## Complex Critical Systems in Event-B

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#### Outline

- Environment Modelling
- 2 Medical Devices
- Formal Development
- 4 Interactive Simulation
- Code Generation



# Safety-Critical System

#### Definition

A life-critical system or safety-critical system is a system whose failure or malfunction may result in serious injuries, loss of life, economical damage and environmental harm.



# Critical System Failure

### Systems Failure

- Therac-25 (1985-1987): six people overexposed through radiation.
- Pacemaker and ICD (1990-2002): 17,323 pacemakers and ICDs were explanted that includes 61 deaths.
- Insulin Infusion Pump (IIP) (2010): 5000 adverse events that includes 30 deaths.
- Missing Malaysian Plane MH370 (8 March, 2014): Unknown.
- Satellite Failure: +150: http://www.sat-nd.com/failures/.



# The Cardiac Pacemaker (Grand Challenge)

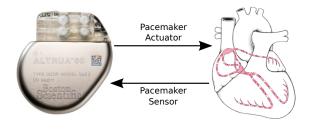


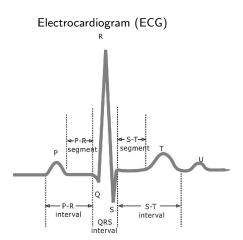
Figure: Closed-loop Modelling of the Heart & Cardiac Pacemaker



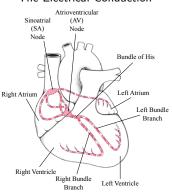
# The Human Heart



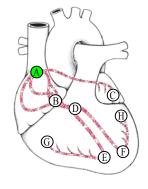
# The Heart Modeling



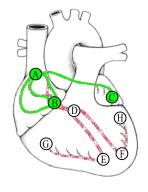
#### The Electrical Conduction



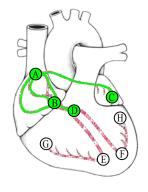




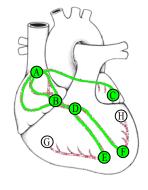




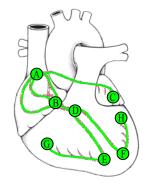














#### Heart Definition

## Definition 1 (The Heart System)

Given a set of nodes N, a transition (conduction) T is a pair (i, j), with  $i, j \in N$ . A transition is denoted by  $i \leadsto j$ . The heart system is a tuple  $HSys = (N, T, N_0, TW_{time}, CW_{speed})$  where:

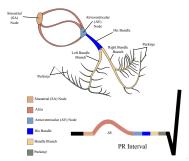
- ullet  $N=\{$  A, B, C, D, E, F, G, H $\}$  is a finite set of landmark nodes;
- $T \subseteq N \times N = \{A \mapsto B, A \mapsto C, B \mapsto D, D \mapsto E, D \mapsto F, E \mapsto G, F \mapsto H\}$  is a set of transitions;
- $N_0 = A$  is the initial landmark node;
- $TW_{time} \in N \rightarrow TIME$  is a weight function as time delay;
- $CW_{speed} \in T \rightarrow SPEED$  is a weight function as impulse propagation speed.



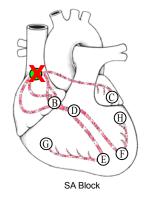
## Time Intervals and Impulse Propagation in the ECG

Location in the heart	Impulse Propagation Time (ms.) Property 1 (TW <sub>time</sub> )	Location in the heart	Impulse Propagation Speed (cm/sec.) Property 2 (CW <sub>speed</sub> )
SA Node (A)	010	$A \mapsto B$	3050
Left atria muscle fibers (C)	7090	$A \mapsto C$	3050
AV Node (B)	5070	$B \mapsto D$	100200
Bundle of His (D)	125160	$D \mapsto E$	100200
Right Bundle Branch (E)	145180	$D \mapsto F$	100200
Left Bundle Branch (F)	145180	$E \mapsto G$	300400
Right Purkinje fibers (G)	150210	$F \mapsto H$	300400
Left Purkinje fibers (H)	150230		

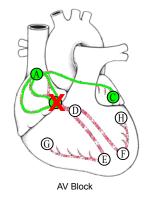
Table: Cardiac Activation Time and Cardiac Velocity



