 1.1.3\(\sigma \text{up}\)

node <1.1.2>
Server Process
bind(service = 98
instance = 1
domain = 0)

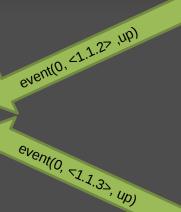
node <1.1.3>
Server Process

bind(service = 98 instance = 2 domain = 0)

Node Subscription

node <1.1.1>

Client Process
subscribe(service = 0
lower = 0
upper = ~0)

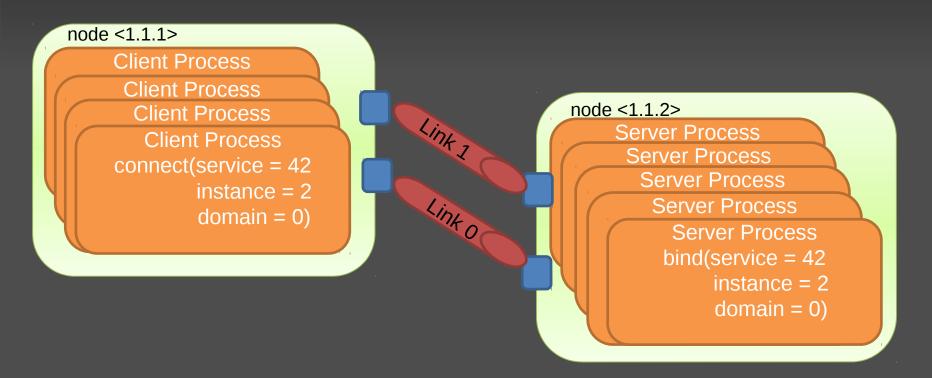


node <1.1.3>

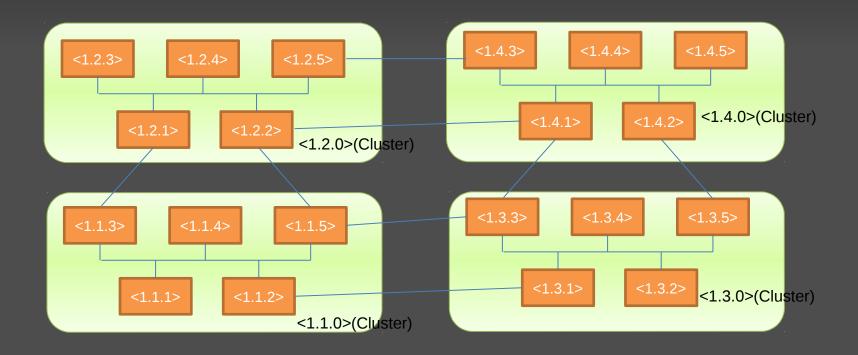
node <1.1.2>



Redundant Dual Links

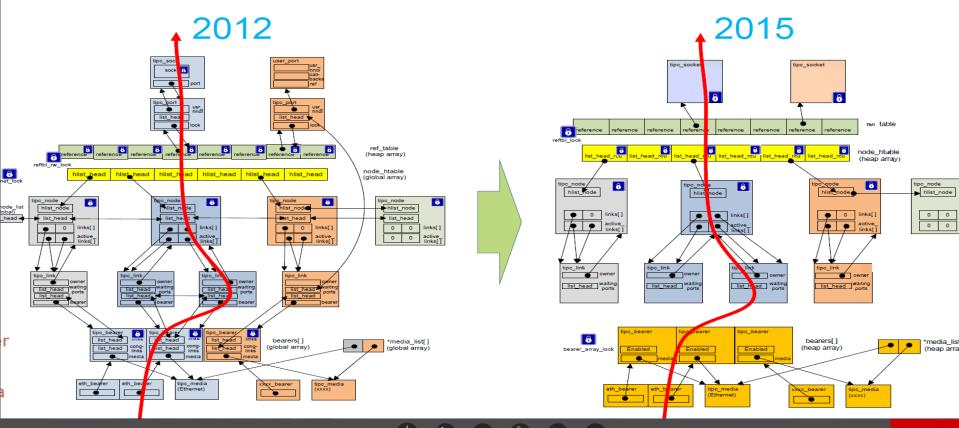


Neighbor Discovery





What Changed?



What Changed?

- Locking Policy
 - Rewrote internal server
 - Decoupled the close relationship between tipc_node and tipc_link structures
 - Introduced reference count for tipc_node structure
 - Replaced RW lock with RCU to protect node list
 - Replaced RW lock with RCU to protect name table
 - Introduced RCU lock to manage bearer list
 - Used RTNL lock to manage the process of configuring both bearer and link
 - Eliminated net lock



What Changed(Cont.)?

- Locking Policy
 - Port structure was purged
 - Introduced rhashtable to manage socket references
- New Important Functionalities
 - Name space was supported
 - Introduced UDP bearer
 - Migrated TIPC configuration tool from tipc_config to ip through netlink



What will be Changed?

- Traffic Control
 - Improve link flow control algorithm, such as adaptable window size and congestion avoidance
 - Datagram and multicast congestion feedback
 - Improve multicast flow control algorithm
- Stability
 - Support full network space
 - Further reduce the area protected by node
 - Further improve network throughout



What will be Changed(Cont.)?

- Optimization
 - Overhaul of broadcast link
 - Introduce new mechanism to take precautions against the overload of socket receive buffer
 - Align code style with Linux



Demo



Questions?

- Ask questions through tipc-discussion mail list:
 - tipc-discussion@lists.sourceforge.net
 - http://sourceforge.net/p/tipc/mailman/tipc-discussion/

