6.034: Lecture 5

Adversarial search & games Can computers play chess?

Professor Robert C. Berwick

Chess (Yes, Chess) Is Now a Streaming Obsession

Hikaru Nakamura Viewers are flocking to games during the pandemic, entranced by a charismatic grandmaster and his lightning-fast play. 19:26:16 Chess.com 6

Wins/Losses/Draws: 286 / 11 / 2





Al Methods

- Problem solving
 - ☐ G+T, search, optimal search, games, constraint satisfaction
- □ Inference
 - □ rule-based systems, Bayesian inference
- ☐ Machine learning
 - k-nearest neighbors, id trees, neural nets, deep neural nets, support vector machines, genetic algorithms, near miss/one-shot
- □ Communication, perception, action
 - □ natural language processing, vision, robotics

Menu for today

Games: Adversarial search

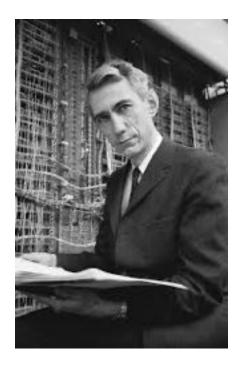
- ☐ Intelligence & Chess
- ☐ Ways to Play Chess
- ☐ How to search: Minimax algorithm
- ☐ Improving search: Alpha-beta pruning
- ☐ Progressive Deepening
- ☐ Reflections on contemporary chess computers & Gold Star ideas

A very brief history

Philosophical Magazine, Ser.7, Vol. 41, No. 314 - March 1950.

XXII. Programming a Computer for Playing Chess¹
By CLAUDE E. SHANNON

Bell Telephone Laboratories, Inc., Murray Hill, N.J.² [Received November 8, 1949]



Claude Shannon (1916-2001)

1. INTRODUCTION

This paper is concerned with the problem of constructing a computing routine or "program" for a modern general purpose computer which will enable it to play chess. Although perhaps of no practical importance, the question is of theoretical interest, and it is hoped that a satisfactory solution of this problem will act as a wedge in attacking other problems of a similar nature and of greater significance. Some possibilities in this direction are: -

The chess machine is an ideal one to start with, since: (1) the problem is sharply defined both in allowed operations (the moves) and in the ultimate goal (checkmate); (2) it is neither so simple as to be trivial nor too difficult for satisfactory solution; (3) chess is generally considered to require "thinking" for skilful play; a solution of this problem will force us either to admit the possibility of a mechanized thinking or to further restrict our concept of "thinking"; (4) the discrete structure of chess fits well into the digital nature of modern computers.

ALCHEMY AND ARTIFICIAL INTELLIGENCE Hubert L. Dreyfus

December 1965

remains unimproved. Burton Bloom at M.I.T. has made the latest attempt to write a chess program; like all the others, it plays a stupid game. In fact, in the nine years since the Los Alamos program beat a weak player, in spite of a great investment of time, energy, and ink, the only improvement seems to be that a machine now plays poorly on an eight-by-eight rather than a six-by-six board. According to Newell, Shaw, and Simon themselves, evaluating the Los Alamos, the IBM, and the NSS programs: "All three programs play roughly the same quality of chess (mediocre) with roughly the same amount of computing time" [20:14]. Still no chess program can play even amateur chess, and the world championship tournament is only two years away.



by RICHARD D. GREENBLATT, DONALD E. EASTLAKE, III,

and

STEPHEN D. CROCKER

Massachussetts Institute of Technology Cambridge, Massachusetts

INTRODUCTION

Since mid-November 1966 a chess program has been under development at the Artificial Intelligence Laboratory of Project MAC at M.I.T. This paper describes the state of the program as of August 1967 and gives some of the details of the heuristics and algorithms employed.

Fall Joint Computer Conference, 1967





MIT AI Lab MacHack IV

H. Dreyfus vs. MacHack IV



THE ARTIFICIAL INTELLIGENCE OF HUBERT L. DREYFUS

A Budget of Pallacies

by

Seymour Papert

January 1968.

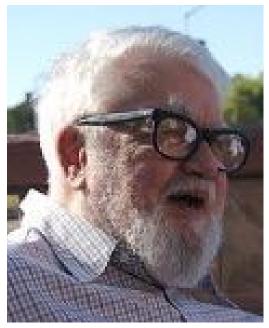
1.5 Computers Can't Play Chess.

1.5.1 Nor Can Dreyfus.

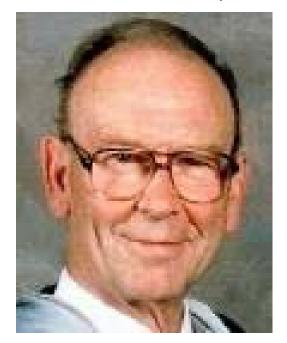
I hope readers will pardon a digression to get a debating point off my chest. They will surely understand why I find it irresistible.