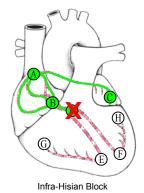




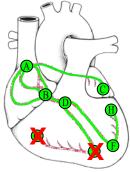






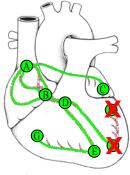






Right Bundle Branch Block





Left Bundle Branch Block

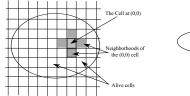


#### 2D Cellular Automata and State Transition Model

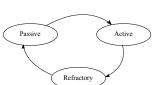
## Definition 3 (State Transition of a Cell)

The heart muscle system is composed of heterogeneous cells, the cellular automata model of the muscle system,  $CAM_{CA}$ , is characterized with no dependencies on the type of cells.  $CAM_{CA}$  is defined as follows:

$$\begin{split} &CAM_{CA} = \langle S, N, T \rangle \\ &S = \{Active, Passive, Refractory\} \\ &N_{m,n} = \{(m,n), (m+1,n), (m-1,n), (m,n+1), (m,n-1)\} \\ &T: S^{|N|} \rightarrow S \\ &s_{t+1}(m,n) = \begin{cases} Refractory & \text{if } s_t(m,n) = Active \\ Passive & \text{if } s_t(m,n) = Refractory \\ Active & \text{if } s_t(m,n) = Passive \text{ and any neighbor is in } Active \text{ state} \\ Passive & \text{if } s_t(m,n) = Passive \text{ and none neighbor is in } Active \text{ state} \end{cases} \end{split}$$



where,  $s_t(m, n)$  denotes the state of the cell located at (m,n).





## Development in Event-B

#### Abstract Model and Chain of Refinements

- Abstract Model
- Refinement 1: Introducing Steps in the Propagation
- Refinement 2: Impulse Propagation
- Refinement 3: Perturbation the Conduction
- Refinement 4: Getting a Cellular Model

#### **Proof Statistics**

Model	Total number	Automatic	Interactive
	of POs	Proof	Proof
Abstract Model	29	22(76%)	7(24%)
First Refinement	9	6(67%)	3(33%)
Second Refinement	159	155(97%)	4(3%)
Third Refinement	10	1(10%)	9(90%)
Fourth Refinement	11	10(91%)	1(9%)
Total	218	194(89%)	24(11%)

# The Cardiac Pacemaker



#### The Cardiac Pacemaker

#### Pacemaker

A pacemaker is an electronic device implanted in a body to regulate the abnormal heart rhythm (bradycardia). Type of pacemakers:1,2 and 3-Electrodes.

#### The Cardiac Pacemaker



## Operating Modes: NASPE/BPEG Generic Code

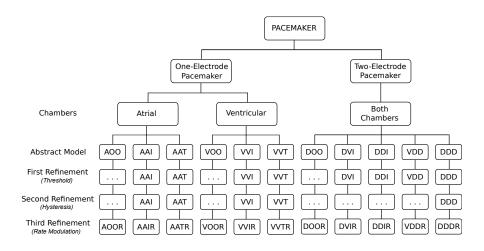
Category	Chambers	Chambers	Response to	Rate Modulation
	Paced	Sensed	Sensing	
Letters	<b>O</b> -None	<b>O</b> -None	<b>O</b> -None	<b>R</b> -Rate Modulation
	<b>A</b> -Atrium	<b>A</b> -Atrium	<b>T</b> -Triggered	
	<b>V</b> -Ventricle	<b>V</b> -Ventricle	I-Inhibited	
	$\mathbf{D}$ -Dual(A+V)	<b>D</b> -Dual(A+V)	D-Dual $(T+I)$	

i.e. AOO, VOO, AAI, AAT, VVI, VVT, AATR, VVTR, AOOR etc...

Periodic stimuli: (AOO, VOO and DOO) Aperiodic stimuli: (AAI, VVI, DDD, DDI, etc.)

