Matrix Forensics

A brief guide to matrix math
and its efficient implementation

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github.com/r-barnes/MatrixForensics

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1 Introduction

Goals: TODO

Contributing: Please contribute on Github at `https://github.com/r-barnes/MatrixForensics` either by opening an issue or making a pull request. If you are not comfortable with this, please send your contribution to `rijard.barnes@gmail.com`.

Contributors: Richard Barnes

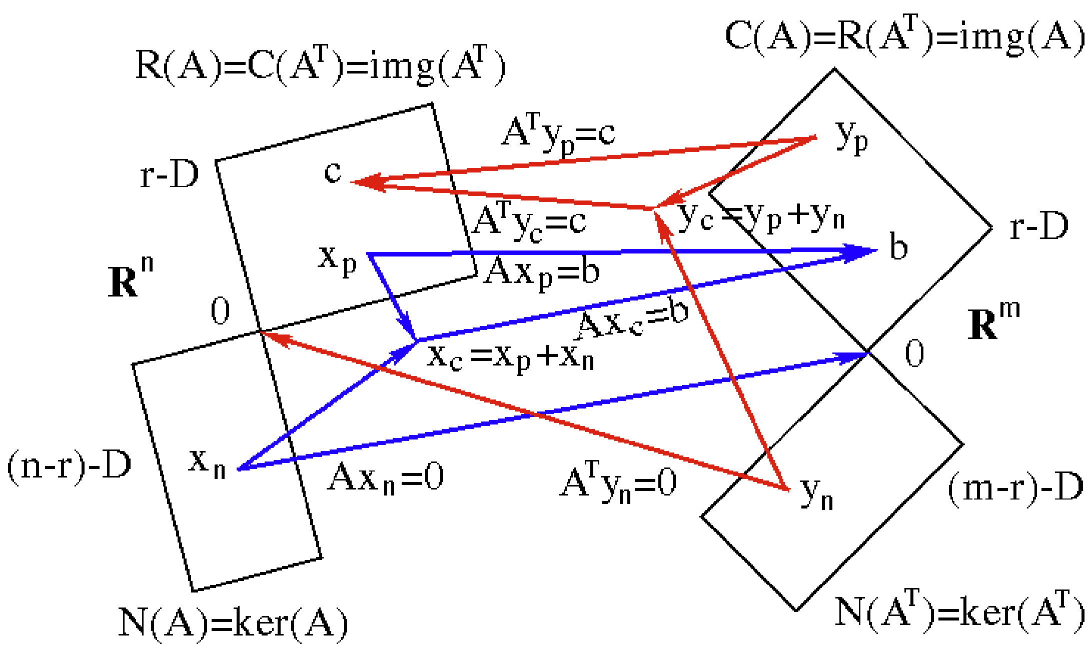
Funding: TODO

2 Nomenclature

$\mathbf{A}$	Matrix.
$\mathbf{a}$	(Column) vector.
a	Scalar.
$\mathbf{A}_{ij}$	Matrix indexed. Returns ith row and jth column.
$\mathbf{A} \circ \mathbf{B}$	Hadamard (element-wise) product of matrices A and B.
$\mathcal{N}(\mathbf{A})$	Nullspace of the matrix $\mathbf{A}$.
$\mathcal{R}(\mathbf{A})$	Range of the matrix $\mathbf{A}$.
$\det(\mathbf{A})$	Determinant of the matrix $\mathbf{A}$.
$\mathrm{eig}(\mathbf{A})$	Eigenvalues of the matrix $\mathbf{A}$.
$\mathbf{A}^H$	Conjugate transpose of the matrix $\mathbf{A}$.
$\mathbf{A}^T$	Transpose of the matrix $\mathbf{A}$.
$\mathbf{A}^+$	Pseudoinverse of the matrix $\mathbf{A}$.
$\mathbf{x} \in \mathbb{R}^n$	The entries of the n-vector $\mathbf{x}$ are all real numbers.
$\mathbf{A} \in \mathbb{R}^{m,n}$	The entries of the matrix $\mathbf{A}$ with m rows and n columns are all real numbers.
$\mathbf{A} \in \mathbb{S}^n$	The matrix $\mathbf{A}$ is symmetric and has n rows and n columns.
$\mathbf{I}_n$	Identity matrix with n rows and n columns.
$\{0\}$	The empty set

3 Basics

3.1 Fundamental Theorem of Linear Algebra



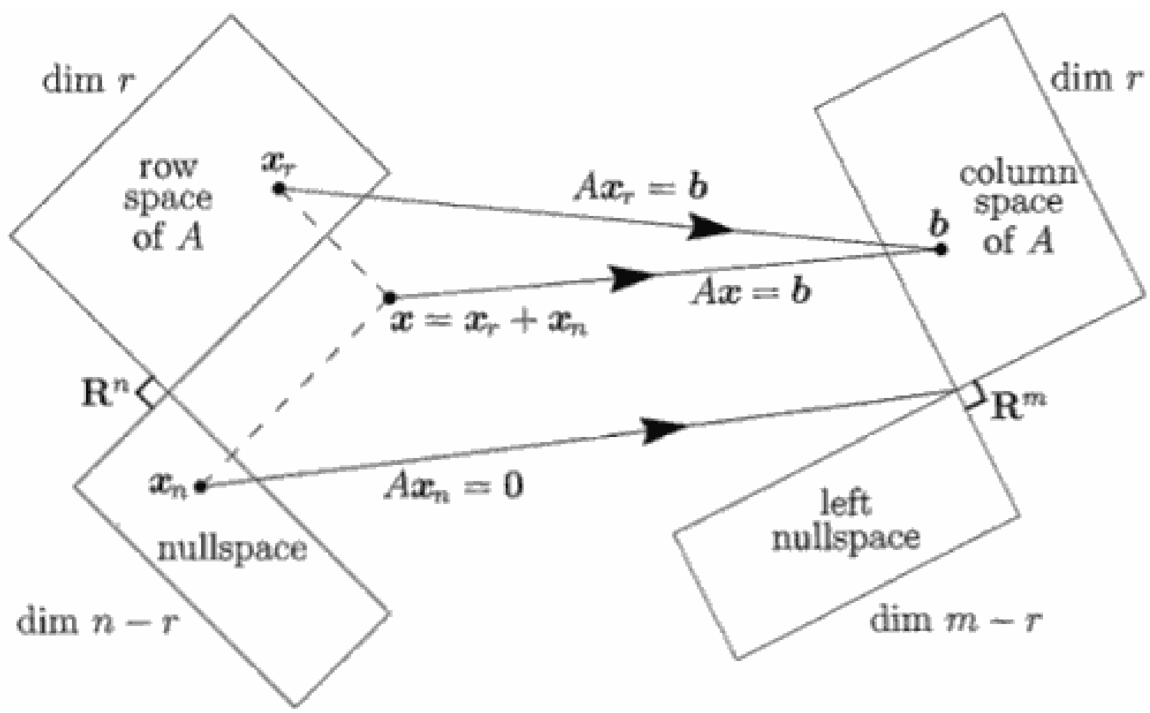
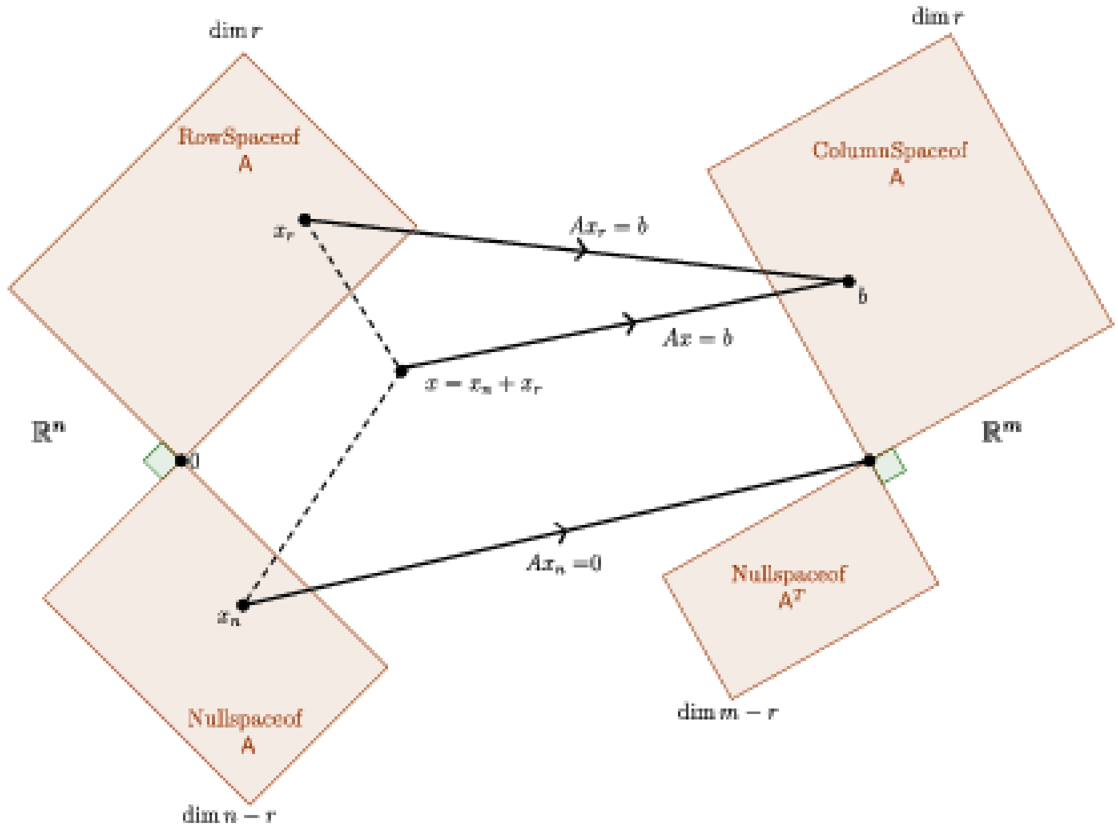
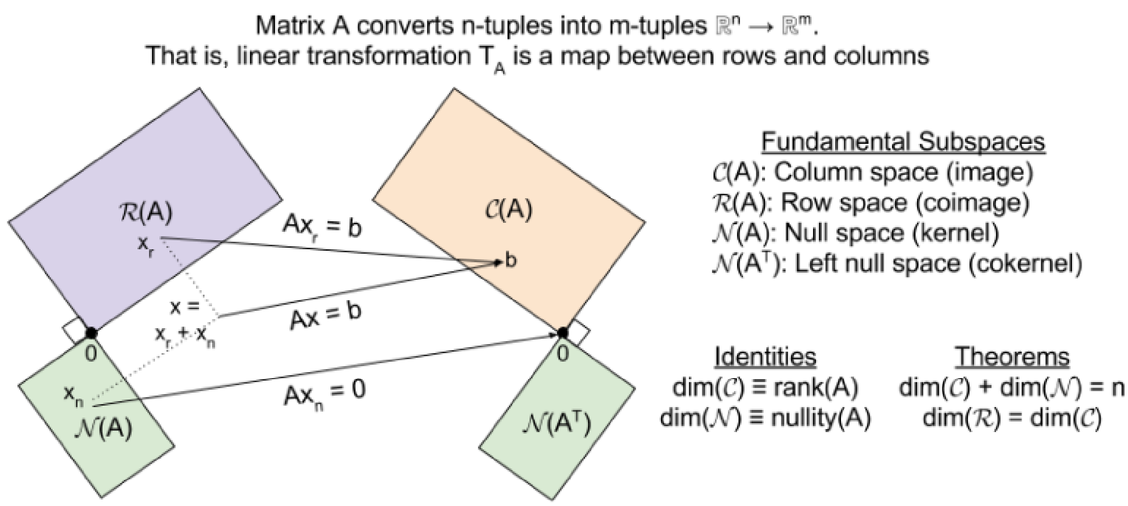


Figure 3.4 The true action $Ax = A(x_{row} + x_{null})$ of any m by n matrix.


dim r
RowSpaceof A
dim r
ColumnSpaceof A
x_r
$Ax_r = b$
b
$Ax = b$
$x = x_n + x_r$
$\mathbb{R}^n$
0
$\mathbb{R}^m$
$Ax_n = 0$
Nullspaceof A^T
x_n
dim $m - r$
Nullspaceof A
dim $n - r$


Matrix A converts n-tuples into m-tuples $\mathbb{R}^n \to \mathbb{R}^m$.
That is, linear transformation T_A is a map between rows and columns
$\mathcal{R}$(A)
x_r
$Ax_r = b$
$\mathcal{C}$(A)
b
x = $x_r + x_n$
Ax = b
0
0
x_n
$Ax_n = 0$
$\mathcal{N}$(A)
$\mathcal{N}(A^T)$
Fundamental Subspaces
$\mathcal{C}$(A): Column space (image)
$\mathcal{R}$(A): Row space (coimage)
$\mathcal{N}$(A): Null space (kernel)
$\mathcal{N}(A^T)$: Left null space (cokernel)
Identities
dim($\mathcal{C}$) ≡ rank(A)
dim($\mathcal{N}$) ≡ nullity(A)
Theorems
dim($\mathcal{C}$) + dim($\mathcal{N}$) = n
dim($\mathcal{R}$) = dim($\mathcal{C}$)

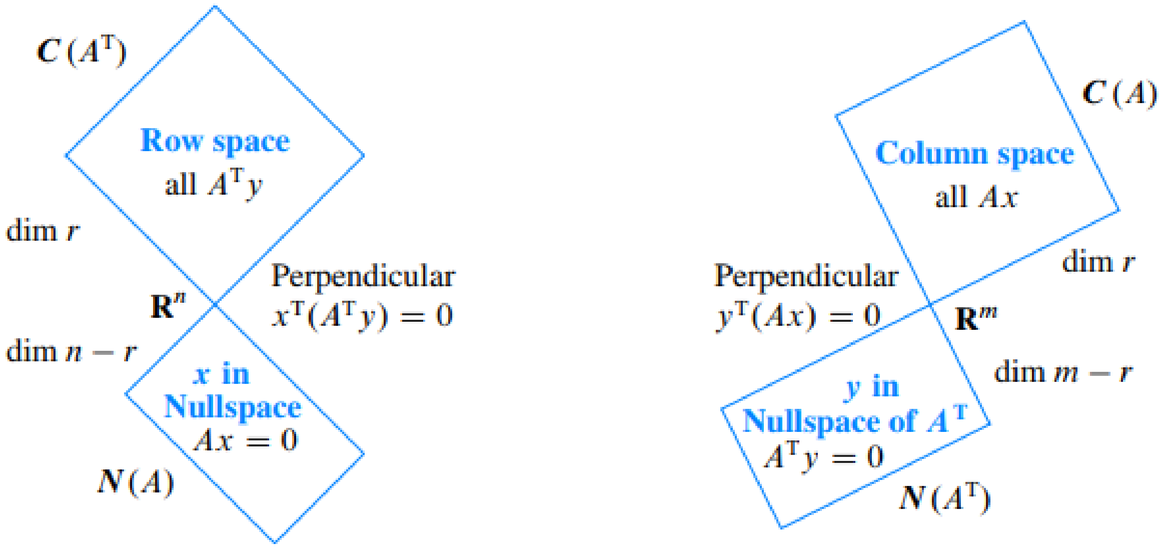


Figure 1: Dimensions and orthogonality for any m by n matrix A of rank r.