In <u>Alchemy</u> and <u>Phenomenology</u> Dreyfus discusses the weakness of chessplaying programs. He plainly gives the impression that they can typically be defeated by human novices. His twice recounted story of how a ten-year-old child defeated a program constructed by Newell, Shaw and Simon was gleefully quoted by the <u>New Yorker</u> and other popular magazines as demonstrating the futility of Artificial Intelligence. While <u>Phenomenology</u> was in press I had the pleasure of arranging for Dreyfus to play against Richard Greenblatt's chess program at M.I.T. and seeing him very roundly trounced. The newsletter SIGART reprinted the game with no comment beyond one phrase from <u>Alchemy</u>:

[&]quot;... no chess program can play even amateur chess." (p. 10)

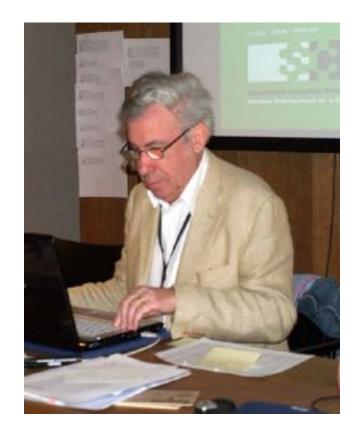


John McCarthy



Donald Michie

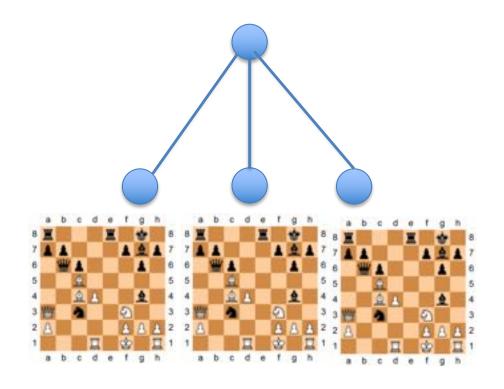
Versus.



Chess Grandmaster David Levy Bet in 1968, 1250 pounds

How could a computer play chess?

- 1. Human like: Describe board via pawn structure, King safety, good time to castle....use tactics, strategy then make a move
- 2. If-then rules make most plausible move
- 3. Look ahead and evaluate



$$S=g(f_1, f_2,...,f_n)$$

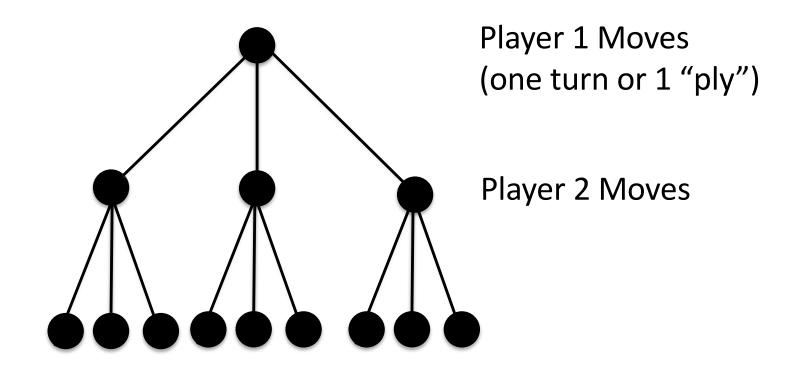
$$= c_1 f_1 + c_2 f_2 + \dots + c_n f_n$$

Linear Scoring Polynomial

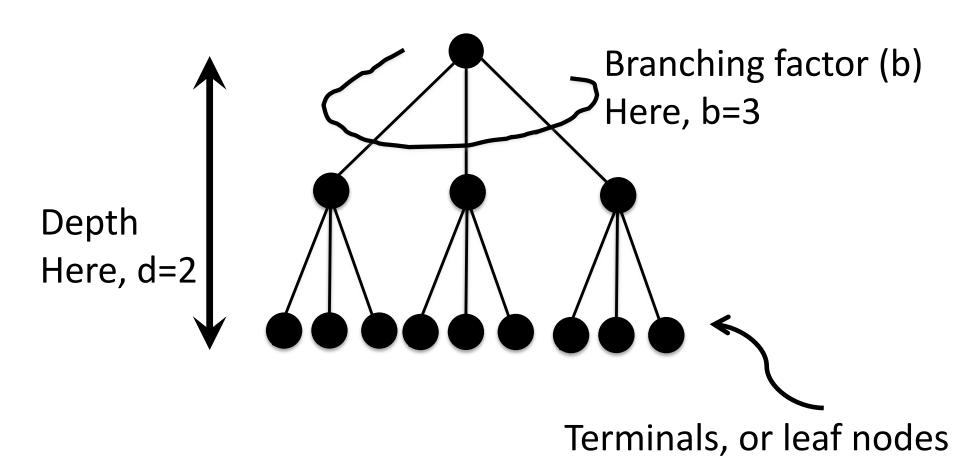
How could a computer play chess?

- 1. Human like
- 2. If-then rules: look & make most plausible move
- 3. Look ahead and evaluate
- 4. British Museum algorithm (exhaustive search)?

