# **REAL TIME SYSTEMS**

3. APEX: ARINC 653 standard

## Introduction

Specify an interface between an OS of an avionics resource and application software: APEX interface (APplication / EXecutive)

Start of work: 1991

Publication of the first ARINC 653 standard: Summer 1996

## **Definitions**

### Application:

- Software realization of an avionics function
- consisting of one or more partitions

**Example: Automatic Pilot** 

#### Partition:

- Entity of execution of an application
- Runs on a single processor
- Deterministic allocation of resources to partitions
- Spatial (memory) and temporal (CPU) segregation unit
- Consisting of one or more processes

Example: FM1, FM2, FE1-COM, FE1-MON, FE2-COM, FE2-MON, FG1-COM, FG1-MON, FG2-COM, FG2-MON

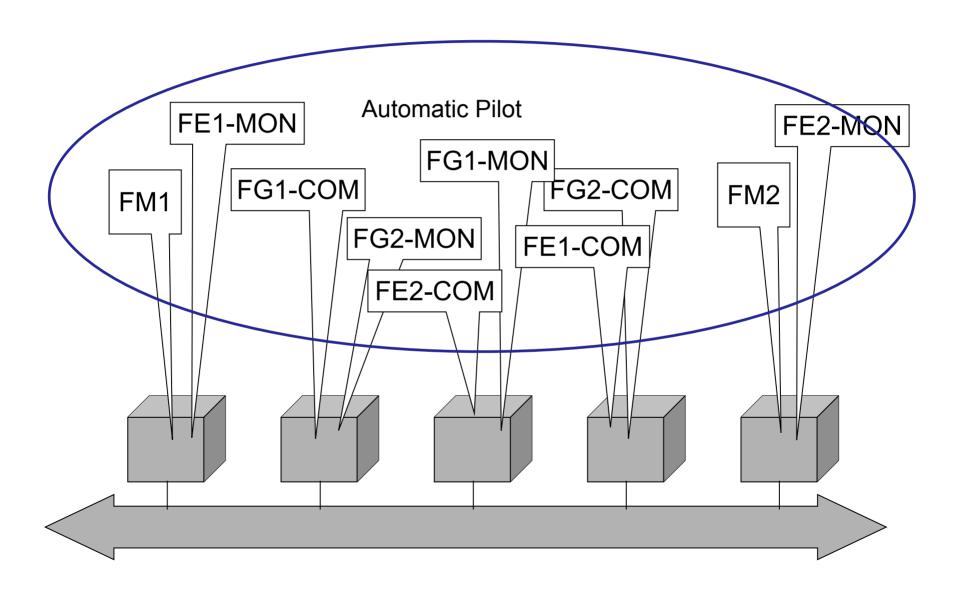
## **Definition**

### **Process:**

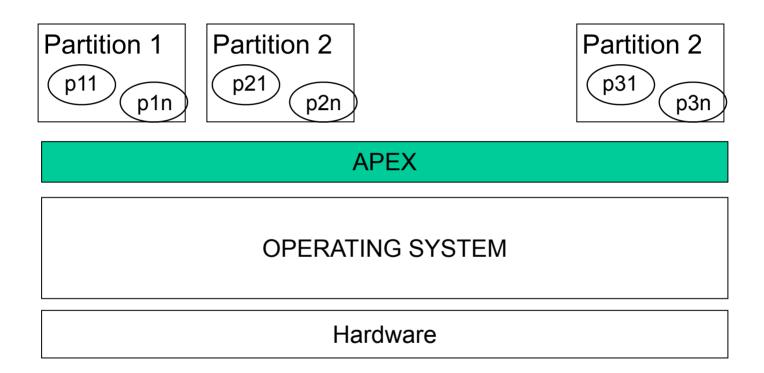
- A program unit that runs in the partition environment.
- Concurrent execution of processes to perform the avionics function

Example: a control law, a control logic...