3.2 Matrix Properties

$$\mathbf{A}(\mathbf{B}+\mathbf{C}) = \mathbf{AB}+\mathbf{AC} \qquad \text{(left distributivity)} \tag{1}$$

$$(\mathbf{B}+\mathbf{C})\mathbf{A} = \mathbf{BA}+\mathbf{CA} \qquad \text{(right distributivity)} \tag{2}$$

$$\mathbf{AB} \neq \mathbf{BA} \qquad \text{(in general)} \tag{3}$$

$$(\mathbf{AB})\mathbf{C} = \mathbf{A}(\mathbf{BC}) \qquad \text{(associativity)} \tag{4}$$

3.3 Rank

If $\mathbf{A} \in \mathbb{R}^{m,n}$ and $\mathbf{B} \in \mathbb{R}^{n,r}$, then

[1] $$\operatorname{rank}(\mathbf{A}) + \operatorname{rank}(\mathbf{B}) - n \le \operatorname{rank}(\mathbf{AB}) \le \min(\operatorname{rank}(\mathbf{A}), \operatorname{rank}(\mathbf{B})) \qquad \text{Sylvester's Inequality} \tag{5}$$

If $\mathbf{AB}$, $\mathbf{ABC}$, $\mathbf{BC}$ are defined, then

[1] $$\operatorname{rank}(\mathbf{AB}) + \operatorname{rank}(\mathbf{BC}) \le \operatorname{rank}(\mathbf{B}) + \operatorname{rank}(\mathbf{ABC}) \qquad \text{Frobenius's inequality} \tag{6}$$

If $\dim(\mathbf{A}) = \dim(\mathbf{B})$, then

$$\operatorname{rank}(\mathbf{A}+\mathbf{B}) \le \operatorname{rank}(\mathbf{A}) + \operatorname{rank}(\mathbf{B}) \qquad \text{Subadditivity} \tag{7}$$

If $\mathbf{A}_1, \mathbf{A}_2, \ldots, \mathbf{A}_l$ have $n_1, n_2, \ldots, n_l$ columns, so that $\mathbf{A}_1\mathbf{A}_2\ldots\mathbf{A}_l$ is well-defined, then

[1] $$\operatorname{rank}(\mathbf{A}_1\mathbf{A}_2\ldots\mathbf{A}_l) \ge \sum_{i=1}^{l-1} \operatorname{rank}(\mathbf{A}_i\mathbf{A}_{i+1}) - \sum_{i=2}^{l-1} \operatorname{rank}(\mathbf{A}_i) \ge \sum_{i=1}^{l} \operatorname{rank}(\mathbf{A}_i) - \sum_{i=1}^{l-1} n_i \tag{8}$$

3.4 Identities

$$\left(\sum_{i=1}^{n} \mathbf{z}_i\right)^2 = \mathbf{z}^T \begin{bmatrix} 1 & \cdots & 1 \\ \vdots & \ddots & \vdots \\ 1 & \cdots & 1 \end{bmatrix} \mathbf{z} \tag{9}$$

3.5 Matrix Multiplication

$$(\mathbf{AB})_{kl} = \sum_m \mathbf{A}_{km}\mathbf{B}_{ml} \quad \mathbf{A} \in \mathbb{R}^{k,m}, \mathbf{B} \in \mathbb{R}^{m,l} \tag{10}$$

3.6 Time Complexities

Operation	**Input**	**Output**	**Algorithm**	**Time**
Matmult	$A, B \in n \times n$	$n \times n$	Schoolbook	$O(n^3)$
			Strassen [2]	$O(n^{2.807})$
			Best	$O(n^\omega)$
Matmult	$A \in n \times m, B \in m \times p$	$n \times p$	Schoolbook	$O(nmp)$
Inversion	$A \in n \times n$	$n \times n$	Gauss–Jordan elimination	$O(n^3)$
			Strassen [2]	$O(n^{2.807})$
			Best	$O(n^\omega)$
SVD	$A \in m \times n$	$m \times m, m \times n, n \times n$ $m \times r, r \times r, n \times r$		$O(mn^2)$ $(m \geq n)$
Determinant	$A \in n \times n$	Scalar	Laplace expansion	$O(n!)$
			Division-free [3]	$O(n!)$
			LU decomposition	$O(n^3)$
			Integer preserving [4]	$O(n^3)$
Back substitution	A triangular	n solutions	Back substitution	$O(n^2)$

A comment on ω

The lower bound on matmult time complexity is $O(n^\omega)$, where ω is an unknown constant bounded by $2 \leq \omega \leq 2.373$. Algorithms achieving lower values of ω tend to be less efficient in practice for all but the largest matrices. Of the algorithm with times of less than $O(n^3)$, only the Strassen algorithm has seen serious attempts at optimized implementation. Most matmult implementations use highly optimized variants of the standard $O(n^3)$ algorithm. At this point, memory and bus speeds dominate the performance of implementations, so simple Big-O notation cannot be used to reliably compare matmult performances.

Name	**Year**	ω
Standard	-	3
Strassen [2]	1969	2.807
Pan [5]	1978	2.796
Bini et al. [6]	1979	2.78
Schönhage [7]	1981	2.548
Schönhage [7]	1981	2.522
Romani [8]	1982	2.517
Coppersmith and Winograd [9]	1982	2.496
Strassen [10]	1986	2.479
Coppersmith and Winograd [11]	1990	2.376
Williams [12]	2012	2.37294
Le Gall [13]	2014	2.3728639
Williams [12]	2012	2.3727

4 Derivatives

4.1 Useful Rules for Derivatives

For general $\mathbf{A}$ and $\mathbf{X}$ (no special structure):

$$\partial\mathbf{A} = 0 \quad \text{where } \mathbf{A} \text{ is a constant} \tag{11}$$
$$\partial(c\mathbf{X}) = c\partial\mathbf{X} \tag{12}$$
$$\partial(\mathbf{X}+\mathbf{Y}) = \partial\mathbf{X} + \partial\mathbf{Y} \tag{13}$$
$$\partial(\mathrm{tr}(\mathbf{X})) = \mathrm{tr}(\partial(\mathbf{X})) \tag{14}$$
$$\partial(\mathbf{X}\mathbf{Y}) = (\partial\mathbf{X})\mathbf{Y} + \mathbf{X}(\partial\mathbf{Y}) \tag{15}$$
$$\partial(\mathbf{X}\circ\mathbf{Y}) = (\partial\mathbf{X})\circ\mathbf{Y} + \mathbf{X}\circ(\partial\mathbf{Y}) \tag{16}$$
$$\partial(\mathbf{X}^{-1}) = -\mathbf{X}^{-1}(\partial\mathbf{X})\mathbf{X}^{-1} \tag{17}$$
$$\partial(\det(\mathbf{X})) = \mathrm{tr}(\mathrm{adj}(\mathbf{X})\partial\mathbf{X}) \tag{18}$$
$$\partial(\det(\mathbf{X})) = \det(\mathbf{X})\,\mathrm{tr}(\mathbf{X}^{-1}\partial\mathbf{X}) \tag{19}$$
$$\partial(\ln(\det(\mathbf{X}))) = \mathrm{tr}(\mathbf{X}^{-1}\partial\mathbf{X}) \tag{20}$$
$$\partial(\mathbf{X}^T) = (\partial\mathbf{X})^T \tag{21}$$
$$\partial(\mathbf{X}^H) = (\partial\mathbf{X})^H \tag{22}$$

4.2 Derivatives of Matrices and Vectors

4.2.1 First-Order

In the following, $\mathbf{J}$ is the Single-Entry Matrix (section 5.11).

$$\frac{\partial\mathbf{x}^T\mathbf{a}}{\partial\mathbf{x}} = \frac{\partial\mathbf{a}^T\mathbf{x}}{\partial\mathbf{x}} = \mathbf{a} \tag{23}$$
$$\frac{\partial\mathbf{a}^T\mathbf{X}\mathbf{b}}{\partial\mathbf{X}} = \mathbf{a}\mathbf{b}^T \tag{24}$$
$$\frac{\partial\mathbf{a}^T\mathbf{X}^T\mathbf{b}}{\partial\mathbf{X}} = \mathbf{b}\mathbf{a}^T \tag{25}$$
$$\frac{\partial\mathbf{a}^T\mathbf{X}\mathbf{a}}{\partial\mathbf{X}} = \frac{\partial\mathbf{a}^T\mathbf{X}^T\mathbf{a}}{\partial\mathbf{X}} = \mathbf{a}\mathbf{a}^T \tag{26}$$
$$\frac{\partial\mathbf{X}}{\partial\mathbf{X}_{ij}} = \mathbf{J}^{ij} \tag{27}$$

4.3 Derivatives of vector norms

$$\frac{\partial}{\partial \mathbf{x}}\|\mathbf{x}-\mathbf{a}\|_2 = \frac{\mathbf{x}-\mathbf{a}}{\|\mathbf{x}-\mathbf{a}\|_2} \tag{28}$$

$$\frac{\partial}{\partial \mathbf{x}}\frac{\mathbf{x}-\mathbf{a}}{\|\mathbf{x}-\mathbf{a}\|_2} = \frac{\mathbf{I}}{\|\mathbf{x}-\mathbf{a}\|_2} - \frac{(\mathbf{x}-\mathbf{a})(\mathbf{x}-\mathbf{a})^T}{\|\mathbf{x}-\mathbf{a}\|_2^3} \tag{29}$$

$$\frac{\partial\|\mathbf{x}\|_2^2}{\partial \mathbf{x}} = \frac{\partial\left\|\mathbf{x}^T\mathbf{x}\right\|_2}{\partial \mathbf{x}} = 2\mathbf{x} \tag{30}$$

4.4 Scalar by Vector

Qualifier	Expression	Numerator layout	Denominator layout
	$\frac{\partial a}{\partial x}$	$\mathbf{0}^T$	$\mathbf{0}$
	$\frac{\partial a u(\mathbf{x})}{\partial \mathbf{x}}$	$a\frac{\partial u}{\partial \mathbf{x}}$	Same
	$\frac{\partial u(\mathbf{x})+v(\mathbf{x})}{\partial \mathbf{x}}$	$\frac{\partial u}{\partial \mathbf{x}}+\frac{\partial v}{\partial \mathbf{x}}$	Same
	$\frac{\partial u(\mathbf{x})v(\mathbf{x})}{\partial \mathbf{x}}$	$u\frac{\partial v}{\partial \mathbf{x}}+v\frac{\partial u}{\partial \mathbf{x}}$	Same
	$\frac{\partial g(u(\mathbf{x}))}{\partial \mathbf{x}}$	$\frac{\partial g(u)}{\partial u}\frac{\partial u}{\partial \mathbf{x}}$	Same
	$\frac{\partial f(g(u(\mathbf{x})))}{\partial \mathbf{x}}$	$\frac{\partial f(g)}{\partial g}\frac{\partial g(u)}{\partial u}\frac{\partial u}{\partial \mathbf{x}}$	Same
	$\frac{\partial \mathbf{u}(\mathbf{x})^T\mathbf{v}(\mathbf{x})}{\partial \mathbf{x}}$	$\mathbf{u}^T\frac{\partial \mathbf{v}}{\partial \mathbf{x}}+\mathbf{v}^T\frac{\partial \mathbf{u}}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{u}}{\partial \mathbf{x}}\mathbf{v}+\frac{\partial \mathbf{v}}{\partial \mathbf{x}}\mathbf{u}$
	$\frac{\partial \mathbf{u}(\mathbf{x})^T\mathbf{A}\mathbf{v}(\mathbf{x})}{\partial \mathbf{x}}$	$\mathbf{u}^T\mathbf{A}\frac{\partial \mathbf{v}}{\partial \mathbf{x}}+\mathbf{v}^T\mathbf{A}^T\frac{\partial \mathbf{u}}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{u}}{\partial \mathbf{x}}\mathbf{A}\mathbf{v}+\frac{\partial \mathbf{v}}{\partial \mathbf{x}}\mathbf{A}^T\mathbf{u}$
	$\frac{\partial^2 f}{\partial \mathbf{x}\partial \mathbf{x}^T}$		$\mathbf{H}$, the Hessian matrix
	$\frac{\partial \mathbf{a}\cdot\mathbf{x}}{\partial \mathbf{x}}=\frac{\partial \mathbf{x}\cdot\mathbf{a}}{\partial \mathbf{x}}$	$\mathbf{a}^T$	$\mathbf{a}$
	$\frac{\partial \mathbf{b}^T\mathbf{A}\mathbf{x}}{\partial \mathbf{x}}$	$\mathbf{b}^T\mathbf{A}$	$\mathbf{A}^T\mathbf{b}$
	$\frac{\partial \mathbf{x}^T\mathbf{A}\mathbf{x}}{\partial \mathbf{x}}$	$\mathbf{x}^T(\mathbf{A}+\mathbf{A}^T)$	$(\mathbf{A}+\mathbf{A}^T)\mathbf{x}$
$\mathbf{A}$ symmetric	$\frac{\partial \mathbf{x}^T\mathbf{A}\mathbf{x}}{\partial \mathbf{x}}$	$2\mathbf{x}^T\mathbf{A}$	$2\mathbf{A}\mathbf{x}$
	$\frac{\partial \mathbf{x}^T\mathbf{A}\mathbf{x}}{\partial \mathbf{x}}$	$\mathbf{A}+\mathbf{A}^T$	Same
$\mathbf{A}$ symmetric	$\frac{\partial \mathbf{x}^T\mathbf{A}\mathbf{x}}{\partial \mathbf{x}}$	$\mathbf{A}$	Same
	$\frac{\partial \mathbf{x}^T\mathbf{x}}{\partial \mathbf{x}}$	$2\mathbf{x}^T$	$2\mathbf{x}$
	$\frac{\partial \mathbf{a}^T\mathbf{u}(\mathbf{x})}{\partial \mathbf{x}}$	$\mathbf{a}^T\frac{\partial \mathbf{u}}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{u}}{\partial \mathbf{x}}\mathbf{a}$
	$\frac{\partial \mathbf{a}^T\mathbf{x}\mathbf{x}^T\mathbf{b}}{\partial \mathbf{x}}$	$\mathbf{x}^T(\mathbf{a}\mathbf{b}^T+\mathbf{b}\mathbf{a}^T)$	$(\mathbf{a}\mathbf{b}^T+\mathbf{b}\mathbf{a}^T)\mathbf{x}$
	$\frac{\partial (\mathbf{A}\mathbf{x}+\mathbf{b})^T\mathbf{C}(\mathbf{D}\mathbf{x}+\mathbf{e})}{\partial \mathbf{x}}$	$(\mathbf{D}\mathbf{x}+\mathbf{e})^T\mathbf{C}^T\mathbf{A}+(\mathbf{A}\mathbf{x}+\mathbf{b})^T\mathbf{C}\mathbf{D}$	$\mathbf{D}^T\mathbf{C}^T(\mathbf{A}\mathbf{x}+\mathbf{b})+\mathbf{A}^T\mathbf{C}(\mathbf{D}\mathbf{x}+\mathbf{e})$
	$\frac{\partial\|\mathbf{x}-\mathbf{a}\|}{\partial \mathbf{x}}$	$\frac{(\mathbf{x}-\mathbf{a})^T}{\|\mathbf{x}-\mathbf{a}\|}$	$\frac{\mathbf{x}-\mathbf{a}}{\|\mathbf{x}-\mathbf{a}\|}$

