	Student information	Date	Number of session
Algorithmics	UO: UO300535	27-02-2025	3
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Activity 1. D&C by Subtraction

After analyzing the complexity of the three previous classes, you are not asked to make the timetables, but to reason whether the times match the theoretical time complexity of each algorithm.

The algorithms in Subtraction1.java and Subtraction2.java are respectively linear and quadratic in their time complexity, although the empirical data is so small in the first one as to make the trend unnoticeable, and Subtraction2.java looks linear in the first few iterations; we'd have to measure them with a lot greater problem sizes for their real growth trends to become apparent.

For what value of n do the Subtraction1 and Subtraction2 classes stop giving times (we abort the algorithm because it exceeds 1 minute)? Why does that happen?

They both seem to overflow once n > 8192, which is way before the times get even close to a minute. This happens because of the massive amount of times the method calls itself, therefore pushing more information to the stack than it can actually hold.

```
terminated > Subtraction1 [Java Application] C:\Program Files\Java\jdk-22\bin\javaw.exe
n=64**TIME=0**cont=1
n=128**TIME=0**cont=1
n=256**TIME=0**cont=1
n=512**TIME=0**cont=1
n=1024**TIME=0**cont=1
n=2048**TIME=1**cont=1
n=4096**TIME=0**cont=1
n=8192**TIME=0**cont=1
                                               Subtraction1.java:10
                                               Subtraction1.java:15
                                               Subtraction1.java:15
                                               Subtraction1.java:15
```

	Student information	Date	Number of session
Algorithmics	UO: UO300535	27-02-2025	3
J	Surname: Cabo Stroup		
	Name: José David		

```
<terminated> Subtraction2 [Java Application] C:\Program Files\Java\jdk-22\bin\javaw.exe
n=64**TIME=0**cont=128
n=128**TIME=0**cont=256
n=256**TIME=0**cont=512
n=512**TIME=1**cont=1024
n=1024**TIME=4**cont=2048
n=2048**TIME=13**cont=4096
n=4096**TIME=52**cont=8192
n=8192**TIME=214**cont=16384
                                              Subtraction2.java:10
                                              Subtraction2.java:16
                                              Subtraction2.java:16
                                              Subtraction2.java:16
```

How many years would it take to complete the Subtraction3 execution for n=80? Reason the answer.

Subtraction3.java, however, is exponential (O(2 n)). Taking [n $_1$ = 30, t $_1$ = 38,051 ms] as a reference point, we get that:

$$t_2 = \frac{2^{n_2}}{2^{n_1}} \times t_1 = 2^{n_2 - n_1} \times t_1 = 2^{80 - 30} \times t_1 = 2^{50} \times 38051 \approx 4.28 \times 10^{19} \, \text{ms}$$

Which is about 1,358.5 million years. Efficient!

- Implement a Subtraction4.java class with a complexity O(n3) and then fill in a table showing the time (in milliseconds) for n=100, 200, 400, 800, ... (until OoT).
- Implement a Subtraction5.java class with a complexity O(3n/2) and then fill in a table showing the time (in milliseconds) for n=30, 32, 34, 36, ... (until OoT).

The times for Subtraction4.java and Subtraction5.java are as follows:

	Student information	Date	Number of session
Algorithmics	UO: UO300535	27-02-2025	3
J	Surname: Cabo Stroup		
	Name: José David		

	Time (ms)		Time (ms)
n	Subtraction4.java	n	Subtraction5.java
100	LoR	30	373
200	LoR	32	1083
400	81	34	3235
800	615	36	9934
1600	4870	38	29184
3200	39568	40	89494
6400	OoT	42	OoT
12800	ОоТ	44	OoT
25600	ОоТ	46	ОоТ
51200	OoT	48	ОоТ

How many years would it take to complete the Subtraction5 execution for n=80?
 Reason the answer.

For n=80, Subtraction5.java would take

$$t_2 = \frac{3^{\frac{n_2}{2}}}{3^{\frac{n_1}{2}}} \times t_1 = 3^{\frac{n_2}{2} - \frac{n_1}{2}} \times t_1 = 3^{\frac{80}{2} - \frac{40}{2}} \times t_1 = 3^{20} \times 89494 \approx 3.12 \times 10^{14} \, ms$$

Or approximately **9,894.9 years.** Hey, at least it's better than Subtraction3.java...

Activity 2. D&C by Division

 After analyzing the complexity of the three previous classes, you are not asked to make the timetables, but to reason whether the times match the theoretical time complexity of each algorithm.

Indeed, they do. Division1.java and Division3.java are linear, as the times would suggest. Division2.java, on the other hand, is O(n*log n), which is also coherent with the data when measured.

	Student information	Date	Number of session
Algorithmics	UO: UO300535	27-02-2025	3
	Surname: Cabo Stroup		
	Name: José David		

- Implement a Division4.java class with a complexity O(n^2) (with a<b^k) and then fill in a table showing the time (in milliseconds) for n=1000, 2000, 4000, 8000, ... (up to OoT).
- Implement a Division5.java class with a complexity O(n^2) (with a>b^k) and then fill in a table showing the time (in milliseconds) for n=1000, 2000, 4000, 8000, ... (up to OoT).