## **Architecture**



## **Architecture**



### Operating system:

Schedules the partitions of the module Schedules the processes in the partition Ensures spatial and temporal segregation (partitioning)

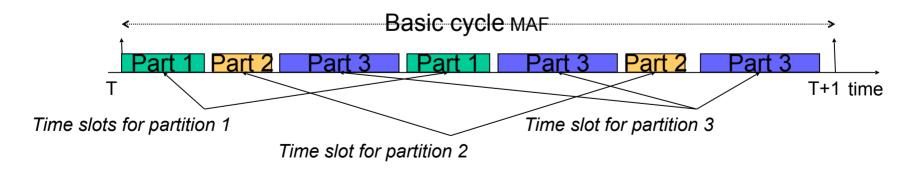
# Resource Sharing

### Spatial sharing:

Predetermined memory area for each partition

### Time sharing (CPU):

- A partition has no priority
- Deterministic and cyclic allocation of the processor to partitions:
  - => The OS repeats a basic cycle (MAjor time Frame: MAF) of fixed duration.
  - => Assign one or more time slots in the MAF to each partition.



=> Resource allocation is defined by configuration and cannot be changed dynamically.

# Partitions management

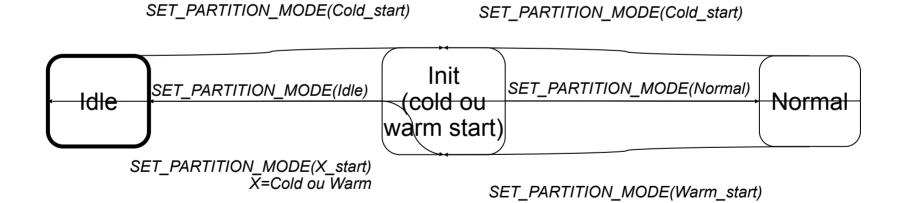
- The OS starts the partition scheduler at the end of module initialization.
- During its time slots, a partition has the following states
  - Idle :
    - No process is running.
  - Cold\_start ou Warm\_start : initialization of the partition
    - Cold start: abstraction of the previous context of the partition
    - Warm start: recovering the previous context of the partition
  - Normal : starting process scheduling
- Outside its time slots a partition is suspended.
- The partition creates all its objects (ports, process, semaphores...) during its initialization phase.
- The partition starts process scheduling at the end of its initialization phase.

# Partitions management

### APEX service for partitions

SET\_PARTITION\_MODE: changes the mode (Idle, Cold or Warm start, Normal)

GET\_PARTITION\_STATUS: provides the status of the partition



SET\_PARTITION\_MODE(Idle)

- 2 process types
  - Periodic: execution in regular intervals
  - Apériodic : execution on event occurrence
- No segregation between processes within the same application
- Processes on one partition are not visible to other processes on other partitions.
- Each process has a priority
- Process scheduling is based on a preemption algorithm by priority level and current process state
- Only one process per partition is "running" at any given time.

### APEX services for processes

#### CREATE\_PROCESS:

Creation of a link between the name of the process and its reservation on partition initialization

Putting the process in a Sleeping state and returning an identifier (id)

#### START:

Initialization of the process

Change from Sleeping to Ready state

#### STOP:

stopping a process by removing access rights to the processor

Change from Ready or Waiting to Sleeping state

#### STOP SELF:

Auto-stop of the running process

Change from Running to Sleeping state

#### SUSPEND:

Suspension of the process until resumed by another process

Change from Running to Waiting state

#### SUSPEND SELF:

self-suspension of the process

Change from Running to Waiting state

### **APEX** service for processes

#### **RESUME:**

Restarting a suspended process

If the process is already waiting for a resource, it remains in the *Waiting* state, otherwise it switches to the *Ready* state

#### **SET PRIORITY:**

Changing the priority of a process

#### LOCK\_PREEMPTION:

Incrementing the preemption counter of the partition and blocking the preemption mechanism

#### **UNLOCK PREEMPTION:**

Decrementing the preemption counter of the partition and unlocking the preemption mechanism if this counter is null

#### GET\_PROCESS\_ID:

returns the identifier of a process by giving its name

#### GET\_PROCESS\_STATUS:

returns status information about a process

### Time management

Time is expressed in real time (absolute).