4.5 Vector by Vector

Qualifier	Expression	Numerator layout	Denominator layout
	$\frac{\partial \mathbf{a}}{\partial \mathbf{x}}$	$\mathbf{0}$	Same
	$\frac{\partial \mathbf{x}}{\partial \mathbf{x}}$	$\mathbf{I}$	Same
	$\frac{\partial \mathbf{A}\mathbf{x}}{\partial \mathbf{x}}$	$\mathbf{A}$	$\mathbf{A}^T$
	$\frac{\partial \mathbf{x}^T\mathbf{A}}{\partial \mathbf{x}}$	$\mathbf{A}^T$	$\mathbf{A}$
	$\frac{\partial a\mathbf{u}(\mathbf{x})}{\partial \mathbf{x}}$	$a\frac{\partial \mathbf{u}}{\partial \mathbf{x}}$	Same
	$\frac{\partial a(\mathbf{x})\mathbf{u}(\mathbf{x})}{\partial \mathbf{x}}$	$a\frac{\partial \mathbf{u}}{\partial \mathbf{x}} + \mathbf{u}\frac{\partial a}{\partial \mathbf{x}}$	$a\frac{\partial \mathbf{u}}{\partial \mathbf{x}} + \frac{\partial a}{\partial \mathbf{x}}\mathbf{u}^T$
	$\frac{\partial \mathbf{A}\mathbf{u}(\mathbf{x})}{\partial \mathbf{x}}$	$\mathbf{A}\frac{\partial \mathbf{u}}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{u}}{\partial \mathbf{x}}\mathbf{A}^T$
	$\frac{\partial(\mathbf{u}(\mathbf{x})+\mathbf{v}(\mathbf{x}))}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{u}}{\partial \mathbf{x}} + \frac{\partial \mathbf{v}}{\partial \mathbf{x}}$	Same
	$\frac{\partial \mathbf{g}(\mathbf{u}(\mathbf{x}))}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{g}(\mathbf{u})}{\partial \mathbf{u}}\frac{\partial \mathbf{u}}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{u}}{\partial \mathbf{x}}\frac{\partial \mathbf{g}(\mathbf{u})}{\partial \mathbf{u}}$
	$\frac{\partial \mathbf{f}(\mathbf{g}(\mathbf{u}(\mathbf{x})))}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{f}(\mathbf{g})}{\partial \mathbf{g}(\mathbf{u})}\frac{\partial \mathbf{g}(\mathbf{u})}{\partial \mathbf{u}}\frac{\partial \mathbf{u}}{\partial \mathbf{x}}$	$\frac{\partial \mathbf{u}}{\partial \mathbf{x}}\frac{\partial \mathbf{g}(\mathbf{u})}{\partial \mathbf{u}}\frac{\partial \mathbf{f}(\mathbf{g})}{\partial \mathbf{g}}$

4.6 Matrix by Scalar

Qualifier	Expression	Numerator layout
	$\frac{\partial a\mathbf{U}(x)}{\partial x}$	$a\frac{\partial \mathbf{U}}{\partial x}$
	$\frac{\partial \mathbf{A}\mathbf{U}(x)\mathbf{B}}{\partial x}$	$\mathbf{A}\frac{\partial \mathbf{U}}{\partial x}\mathbf{B}$
	$\frac{\partial(\mathbf{U}(x)+\mathbf{V}(x))}{\partial x}$	$\frac{\partial \mathbf{U}}{\partial x} + \frac{\partial \mathbf{V}}{\partial x}$
	$\frac{\partial(\mathbf{U}(x)\mathbf{V}(x))}{\partial x}$	$\mathbf{U}\frac{\partial \mathbf{V}}{\partial x} + \frac{\partial \mathbf{U}}{\partial x}\mathbf{V}$
	$\frac{\partial(\mathbf{U}(x)\otimes\mathbf{V}(x))}{\partial x}$	$\mathbf{U}\otimes\frac{\partial \mathbf{V}}{\partial x} + \frac{\partial \mathbf{U}}{\partial x}\otimes\mathbf{V}$
	$\frac{\partial(\mathbf{U}(x)\circ\mathbf{V}(x))}{\partial x}$	$\mathbf{U}\circ\frac{\partial \mathbf{V}}{\partial x} + \frac{\partial \mathbf{U}}{\partial x}\circ\mathbf{V}$
	$\frac{\partial \mathbf{U}^{-1}(x)}{\partial x}$	$-\mathbf{U}^{-1}\frac{\partial \mathbf{U}}{\partial x}\mathbf{U}^{-1}$
	$\frac{\partial^2 \mathbf{U}^{-1}}{\partial x \partial y}$	$\mathbf{U}^{-1}\left(\frac{\partial \mathbf{U}}{\partial x}\mathbf{U}^{-1}\frac{\partial \mathbf{U}}{\partial y} - \frac{\partial^2 \mathbf{U}}{\partial x\partial y} + \frac{\partial \mathbf{U}}{\partial y}\mathbf{U}^{-1}\frac{\partial \mathbf{U}}{\partial x}\right)\mathbf{U}^{-1}$
	$\frac{\partial e^{x\mathbf{A}}}{\partial x}$	$\mathbf{A}e^{x\mathbf{A}} = e^{x\mathbf{A}}\mathbf{A}$

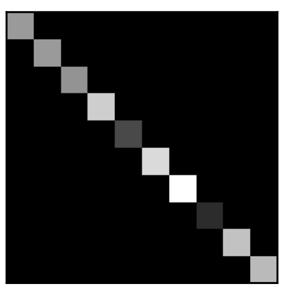
5 | Matrix Rogue Gallery

5.1 Non-Singular vs. Singular Matrices

For $\mathbf{A} \in \mathbb{R}^{n,n}$ (initially drawn from [14, p. 574]):

Non-Singular	**Singular**
$\mathbf{A}$ is invertible	$\mathbf{A}$ is not invertible
The columns are independent	The columns are dependent
The rows are independent	The rows are dependent
$\det(\mathbf{A}) \neq 0$	$\det(\mathbf{A}) = 0$
$\mathbf{Ax} = 0$ has one solution: $\mathbf{x} = 0$	$\mathbf{Ax} = 0$ has infinitely many solutions
$\mathbf{Ax} = \mathbf{b}$ has one solution: $\mathbf{x} = \mathbf{A}^{-1}\mathbf{b}$	$\mathbf{Ax} = \mathbf{b}$ has either no or infinitely many solutions
$\mathbf{A}$ has n nonzero pivots	$\mathbf{A}$ has $r < n$ pivots
$\mathbf{A}$ has full rank $r = n$	$\mathbf{A}$ has rank $r < n$
The reduced row echelon form is $\mathbf{R} = \mathbf{I}$	$\mathbf{R}$ has at least one zero row
The column space is all of $\mathbb{R}^n$	The column space has dimension $r < n$
The row space is all of $\mathbb{R}^n$	The row space has dimension $r < n$
All eigenvalues are nonzero	Zero is an eigenvalue of $\mathbf{A}$
$\mathbf{A}^T\mathbf{A}$ is symmetric positive definite	$\mathbf{A}^T\mathbf{A}$ is only semidefinite
$\mathbf{A}$ has n positive singular values	$\mathbf{A}$ has $r < n$ singular values

5.2 Diagonal Matrix



$$A = \operatorname{diag}(a_1, \ldots, a_n) = \begin{bmatrix} a_1 & & \\ & \ddots & \\ & & a_n \end{bmatrix} \tag{31}$$

Square matrix. Entries above diagonal are equal to entries below diagonal.

Number of “free entries”: $\frac{n(n+1)}{2}$.

Special Properties

$$\mathrm{eig}(A) = a_1, \dots, a_n \tag{32}$$

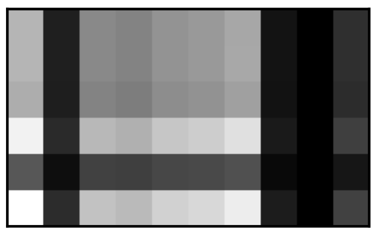
$$\det(A) = \prod_i a_i \tag{33}$$

$$A^{-1} = \begin{bmatrix} \frac{1}{a_1} & & \\ & \ddots & \\ & & \frac{1}{a_n} \end{bmatrix} \tag{34}$$

$$\mathbf{x}^T\mathbf{A}\mathbf{x} = \sum_i a_i \mathbf{x}_i^2 \tag{35}$$

$$\tag{36}$$

5.3 Dyads



$\mathbf{A} \in \mathbb{R}^{m,n}$ is a dyad if it can be written as

$$\mathbf{A} = \mathbf{u}\mathbf{v}^T \quad \mathbf{u} \in \mathbb{R}^m, \mathbf{v} \in \mathbb{R}^n \tag{37}$$

Special Properties

- The columns of $\mathbf{A}$ are copies of $\mathbf{u}$ scaled by the values of $\mathbf{v}$.
- The rows of $\mathbf{A}$ are copies of $\mathbf{u}^T$ scaled by the values of $\mathbf{v}$.
- If $\mathbf{A}$ is a dyad, it acts on a vector $\mathbf{x}$ as $\mathbf{A}\mathbf{x} = (\mathbf{u}\mathbf{v}^T)\mathbf{x} = (\mathbf{v}^T\mathbf{u})\mathbf{x}$.
- $\mathbf{A}\mathbf{x} = c\mathbf{u}$ ($\mathbf{A}$ scales $\mathbf{x}$ and points it along $\mathbf{u}$).
- $\mathbf{A}_{ij} = \mathbf{u}_i\mathbf{v}_j$.
- If $\mathbf{u}, \mathbf{v} \neq 0$, then $\mathrm{rank}(\mathbf{A}) = 1$.
- If $m = n$, $\mathbf{A}$ has one eigenvalue $\lambda = \mathbf{v}^T\mathbf{u}$ and eigenvector $\mathbf{u}$.
- A dyad can always be written in a normalized form $c\tilde{\mathbf{u}}\tilde{\mathbf{v}}^T$.

5.4 Hermitian Matrix

$\mathbf{H} \in \mathbb{C}^{m,n}$ is Hermitian iff

$$\mathbf{H} = \mathbf{H}^H \tag{38}$$

where $\mathbf{H}^H$ is the conjugate transpose of $\mathbf{H}$.

For $\mathbf{H} \in \mathbb{R}^{m,n}$, Hermitian and symmetric matrices are equivalent.

Special Properties

$$\mathbf{H}_{ii} \in \mathbb{R} \tag{39}$$
$$\mathbf{H}\mathbf{H}^H = \mathbf{H}^H\mathbf{H} \tag{40}$$
$$\mathbf{x}^H\mathbf{H}\mathbf{x} \in \mathbb{R} \quad \forall \mathbf{x} \in \mathbb{C} \tag{41}$$
$$\mathbf{H}_1 + \mathbf{H}_2 = \text{Hermitian} \tag{42}$$
$$\mathbf{H}^{-1} = \text{Hermitian} \tag{43}$$
$$\mathbf{A} + \mathbf{A}^H = \text{Hermitian} \tag{44}$$
$$\mathbf{A} - \mathbf{A}^H = \text{Skew-Hermitian} \tag{45}$$
$$\mathbf{AB} = \text{Hermitian iff } \mathbf{AB} = \mathbf{BA} \tag{46}$$
$$\det(\mathbf{H}) \in \mathbb{R} \tag{47}$$
$$\text{eig}(\mathbf{H}) \in \mathbb{R} \tag{48}$$

5.5 Idempotent Matrix

A matrix $\mathbf{A}$ is idempotent iff

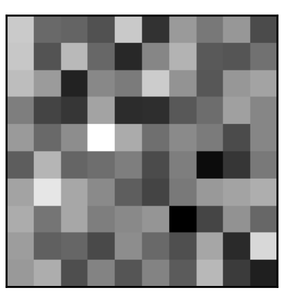
$$\mathbf{AA} = \mathbf{A} \tag{49}$$

Special Properties

$$\mathbf{A}^n = A \quad \forall n \tag{50}$$
$$\mathbf{I} - \mathbf{A} \quad \text{is idempotent} \tag{51}$$
$$\mathbf{A}^H \quad \text{is idempotent} \tag{52}$$
$$\mathbf{I} - \mathbf{A}^H \quad \text{is idempotent} \tag{53}$$
$$\text{rank}(\mathbf{A}) = \text{tr}(\mathbf{A}) \tag{54}$$
$$\mathbf{A}(I - \mathbf{A}) = 0 \tag{55}$$
$$\mathbf{A}^+ = \mathbf{A} \tag{56}$$
$$f(s\mathbf{I} + t\mathbf{A}) = (\mathbf{I} - \mathbf{A})f(s) + \mathbf{A}f(s+t) \tag{57}$$
$$\mathbf{AB} = \mathbf{BA} \implies \mathbf{AB} \text{ is idempotent} \tag{58}$$
$$\text{eig}(\mathbf{A})_i \in \{0, 1\} \tag{59}$$
$$\mathbf{A} \text{ is always diagonalizable} \tag{60}$$

$\mathbf{A} - \mathbf{I}$ may not be idempotent.

