Language Design meets Verifying Compilers

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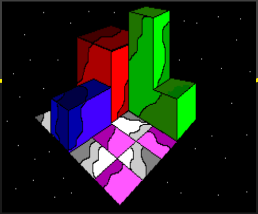
ConsenSys

@WhileyDave
whileydave.com

```
class sprlib {
public:
   sprlib(char *,int = 0);
   ~sprlib();
   void setscreenptr(word);
   void drawspr(int,int,int);
   void xchgspr(int,int,int,int);
   ...
};
```



(Circa 1995)

History

Friday, 24th, June

Checking a large routine. by Dr. A. Turing.

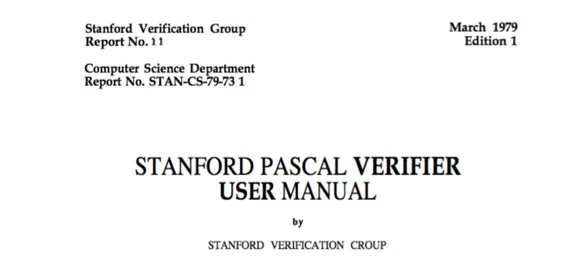
How can one check a routine in the sense of making sure that it is right?

In order that the man who checks may not have too difficult a task the programmer should make a number of definite assertions which can be checked individually, and from which the correctness of the whole program easily follows.

–Alan Turing, 1949

"it was the **Stanford Pascal Verifier** *project that produced the first verification system to target a real programming language"*

–Ireland'04



Stanford Verification Group
Report No. 11

March 1979
Edition 1

Computer Science Department
Report No. STAN-CS-79-731

STANFORD PASCAL VERIFIER
USER MANUAL

by

STANFORD VERIFICATION CROUP

The Verifying Compiler: A Grand Challenge for Computing Research

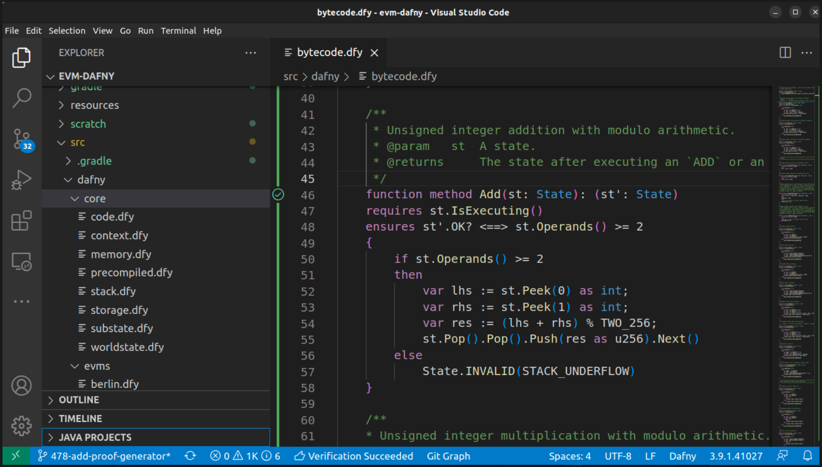
Abstract. This contribution proposes a set of criteria that distinguish a grand challenge in science or engineering from the many other kinds of short-term or long-term research problems that engage the interest of scientists and engineers. As an example drawn from Computer Science, it revives an old challenge: **the construction and application of a verifying compiler that guarantees correctness of a program before running it**.

–Hoare'03

```
deposit(...)
{
  while C1 {
    if C2 { return; }
    ...
  }
  // As the loop should always end prematurely with the 'return'
  // statement, this code should be unreachable. We assert 'false'
  // just to be safe.
  assert (false);
}
```

–Cassez, *et al.*, FM'21

(contract currently holds around 9million ETH)


bytecode.dfy - evm-dafny - Visual Studio Code
File Edit Selection View Go Run Terminal Help
EXPLORER
EVM-DAFNY
resources
scratch
src
.gradle
dafny
core
code.dfy
context.dfy
memory.dfy
precompiled.dfy
stack.dfy
storage.dfy
substate.dfy
worldstate.dfy
evms
berlin.dfy
OUTLINE
TIMELINE
JAVA PROJECTS
bytecode.dfy
src > dafny > bytecode.dfy
/**
* Unsigned integer addition with modulo arithmetic.
* @param st A state.
* @returns The state after executing an `ADD` or an
*/
function method Add(st: State): (st': State)
requires st.IsExecuting()
ensures st'.OK? <==> st.Operands() >= 2
{
if st.Operands() >= 2
then
var lhs := st.Peek(0) as int;
var rhs := st.Peek(1) as int;
var res := (lhs + rhs) % TWO_256;
st.Pop().Pop().Push(res as u256).Next()
else
State.INVALID(STACK_UNDERFLOW)
}
/**
* Unsigned integer multiplication with modulo arithmetic.
478-add-proof-generator*
0 1K 6
Verification Succeeded
Git Graph
Spaces: 4
UTF-8
LF
Dafny
3.9.1.41027