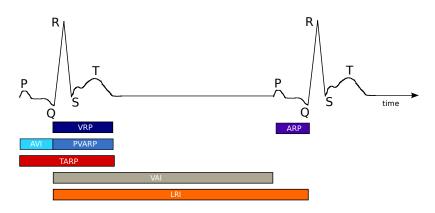


# Hierarchical Development of Operating Modes





## Timing Cycles



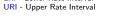
VRP - Ventricular Refractory Period ARP - Atrial Refractory Period

AVI - Atrioventricular Interval

PVARP - Post Ventricular Atrial Refractory Period

TARP - Total Atrial Refractory Period

VAI - Ventriculoatrial Interval LRI - Lower Rate Interval



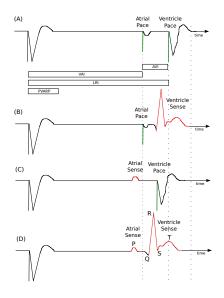


# General Requirements (Boston Scientific)

- The bradycardia operating modes (AOO, AAT, VVI, ..., DDDR) shall be programmable.
- 2 The Pacemaker shall support single and dual chamber pacing modes.
- The device shall actuate for pacing in the heart with programmable voltages and widths.
- The pacing pulse amplitudes and pacing width shall be programmable for each chamber (e.g. atrial, ventricular).
- The device shall support rate adaptive pacing in order to meet the physiological needs.



# The DDD Pacing Scenarios





# DOO and DDD Modes Requirements

- The pacemaker in DOO mode must pace in both atria and ventricular chambers without sensing any intrinsic activities from both chambers according to the programmable parameters.
- In DDD mode, when the pacing system senses intrinsic activities in both chambers, the pacemaker must be triggered or inhibited to pace in both chambers according to the programmable parameters.

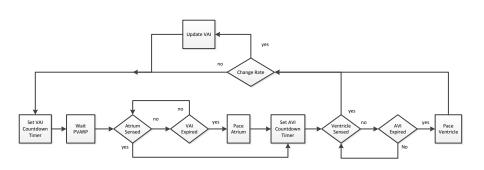


# Rate Adaptive Modes (AOOR, DOOR, ..., DDDR) Requirements

• The pacemaker in rate adaptive mode must pace and sense in any chamber according to the selected mode by changing the rate to meet the physical needs according to the programmable parameters.



# Flowchart for DDDR Pacing Cycle



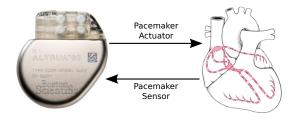


# Formal Development



# Closed-loop Model (Heart & Cardiac Pacemaker)

- Simple Pacemaker Model
- Closed-loop Model



- Abstract Model: Pacing and Sensing Activities + Normal and Abnormal Heart Behaviour.
- Refinement 1: Threshold + Impulse Propagation .
- Refinement 2: Hysteresis + Perturbation the Conduction.
- Refinement 3: Rate Modulation + Cellular Model.



## Formalization of the Closed-loop model: Abstract Model

```
\begin{array}{l} \textit{axm1}: \textit{partition}(\textit{ConductionNode}, \{A\}, \{B\}, \{C\}, \{D\}, \{E\}, \{F\}, \{G\}, \{H\})) \\ \textit{axm2}: \textit{ConductionTime} \in \textit{ConductionNode} \rightarrow \mathbb{P}(0 \dots 230) \\ \textit{axm3}: \textit{ConductionPath} \subseteq \textit{ConductionNode} \times \textit{ConductionNode} \\ \textit{axm4}: \textit{ConductionSpeed} \in \textit{ConductionPath} \rightarrow \mathbb{P}(5 \dots 400) \\ \textit{axm5}: \textit{LRL} \in 30 \dots 175 \land \textit{URL} \in 50 \dots 175 \land \textit{PVARP} \in 150 \dots 500 \\ \textit{axm6}: \textit{ARP} \in 150 \dots 500 \land \textit{VRP} \in 150 \dots 500 \land \textit{status} = \{\textit{ON}, \textit{OFF}\} \\ \end{array}
```

```
inv1: ConductionNodeState \in ConductionNode \rightarrow BOOL
inv2: CConductionTime \in ConductionNode \rightarrow 0...300
inv3: CConductionSpeed \in ConductionPath \rightarrow 0..500
inv4 : HeartState ∈ BOOL
inv5: PM\_Actuator\_A \in status \land PM\_Actuator\_V \in status
inv6: PM\_Sensor\_A \in status \land PM\_Sensor\_V \in status
inv7: Pace_Int \in URI ... LRI \land sp \in 1... Pace_Int
inv8: sp < VRP \land sp < PVARP
     \Rightarrow
     PM Actuator V = OFF \land PM Sensor A = OFF \land
     PM Sensor V = OFF \land PM Actuator A = OFF
inv9: PM\_Actuator\_V = ON \Rightarrow sp = Pace\_Int
inv10 : PM\_Actuator\_A = ON \Rightarrow (sp \ge Pace\_Int - FixedAV)
. . .
```