5.6 Orthogonal Matrix



(Not much visible structure)

$$U = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix} \tag{61}$$

A matrix $\mathbf{U}$ is orthogonal iff:

$$\mathbf{U}^T\mathbf{U} = \mathbf{U}\mathbf{U}^T = I \tag{62}$$

Square matrix. The columns form an orthonormal basis of $\mathbb{R}^n$.

Special Properties

- The eigenvalues of $\mathbf{U}$ are placed on the unit circle.
- The eigenvectors of $\mathbf{U}$ are unitary (have length one).
- $\mathbf{U}^{-1}$ is orthogonal.
- The product of two orthogonal matrices is itself orthogonal.

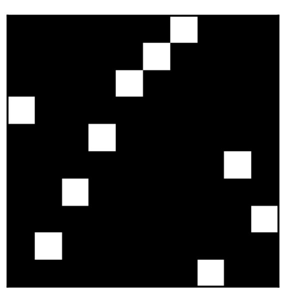
$$\begin{aligned} \mathbf{U}^T &= \mathbf{U}^{-1} && (63) \\ \mathbf{U}^{-T} &= \mathbf{U} && (64) \\ \mathbf{U}^T\mathbf{U} &= \mathbf{I} && (65) \\ \mathbf{U}\mathbf{U}^T &= \mathbf{I} && (66) \\ \det(\mathbf{U}) &= \pm 1 && (67) \end{aligned}$$

Orthogonal matrices preserve the lengths and angles of the vectors they operator on. The converse is true: any matrix which preserves lengths and angles is orthogonal.

$$\|\mathbf{U}\mathbf{x}\|_2^2 = (\mathbf{U}\mathbf{x})^T(\mathbf{U}\mathbf{x}) = \mathbf{x}^T\mathbf{U}^T\mathbf{U}\mathbf{x} = \mathbf{x}^T\mathbf{x} = \|\mathbf{x}\|_2^2 \quad \forall \mathbf{x} \tag{68}$$

$$\|\mathbf{U}\mathbf{A}\mathbf{V}\|_F = \|\mathbf{A}\|_F \quad \forall \mathbf{A}, \mathbf{U}, \mathbf{V} \text{ with } \mathbf{U}, \mathbf{V}\text{orthogonal} \tag{69}$$

5.7 Permutation Matrix



TODO

5.8 Positive Definite

$\mathbf{A} \in \mathbb{S}^n$ is positive definite (denoted $\mathbf{A} \succ 0$) if any of the following are true:

- $\mathbf{x}^T\mathbf{A}\mathbf{x} > 0, \forall \mathbf{x} \in \mathbb{R}^n$.
- $\mathrm{eig}(\mathbf{A}) > 0$

Special Properties

- If $\mathbf{A}$ is PD and invertible, $\mathbf{A}^{-1}$ is also PD.
- If $\mathbf{A}$ is PD and $c \in \mathbb{R}$ then $c\mathbf{A}$ is PD.
- The diagonal entries $\mathbf{A}_{ii}$ are real and non-negative, so $\mathrm{tr}(\mathbf{A}) \geq 0$.
- $\det(\mathbf{A}) > 0$
- For $\mathbf{A} \in \mathbb{R}^{m,n}$, $\mathbf{A}^T\mathbf{A} \succ 0 \iff \mathbf{A}$ is full-column rank ($\mathrm{rank}(\mathbf{A}) = n$)
- For $\mathbf{A} \in \mathbb{R}^{m,n}$, $\mathbf{A}\mathbf{A}^T \succ 0 \iff \mathbf{A}$ is full-row rank ($\mathrm{rank}(\mathbf{A}) = m$)
- $\mathbf{P} \succ 0$ defines a full-dimensional, bounded ellipsoid centered at the origin and defined by the set $\mathcal{E} = \{\mathbf{x} \in \mathbb{R}^n : x^T\mathbf{P}^{-1}x \leq 1\}$. The eigenvalues λ_i and eigenvectors u_i of $\mathbf{P}$ define the orientation and shape of the ellipsoid. u_i are the semi-axes while the lengths of the semi-axes are given by $\sqrt{\lambda_i}$. Using the Cholesky decomposition, $\mathbf{P}^{-1} = \mathbf{A}^T\mathbf{A}$, an equivalent definition of the ellipsoid is $\mathcal{E} = \{\mathbf{x} \in \mathbb{R}^n : \|\mathbf{A}\mathbf{x}\|_2 \leq 1\}$.

5.9 Positive Semi-Definite

$\mathbf{A}$ is positive semi-definite (denoted $\mathbf{A} \succeq 0$) if any of the following are true:

- $\mathbf{x}^T\mathbf{A}\mathbf{x} \geq 0, \forall \mathbf{x} \in \mathbb{R}^n$.
- $\mathrm{eig}(\mathbf{A}) \geq 0$

Special Properties

- For $\mathbf{A} \in \mathbb{R}^{m,n}$, $\mathbf{A}^T\mathbf{A} \succeq 0$
- For $\mathbf{A} \in \mathbb{R}^{m,n}$, $\mathbf{A}\mathbf{A}^T \succeq 0$
- The positive semi-definite matrices $\mathbb{S}^n_+$ form a convex cone. For any two PSD matrices $\mathbf{A}, \mathbf{B} \in \mathbb{S}^n_+$ and some $\alpha \in [0, 1]$:

$$\mathbf{x}^T(\alpha\mathbf{A} + (1-\alpha)\mathbf{B})\mathbf{x} = \alpha\mathbf{x}^T\mathbf{A}\mathbf{x} + (1-\alpha)\mathbf{x}^T\mathbf{B}\mathbf{x} \geq 0 \quad \forall \mathbf{x} \tag{70}$$

$$\alpha\mathbf{A} + (1-\alpha)\mathbf{B} \in \mathbb{S}^n_+ \tag{71}$$

- For $\mathbf{A} \in \mathbb{S}^n_+$ and $\alpha \geq 0$, $\alpha\mathbf{A} \succeq 0$, so $\mathbb{S}^n_+$ is a cone.
- $\mathbf{A} \succeq 0$ has a unique PSD matrix $\mathbf{S}^{1/2}$ such that $\mathbf{S}^{1/2}\mathbf{S}^{1/2} = \mathbf{A}$

5.9.1 Loewner order

If $\mathbf{A} - \mathbf{B} \succeq 0$, then we say $\mathbf{A} \succeq \mathbf{B}$. A sufficient condition for this is that $\lambda_n(\mathbf{A}) \geq \lambda_1(\mathbf{B})$.

5.10 Projection Matrix

A square matrix $\mathbf{P}$ is a projection matrix that projects onto a vector space $\mathcal{S}$ iff

$$\mathbf{P} \text{ is idempotent} \tag{72}$$
$$\mathbf{Px} \in \mathcal{S} \quad \forall \mathbf{x} \tag{73}$$
$$\mathbf{Pz} = \mathbf{z} \quad \forall \mathbf{z} \in \mathcal{S} \tag{74}$$

5.11 Single-Entry Matrix

$$\mathbf{J}^{2,3} = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} \tag{75}$$

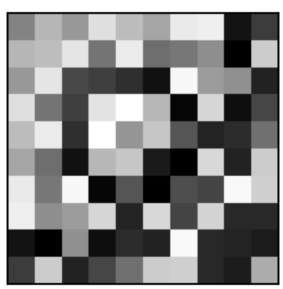
The single-entry matrix $\mathbf{J}^{iJ} \in \mathbb{R}^{n,n}$ is defined as the matrix which is zero everywhere except for the entry (i, j), which is 1.

5.12 Singular Matrix

A square matrix that is not invertible.

$\mathbf{A} \in \mathbb{R}^{n,n}$ is singular iff $\det \mathbf{A} = 0$ iff $\mathcal{N}(A) \neq \{0\}$.

5.13 Symmetric Matrix



$\mathbf{A} \in \mathbb{S}^n$ is a symmetric matrix if $\mathbf{A} = \mathbf{A}^T$ (entries above diagonal are equal to entries below diagonal).

$$\begin{bmatrix} a & b & c & d & e & f \\ b & g & l & m & o & p \\ c & l & h & n & q & r \\ d & m & n & i & s & t \\ e & o & q & s & j & u \\ f & p & r & t & u & k \end{bmatrix} \tag{76}$$

Special Properties

$$\mathbf{A} = \mathbf{A}^T \tag{77}$$

$$\text{eig}(A) \in \mathbb{R}^n \tag{78}$$

$$\text{Number of “free entries”} = \frac{n(n+1)}{2} \tag{79}$$

If $\mathbf{A}$ is real, it can be decomposed into $\mathbf{A} = \mathbf{Q}^T\mathbf{D}\mathbf{Q}$ where $\mathbf{Q}$ is a real orthogonal matrix (the columns of which are eigenvectors of $\mathbf{A}$) and $\mathbf{D}$ is real and diagonal containing the eigenvalues of $\mathbf{A}$.

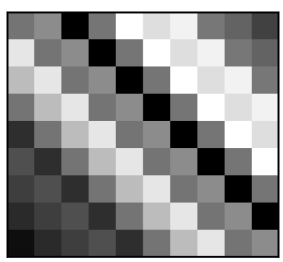
For a real, symmetric matrix with non-negative eignevalues, the eigenvalues and singular values coincide.

5.14 Skew-Hermitian

A matrix $\mathbf{H} \in \mathbb{C}^{m,n}$ is Skew-Hermitian iff

$$\mathbf{H} = -\mathbf{H}^H \tag{80}$$

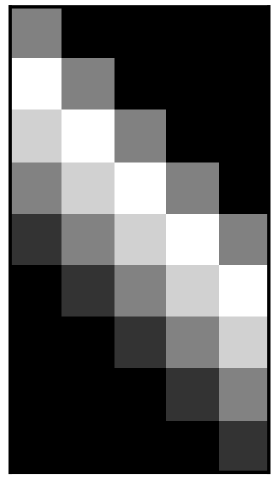
5.15 Toeplitz Matrix, General Form



Constant values on descending diagonals.

$$\begin{bmatrix} a_0 & a_{-1} & a_{-2} & \cdots & \cdots & a_{-(n-1)} \\ a_1 & a_0 & a_{-1} & \ddots & & \vdots \\ a_2 & a_1 & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & a_{-1} & a_{-2} \\ \vdots & & \ddots & a_1 & a_0 & a_{-1} \\ a_{n-1} & \cdots & \cdots & a_2 & a_1 & a_0 \end{bmatrix} \tag{81}$$

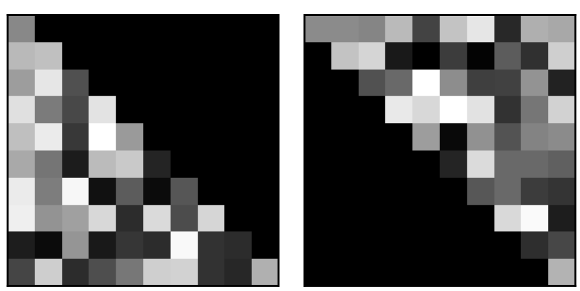
5.16 Toeplitz Matrix, Discrete Convolution



Constant values on main and subdiagonals.

$$
\begin{bmatrix}
h_m & 0 & 0 & \dots & 0 & 0 \\
\vdots & h_m & 0 & \dots & 0 & 0 \\
h_1 & \vdots & h_m & \dots & 0 & 0 \\
0 & h_1 & \ddots & \ddots & 0 & 0 \\
0 & 0 & h_1 & \ddots & h_m & 0 \\
0 & 0 & 0 & \ddots & \vdots & h_m \\
0 & 0 & 0 & \dots & h_1 & \vdots \\
0 & 0 & 0 & \dots & 0 & h_1
\end{bmatrix}
\tag{82}
$$

5.17 Triangular Matrix



$$
\begin{bmatrix}
a & b & c & d & e & f \\
 & g & h & i & j & k \\
 & & l & m & n & o \\
 & & & p & q & r \\
 & & & & s & t \\
 & & & & & u
\end{bmatrix}
\begin{bmatrix}
a & & & & & \\
b & g & & & & \\
c & h & l & & & \\
d & i & m & p & & \\
e & j & n & q & s & \\
f & k & o & r & t & u
\end{bmatrix}
\tag{83}
$$

Square matrices in which all elements either above or below the main diagonal are zero. An upper (left) and a lower (right) triangular matrix are shown above.

For an upper triangular matrix $A_{ij} = 0$ whenever $i > j$; for a lower triangular matrix $A_{ij} = 0$ whenever $i < j$.

Special Properties

$$\operatorname{eig}(A) = \operatorname{diag}(A) \tag{84}$$

$$\det(A) = \prod_i \operatorname{diag}(A)_i \tag{85}$$

The product of two upper (lower) triangular matrices is still upper (lower) triangular.

The inverse of a nonsingular upper (lower) triangular matrix is still upper (lower) triangular.

5.18 Vandermonde Matrix

$$V = \begin{bmatrix} 1 & \alpha_1 & \alpha_1^2 & \dots & \alpha_1^{n-1} \\ 1 & \alpha_2 & \alpha_2^2 & \dots & \alpha_2^{n-1} \\ 1 & \alpha_3 & \alpha_3^2 & \dots & \alpha_3^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \alpha_m & \alpha_m^2 & \dots & \alpha_m^{n-1} \end{bmatrix} \tag{86}$$

Alternatively,

$$V_{i,j} = \alpha_i^{j-1} \tag{87}$$

Uses

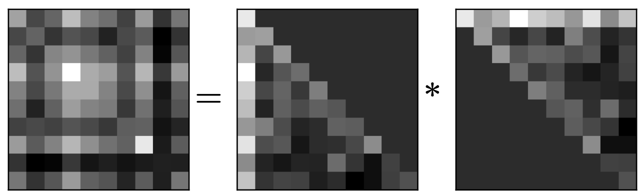
Polynomial interpolation of data.

