All PDA's are defined as the 6-tuple: $(Q, \Sigma, \Gamma, \delta, q_0, F)$, where:

Q: State Set

 Σ : Automaton Alphabet

Γ: Stack Alphabet

 $\delta \colon Q \times \Sigma_{\epsilon} \times \Gamma_{\epsilon} \to P(Q \times \Gamma_{\epsilon})$: Transition Function

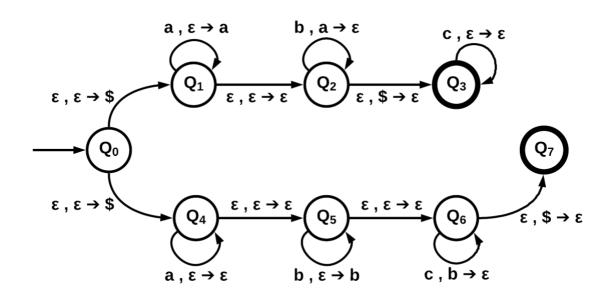
q₀: Starting StateF: Accepting State Set

Exercise 1:

Draw the PDA for the language:

$$L_1 = \{a^i b^j c^k | i = j \lor j = k, \forall i, j, k \ge 0\}$$

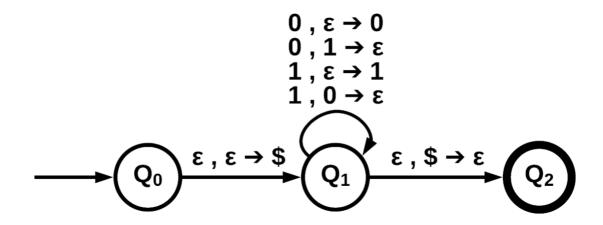
$$\begin{split} Q &= \{Q_0, \dots, Q_7\} \\ \Sigma &= \{a, b, c\} \\ \Gamma &= \{a, b\} \\ Q_0 \text{ is the starting state } \\ F &= \{Q_3, Q_7\} \end{split}$$



Exercise 2:

Draw the PDA for the language: $L_2 = \{w | w \text{ has as many } 1s \text{ as } 0s\}$

$$\begin{aligned} Q &= \{Q_0,Q_1,Q_2\} \\ \Sigma &= \Gamma = \{0,1\} \\ Q_0 \text{ is the starting state } \\ F &= \{Q_2\} \end{aligned}$$

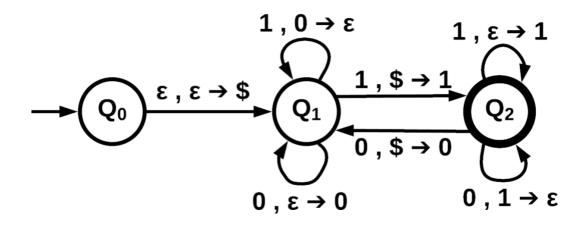


Exercise 3:

Draw the PDA for the language:

 $L_3 = \{w | w \text{ has more 1s than 0s}\}$

$$\begin{aligned} Q &= \{Q_0,Q_1,Q_2\} \\ \Sigma &= \Gamma = \{0,1\} \\ Q_0 \text{ is the starting state } \\ F &= \{Q_2\} \end{aligned}$$

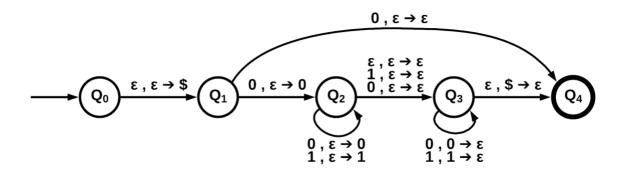


Exercise 4:

Draw the PDA for the language:

 $L_4 = \{w | w \text{ is a palindrome and starts with } 0\}$

$$\begin{split} Q &= \{Q_0,Q_1,Q_2,Q_3,Q_4\} \\ \Sigma &= \Gamma = \{0,1\} \\ Q_0 \text{ is the starting state } \\ F &= \{Q_4\} \end{split}$$



Exercise 5:

Draw the PDA for the language:

 $L_5 = \{w | w \text{ has odd length and the middle symbol is a 0}\}$

$$\begin{split} Q &= \{Q_0,Q_1,Q_2,Q_3\} \\ \Sigma &= \{0,1\} \\ \Gamma &= \{1\} \\ Q_0 \text{ is the starting state } \\ F &= \{Q_3\} \end{split}$$