Vocabulary for game trees

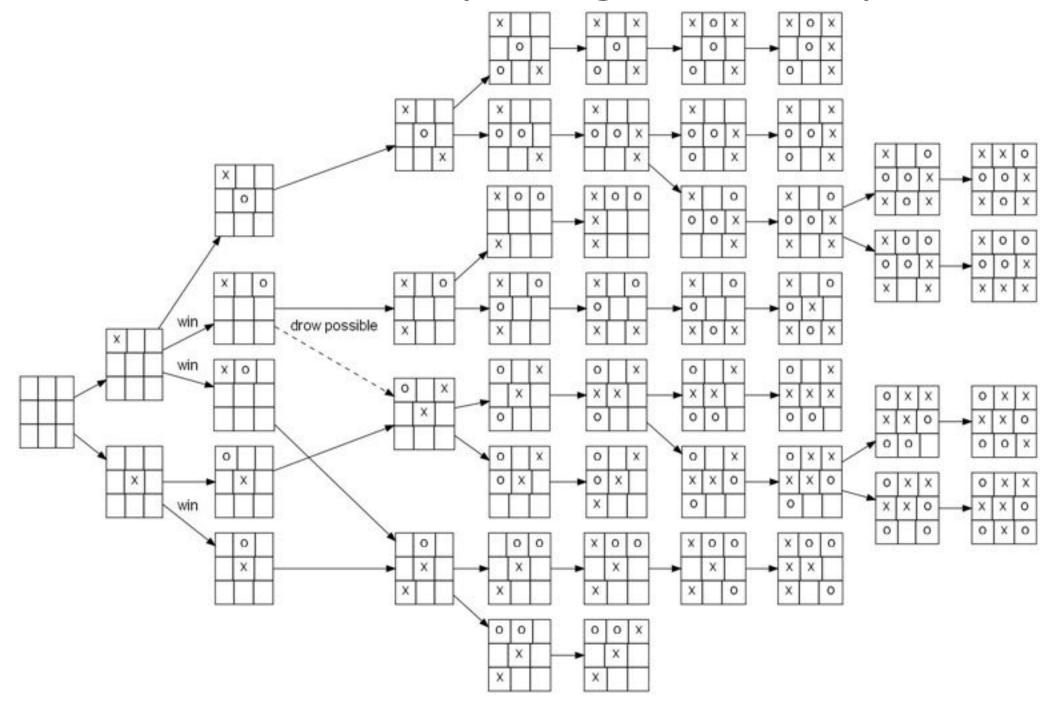


Vocabulary for game trees



 $b^d = 3^2 = 9$

Tic-tac-toe: complete game tree space



Could we use cloud computing to evaluate all possible chess moves in the same way?

10¹²⁰ moves

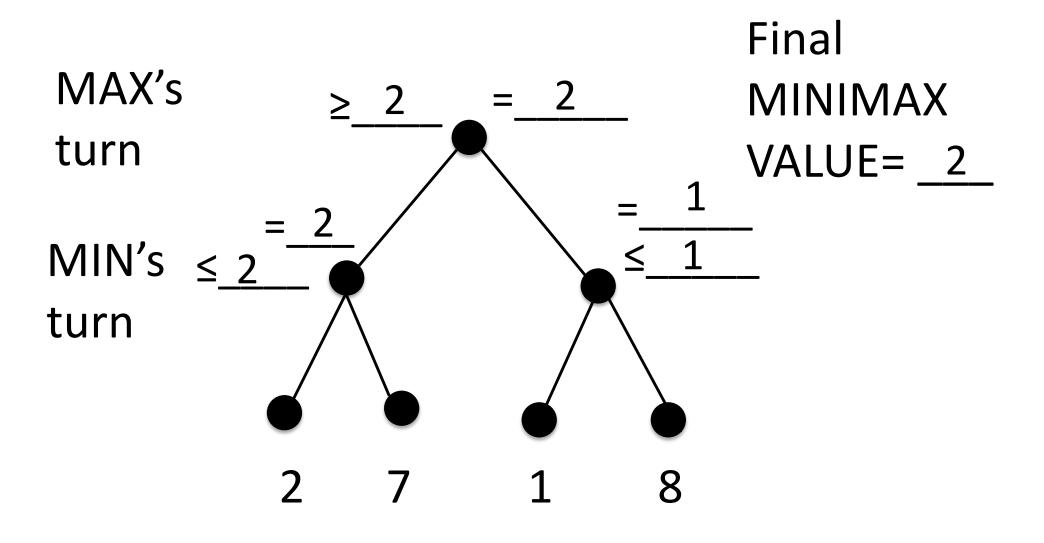
Could we use cloud computing to evaluate all possible chess moves in the same way?

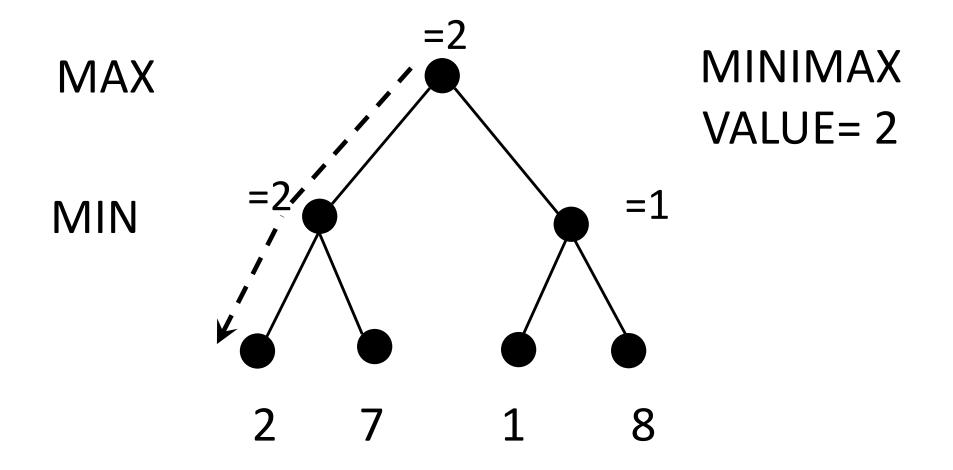
> ≈10²⁶ nanoseconds x 10⁸⁰ atoms ≈10¹⁰⁶ ops, each atom operating @ nanosecond speeds Not enough compute power for 10¹²⁰ moves!