Special Properties

- $\det(V) = \prod_{1 \le i < j \le n} (x_j - x_i)$

6 Matrix Decompositions

6.1 LLT/UTU: Cholesky Decomposition

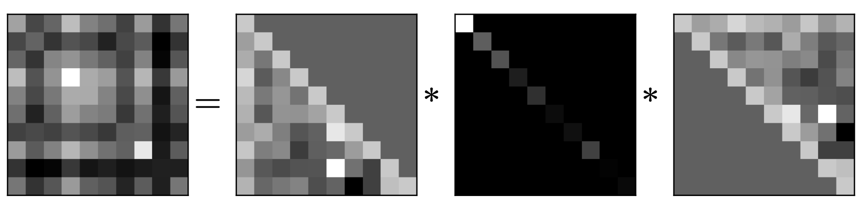


If $\mathbf{A}$ is symmetric, positive definite, square, then

$$\mathbf{A} = \mathbf{U}^T\mathbf{U} = \mathbf{L}\mathbf{L}^T \tag{88}$$

where $\mathbf{U}$ is a unique upper triangular matrix and $\mathbf{L}$ is a unique lower-triangular matrix.

6.2 LDLT



If $\mathbf{A}$ is a non-singular symmetric definite square matrix, then

$$\mathbf{A} = \mathbf{L}\mathbf{D}\mathbf{L}^T = \mathbf{L}^T\mathbf{D}\mathbf{L} \tag{89}$$

where $\mathbf{L}$ is a unit lower triangular matrix and $\mathbf{D}$ is a diagonal matrix. If $\mathbf{A} \succ 0$, then $\mathbf{D}_{ii} > 0$.

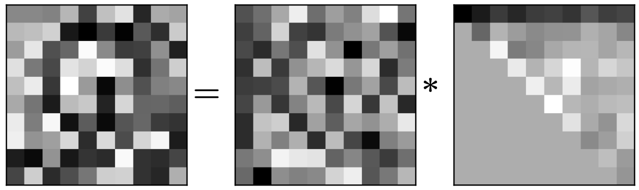
6.3 PCA: Principle Components Analysis

Find normalized directions in data space such that the variance of the projections of the centered data points is maximal. For centered data $\tilde{\mathbf{X}}$, the mean-square variation of data along a vector $\mathbf{x}$ is $\mathbf{x}^T\tilde{\mathbf{X}}\tilde{\mathbf{X}}^T\mathbf{x}$.

$$\max_{\mathbf{x}\in\mathbb{R}^n,\ \|\mathbf{x}\|_2=1} \mathbf{x}^T\tilde{\mathbf{X}}\tilde{\mathbf{X}}^T\mathbf{x} \tag{90}$$

Taking an SVD of $\tilde{\mathbf{X}}\tilde{\mathbf{X}}^T$ gives $H = \mathbf{U}_r\mathbf{D}^2\mathbf{U}^T$, which is maximized by taking $\mathbf{x} = \mathbf{u}_1$. By repeatedly removing the first principal components and recalculating, all the principal axes can be found.

6.4 QR: Orthogonal-triangular

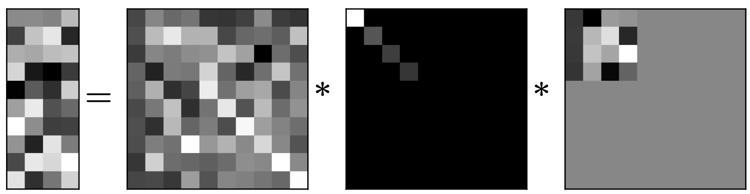


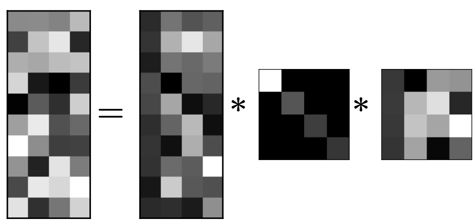
For $\mathbf{A} \in \mathbb{R}^{n,n}$, $\mathbf{A} = \mathbf{QR}$ where $\mathbf{Q}$ is orthogonal and $\mathbf{R}$ is an upper triangular matrix. If $\mathbf{A}$ is non-singular, then $\mathbf{Q}$ and $\mathbf{R}$ are uniquely defined if $\mathrm{diag}(\mathbf{R})$ are imposed to be positive.

Algorithms

Gram-Schmidt.

6.5 SVD: Singular Value Decomposition





Any matrix $\mathbf{A} \in \mathbb{R}^{m,n}$ can be written as

$$\mathbf{A} = \mathbf{UDV}^T = \sum_{i=1}^{r} \sigma_i u_i v_i^T \tag{91}$$

where

$$U = \text{eigenvectors of } \mathbf{AA}^T \qquad \mathbb{R}^{m,m} \tag{92}$$
$$D = \mathrm{diag}(\sigma_i) = \sqrt{\mathrm{diag}(\mathrm{eig}(\mathbf{AA}^T))} \qquad \mathbb{R}^{n,m} \tag{93}$$
$$V = \text{eigenvectors of } \mathbf{A}^T\mathbf{A} \qquad \mathbb{R}^{n,n} \tag{94}$$

Let σ_i be the non-zero singular values for $i = 1, \ldots, r$ where r is the rank of $\mathbf{A}$; $\sigma_1 \geq \ldots \geq \sigma_r$.

We also have that

$$\mathbf{Av}_i = \sigma_i \mathbf{u}_i \tag{95}$$
$$\mathbf{A}^T\mathbf{u}_i = \sigma_i \mathbf{v}_i \tag{96}$$
$$\mathbf{U}^T\mathbf{U} = \mathbf{I} \tag{97}$$
$$\mathbf{V}^T\mathbf{V} = \mathbf{I} \tag{98}$$

$\mathbf{D}$ can be written in an expanded form:

$$\tilde{\mathbf{D}} = \begin{bmatrix} \mathbf{D} & 0_{r,n-r} \\ 0_{m-r,r} & 0_{m-r,n-r} \end{bmatrix} \tag{99}$$

The final $n - r$ columns of $\mathbf{V}$ give an orthonormal basis spanning $\mathcal{N}(\mathbf{A})$. An orthonormal basis spanning the range of $\mathbf{A}$ is given by the first r columns of $\mathbf{U}$.

$$\|\mathbf{A}\|_F^2 = \text{Frobenius norm} = \text{tr}(\mathbf{A}^T\mathbf{A}) = \sum_{i=1}^{r} \sigma_i^2 \tag{100}$$

$$\|\mathbf{A}\|_2^2 = \sigma_1^2 \tag{101}$$

$$\|\mathbf{A}\|_* = \text{nuclear norm} = \sum_{i=1}^{r} \sigma_i \tag{102}$$

The **condition number** κ of an invertible matrix $\mathbf{A} \in \mathbb{R}^{n,n}$ is the ratio of the largest and smallest singular value. Matrices with large condition numbers are closer to being singular and more sensitive to changes.

$$\kappa(\mathbf{A}) = \frac{\sigma_1}{\sigma_n} = \|A\|_2 \cdot \left\|A^{-1}\right\|_2 \tag{103}$$

Low-Rank Approximation

Approximating $\mathbf{A} \in \mathbb{R}^{m,n}$ by a matrix $\mathbf{A}_k$ of rank $k > 0$ can be formulated as the optimization probem

$$\min_{\mathbf{A}_k \in \mathbb{R}^{m,n}} \|\mathbf{A} - \mathbf{A}_k\|_F^2 : \text{rank}\, \mathbf{A}_k = k, 1 \leq k \leq \text{rank}(\mathbf{A}) \tag{104}$$

The optimal solution of this problem is given by

$$\mathbf{A}_k = \sum_{i=1}^{k} \sigma_i \mathbf{u}_i \mathbf{v}_i^T \tag{105}$$

where

$$\frac{\|\mathbf{A}_k\|_F^2}{\|\mathbf{A}\|_F^2} = \frac{\sigma_1^2 + \ldots + \sigma_k^2}{\sigma_1^2 + \ldots + \sigma_r^2} \tag{106}$$

$$1 - \frac{\|\mathbf{A}_k\|_F^2}{\|\mathbf{A}\|_F^2} = \frac{\sigma_{k+1}^2 + \ldots + \sigma_r^2}{\sigma_1^2 + \ldots + \sigma_r^2} \tag{107}$$

is the fraction of the total variance in $\mathbf{A}$ explained by the approximation $\mathbf{A}_k$.

Range and Nullspace

$$\mathcal{N}(\mathbf{A}) = \mathcal{R}(\mathbf{V}_{nr}) \tag{108}$$

$$\mathcal{N}(\mathbf{A})^\perp \equiv \mathcal{R}(\mathbf{A}^T) = \mathcal{R}(\mathbf{V}_r) \tag{109}$$

$$\mathcal{R}(\mathbf{A}) = \mathcal{R}(\mathbf{U}_r) \tag{110}$$

$$\mathcal{R}(\mathbf{A})^\perp \equiv \mathcal{N}(\mathbf{A}^T) = \mathcal{R}(\mathbf{U}_{nr}) \tag{111}$$

where $\mathbf{V}_r$ is the first r columns of V and $V_n r$ are the last $[r+1, n]$ columns; similarly for $\mathbf{U}$.

Projectors

The projection of $\mathbf{x}$ onto $\mathcal{N}(\mathbf{A})$ is $(\mathbf{V}_{nr}\mathbf{V}_{nr}^T)\mathbf{x}$. Since $\mathbf{I}_n = \mathbf{V}_r\mathbf{V}_r^T + \mathbf{V}_{nr}\mathbf{V}_{nr}^T$, $(\mathbf{I}_n - \mathbf{V}_r\mathbf{V}_r^T)\mathbf{x}$ also works. The projection of $\mathbf{x}$ onto $\mathcal{R}(\mathbf{A})$ is $(\mathbf{U}_r\mathbf{U}_r^T)\mathbf{x}$.

If $\mathbf{A} \in \mathbb{R}^{m,n}$ is full row rank ($\mathbf{A}\mathbf{A}^T \succ 0$), then the minimum distance to an affine set $\{x : \mathbf{A}\mathbf{x} = \mathbf{b}\}, \mathbf{b} \in \mathbb{R}^m$ is given by $\mathbf{x}^* = \mathbf{A}^T(\mathbf{A}\mathbf{A}^T)^{-1}\mathbf{b}$.

If $\mathbf{A} \in \mathbb{R}^{m,n}$ is full column rank ($\mathbf{A}^T\mathbf{A} \succ 0$), then the minimum distance to an affine set $\{x : \mathbf{A}\mathbf{x} = \mathbf{b}\}, \mathbf{b} \in \mathbb{R}^m$ is given by $\mathbf{x}^* = \mathbf{A}(\mathbf{A}^T\mathbf{A})^{-1}\mathbf{A}^T\mathbf{b}$.

Computational Notes

A *numerical rank* can be estimated for the matrix as the largest k such that $\sigma_k > \epsilon\sigma_1$ for $\epsilon \geq 0$.

6.6 Eigenvalue Decomposition for Diagonalizable Matrices

For a square, diagonalizable matrix $\mathbf{A} \in \mathbb{R}^{n,n}$

$$\mathbf{A} = U\Lambda U^{-1} \tag{112}$$

where $U \in \mathbb{C}^{n,n}$ is an invertible matrix whose columns are the eigenvectors of $\mathbf{A}$ and Λ is a diagonal matrix containing the eigenvalues $\lambda_1, \ldots, \lambda_n$ of $\mathbf{A}$ in the diagonal.

The columns $\mathbf{u}_1, \ldots, \mathbf{u}_n$ satisfy

$$\mathbf{A}\mathbf{u}_i = \lambda_i\mathbf{u}_i \quad i = 1, \ldots, n \tag{113}$$

6.7 Eigenvalue (Spectral) Decomposition for Symmetric Matrices

A symmetric matrix $\mathbf{A} \in \mathbb{R}^{n,n}$ can be factored as

$$\mathbf{A} = U\Lambda U^T = \sum_i^n \lambda_i\mathbf{u}_i\mathbf{u}_i^T \tag{114}$$

where $U \in \mathbb{R}^{n,n}$ is an orthogonal matrix whose columns $\mathbf{u}_i$ are the eigenvectors of $\mathbf{A}$ and Λ is a diagonal matrix containing the eigenvalues $\lambda_1 \geq \ldots \geq \lambda_n$ of $\mathbf{A}$ in the diagonal. These eigenvalues are always real. The eigenvectors can always be chosen to be real and to form an orthonormal basis.