Dafny

```
function abs(x:int) : (r:int)
ensures r >= 0
ensures (x == r) || (-x == r) {
    if x >= 0 then x else -x
}
```

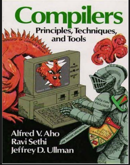
Whiley

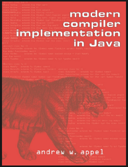
```
function abs(int x) -> (int r)
ensures r >= 0
ensures (r == x) || (r == -x):
  if x >= 0:
    return x
  else:
    return -x
```

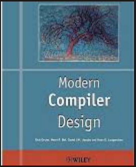
(Verifying) Compilers

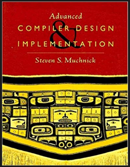
"In computing, a **compiler** is a computer program that translates computer code written in one programming language (the source language) into another language (the target language)."

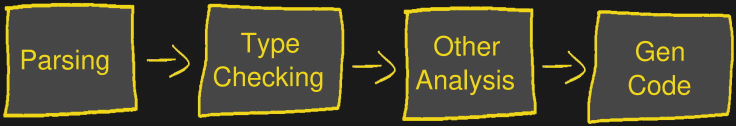
—Wikipedia










Parsing
Type Checking
Other Analysis
Gen Code


Parsing
Verification
Gen
Code


Parsing
Type Checking
Other Analysis
Verification
Gen Code

Dafny

```
datatype Option = Some(value:int) | None

method unboxer(x:int, p:Option) returns (r:int)
requires x >= 0 ==> p.Some? {
  if x >= 0 {
    return p.value;
  } else {
    return x;
  }
}
```

Whiley

```
type Some is {int value}
type Option is Some | null

function unboxer(int x, Option p) -> (int r)
requires (x >= 0) ==> (p is Some):
  if x >= 0:
    // Error!
    return p.value
  else:
    return x
```

Whiley

```
type Some is {int value}
type Option is Some | null

function unboxer(int x, Option p) -> (int r)
requires (x >= 0) ==> (p is Some):
  if x >= 0:
    assert p is Some
    return p.value
  else:
    return x
```

Dafny

```
method maxer(x:int, y:int) returns (r:int)
requires x != y
ensures (r == x) || (r == y) {
    if x > y {
        return x;
    } else if x < y {
        return y;
    }
}
```

Whiley

```
function maxer(int x, int y) -> (int r)
requires x != y:
  if x > y:
    return 1
  else if x < y:
    return 0
  else:
    fail
```

Flow Typing

> “**Flow typing** offers an alternative where a variable may have different types at different program points. The technique is adopted from flow-sensitive program analysis ...”
>
> —Pearce’13

Dafny

```
method iof(xs:seq<int>, x:nat) returns (r:Opt<nat>)
// If valid index returned, element matches item
ensures r.Some? ==> (r.v < |xs| && xs[r.v] == x)
```

Whiley

```
method iof(int[] xs, int x) -> (int|null r)
// If valid index returned, element matches item
ensures (r is int) ==> (xs[r] == x)
```

Functional Purity

"To be functionally pure, a method must satisfy two critical properties: **First**, it must have no side effects. ... The **second** property is functional determinism."

—Finifter *et al.*, 2008

Dafny

```
function max(x:int, y:int) : (r:int)
ensures (r == x) || (r == y)
ensures (r >= x) && (r >= y) {
    ...
}
```

Whiley

```
function max(int x, int y) -> (int r)
ensures (r == x) || (r == y)
ensures (r >= x) && (r >= y):
    ...
```

"Unlike pure functional programming, however, **mutable value semantics** allows part-wise in-place mutation, thereby eliminating the memory traffic usually associated with functional updates of immutable data"

—Racordon *et al.*, 2022

Dafny

```
function fill(xs:seq<int>, n:nat, x:int) : seq<int>
requires n <= |xs|
{
   if n == 0 then xs
   else [x] + fill(xs[1..],n-1,x)
}
```

Whiley

```
function fill(int[] xs, uint n, int x) -> (int[] rs)
requires n <= |xs|:
   for i in 0..n:
     xs[i] = x
   return xs
```

(Un)interpreted Functions

"Normally function bodies are transparent and available for constructing proofs of assertions that use those functions. However, sometimes it is helpful to mark a function `{:opaque}` and treat it as an **uninterpreted function**, whose properties are just its specifications."