How could a computer play chess?

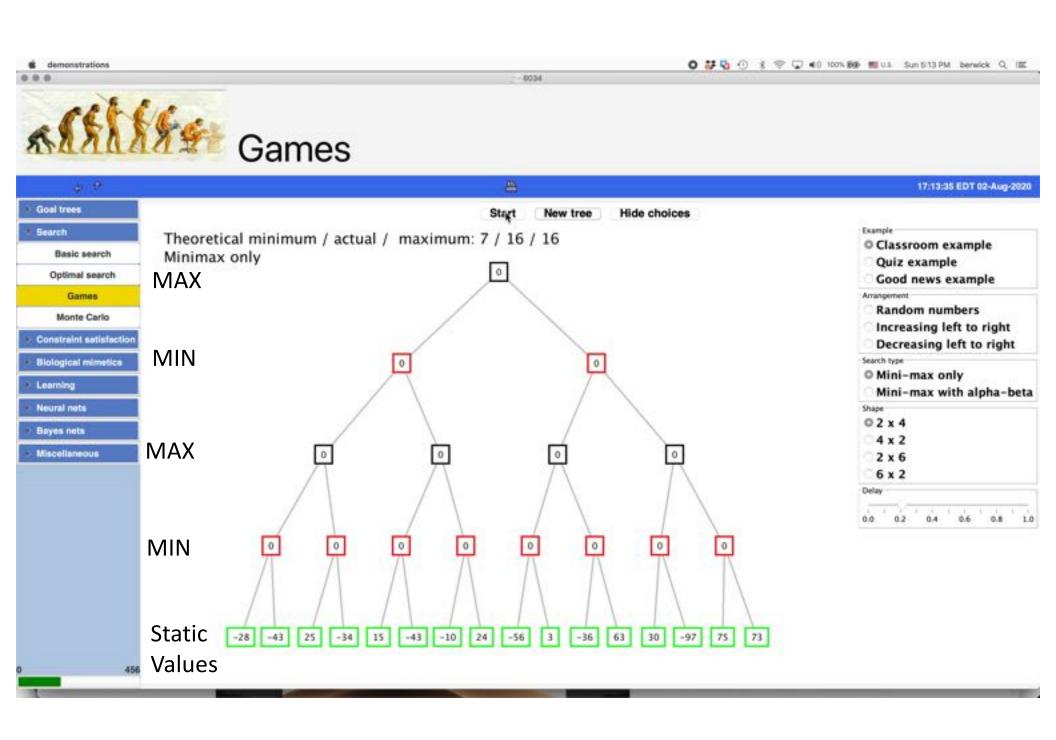
- 1. Human like
- 2. If-then— make most plausible move
- 3. Look ahead and evaluate
- 4. British Museum algorithm (exhaustive search)
- 5. Look ahead as far as possible

Minimax algorithm for searching game tree to find optimal move sequence

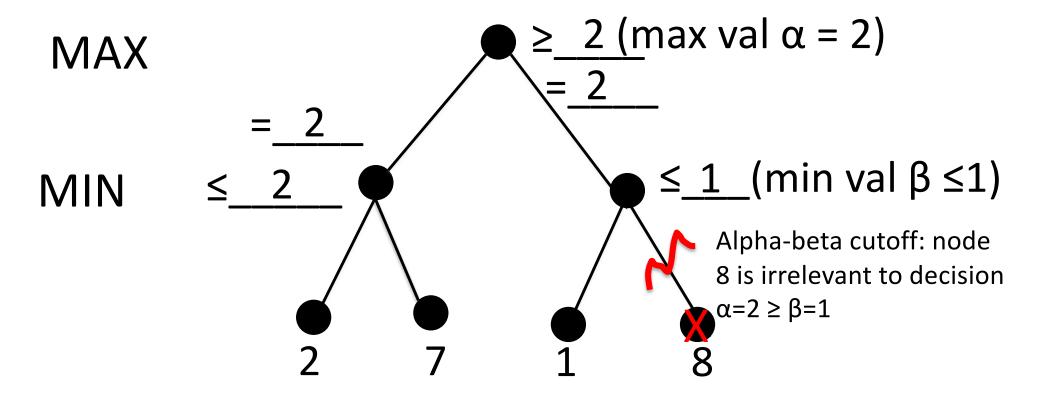




Note that the best minimax path is preserved under any transform of the leaf scores that preserves their rank ordering