The OS ensures time segregation between partitions

APEX services for time management :

TIMED\_WAIT:

Suspension of a process for a given minimum amount of time Change from Running to Waiting state

A time equal to zero performs a "Round Robin" on processes with the same priority.

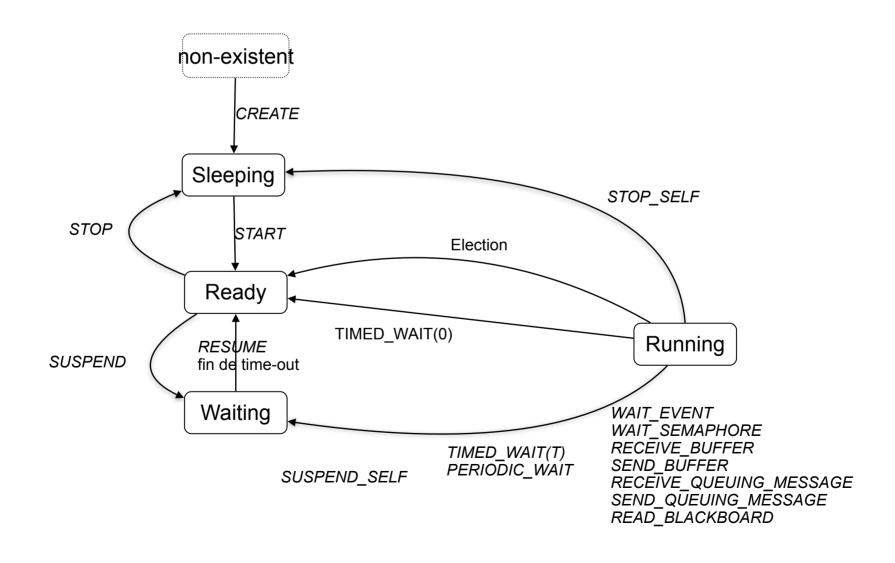
PERIODIC\_WAIT

Suspension of a periodic process until the next period reallocation of the process time budget at the beginning of the period

GET\_TIME

returns time (absolute) and local time to the module

### State graph of the processes



Processes on the same partition can communicate and synchronize through

- communication via Buffer or Blackboard type mailboxes
- use of global variables to the partition
- synchronization by semaphore and event mechanisms

A mailbox allows the exchange of messages between several processes without indicating the names of the sending and receiving processes.

A memory area is reserved for the initialization of the partition for these communication / synchronization objects.

### The "Buffer" type mailboxes:

communication by message that can carry different data a message occurrence does not overwrite previous occurrences messages are stored in FIFO message queues the message tails are limited in length a message queue can be full processes trying to read from an empty message queue are put on hold processes seeking to transmit in a full message queue are put on hold

#### Corresponding APEX services:

```
CREATE_BUFFER

SEND_BUFFER

RECEIVE_BUFFER

GET_BUFFER_ID

GET_BUFFER_STATUS

a timeout can be specified
```

### "Blackboard" type mailboxes:

data write communication

- a data entry overwrites the previous data
- a written data remains displayed until the next data is written or a delete request is made.
- all processes waiting on an empty "Blackboard" are woken up when writing data (change from *Waiting* to *Ready* status)

#### Corresponding APEX services:

```
CREATE_BLACKBOARD
DISPLAY_ BLACKBOARD
READ_ BLACKBOARD
CLEAR_ BLACKBOARD ______ a timeout can be specified
GET_ BLACKBOARD _ID
GET_ BLACKBOARD STATUS
```

### Smaphores:

```
=> semaphores with counter
```

resource and signals the semaphore when it releases the resource

if counter > 0: the value represents the number of processes that can access the resource.