—Dafny Reference Manual

"In mathematical logic, an **uninterpreted function** or function symbol is one that has no other property than its name and n-ary form."

—Wikipedia

"**Uninterpreted functions** are used for abstracting, or generalizing, theorems. Unlike other function symbols, they should not be interpreted as part of a model of a formula."

—Kroening & Strichman

Dafny

```
function zero_f(xs:seq<int>, n:nat) : (r:seq<int>)
requires n <= |xs| { ... }

method zero_m(xs:seq<int>,n:nat) returns(r:seq<int>)
requires n <= |xs| { ... }
```

```
assert zero_f([1,2,3],2) == [0,0,3];

var r := zero_m([1,2,3],2);
assert r == [0,0,3];
```

Whiley

```
property zero_p(int[] xs, uint n) -> (int[] rs)
requires n <= |xs|:
    ...

function zero_f(int[] xs, uint n) -> (int[] rs)
requires n <= |xs|:
    ...

method zero_m(int[] xs, uint n) -> (int[] rs)
requires n <= |xs|:
    ...
```

```
assert zero_p([1,2,3],2) == [0,0,3]
assert zero_f([1,2,3],2) == [0,0,3]
int[] rs = zero_m([1,2,3],2)
assert rs == [0,0,3]
```

Inference

Dafny

```
function contains(xs:seq<int>, x:int) : bool {
    exists k:nat | k < |xs| :: xs[k] == x
}

method find(xs:seq<int>, x:int) returns (r:nat)
requires contains(xs,x)
ensures xs[r] == x
{
    for i := 0 to |xs|
    invariant contains(xs[i..],x) {
        if xs[i] == x { return i; }
    }
    assert false;
}
```

Dafny

```
function contains(xs:seq<int>, x:int) : bool {
    exists k:nat | k < |xs| :: xs[k] == x
}

method find(xs:seq<int>, x:int) returns (r:nat)
requires contains(xs,x)
ensures  r < |xs| && xs[r] == x
{
    for i := 0 to |xs|
    invariant contains(xs[i..],x) {
        if xs[i] == x { return i; }
    }
    assert false;
}
```

Dafny

```
method indexOf(xs: seq<int>, x:int) returns (r:nat)
ensures r <= |xs|
ensures r < |xs| ==> xs[r] == x {
    var i : int := 0;

    while i < |xs|
    invariant i <= |xs| {
        if xs[i] == x { break; }
        i := i + 1;
    }
    return i;
}
```

Whiley

```
function indexOf(int[] xs, int x) -> (int r)
ensures (r != |xs|) ==> (xs[r] == x):
    int i = 0

    while i < |xs|
    where i <= |xs|:
        if xs[i] == x:
            break
        i = i + 1
    return i
```

Whiley

```
function indexOf(int[] xs, int x) -> (int r)
ensures (r != |xs|) ==> (xs[r] == x):
    int i = 0

    while i < |xs|
    where i >= 0 && i <= |xs|:
        if xs[i] == x:
            break
        i = i + 1
    return i
```

Language Features?

Whiley

```
function copy(int[] xs, uint n) -> (int[] ys)
ensures |ys| == n:
    // Create array of given size
    ys = [0; n]
    // Copy over what we can
    for i in 0..min(n,|xs|)
    where n == |ys|:
        ys[i] = xs[i]
    // Done
    return ys
```

Whiley

```
type Box<T> is &T

method destroy(Box<T> p):
    //
    delete p
```

Whiley

```
type Box<T> is &T

method destroy(Box<T> p)
requires #p == 1:
    //
    delete p
```

Whiley

```
type Box<T> is &T where #p == 1

method destroy(Box<T> p):
    //
    delete p
```

ShriramKrishnamurthi
@ShriramKMurthi

10/ The next generation of computing problems will not be about writing 80s style 5-line for-loops. It'll be about properties, specification, reasoning, verification, prompt eng, synthesis, etc. How will we get there?

And no, I will not be taking questions. (-:

http://whiley.org

@WhileyDave
http://github.com/Whiley

"It's this question that leads to the **Closet Contract Conjecture***: are the contracts of Eiffel libraries a figment of the Eiffel programmer's obsession with this mechanism? Or are they present anyway, hidden, in non-Eiffel libraries as well?"*

Arnout & Meyer, 2002