Minimax with alpha-beta (α - β) pruning



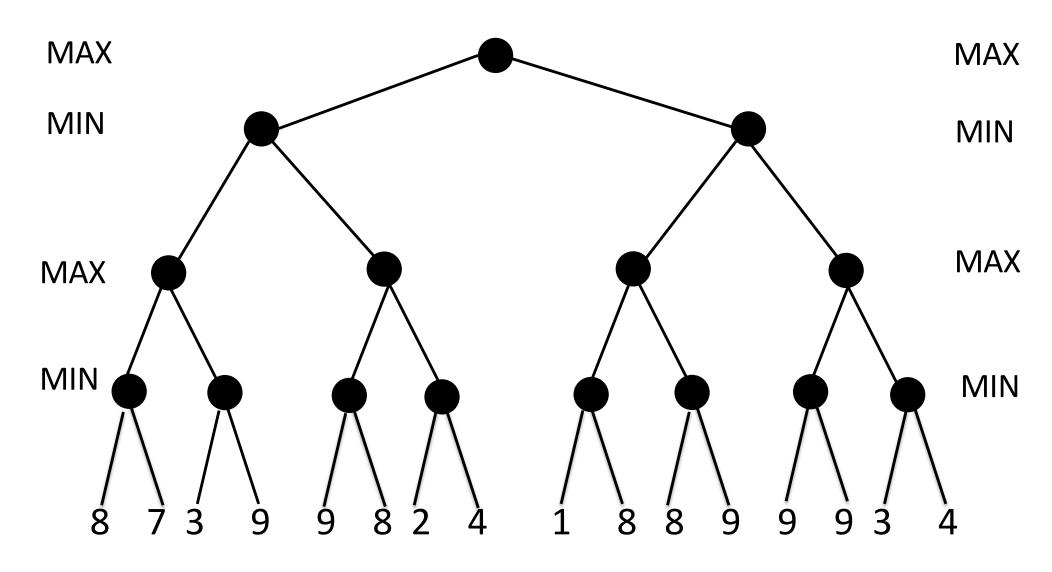
 α value = MAX is guaranteed to gain at least this much, or $\geq \alpha$; a floor on MAX's gain.

 β value = MIN is guaranteed to lose at most this much, or $\leq \beta$; a ceiling on MIN's loss. Note: $[\alpha < \beta]$, or floor < ceiling. Why?

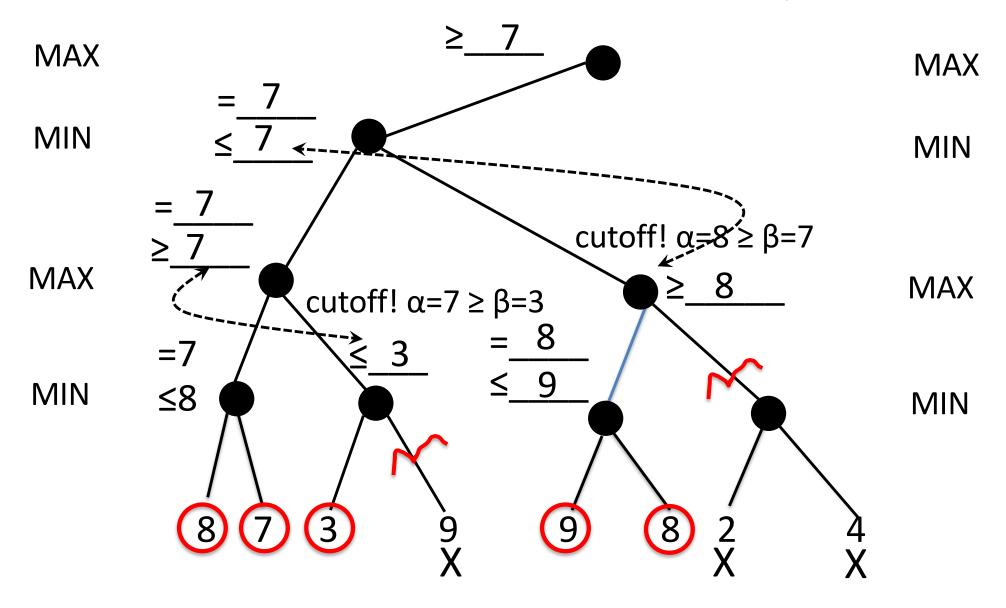
Q: Alpha-Beta pruning value = Minimax value w/o pruning?

Ans: 2

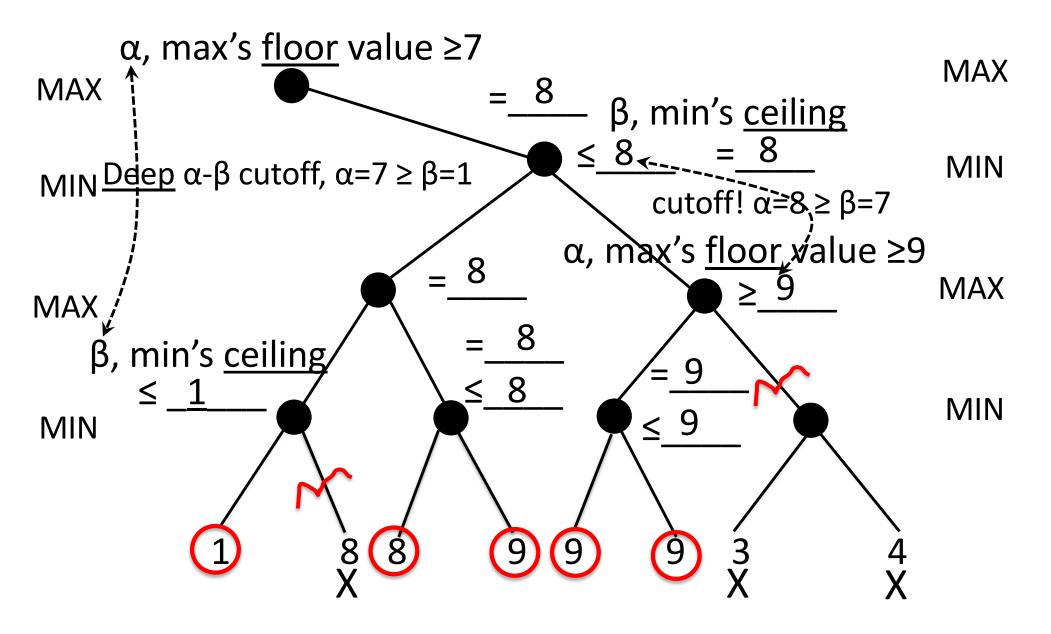
A deeper game tree illustrates how much pruning alpha-beta can do (b = 2, d = 3)