The columns $\mathbf{u}_1, \ldots, \mathbf{u}_n$ satisfy

$$\mathbf{A}\mathbf{u}_i = \lambda_i\mathbf{u}_i \quad i = 1, \ldots, n \tag{115}$$

6.8 Schur Complements

For $\mathbf{A} \in \mathbb{S}^n$, $\mathbf{B} \in \mathbb{S}^n$, $\mathbf{X} \in \mathbb{R}^{n,m}$ with $\mathbf{B} \succ 0$ and the block matrix

$$\mathbf{M} = \begin{bmatrix} \mathbf{A} & \mathbf{X} \\ \mathbf{X}^T & \mathbf{B} \end{bmatrix} \tag{116}$$

and the Schur complement of $\mathbf{A}$ in $\mathbf{M}$

$$S = \mathbf{A} - \mathbf{X}\mathbf{B}^{-1}\mathbf{X}^T \tag{117}$$

Then

$$\mathbf{M} \succeq 0 \iff S \succeq 0 \tag{118}$$

$$\mathbf{M} \succ 0 \iff S \succ 0 \tag{119}$$

7 | Transpose Properties

$$(\mathbf{AB})^T = \mathbf{B}^T \mathbf{A}^T \tag{120}$$

$$(\mathbf{A} + \mathbf{B})^T = \mathbf{A}^T + \mathbf{B}^T \tag{121}$$

$$(\mathbf{A}^{-1})^T = (\mathbf{A}^T)^{-1} \tag{122}$$

8 Determinant Properties

Geometrically, if a unit volume is acted on by $\mathbf{A}$, then $|\det(\mathbf{A})|$ indicates the volume after the transformation.

$$\det(I_n) = 1 \tag{123}$$

$$\det(\mathbf{A}^T) = \det(\mathbf{A}) \tag{124}$$

$$\det(\mathbf{A}^{-1}) = \frac{1}{\det(\mathbf{A})} = \det(\mathbf{A})^{-1} \tag{125}$$

$$\det(\mathbf{AB}) = \det(\mathbf{BA}) \tag{126}$$

$$\det(\mathbf{AB}) = \det(\mathbf{A})\det(\mathbf{B}) \qquad \mathbf{A}, \mathbf{B} \in \mathbb{R}^{n,n} \tag{127}$$

$$\det(c\mathbf{A}) = c^n \det(\mathbf{A}) \qquad \mathbf{A} \in \mathbb{R}^{n,n} \tag{128}$$

$$\det(\mathbf{A}) = \prod \operatorname{eig}(\mathbf{A}) \tag{129}$$

For $\mathbf{A} \in \mathbb{R}^{m,n}, \mathbf{B} \in \mathbb{R}^{n,m}$

$$\det(\mathbf{I}_m + \mathbf{AB}) = \det(\mathbf{I}_n + \mathbf{BA}) \qquad \text{Sylvester's determinant identity} \tag{130}$$ [15]

9 Trace Properties

$$\operatorname{tr}(\mathbf{A}) = \sum_{i=1}^{n} \mathbf{A}_{ii} \qquad \mathbf{A} \in \mathbb{R}^{n,n} \tag{131}$$

$$\operatorname{tr}(\mathbf{A} + \mathbf{B}) = \operatorname{tr}(\mathbf{A}) + \operatorname{tr}(\mathbf{B}) \tag{132}$$

$$\operatorname{tr}(c\mathbf{A}) = c \operatorname{tr}(\mathbf{A}) \tag{133}$$

$$\operatorname{tr}(\mathbf{A}) = \operatorname{tr}(\mathbf{A}^T) \tag{134}$$

For $\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D}$ of compatible dimensions,

$$\operatorname{tr}(\mathbf{A}^T\mathbf{B}) = \operatorname{tr}(\mathbf{A}\mathbf{B}^T) = \operatorname{tr}(\mathbf{B}^T\mathbf{A}) = \operatorname{tr}(\mathbf{B}\mathbf{A}^T) \tag{135}$$

$$\operatorname{tr}(\mathbf{ABCD}) = \operatorname{tr}(\mathbf{BCDA}) = \operatorname{tr}(\mathbf{CDAB}) = \operatorname{tr}(\mathbf{DABC}) \tag{136}$$

(Invariant under cyclic permutations)

10 Inverse Properties

The inverse of $\mathbf{A} \in \mathbb{C}^{n,n}$ is denoted $\mathbf{A}^{-1}$ and defined such that

$$\mathbf{A}\mathbf{A}^{-1} = \mathbf{A}^{-1}\mathbf{A} = \mathbf{I}_n \tag{137}$$

where $\mathbf{I}_n$ is the $n \times n$ identity matrix. $\mathbf{A}$ is nonsingular if $\mathbf{A}^{-1}$ exists; otherwise, $\mathbf{A}$ is singular.

If individual inverses exist

$$(\mathbf{A}\mathbf{B})^{-1} = \mathbf{B}^{-1}\mathbf{A}^{-1} \tag{138}$$

more generally

$$(\mathbf{A}_1\mathbf{A}_2 \ldots \mathbf{A}_n)^{-1} = \mathbf{A}_n^{-1} \ldots \mathbf{A}_2^{-1}\mathbf{A}_1^{-1} \tag{139}$$

$$(\mathbf{A}^{-1})^T = (\mathbf{A}^T)^{-1} \tag{140}$$

11 Pseudo-Inverse Properties

For $\mathbf{A} \in \mathbb{R}^{m,n}$, a pseudoinverse satisfies:

$$\mathbf{A}\mathbf{A}^{+}\mathbf{A} = \mathbf{A} \tag{141}$$
$$\mathbf{A}^{+}\mathbf{A}\mathbf{A}^{+} = \mathbf{A}^{+} \tag{142}$$
$$(\mathbf{A}\mathbf{A}^{+})^{T} = \mathbf{A}\mathbf{A}^{+} \tag{143}$$
$$(\mathbf{A}^{+}\mathbf{A})^{T} = \mathbf{A}^{+}\mathbf{A} \tag{144}$$

11.1 Moore-Penrose Pseudoinverse

$$\mathbf{A}^{+} = \mathbf{V}\mathbf{D}^{-1}\mathbf{U}^{T} \tag{145}$$

where the foregoing comes from a singular-value decomposition and $\mathbf{D}^{-1} = \mathrm{diag}(\frac{1}{\sigma_1}, \ldots, \frac{1}{\sigma_r})$

Special Properties

- $\mathbf{A}^{+} = \mathbf{A}^{-1}$ if $\mathbf{A} \in \mathbb{R}^{n,n}$ and $\mathbf{A}$ is square and nonsingular.
- $\mathbf{A}^{+} = (\mathbf{A}^{T}\mathbf{A})^{-1}\mathbf{A}^{T}$, if $\mathbf{A} \in \mathbb{R}^{m,n}$ is full column rank ($r = n \leq m$). $\mathbf{A}^{+}$ is a left inverse of $\mathbf{A}$, so $\mathbf{A}^{+}\mathbf{A} = \mathbf{V}_r\mathbf{V}_r^{T} = \mathbf{V}\mathbf{V}^{T} = \mathbf{I}_n$.
- $\mathbf{A}^{+} = \mathbf{A}^{T}(\mathbf{A}\mathbf{A}^{T})^{-1}$, if $\mathbf{A} \in \mathbb{R}^{m,n}$ is full row rank ($r = m \leq n$). $\mathbf{A}^{+}$ is a right inverse of $\mathbf{A}$, so $\mathbf{A}\mathbf{A}^{+} = \mathbf{U}_r\mathbf{U}_r^{T} = \mathbf{U}\mathbf{U}^{T} = \mathbf{I}_m$.

12 Hadamard Identities

$$(\mathbf{A} \circ \mathbf{B})_{ij} = A_{ij} B_{ij} \; \forall \; i, j \tag{146}$$

$$\mathbf{A} \circ \mathbf{B} = \mathbf{B} \circ \mathbf{A} \tag{147}$$ [16]

$$\mathbf{A} \circ (\mathbf{B} \circ \mathbf{C}) = (\mathbf{A} \circ \mathbf{B}) \circ \mathbf{C} \tag{148}$$

$$\mathbf{A} \circ (\mathbf{B} + \mathbf{C}) = \mathbf{A} \circ \mathbf{B} + \mathbf{A} \circ \mathbf{C} \tag{149}$$ [16]

$$a(\mathbf{A} \circ \mathbf{B}) = (a\mathbf{A}) \circ \mathbf{B} = \mathbf{A} \circ (a\mathbf{B}) \tag{150}$$ [16]

$$(\mathbf{A}^T \circ \mathbf{B}^T) = (\mathbf{A} \circ \mathbf{B})^T \tag{151}$$

$$(\mathbf{A}^T \circ \mathbf{B}^T) = (\mathbf{A} \circ \mathbf{B})^T \tag{152}$$

$$(\mathbf{x}^T \mathbf{A} \mathbf{x}) = \sum_{i,j} \big((\mathbf{x}\mathbf{x}^T) \circ \mathbf{A} \big) \tag{153}$$

$$\mathbf{x}^T (\mathbf{A} \circ \mathbf{B}) \mathbf{y} = \mathrm{tr}((\mathrm{diag}(\mathbf{x})\mathbf{A})^T \mathbf{B} \, \mathrm{diag}(\mathbf{y})) \quad \mathbf{A}, \mathbf{B} \in \mathbb{R}^{m,n} \tag{154}$$ [17]

$$\mathrm{tr}(\mathbf{A}^T \mathbf{B}) = \mathbf{1}^T (\mathbf{A} \circ \mathbf{B}) \mathbf{1} \tag{155}$$

13 Eigenvalue Properties

$\lambda \in \mathbb{C}$ is an eigenvalue of $\mathbf{A} \in \mathbb{R}^{n,n}$ and $u \in \mathbb{C}^n$ is a corresponding eigenvector if $\mathbf{Au} = \lambda\mathbf{u}$ and $\mathbf{u} \neq 0$. Equivalantly, $(\lambda\mathbf{I}_n - \mathbf{A})\mathbf{u} = 0$ and $\mathbf{u} \neq 0$. Eigenvalues satisfy the equation $\det(\lambda\mathbf{I}_n - \mathbf{A}) = 0$.

Any matrix $\mathbf{A} \in \mathbb{R}^{n,n}$ has n eigenvalues, though some may be repeated. λ_1 is the largest eigenvalue and λ_n the smallest.

$$\mathrm{eig}(\mathbf{A}\mathbf{A}^T) = \mathrm{eig}(\mathbf{A}^T\mathbf{A}) \tag{156}$$

(Note that the number of entries in $\mathbf{A}\mathbf{A}^T$ and $\mathbf{A}^T\mathbf{A}$ may differ significantly leading to different compute times.)

$$\mathrm{eig}(\mathbf{A}^T\mathbf{A}) \geq 0 \tag{157}$$

$$\lambda_{\min}(\mathbf{A}) \leq \frac{\mathbf{x}^T\mathbf{A}\mathbf{x}}{\mathbf{x}^T\mathbf{x}} \leq \lambda_{\max}(\mathbf{A}) \quad \mathbf{x} \neq 0 \tag{158}$$

13.0.1 Weyl's Inequality

If $\mathbf{M}, \mathbf{H}, \mathbf{P} \in \mathbb{R}^{n,n}$ are Hermitian matrices and $\mathbf{M} = \mathbf{H} + \mathbf{P}$ ($\mathbf{H}$ is perturbed by $\mathbf{P}$) and $\mathbf{M}$ has eigenvalues $\mu_1 \geq \cdots \geq \mu_n$, $\mathbf{H}$ has eigenvalues $\nu_1 \geq \cdots \geq \nu_n$, and $\mathbf{P}$ has eigenvalues $\rho_1 \geq \cdots \geq \rho_n$, then

$$\nu_i + \rho_n \leq \mu_i \leq \nu_i + \rho_1 \ \forall i \tag{159}$$

If $j + k - n \geq i \geq r + s - 1$, then

$$\nu_j + \rho_k \leq \mu_i \leq \nu_r + \rho_s \tag{160}$$

If $\mathbf{P} \succeq 0$, then $\mu_i > \nu_i \ \forall i$.

14 Norms

14.1 General Properties

Matrix norms satisfy some properties:

$$f(\mathbf{A}) \geq 0 \tag{161}$$
$$f(\mathbf{A}) = 0 \iff \mathbf{A} = 0 \tag{162}$$
$$f(c\mathbf{A}) = |c| f(\mathbf{A}) \tag{163}$$
$$f(\mathbf{A} + \mathbf{B}) \leq f(\mathbf{A}) + f(\mathbf{B}) \tag{164}$$

Many popular norms also satisfy "sub-multiplicativity": $f(\mathbf{AB}) \leq f(\mathbf{A}) f(\mathbf{B})$.

14.2 Matrices

14.2.1 Frobenius norm

$$\|\mathbf{A}\|_F = \sqrt{\operatorname{tr} \mathbf{A}\mathbf{A}^H} \tag{165}$$
$$= \sqrt{\sum_{i=1}^{m} \sum_{j=1}^{n} |\mathbf{A}_{ij}|^2} \tag{166}$$
$$= \sqrt{\sum_{i=1}^{m} \operatorname{eig}(A^H A)_i} \tag{167}$$

Special Properties

$$\|\mathbf{Ax}\|_2 \leq \|\mathbf{A}\|_F \|\mathbf{x}\|_2 \quad \mathbf{x} \in \mathbb{R}^n \tag{168}$$
$$\|\mathbf{AB}\|_F \leq \|\mathbf{A}\|_F \|\mathbf{B}\|_F \tag{169}$$
$$\left\|\mathbf{C} - \mathbf{x}\mathbf{x}^T\right\|_F^2 = \|\mathbf{C}\|_F^2 + \|\mathbf{x}\|_2^4 - 2\mathbf{x}^T \mathbf{C} \mathbf{x} \tag{170}$$

14.2.2 Operator Norms

For $p = 1, 2, \infty$ or other values, an operator norm indicates the maximum input-output gain of the matrix.

$$\|\mathbf{A}\|_p = \max_{\|\mathbf{u}\|_p = 1} \|\mathbf{Au}\|_p \tag{171}$$

$$\|\mathbf{A}\|_1 = \max_{\|\mathbf{u}\|_1 = 1} \|\mathbf{Au}\|_1 \tag{172}$$
$$= \max_{j=1,\ldots,n} \sum_{i=1}^{m} |\mathbf{A}_{ij}| \tag{173}$$
$$= \text{Largest absolute column sum} \tag{174}$$

$$\|\mathbf{A}\|_\infty = \max_{\|\mathbf{u}\|_\infty=1} \|\mathbf{A}\mathbf{u}\|_\infty \tag{175}$$
$$= \max_{j=1,\dots,m} \sum_{i=1}^{n} |\mathbf{A}_{ij}| \tag{176}$$
$$= \text{Largest absolute row sum} \tag{177}$$

$$\|\mathbf{A}\|_2 = \text{“spectral norm”} \tag{178}$$
$$= \max_{\|\mathbf{u}\|_2=1} \|\mathbf{A}\mathbf{u}\|_2 \tag{179}$$
$$= \sqrt{\max(\text{eig}(\mathbf{A}^T\mathbf{A}))} \tag{180}$$
$$= \text{Square root of largest eigenvalue of } \mathbf{A}^T\mathbf{A} \tag{181}$$

Special Properties

$$\|\mathbf{A}\mathbf{u}\|_p \le \|\mathbf{A}\|_p \|\mathbf{u}\|_p \tag{182}$$
$$\|\mathbf{A}\mathbf{B}\|_p \le \|\mathbf{A}\|_p \|\mathbf{B}\|_p \tag{183}$$

14.2.3 Spectral Radius

Not a proper norm.

$$\rho(\mathbf{A}) = \text{spectral radius}(\mathbf{A}) = \max_{i=1,\dots,n} |\,\text{eig}(\mathbf{A})_i| \tag{184}$$

