Math 239 - Lecture # 29

Definition:

A connected planar graph is platonic if every vertex has an embedding where every vertex has the same degree (>3) and every face has the same degree (>3).

We can turn a planar embedding into a polyhedron by drawing it on a sphere and cut off the faces.

Suppose a platonic graph has vertex deg do 33 and face deg do 33, a vertices, m edges, s faces.

- 1 Hondshaking Lemma: nodu = zm = n = zm
- 2) HSL for faces: S. d= Zm = S = Zm df
- 3 Euler's Formula: n-m+s=2
 - > 2m m + 2m = Z, mult by duds
 - => 2mdf mdrdf + 2mdr = Zdrdf > 0, divide out m
 - => Zdf drdf + Zdr > 0
 - => 2df drdf + 2dr + 4-4>0 factor
 - =) (dv-Z)(df-Z)+4 >0
 - => (dn-Z)(dx-Z)<4 * (dn,dx3,3)

Possible (du, df) pairs

-	A STATE OF THE PARTY OF THE PAR		5	These are all the
df	3,4,5	3	3	possible pairs for
	4> df=2	Z> d=-Z	43 > df-Z	Platonic graphs



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Platonic Graphs

Platonic: (1) dr=3, dr=3

Tetrahedron Am

2) dv = 3, df = 4

Cube/Hexahedran
(d6)

n=8 m=12

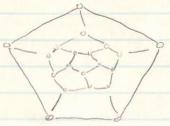
F= 6

(3) dr=4, df=3

n=6 m=12S=8

Octahedron (18)

(4) dr=3, dr=5

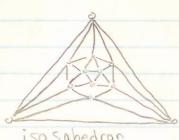


Dodedhedron

N = 20 M = 30S = 12

(212)

(5) dr=5, dq=3



iso sahedran

n=12 m=30S=20

Non-Planar Graphs

Theorem: When n>3, a connected planar graph on n vertices has at most 3n-6 edges.

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Proof:

Consider a connected embedding with a vertices, medges, s faces. We claim that every face has deg at least 3.

If G does not have a cycle, then G is a tree, which has only I face since n=3, G has at least 2 edges, so this face has deg at least 4.

If F contains a cycle, then every face must contain a cycle on its boundary (to seperale it from the other faces). So each face has deg > 3.

Using HSLFF, $z_m = \sum_{f \in F} deg(f) > 35$

=3(2-n+3) by =F(n-m+s=2)

= 6 - 3n + 3m

s. m < 3n-6. 0