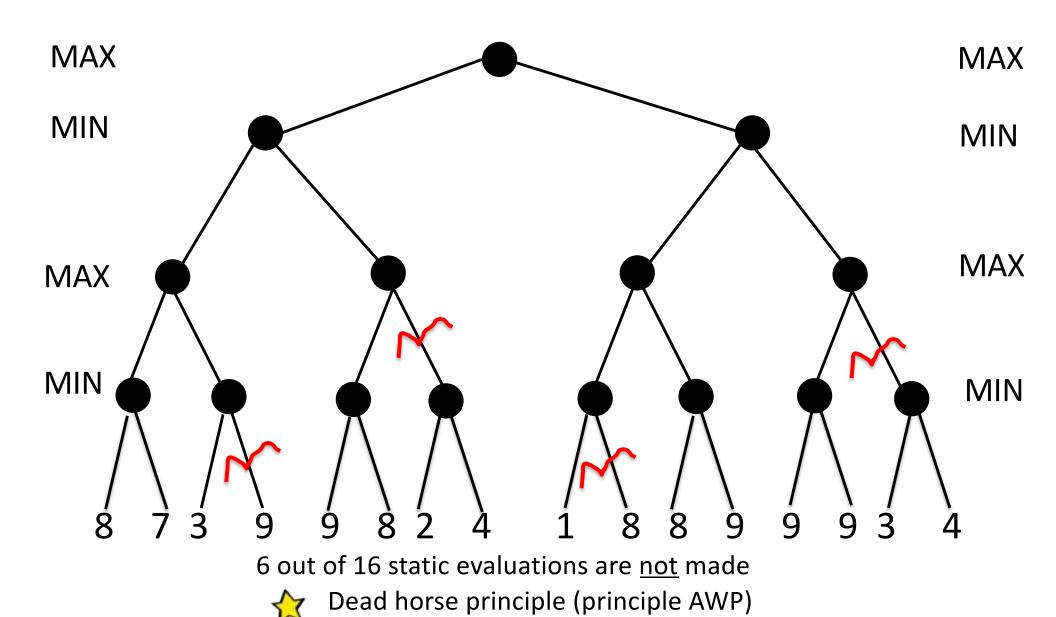
A deeper game tree illustrates how much pruning alpha-beta can do First half of tree (max-α values; min-β values)

