

Math 239 - Lecture #13

Recall ambiguity within the context of binary strings.

Three unambiguous expressions for the set of all strings:

All:
binstr

(1) $\{0,1\}^*$ Every string can be uniquely generated by decomposing after each bit. I.E. $\frac{0/1/0/1}{\in \{0,1\}}$

(2) $\{0\}^* (\{1\} \{0\}^*)^*$ Every string with k 1's has a form $0^{a_1} 1 0^{a_2} 1 \dots 1 0^{a_{k+1}}$; break just before each 1.
 $\underbrace{0^{a_1}}_{\in \{0\}^*} \underbrace{1 0^{a_2}}_{\in \{1\} \{0\}^*} \dots \underbrace{1 0^{a_{k+1}}}_{\in \{1\} \{0\}^*}$

(3) $\{0\}^* (\{1\} \{1\}^* \{0\} \{0\}^*)^* \{1\}^*$ block decomposition
 $\underbrace{000}_{\in \{0\}^*} \underbrace{111100}_{\in \{1\} \{1\}^* \{0\} \{0\}^*} \underbrace{100}_{\in \{1\} \{1\}^* \{0\} \{0\}^*} \underbrace{1110}_{\in \{1\} \{1\}^* \{0\} \{0\}^*} \underbrace{10011}_{\in \{1\} \{1\}^* \{0\} \{0\}^*}$

- Break up each block of 1's and 0's.
- Leading 0's and trailing 1's could be empty.

Problems on restriction of substrings. Take one of the unambiguous decompositions for all strings, and remove/restrict parts of it. The resulting expression is still unambiguous.

Example: Let S be the set of all strings with no 3 consecutive 0's. (no 000 as a substring)

Start with $\{0\}^* (\{1\} \{0\}^*)^*$

- Where can we find '000' in this expression?

Remove instances of '000' from head. So

$$\{0\}^* = \{\epsilon, 0, 00, 000, 0000, 00000, \dots\}$$

So replace $\{0\}^*$ with $\{\epsilon, 0, 00\}$.

Do the exact same thing to the inner $\{0\}^*$.

$$\therefore S = \{\epsilon, 0, 00\} (\{1\} \{\epsilon, 0, 00\})^*$$

Cont

$$\Phi_{\{\epsilon, 0, 00\}}(x) = 1 + x + x^2, \quad \Phi_{\{1\}}(x) = x$$

$$\Phi_S = (1 + x + x^2) \cdot \frac{1}{1 - (x \cdot (1 + x + x^2))} = \frac{1 + x + x^2}{1 - x - x^2 - x^3}$$

The # of str of length n with no 000 is $[x^n] \frac{1 + x + x^2}{1 - x - x^2 - x^3}$

Example: Let S be the set of all strings where every block has length of at least 2.

Start with $\{0\}^* (\{1\}\{1\}^* \{0\}\{0\}^*)^* \{1\}^*$.

$\{0\}^* = \{\epsilon, 0, 00, 000, 0000, \dots\}$ remove blocks length < 2 .

- We keep ϵ since it isn't a block, it's just empty!
- Similar for tail $\{1\}^*$.

$\{1\}\{1\}^* = \{1, 11, 111, 1111, \dots\}$ just remove 1. No epsilon!

- Similar for $\{0\}\{0\}^*$.

$$\text{So } S = (\{\epsilon\} \cup \{00\}\{0\}^*) (\{11\}\{1\}^* \{00\}\{0\}^*)^* (\{1\}\{1\}^*)$$

$$\begin{aligned} \therefore \Phi_S(x) &= \left(1 + x^2 \cdot \frac{1}{1-x}\right) \left(\frac{1}{1 - (x^2 \frac{1}{1-x} \cdot x^2 \frac{1}{1-x})}\right) \left(1 + x^2 \cdot \frac{1}{1-x}\right) \\ &= \frac{1 - x + x^2}{1 - x - x^2} \end{aligned}$$

Example: Let S be the set of all str where an even block of 0's cannot be followed by an odd block of 1's.

Cont next page

Math 239 - Lecture #13 Cont.

Example:ContStart with $\{1\}^* (\{0\}\{0\}^* \{1\}\{1\}^*)^* \{0\}^*$

- Modified block - swapped the 0's and 1's.

We can only find what we're looking to exclude in the middle partition. Break into 2 cases.

(1) Even # of 0's : $\{00\}\{00\}^* + \text{even 1's, } \{11\}\{11\}^*$ (2) Odd # of 0's : $\{0\}\{00\}^* + \text{any 1's, } \{1\}\{1\}^*$

$$\text{So } S = \{1\}^* \left(\{00\}\{00\}^* \{11\}\{11\}^* \cup \{0\}\{00\}^* \{1\}\{1\}^* \right)^* \{0\}^*$$

$$\begin{aligned} \overline{\Phi}_S(x) &= \frac{1}{1-x} \left(\frac{1}{1 - \left(x^2 \frac{1}{1-x^2} \cdot x^2 \frac{1}{1-x^2} + x \frac{1}{1-x^2} \cdot x \frac{1}{1-x} \right)} \right) \frac{1}{1-x} \\ &= \frac{1+2x+x^2}{1-3x^2-x^3} \end{aligned}$$

$$x - x^3 + x^2 - x^3 = \frac{x + x^2 - 2x^3}{(1-x)(1-x^2)}$$

$$\frac{1-x-x^2+x^3-x-x^2+2x^3}{(1-x)(1-x^2)}$$

$$= \frac{1-2x-2x^2-2x^3}{(1-x)(1-x^2)}$$