Math 239 - Lecture #30

Recall: If n33, then any connected planar graph has at most 3n-6 edges.

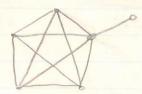
A) The converse of this is not true. A graph with &3n-6 edges might not be planar. (Example: Peterson) Corollary. Ks is not planar

Proof:

Ks has 5 vertices, and any planar graph on 5 vertices has at most (3.5)-6=9 edges. But Kz has 10 edges, so Kz is not planar. 1



(A) Counterexample:



6 vertices 11 edges 5 3n-6 6.3-6=12; not planar



6 vertices 9 odges < 3n-6=12

Theorem.

Proof_

when 123, a connected planar bipartite graph on n vertices has at most 2n-4 edges. Consider a planar embedding of a bipartite graph G with n vertices, m edges, s faces. We claim that every face has deg > 4. We proved the tree case in previous lesson.

If there is a cycle, then all face boundaries contain cycles of length >4 (since no is exist in a bipartite graph). So every face has deg 34.

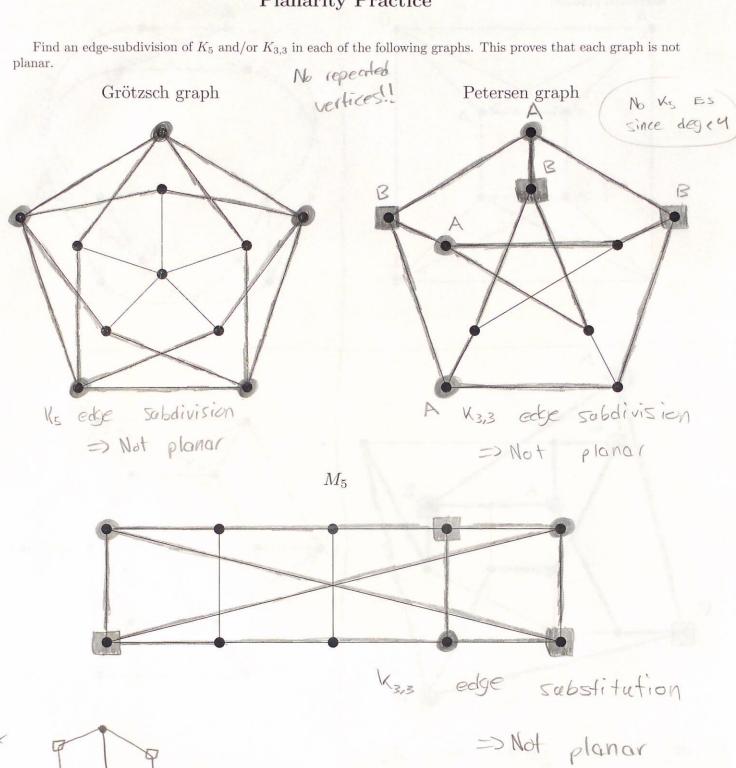
By HSLFF, 2m = \(\frac{1}{2} \deg(f) \) \(\gamma \text{US} \)

= 4(2-n+m) by EF (n-m+s=2) = 8- Un+UM

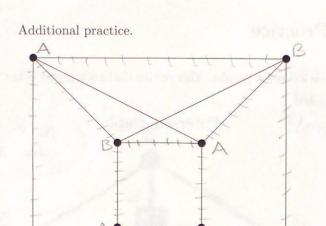
So 2n-4 > m. 0

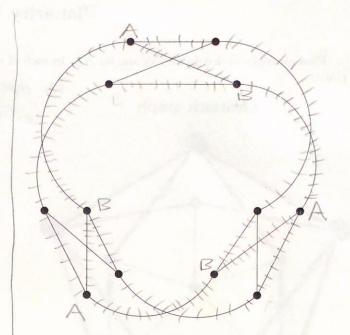
Corollary: K3.3 is not planar. K3,3 has 6 vertices. Any planar bipartite graph with Proof: 6 vertices has at most (2.6)-4=8 edges. But K3,3 has 9 edges, so it is not planar. I Kuratows Ki's Theorem Adding vertices of deg 2 doesn't after planarity Def: An edge subdivision of F is obtained by introducing some new vertices of degree 2 to the edges of G.
Could be O new vertices . We are replacing each edge with a new path A graph is not planar if and only if any edge subdivision of the graph is not planar. Kuratowski's theorem says that a graph is planar if and only if it does not contain an edge subdivision of Ks or Ks.3 as a subgraph.

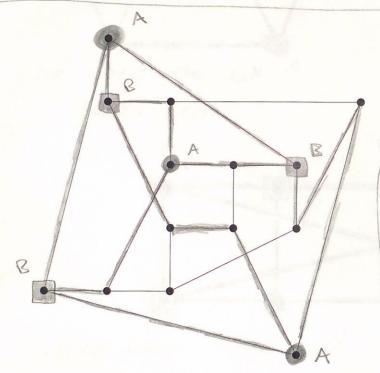
Planarity Practice

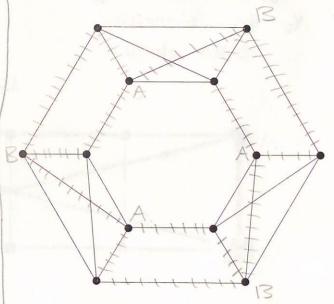


K3.3 also!









All deg 3, not gonna have Ks

K3,3 edge substitution

Not Planar