	Time (ms)		Time (ms)
n	Division4.java	n	Division5.java
100	LoR	10	D LoR
200	LoR	20	98
400	66	400	390
800	253	80	1510
1600	1042	160	6046
3200	4059	320	23673
6400	16387	640	O oT
12800	OoT	1280	O OT
25600	OoT	2560	O OT
51200	OoT	5120	O OT
25600	ОоТ	2560	O OT

	Student information	Date	Number of session
Algorithmics	UO: UO300535	27-02-2025	3
Ü	Surname: Cabo Stroup		
	Name: José David		

	Fibonacci Times (ms)					Fibonacci 7	Times (ms)			
n		fib1	fib2	fib3	fib4	n	fib1	fib2	fib3	fib4
	10	0,000091	0,000127	0,000209	0,00261	35	0,000187	0,000398	0,000602	433
	11	0,000089	0,000138	0,000216	0,00417	36	0,0002	0,000432	0,000651	696
	12	0,000091	0,000148	0,000232	0,00679	37	0,000197	0,00044	0,000637	1129
	13	0,000097	0,000158	0,000249	0,01148	38	0,0002	0,000429	0,000646	1857
	14	0,000099	0,000168	0,000264	0,01846	39	0,000213	0,000446	0,000669	3000
	15	0,000107	0,000178	0,000278	0,03078	40	0,000215	0,000449	0,000683	4920
	16	0,000109	0,000188	0,000293	0,04917	41	0,000221	0,00046	0,000703	8030
	17	0,000111	0,000199	0,000312	0,07865	42	0,000224	0,000476	0,000719	12737
	18	0,000117	0,000209	0,000346	0,12842	43	0,000235	0,000487	0,000733	20344
	19	0,00012	0,000221	0,000371	0,2027	44	0,000232	0,000493	0,000755	33088
	20	0,000128	0,000233	0,000361	0,318	45	0,000227	0,000508	0,000753	54959
	21	0,00013	0,000242	0,000389	0,516	46	0,000234	0,000518	0,000774	OoT
	22	0,000132	0,000255	0,000401	0,829	47	0,000242	0,000531	0,000788	OoT
	23	0,000137	0,000266	0,000409	1,364	48	0,000238	0,000542	0,000808	OoT
	24	0,00014	0,000278	0,000425	2,196	49	0,000245	0,000577	0,00083	OoT
	25	0,000158	0,000289	0,000434	3,561	50	0,000266	0,000572	0,000843	OoT
	26	0,000148	0,000299	0,000451	5,754	51	0,000253	0,000589	0,000849	OoT
	27	0,000152	0,00031	0,000505	9,204	52	0,000257	0,000586	0,000867	OoT
	28	0,000157	0,00032	0,000486	14,901	53	0,00026	0,00059	0,000875	OoT
	29	0,000162	0,00033	0,000501	24,103	54	0,000266	0,000601	0,00091	OoT
	30	0,000162	0,000342	0,000524	38,866	55	0,000269	0,000606	0,000945	OoT
	31	0,000172	0,000357	0,000532	63,765	56	0,000273	0,000616	0,000934	OoT
	32	0,000173	0,000364	0,00055	102	57	0,000277	0,000627	0,000947	OoT
	33	0,000178	0,000383	0,000577	167	58	0,00028	0,000643	0,000953	OoT
	34	0,00018	0,000392	0,000588	266	59	0,000287	0,000651	0,000971	OoT

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	VectorSum Times (ms)						
n	sum1	sum2	sum3				
3	0,0000395	0,000077	0,000088				
6	0,0000623	0,000117	0,00016				
12	0,0000856	0,000234	0,000345				
24	0,0001314	0,000427	0,000704				
48	0,0002226	0,000811	0,001448				
96	0,0004046	0,001598	0,002894				
192	0,0007699	0,003182	0,005747				
384	0,0015166	0,006476	0,011674				
768	0,00296	0,012749	0,023393				
1536	0,00586	0,025404	0,0473				
3072	0,01175	0,0496	0,0977				
6144	0,02389	0,0992	0,1881				
12288	0,04725	OVERFLOW	0,3753				
24576	0,09443	OVERFLOW	0,7499				
49152	0,18808	OVERFLOW	1,5232				
98304	0,37758	OVERFLOW	3,0049				

After analyzing the complexity of the various algorithms within the two classes, executing them and afterwards putting the times obtained in a table, compare the efficiency of each algorithm.

For Fibonacci, the algorithms are of decreasing efficiency: the first three are O(n), whereas the last one is $O(1.6^n)$. The VectorSum algorithms, on the other hand, are all linear, with a clear preference for the first one. The second one might be better than the third one if it didn't overflow, but alas, I guess we'll never know.