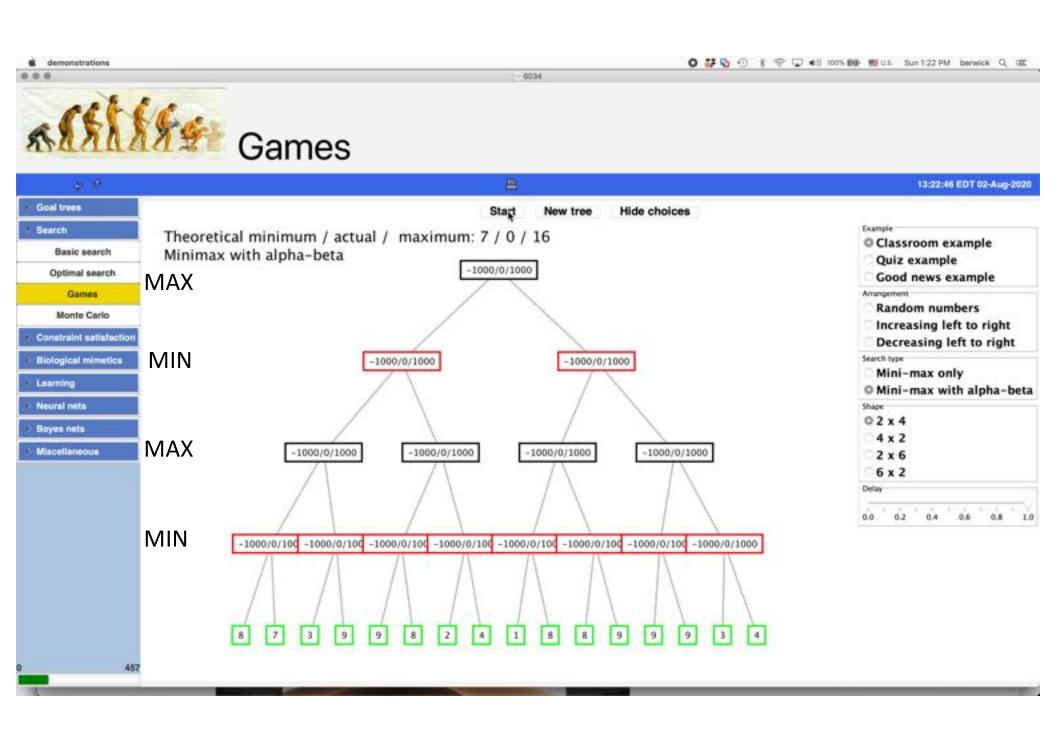


Second half of tree: cutoff <u>whenever</u> α≥ β



A deeper game tree

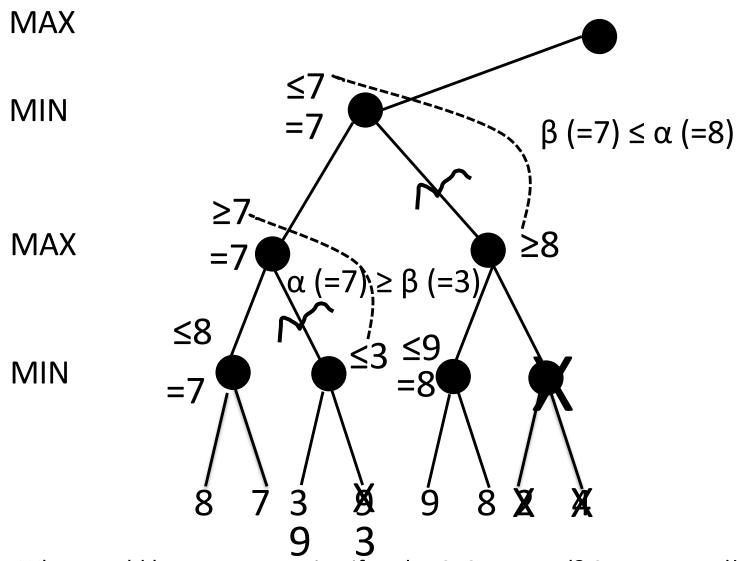




Ordering of static values can greatly affect alpha-beta pruning

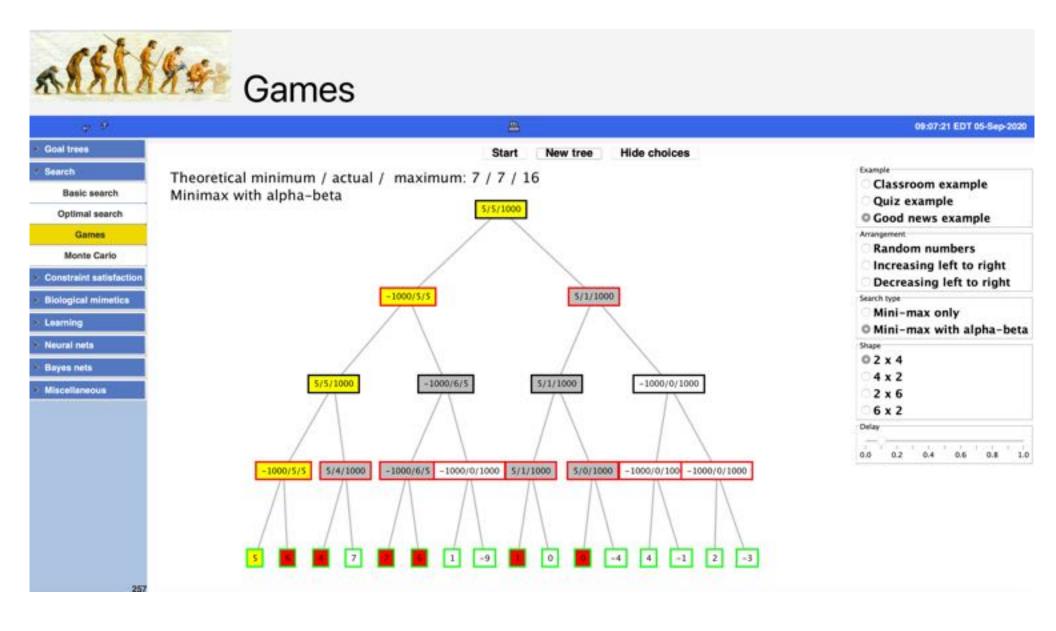
- If the most favorable successor nodes for <u>both</u> MAX and MIN are on the <u>LEFT</u> so we explore them <u>FIRST</u>, then this leads to maximal pruning
- If the most favorable successor nodes for <u>both</u> MAX and MIN are on the <u>RIGHT</u> so they are explored <u>LAST</u>, then there is less pruning, possibly none at all
- Maximal pruning (in terms of branching factor b and tree depth d is approx: $\frac{2 b^{d/2}}{}$, so can search down 2x as far using α - β search in this "good news" optimal case, compared to std minimax

Prune whenever $\alpha \geq \beta$



What would happen to pruning if nodes 3, 9 reversed? 3 Not pruned!