## **DDD Pacing Modes**

# EVENT Actuator\_ON\_V WHEN

 $grd1: PM\_Actuator\_V = OFF$ 

 $grd2: sp = Pace\_Int$ 

 $grd3: sp \geq VRP \land sp \geq PVARP$ 

THEN

 $act1 : PM\_Actuator\_V := ON$ 

 $act2: last\_sp := sp$ 

END

# EVENT Actuator\_OFF\_V WHEN

 $grd1: PM\_Actuator\_V = ON$ 

 $grd2: sp = Pace\_Int$ 

 $grd3 : PM\_Actuator\_A = OFF$  $grd4 : PM\_Sensor\_A = OFF$ 

THEN

 $act1 : PM\_Actuator\_V := OFF$ 

 $act2 : AV_{-}Count := 0$ 

 $act3: PM\_Sensor\_V := OFF$ 

act4: sp := 1

END

# EVENT tic WHEN

 $grd1: sp < Pace\_Int$ 

THEN

 $\mathsf{act1}: \mathit{sp} := \mathit{sp} + 1$ 

**END** 



#### Continue...

# EVENT Sensor\_ON\_V WHEN

 $\begin{array}{l} \operatorname{grd1}: PM\_Sensor\_V = OFF \\ \operatorname{grd2}: (sp \geq PVARP \land sp > VRP \\ \operatorname{grd3}: sp > Pace\_Int - FixedAV \\ \operatorname{grd4}: PM\_Actuator\_A = OFF \\ \mathbf{THFN} \end{array}$ 

act1 : PM\_Sensor\_V := ON END

# EVENT Sensor\_OFF\_V WHEN

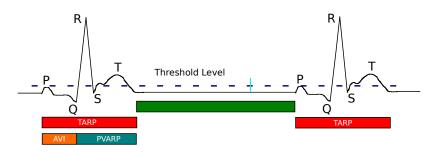
 $grd1: PM\_Sensor\_V = ON$   $grd2: sp \ge VRP \land sp \ge PVARP$   $grd3: sp \ge Pace\_Int - FixedAV$  $grd4: PM\_Actuator\_V = OFF$ 

THEN

act1 : PM\_Sensor\_V := OFF END



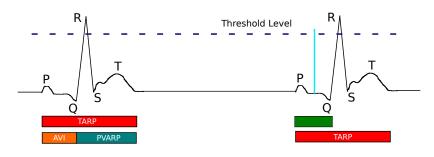
# First Refinement (Threshold): Sensor Activity in DDD



```
inv1: sp > PVARP \land sp < Pace\_Int - FixedAV \Rightarrow PM\_Sensor\_A = ON \land PM\_Sensor\_V = OFF \land PM\_Actuator\_A = OFF \land PM\_Actuator\_V = OFF
```



## First Refinement (Threshold): Sensor Activity in DDD



```
inv1: sp > Pace\_Int - FixedAV \land sp < Pace\_Int \Rightarrow PM\_Sensor\_V = ON \land PM\_Sensor\_A = OFF \land PM\_Actuator\_A = OFF \land PM\_Actuator\_V = OFF
```



## Second Refinement: Hysteresis

### What is Hysteresis?

Hysteresis is a programmed feature whereby the pacemaker paces at a faster rate than the sensing rate and it provides consistent pacing to the atrial or ventricle, or prevents the constant pacing of the atrial or ventricle. The main purpose of hysteresis is to allow the patient to have his or her own underlying rhythm as much as possible.

```
EVENT Hyt_Pace_Updating Refines Change_Pace_Int
ANY
Hyt_Pace_Int
WHERE
grd1: Pace_Int_flag = TRUE
grd2: Hyt_Pace_Int_flag = TRUE
grd3: Hyt_Pace_Int ∈ Pace_Int .. LRI
THEN
act1: Pace_Int := Hyt_Pace_Int
act2: Hyt_Pace_Int_flag := FALSE
act3: HYT_State := TRUE
END
```