Special Properties

$$\rho(\mathbf{A}) \le \|\mathbf{A}\|_p \tag{185}$$
$$\rho(\mathbf{A}) \le \min(\ \|\mathbf{A}\|_1, \|\mathbf{A}\|_\infty) \tag{186}$$

14.3 Vectors

$$\|\mathbf{x}\|_1 = \sum_i |\mathbf{x}_i| \qquad \text{L1-norm} \tag{187}$$
$$\|\mathbf{x}\|_p = (\sum_i |\mathbf{x}_i|^p)^{1/p} \qquad \text{P-norm} \tag{188}$$
$$\|\mathbf{x}\|_\infty = \max_i |\mathbf{x}_i| \qquad \text{L}\infty\text{-norm, L-infinity norm} \tag{189}$$

14.3.1 Identities

$$2\|\mathbf{u}\|_2^2 + 2\|\mathbf{v}\|_2^2 = \|\mathbf{u}+\mathbf{v}\|_2^2 + \|\mathbf{u}-\mathbf{v}\|_2^2 \qquad \text{Polarization Identity} \tag{190}$$
$$<\mathbf{x},\mathbf{y}> = \frac{1}{4}\left(\|\mathbf{x}+\mathbf{y}\|_2^2 - \|\mathbf{x}-\mathbf{y}\|_2^2\right) \quad \forall \mathbf{x},\mathbf{y} \in \mathcal{V} \qquad \text{Polarization Identity} \tag{191}$$

14.3.2 Bounds

$$|\mathbf{x}^T\mathbf{y}| \leq \|\mathbf{x}\|_2\|\mathbf{y}\|_2 \qquad \text{Cauchy-Schwartz Inequality} \tag{192}$$

$$|\mathbf{x}^T\mathbf{y}| \leq \sum_{k=1}^{n} |\mathbf{x}_k\mathbf{y}_k| \leq \|\mathbf{x}\|_p\|\mathbf{x}\|_q \quad \forall p, q \geq 1 : 1/p + 1/q = 1 \qquad \text{Hölder Inequality} \tag{193}$$

For $\mathbf{x} \in \mathbb{R}^n$

$$\frac{1}{\sqrt{n}}\|\mathbf{x}\|_2 \leq \|\mathbf{x}\|_\infty \leq \|\mathbf{x}\|_2 \leq \|\mathbf{x}\|_1 \leq \sqrt{\mathrm{card}(\mathbf{x})}\|\mathbf{x}\|_2 \leq \sqrt{n}\|\mathbf{x}\|_2 \leq n\|\mathbf{x}\|_\infty \tag{194}$$

For any $0 < p < q$ we have that $\|\mathbf{x}\|_q \leq \|\mathbf{x}\|_p$.

15 Bounds

15.1 Matrix Gain

$$\lambda_{\min}(\mathbf{A}^T\mathbf{A}) \leq \frac{\|\mathbf{A}\mathbf{x}\|_2^2}{\|\mathbf{x}\|_2^2} \leq \lambda_{\max}(\mathbf{A}^T\mathbf{A}) \tag{195}$$

$$\max_{\mathbf{x}\neq 0} \frac{\|\mathbf{A}\mathbf{x}\|_2}{\|\mathbf{x}\|_2} = \|\mathbf{A}\|_2 = \sqrt{\lambda_{\max}(\mathbf{A}^T\mathbf{A})} \implies \mathbf{x} = u_1 \tag{196}$$

$$\min_{\mathbf{x}\neq 0} \frac{\|\mathbf{A}\mathbf{x}\|_2}{\|\mathbf{x}\|_2} = \sqrt{\lambda_{\min}(\mathbf{A}^T\mathbf{A})} \implies \mathbf{x} = u_n \tag{197}$$

15.2 Rayleigh quotients

The Rayleigh quotient of $\mathbf{A} \in \mathbb{S}^n$ is given by

$$\frac{\mathbf{x}^T\mathbf{A}\mathbf{x}}{\mathbf{x}^T\mathbf{x}} \quad \mathbf{x} \neq 0 \tag{198}$$

$$\lambda_{\min}(\mathbf{A}) \leq \frac{\mathbf{x}^T\mathbf{A}\mathbf{x}}{\mathbf{x}^T\mathbf{x}} \leq \lambda_{\max}(\mathbf{A}) \quad \mathbf{x} \neq 0 \tag{199}$$

$$\lambda_{\max}(A) = \max_{\mathbf{x}:\|\mathbf{x}\|_2=1} \mathbf{x}^T\mathbf{A}\mathbf{x} = u_1 \tag{200}$$

$$\lambda_{\min}(A) = \min_{\mathbf{x}:\|\mathbf{x}\|_2=1} \mathbf{x}^T\mathbf{A}\mathbf{x} = u_n \tag{201}$$

where u_1 and u_n are the eigenvectors associated with $\lambda_{\max}$ and $\lambda_{\min}$, respectively.

16 Linear Equations

The linear equation $\mathbf{Ax} = \mathbf{y}$ with $\mathbf{A} \in \mathbb{R}^{m,n}$ admits a solution iff $\text{rank}([\mathbf{Ay}]) = \text{rank}(\mathbf{A})$. If this is satisfied, the set of all solutions is an affine set $\mathcal{S} = \{\mathbf{x} = \bar{\mathbf{x}} + z : z \in \mathcal{N}(\mathbf{A})\}$ where $\bar{\mathbf{x}}$ is any vector such that $\mathbf{A}\bar{\mathbf{x}} = \mathbf{y}$. The solution is unique if $\mathcal{N}(\mathbf{A}) = \{0\}$.

$\mathbf{Ax} = \mathbf{y}$ is *overdetermined* if it is tall/skinny ($m > n$); that is, if there are more equations than unknowns. If $\text{rank}(\mathbf{A}) = n$ then $\dim \mathcal{N}(\mathbf{A}) = 0$, so there is either no solution or one solution. Overdetermined systems often have no solution ($\mathbf{y} \notin \mathcal{R}(\mathbf{A})$), so an approximate solution is necessary. See section 16.1.

$\mathbf{Ax} = \mathbf{y}$ is *underdetermined* if it is short/wide ($n > m$); that is, if has more unknowns than equations. If $\text{rank}(\mathbf{A}) = m$ then $\mathcal{R}(\mathbf{A}) = \mathbb{R}^m$, so $\dim \mathcal{N}(\mathbf{A}) = n - m > 0$, so the set of solutions is infinite. Therefore, finding a single solution that optimizes some quantity is of interest.

$\mathbf{Ax} = \mathbf{y}$ is *square* if $n = m$. If $\mathbf{A}$ is invertible, then the equations have the unique solution $\mathbf{x} = \mathbf{A}^{-1}\mathbf{y}$. See section 16.2.

16.1 Least-Squares

For an overdetermined system we wish to find:

$$\min_{\mathbf{x}} \|\mathbf{Ax} - \mathbf{y}\|_2^2 \tag{202}$$

Since $\mathbf{Ax} \in \mathcal{R}(\mathbf{A})$, we need a point $\tilde{\mathbf{y}} = \mathbf{Ax}^* \in \mathcal{R}(\mathbf{A})$ closest to $\mathbf{y}$. This point lies in the nullspace of $\mathbf{A}^T$, so we have $\mathbf{A}^T(\mathbf{y} - \mathbf{Ax}^*) = 0$. There is always a solution to this problem and, if $\text{rank}(\mathbf{A}) = n$, it is unique [18, p. 161]

$$\mathbf{x}^* = (\mathbf{A}^T\mathbf{A})^{-1}\mathbf{A}^T\mathbf{y} \tag{203}$$

16.1.1 Regularized least-squares with low-rank data

For $\mathbf{A} \in \mathbb{R}^{m,n}$, $\mathbf{y} \in \mathbb{R}^m$, $\lambda \geq 0$, the regularized least-squares problem

$$\text{argmin}_{\mathbf{x}} \|\mathbf{Ax} - \mathbf{y}\|_2^2 + \lambda\|\mathbf{x}\|_2^2 \tag{204}$$

has a closed form solution

$$\mathbf{x} = (\mathbf{A}^T\mathbf{A} + \lambda\mathbf{I})^{-1}\mathbf{A}^T\mathbf{y} \tag{205}$$

However, if $\mathbf{A}$ has a rank $r \ll \min(n, m)$ and a known low-rank decomposition $\mathbf{A} = \mathbf{LR}^T$ with $\mathbf{L} \in \mathbb{R}^{m,r}$ and $\mathbf{R} \in \mathbb{R}^{n,r}$, then we can rewrite Equation 205 as

$$\mathbf{x} = (\mathbf{R}^T\mathbf{R}\mathbf{L}^T\mathbf{L} + \lambda\mathbf{I})^{-1}\mathbf{L}^T\mathbf{y} \tag{206}$$

This decreases the time complexity from $O(mn^2 + n^\omega)$ to $O(nr^2 + mr^2)$.

16.2 Minimum Norm Solutions

For undertermined systems in which $\mathbf{A} \in \mathbb{R}^{m,n}$ with $m < n$. We wish to find

$$\min_{\mathbf{x}:\mathbf{Ax}=\mathbf{y}} \|\mathbf{x}\|_2 \tag{207}$$

The solution $\mathbf{x}^*$ must be orthogonal to $\mathcal{N}(\mathbf{A})$, so $\mathbf{x}^* \in \mathcal{R}(\mathbf{A}^T)$, so $\mathbf{x}^* = \mathbf{A}^T c$ for some c. Substituting into $\mathbf{A}\mathbf{x} = \mathbf{y}$ gives $\mathbf{A}\mathbf{A}^T c = \mathbf{y}$, therefore [18, p. 162]:

$$\mathbf{x}^* = \mathbf{A}^T(\mathbf{A}\mathbf{A}^T)^{-1}\mathbf{y} \tag{208}$$

17 | Updates

17.1 Removing a row from $\mathbf{A}^T\mathbf{A}$ ($\mathbf{A}^T\mathbf{A} \rightarrow \mathbf{A}_{\setminus i}^T\mathbf{A}_{\setminus i}$)

Plain English: Matrix times its transpose after eliminating row i from the matrix

Inputs: $\mathbf{A} \in \mathbb{R}^{k,m}, \mathbf{u} \in \mathbb{R}^m, \mathbf{v} \in \mathbb{R}^n$ and i, the row to remove from $\mathbf{A}$

Reduces to: $\mathbf{C} \in \mathbb{R}^{k,l}$

Algorithm:

$$\mathbf{A}_{\setminus i}^T\mathbf{A}_{\setminus i} = \mathbf{A}^T\mathbf{A} - \mathbf{A}_{*i}\mathbf{A}_{*i}^T \tag{209}$$

Similarly:

$$\mathbf{A}_{\setminus i}^T y_{\setminus i} = \mathbf{A}^T y - \mathbf{A}_{*i} y_i^T \tag{210}$$

17.2 $\mathbf{1}_r^T\mathbf{A}\mathbf{1}_c$

Plain English: The sum of the elements of the matrix.

Reduces to: Scalar

Notation: For $\mathbf{A} \in \mathbb{R}^{r\times c}$, $\mathbf{1}_r$ is in $\mathbb{R}^{r\times 1}$ and $\mathbf{1}_c$ is in $\mathbb{R}^{c\times 1}$.

Algorithm: Traverse all the element of the matrix in order keeping track of the sum. For applications where accuracy is important and the matrices have a large dynamic range, Kahan summation or a similar technique should be used.

Update Algorithm: If an entry changes, subtract its old value from the sum and add its new value to the sum.

17.3 $\mathbf{x}^T\mathbf{A}\mathbf{x}$

Plain English: TODO

Reduces to: Scalar

Notation: $\mathbf{A}$ must be in $\mathbb{R}^{i\times i}$. $\mathbf{x}$ is in $\mathbb{R}^{i\times 1}$.

Algorithm: TODO

Update Algorithm: We make use of the identity $(\mathbf{x}^T\mathbf{A}\mathbf{x}) = \sum_{i,j}\big((\mathbf{x}\mathbf{x}^T)\circ\mathbf{A}\big)$. If an entry $\mathbf{A}_{i,j}$ in the matrix changes subtract its old value $\mathbf{x}_i\mathbf{x}_j\mathbf{A}_{ij}$ and add the new value $\mathbf{x}_i\mathbf{x}_j\mathbf{A}'_{ij}$. If an entry $\mathbf{x}_i$ changes TODO.

18 Optimization

18.1 Standard Forms

Least Squares

$$\min_{\mathbf{x}\in\mathbb{R}^n} \|\mathbf{y}-\mathbf{A}\mathbf{x}\|_2 \tag{211}$$

LASSO

$$\min_{\mathbf{b}\in\mathbb{R}^n} \left(\frac{1}{N}\|\mathbf{y}-\mathbf{X}\mathbf{b}\|_2^2 + \lambda\|\mathbf{b}\|_1\right) \tag{212}$$

LP: Linear program

$$\begin{aligned} \underset{\mathbf{x}}{\text{minimize}} \quad & \mathbf{c}^T\mathbf{x} && (213a)\\ \text{subject to} \quad & \mathbf{A}_{\text{eq}}\mathbf{x} = \mathbf{b}_{\text{eq}}, && (213b)\\ & \mathbf{A}\mathbf{x} \le \mathbf{b} && (213c) \end{aligned}$$

Linear Fractional Program

$$\begin{aligned} \underset{\mathbf{x}}{\text{maximize}} \quad & \frac{\mathbf{c}^T\mathbf{x}+a}{\mathbf{d}^T\mathbf{x}+b} && (214a)\\ \text{subject to} \quad & \mathbf{A}\mathbf{x} \le \mathbf{b} && (214b) \end{aligned}$$

Additional constraints must ensure $\mathbf{d}^T\mathbf{x}+b$ has the same sign throughout the entire feasible region.

QCQP: Quadratic Constrainted Quadratic Programs

$$\begin{aligned} \underset{\mathbf{x}}{\text{minimize}} \quad & \mathbf{x}^T\mathbf{H}_0\mathbf{x} + 2\mathbf{c}_0^T\mathbf{x} + \mathbf{d}_0 && (215a)\\ \text{subject to} \quad & \mathbf{x}^T\mathbf{H}_i\mathbf{x} + 2\mathbf{c}_i^T\mathbf{x} + \mathbf{d}_i \le 0 \quad i\in\mathcal{I}, && (215b)\\ & \mathbf{x}^T\mathbf{H}_j\mathbf{x} + 2\mathbf{c}_j^T\mathbf{x} + \mathbf{d}_j = 0 \quad j\in\mathcal{E} && (215c) \end{aligned}$$

If $\mathbf{H}_i \succeq 0\ \forall i$, then the program is convex. In general, QCQPs are NP-Hard.