if counter = 0: resource unavailable

processes waiting on a semaphore are ordered either by FIFO or by priority

#### Corresponding APEX services:

```
CREATE_SEMAPHORE
```

SIGNAL\_SEMAPHORE

GET SEMAPHORE ID

GET\_ SEMAPHORE \_STATUS

#### **Events**

```
=> event at level (UP and DOWN)
```

allows to synchronize processes on the partition

when an event is set to UP, all processes pending on that event are set to *Ready* and a new process scheduling takes place

when a process is waiting for an event set to DOWN, it goes to the *Waiting* state with possibly waiting for a time out.

```
Services APEX correspondants:
```

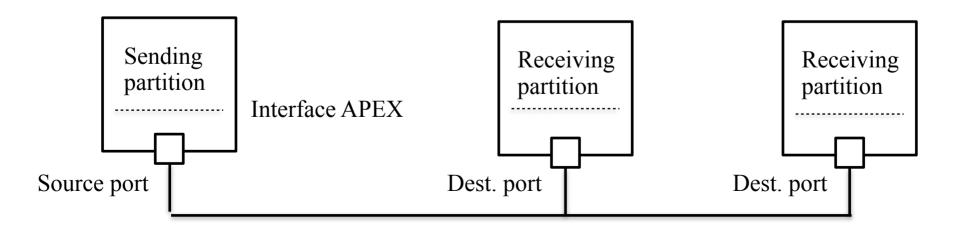
Communication between partitions (of the same module or not) is only done by message exchange.

A transmitter is connected to one or more receivers by a logical channel called "Channel"

The transmitting or receiving entities correspond to the partitions

### A partition accesses a channel via a "Port

=> APEX service for sending and receiving via ports



- The transmitting partition does not know the name and location of the receiving partitions.
- The physical connection between the sending and receiving ports is established by static network configuration.
- The communication protocol between transmitter and receiver must be the same.
  - (same message format, same acknowledgement policy...)
- The periodic call of an APEX transmission service generates a periodic message, and conversely, an aperiodic call generates an aperiodic message.
  - => the periodicity is not an attribute of the port

### Transfert protocol:

- Sampling mode:
  - the message always conveys the same updated data
  - each occurrence of a message overwrites the previous one
  - messages are of fixed size
  - the messages are not segmentable by the OS
- Queuing mode:
  - the message can convey different data
  - occurrences are saved in a FIFO
  - messages can be of variable length
  - messages are segmentable by the OS

#### Ports attributes

- Identifier (id): value returned when creating the port
- Name
- Transfer protocol: sampling or queuing
- Direction: input or outpout
- Maximum message size
- Port size
  - identical to the message size for sampling ports
  - FIFO depth for queuing ports
- Refresh period for sampling ports
  - allows to monitor the freshness of the last received message
- Queuing mode waiting policy:
  - FIFO or priority

### Sampling APEX services

```
CREATE_ SAMPLING_PORT
```

WRITE\_SAMPLING\_MESSAGE

READ\_ SAMPLING\_MESSAGE

reads the last message received in the port and indicates if the age of the message is consistent with the Refresh period attribute of the port

```
GET_ SAMPLING _ID
```

GET\_ SAMPLING \_STATUS

## Service APEX queuing

```
CREATE_ QUEUING_PORT
```

SEND\_ QUEUING \_MESSAGE

RECEIVE\_QUEUING \_MESSAGE ————— a timeout can be specified

GET\_ QUEUING \_ID

**GET\_ QUEUING \_STATUS** 

## A conclusion

#### Summary:

- A two-level OS:
  - partition + static cyclic scheduling
  - process + dynamic scheduling by priority
- Intuitive objective: emulate classic federated architectures where each function has its own resources and manages its own processes as it sees fit
- Partition level required for certification requirements
  - => guarantee that a faulty function will never be able to disrupt another one
  - => notion of strict partitioning

Question: Wouldn't it have been simpler to remove the partition level and multiply the physical resources?

- => Non-Integrated Modular Avionics
- => Only resource sharing = the network