#### Third Refinement: Rate Modulation

#### What is Rate Modulation?

```
Increase_Interval Refines Change_Pace_Int
WHEN
grd1: grd1: Pace_Int_flag = TRUE
grd1: aclor sonsed > threshold
```

grd1 : acler\_sensed ≥ threshold grd1 : HYT\_State = FALSE

#### THEN

act1 : Pace\_Int := 60000/MSR act1 : acler\_sensed\_flag := TRUE END

```
inv3 : acler\_sensed < acc\_thr \land acler\_sensed\_flag = TRUE \Rightarrow Pace\_Int = 60000/LRL inv4 : acler\_sensed \geq acc\_thr \land acler\_sensed\_flag = TRUE \Rightarrow Pace\_Int = 60000/MSR
```



# Model Validation & Analysis

#### ProB Model Checker

The ProB model checker is used for,

- to verify the developed system requirements;
- to check the required behaviour of the pacemaker for each operating mode;
- deadlock checking;
- to discover the counter examples;



# Simple Model Vs. Closed-loop Model

Model	Total number	Automatic	Interactive			
	of POs Proof		Proof			
Simple One-electrode pacemaker						
Abstract Model	203	199(98%)	4(2%)			
First Refinement	48	44(91%)	4(9%)			
Second Refinement	12	8(66%)	4(34%)			
Third Refinement	105	99(94%)	6(6%)			
Simple Two-electrode pacemaker						
Abstract Model	204	195(95%)	9(5%)			
First Refinement	234	223(95%)	11(5%)			
Second Refinement	3	3(100%)	0(0%)			
Third Refinement	83	74(89%)	9(11%)			
Total	892	845(94%)	47(6%)			

	T					
Model	Total number	Automatic	Interactive			
	of POs	Proof	Proof			
Closed-loop model of One-electrode pacemaker						
Abstract Model	304	258(85%)	46(15%)			
First Refinement	1015	730(72%)	285(28%)			
Second Refinement	72 8(11%)		64(89%)			
Third Refinement	153	79(52%)	74(48%)			
Closed-loop model of Two-electrode pacemaker						
Abstract Model	291	244(84%)	47(16%)			
First Refinement	1039	766(74%)	273(26%)			
Second Refinement	53	2(4%)	51(96%)			
Third Refinement	122	60(49%)	62(51%)			
Total	3049	2147(70%)	902(30%)			



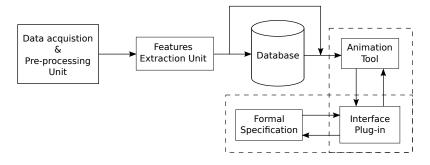
## Interactive Simulation



#### An Architecture of Real-time Animator

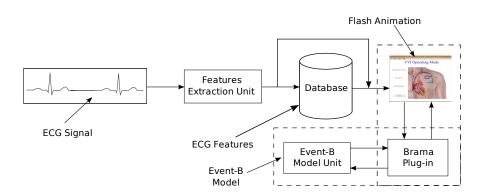
#### Real-Time Animator

Visual representation of a proved formal model using real-time data set to show the system behaviour to domain experts.





## Real-time Animation: Pacemaker





#### ECG Data and Features

- MIT-BIH Database Distribution
- Algorithms to calculate ECG Features (http://ecg.mit.edu/index.html)



## Demo



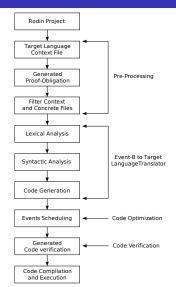
# EB2ALL: Automatic Code Generation from Event-B Models



#### EB2ALL Code Generation tool

## EB2ALL (Event-B to All)

- EB2C
- EB2C++
- EB2J
- EB2C#





## Pre-Processing

- Objective : Make a system deterministic
- Clean termination approach
- Introduction of new Context Files

Event-B type	Formal Range	C & C++ type	Java type	C# type
tl_int16	$-2^{15}2^{15}-1$	int	short	short
tl_uint16	$02^{16} - 1$	unsigned int	-	ushort
tl_int32	$-2^{31}2^{31}-1$	long int	int	int
tl_uint32	$02^{32}-1$	unsigned long int	-	uint
tl_int64	$-2^{63}2^{63}-1$	-	long	long
tl_uint64	$02^{64} - 1$	-	-	ulong