QP: Quadratic Program

$$\begin{aligned} \underset{\mathbf{x}}{\text{minimize}} \quad & \frac{1}{2}\mathbf{x}^T\mathbf{H}_0\mathbf{x} + \mathbf{c}_0^T\mathbf{x} && (216a)\\ \text{subject to} \quad & \mathbf{A}_{\text{eq}}\mathbf{x} = \mathbf{b}_{\text{eq}}, && (216b)\\ & \mathbf{A}\mathbf{x} \le \mathbf{b} && (216c) \end{aligned}$$

If $\mathbf{H}_0 \succ 0$, then the program is convex.

If only equality constraints are present, then the solution is the linear system:

$$\begin{bmatrix} \mathbf{H}_0 & \mathbf{A}^T \\ \mathbf{A} & 0 \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ \lambda \end{bmatrix} = \begin{bmatrix} -\mathbf{c}_0 \\ \mathbf{b} \end{bmatrix} \tag{217}$$

where λ is a set of Lagrange multipliers.

For $\mathbf{H}_0 \succ 0$, the ellipsoid method solves the problem in polynomial time. [19] If, $\mathbf{H}_0$ is indefinite, then the problem is NP-hard [20], even if $\mathbf{H}_0$ has only one negative eigenvalue [21].

SOCP: Second Order Cone Program (Standard Form)

$$\min_{\mathbf{x}} \ \mathbf{c}^T\mathbf{x} \tag{218}$$

$$\text{s.t. } \|\mathbf{A}_i\mathbf{x}+\mathbf{b}_i\|_2 \le \mathbf{c}_i^T\mathbf{x}+\mathbf{d}_i, \quad i=1,\dots,m \tag{219}$$

SOCP: Second Order Cone Program (Conic Standard Form)

$$\min_{\mathbf{x}} \ \mathbf{c}^T\mathbf{x} \tag{220}$$

$$\text{s.t. } (\mathbf{A}_i\mathbf{x}+\mathbf{b}_i, \mathbf{c}_i^T\mathbf{x}+\mathbf{d}_i) \in \mathcal{K}_{m_i} \quad i=1,\dots,m \tag{221}$$

18.2 Transformations

18.2.1 Linear-Fractional to Linear

We transform a Linear-Fractional Program

$$\underset{\mathbf{x}}{\text{maximize}} \quad \frac{\mathbf{c}^T\mathbf{x}+a}{\mathbf{d}^T\mathbf{x}+b} \tag{222a}$$

$$\text{subject to} \quad \mathbf{A}\mathbf{x} \le \mathbf{b} \tag{222b}$$

where $\mathbf{d}^T\mathbf{x}+b$ has the same sign throughout the entire feasible region to a linear program using the Charnes–Cooper transformation [22] by defining

$$\mathbf{y} = \frac{1}{\mathbf{d}^T\mathbf{x}+b}\cdot\mathbf{x} \tag{223}$$

$$t = \frac{1}{\mathbf{d}^T\mathbf{x}+b} \tag{224}$$

to form the equivalent program

$$\underset{\mathbf{y},t}{\text{maximize}} \quad \mathbf{c}^T\mathbf{y}+at \tag{225a}$$

$$\text{subject to} \quad \mathbf{A}\mathbf{y} \le \mathbf{b}t, \tag{225b}$$

$$\mathbf{d}^T\mathbf{y}+bt = 1, \tag{225c}$$

$$t \ge 0 \tag{225d}$$

We then have $\mathbf{x}^* = \frac{1}{t}\mathbf{y}$.

18.2.2 LP as SOCP

The linear program

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{226a}$$

$$\text{subject to} \quad \mathbf{A}\mathbf{x} \le \mathbf{b} \tag{226b}$$

becomes can be cast as an SOCP:

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{227a}$$

$$\text{subject to} \quad \|\mathbf{C}_i\mathbf{x}+\mathbf{d}_i\|_2 \le \mathbf{b}_i - \mathbf{a}_i^T\mathbf{x}\forall i \tag{227b}$$

where $\mathbf{C}_i = 0, d_i = 0 \ \forall i$.

18.2.3 QCQP as SOCP

The quadratic constrainted quadratic program

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{x}^T\mathbf{Q}_0\mathbf{x} + \mathbf{a}_0^T\mathbf{x} \tag{228a}$$
$$\text{subject to} \quad \mathbf{x}^T\mathbf{Q}_i\mathbf{x} + \mathbf{a}_i^T\mathbf{x} \le b_i \quad i = 1, \dots, m \tag{228b}$$

with $\mathbf{Q}_i = \mathbf{Q}_i^T \succeq 0$, $i = 0, \dots, m$ can be cast as an SOCP:

$$\underset{\mathbf{x}, t}{\text{minimize}} \quad \mathbf{a}_0^T\mathbf{x} + t \tag{229a}$$
$$\text{subject to} \quad \left\| \begin{bmatrix} 2\mathbf{Q}_0^{1/2}\mathbf{x} \\ t-1 \end{bmatrix} \right\|_2 \le t + 1, \tag{229b}$$
$$\left\| \begin{bmatrix} 2\mathbf{Q}_i^{1/2}\mathbf{x} \\ b_i - \mathbf{a}_i^T\mathbf{x} - 1 \end{bmatrix} \right\|_2 \le b_i - \mathbf{a}_i^T\mathbf{x} + 1 \quad i = 1, \dots, m \tag{229c}$$

18.2.4 QP as SOCP

The quadratic program

$$\underset{\mathbf{x}}{\text{minimize}} \quad \frac{1}{2}\mathbf{x}^T\mathbf{Q}\mathbf{x} + \mathbf{c}^T\mathbf{x} \tag{230a}$$
$$\text{subject to} \quad \mathbf{a}_i^T\mathbf{x} \le \mathbf{b}_i \tag{230b}$$

with $\mathbf{Q} = \mathbf{Q}^T \succeq 0$ can be cast as an SOCP:

$$\underset{\mathbf{x}, \mathbf{y}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} + y \tag{231a}$$
$$\text{subject to} \quad \left\| \begin{bmatrix} 2\mathbf{Q}^{1/2}\mathbf{x} \\ y-1 \end{bmatrix} \right\|_2 \le y + 1, \tag{231b}$$
$$\mathbf{a}_i^T\mathbf{x} \le \mathbf{b}_i \quad \forall i \tag{231c}$$

18.2.5 Sum of L2 Norms to SOCP

$$\underset{\mathbf{x}}{\text{minimize}} \quad \sum_{i=1}^{p} \|\mathbf{A}_i\mathbf{x} - \mathbf{b}_i\|_2 \tag{232a}$$

becomes

$$\underset{\mathbf{x}, y}{\text{minimize}} \quad \sum_{i=1}^{p} y_i \tag{233a}$$
$$\text{subject to} \quad \|\mathbf{A}_i\mathbf{x} - \mathbf{b}_i\|_2 \le y_i \quad i = 1, \dots, p \tag{233b}$$

18.2.6 Minimax of L2 Norms to SOCP

$$\underset{\mathbf{x}}{\text{minimize}} \quad \max_{i=1,\ldots,p} \|\mathbf{A}_i\mathbf{x} - \mathbf{b}_i\|_2 \tag{234a}$$

becomes

$$\begin{aligned} &\underset{\mathbf{x}, y}{\text{minimize}} && y && \text{(235a)} \\ &\text{subject to} && \|\mathbf{A}_i\mathbf{x} - \mathbf{b}_i\|_2 \le y \quad i = 1, \ldots, p && \text{(235b)} \end{aligned}$$

18.2.7 Hyperbolic Constraints to SOCP

For scalar w, a constraint of the form

$$w^2 \le xy, \quad x \ge 0, \quad y \ge 0 \tag{236}$$

can be transformed into the SOCP constraint

$$\left\| \begin{bmatrix} 2w \\ x - y \end{bmatrix} \right\|_2 \le x + y \tag{237}$$ [23]

For vector $\mathbf{w}$, a constraint of the form

$$\mathbf{w}^T\mathbf{w} = \|\mathbf{w}\|_2^2 \le xy, \quad x \ge 0, \quad y \ge 0 \tag{238}$$

can be transformed into the SOCP constraint

$$\left\| \begin{bmatrix} 2\mathbf{w} \\ x - y \end{bmatrix} \right\|_2 \le x + y \tag{239}$$ [23]

18.2.8 Matrix Fractional to SOCP

The problem

$$\begin{aligned} &\underset{\mathbf{x}}{\text{minimize}} && (\mathbf{F}\mathbf{x} + \mathbf{g})^T(\mathbf{P}_0 + \mathbf{x}_1\mathbf{P} + \ldots + \mathbf{x}_p\mathbf{P}_P)^{-1}(\mathbf{F}\mathbf{x} + \mathbf{g}) && \text{(240a)} \\ &\text{subject to} && \mathbf{P}_0 + \mathbf{x}_1\mathbf{P} + \ldots + \mathbf{x}_p\mathbf{P}_P > 0, && \text{(240b)} \\ &&& \mathbf{x} \ge 0 && \text{(240c)} \end{aligned}$$

where $\mathbf{P}_i = \mathbf{P}_i^T \in \mathbb{R}^{n,n}$, $\mathbf{F} \in \mathbb{R}^{n,p}$, $\mathbf{g} \in \mathbb{R}^n$, and $\mathbf{x} \in \mathbb{R}^p$ can be transformed into the SOCP where $t_i \in \mathbb{R}, \mathbf{y}_i \in \mathbb{R}^n$:

$$\begin{aligned} &\underset{\mathbf{x}, t}{\text{minimize}} && t_0 + \ldots + t_p && \text{(241a)} \\ &\text{subject to} && \mathbf{P}_0^{1/2}\mathbf{y}_0 + \ldots + \mathbf{P}_p^{1/2}\mathbf{y}_p = \mathbf{F}\mathbf{x} + \mathbf{g}, && \text{(241b)} \\ &&& \left\| \begin{bmatrix} 2\mathbf{y}_0 \\ t_0 - 1 \end{bmatrix} \right\|_2 \le t_0 + 1, && \text{(241c)} \\ &&& \left\| \begin{bmatrix} 2\mathbf{y}_i \\ t_i - x_i \end{bmatrix} \right\|_2 \le t_i + x_i \quad i = 1, \ldots, p && \text{(241d)} \end{aligned}$$ [23]

18.2.9 Fractional Objective to SOCP

Convert

$$\underset{\mathbf{x}}{\text{minimize}} \quad \frac{f(x)^2}{g(x)} \tag{242a}$$

$$\text{subject to} \quad g(x) > 0 \tag{242b}$$

to

$$\underset{\mathbf{x},t}{\text{minimize}} \quad t \tag{243a}$$

$$\text{subject to} \quad f(x)^2 \le tg(y), \tag{243b}$$

$$g(y) > 0, \tag{243c}$$

$$t \ge 0 \tag{243d}$$

and apply Equation 239.

18.2.10 Chance-Constrained LP to SOCP

The problem

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{244a}$$

$$\text{subject to} \quad \text{Prob}\{\mathbf{a}_i^T\mathbf{x} \le \mathbf{b}_i\} \ge p_i \quad i = 1, \dots, m \tag{244b}$$

where $p_i > 0.5$ and all $\mathbf{a}_i$ are independent normal random vectors with expected values $\bar{\mathbf{a}}_i$ and covariance matrices $\Sigma_i \succ 0$, can be transformed into the SOCP:

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{245a}$$

$$\text{subject to} \quad \bar{\mathbf{a}}_i^T\mathbf{x} \le b_i - \Phi^{-1}(p_i)\left\|\Sigma_i^{1/2}\mathbf{x}\right\|_2 \quad i = 1, \dots, m \tag{245b}$$

where $\Phi^{-1}(p)$ is the inverse cumulative probability distribution of a standard normal variable.

18.2.11 Robust LP with Box Uncertainty as LP

The problem

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{246a}$$

$$\text{subject to} \quad \mathbf{a}_i^T\mathbf{x} \le b_i \quad \forall \mathbf{a}_i \in \{\hat{\mathbf{a}}_i + \rho_i\mathbf{u} : \|\mathbf{u}\|_\infty \le 1\} \quad i = 1, \dots, m \tag{246b}$$

is equivalent to

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{247a}$$

$$\text{subject to} \quad \hat{\mathbf{a}}_i^T\mathbf{x} + \rho_i\|\mathbf{x}\|_1 \le b_i \quad i = 1, \dots, m \tag{247b}$$

which is equivalent to:

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{248a}$$

$$\text{subject to} \quad \hat{\mathbf{a}}_i^T\mathbf{x} + \rho_i\sum_{j=1}^n \mathbf{u}_j \le b_i \qquad i = 1, \dots, m, \tag{248b}$$

$$-\mathbf{u}_j \le \mathbf{x}_j \le \mathbf{u}_j \quad j = 1, \dots, n \tag{248c}$$

18.2.12 Robust LP with Ellipsoidal Uncertainty as SOCP

The problem

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{249a}$$
$$\text{subject to} \quad \mathbf{a}_i^T\mathbf{x} \leq b_i \quad \forall \mathbf{a}_i \in \{\hat{\mathbf{a}}_i + \mathbf{R}_i\mathbf{u} : \|\mathbf{u}\|_2 \leq 1\} \quad i = 1, \dots, m \tag{249b}$$

is equivalent to

$$\underset{\mathbf{x}}{\text{minimize}} \quad \mathbf{c}^T\mathbf{x} \tag{250a}$$
$$\text{subject to} \quad \hat{\mathbf{a}}_i^T\mathbf{x} + \left\|\mathbf{R}_i^T\mathbf{x}\right\|_2 \leq b_i \quad i = 1, \dots, m \tag{250b}$$

18.3 Useful Problems

$$\text{average}(\mathbf{v}) = \min_{x\in\mathbb{R}} \|\mathbf{v} - x\mathbf{1}\|_2^2 \tag{251}$$
$$\text{median}(\mathbf{v}) = \min_{x\in\mathbb{R}} \|\mathbf{v} - x\mathbf{1}\|_1 \tag{252}$$

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19.1 Gram-Schmidt

TODO

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