Table: Integer bounded data type declaration in different context files

#### Example:

 $\mathsf{I}d \in \mathbb{N}_1$   $\mathsf{I}d \in \mathit{tl\_uint}16$ 



## Basic Principles of Code Generation

	10101010	
Event-B	'C' & 'C++' Language	Comment
nm	int	Interger type
$x \in Y$	Y x;	Scaler declaration
x ∈ tl_int16	int x;	'C' & 'C++' Contexts
$x \in nm \rightarrow Y$	$Y \times [m+1];$	Array declaration
× :∈ Y	/* No Action */	Indeterminate Init.
x:   Y	/* No Action */	Indeterminate Init.
x = y	if(x==y) {	Conditional
$x \neq y$	if(x!=y) {	Conditional
x < y	if(x <y) td="" {<=""><td>Conditional</td></y)>	Conditional
$x \le y$	if(x<=y) {	Conditional
x > y	if(x>y) {	Conditional
$x \ge y$	if(x>=y) {	Conditional
(x>y) ∧ (x≥z)	if ((x>y) && (x>=z) {	Conditional
(x>y) ∨ (x≥z)	if $((x>y) \parallel (x>=z) $ {	Conditional
x := y + z	x = y + z;	Arithmetic assignment
x := y - z	x = y - z;	Arithmetic assignment
x := y * z	x = y * z;	Arithmetic assignment
$x := y \div z$	x = y / z;	Arithmetic assignment
x := F(y)	x = F(y);	Function assignment
$a := F(x \mapsto y)$	a = F(x, y);	Function assignment
x := a(y)	x = a[y];	Array assignment
x := y	x = y;	Scalar action
$a := a \Leftrightarrow \{x \mapsto y\}$	a[x] = y;	Array action
$a := a \Leftrightarrow \{x \mapsto y\} \Leftrightarrow \{i \mapsto j\}$	a[x]=y; a[i]=j;	Array action
X⇒Y	if(!X    Y){	Logical Implication
X⇔Y	if((!X    Y) && (!Y    X)){	Logical Equivalence
¬x <y< td=""><td>if(!(x<y)){< td=""><td>Logical not</td></y)){<></td></y<>	if(!(x <y)){< td=""><td>Logical not</td></y)){<>	Logical not
$x \in \mathbb{N}$	unsigned long int x	Natural numbers
$x \in \mathbb{Z}$	signed long int x	Integer numbers
A	/* No Action */	Quantifier
∃	/* No Action */	Quantifier



#### Tool Installation

Web: http://eb2all.loria.fr/(Until Rodin 2.8)

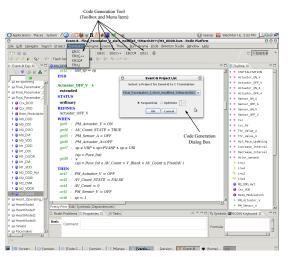


Figure: Screen shots of the Code Generation Tool: EB2ALL



## An Example - Code Generation

```
EVENT Actuator_ON_VWHENActuator_ON_V.Guard1 : PM\_Actuator\_V = OFFActuator_ON_V.Guard2 : (sp = Pace\_Int)\langle sp < Pace\_Int \wedge\langle AV\_Count > V\_Blank \wedge\langle AV\_Count > FixedAV \rangleActuator_ON_V.Guard3 : sp \geq VRP \wedge sp \geq PVARPTHENActuator_ON_V.Action1 : PM\_Actuator\_V := ONActuator_ON_V.Action2 : last\_sp := spEND
```

```
...

BOOL Actuator\_ON\_V(void)

{

    /* Guards No. 1*/
    if(PM_Actuator_V == OFF) {
        /* Guards No. 2*/
    if((sp == Pace_Int) || ((sp < Pace_Int) \&\&
        (AV_Count > V_Blank) && (AV_Count >= FixedAV))) {
        /* Guards No. 3*/
        if((sp >= VRP) && (sp >= PVARP) && (sp >= URI)) {
        /* Actions */
        PM_Actuator_V = ON;
        last_sp = sp;
        return TRUE;
    }}
    return FALSE;
}
```



## Demo



Neeraj Kumar Singh Using Event-B for Critical **Device Software** Systems 2 Springer