Optimal (good news) α - β game tree pruning



 $\approx 2b^{d/2}$

≈ maximum savings using α-β <u>if</u> optimally ordered game tree

But suppose we run out of compute time? The branching factor depends on the game & board state – won't know for sure how deep we can go in 1 minute....

Suppose we are playing blitz chess?
Will we always have a (good) move at hand?
Let's combine this thought with the idea of optimally ordering the node evaluations...

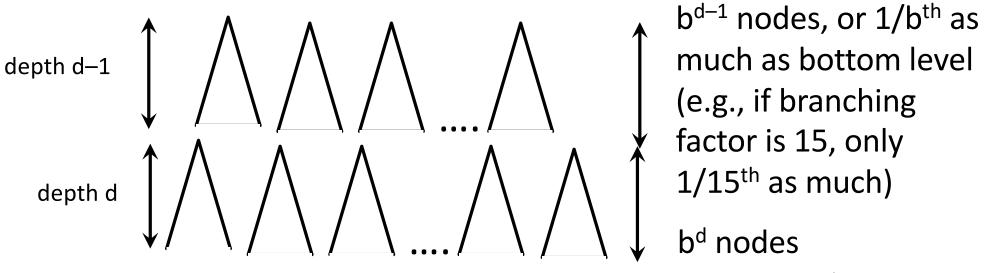
Anytime algorithm



How not to run out of time



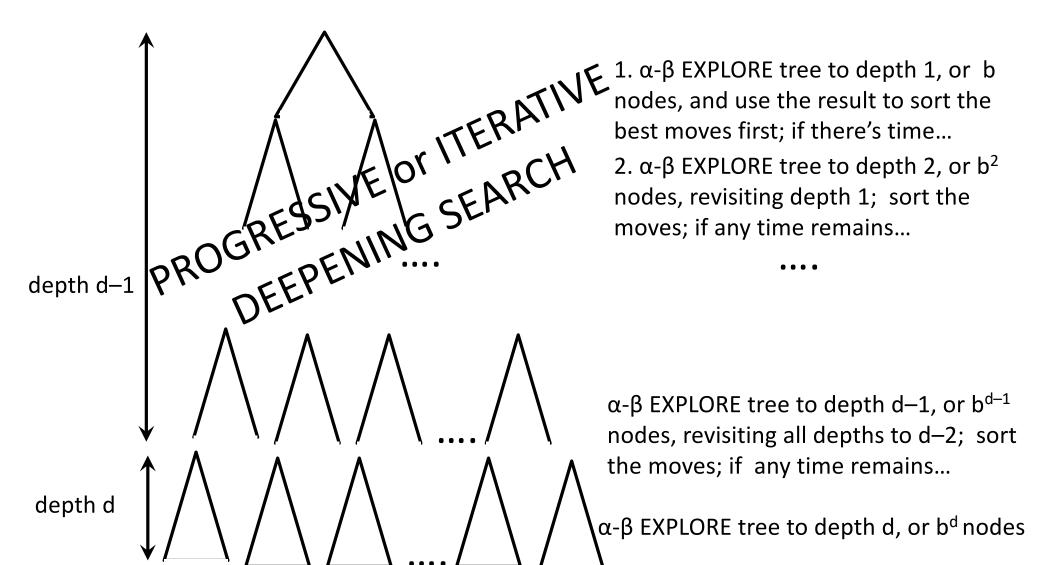
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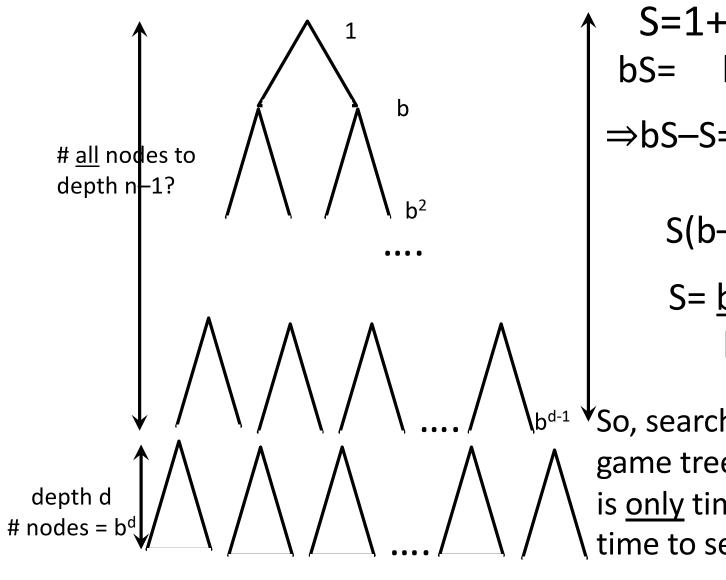
So, we can have a move in hand at level d–1 for only 1/bth of computation cost required to go all the way down to level d

We can do this recursively to great benefit...

How not to run out of time (ignoring α - β for now)



Computational cost of progressive deepening? How many total nodes S in game tree to depth d-1?



S=1+b+b²+...+b^{d-1}
bS= b+b²+...+b^d

⇒bS-S= b^d-1

S(b-1)= b^d-1

S=
$$\frac{b^d-1}{b-1}$$
 ≈ b^{d-1}

VSo, search through <u>all</u> the game tree up thru depth d−1 is <u>only</u> time b^{d−1}, or 1/bth of time to search to full depth d

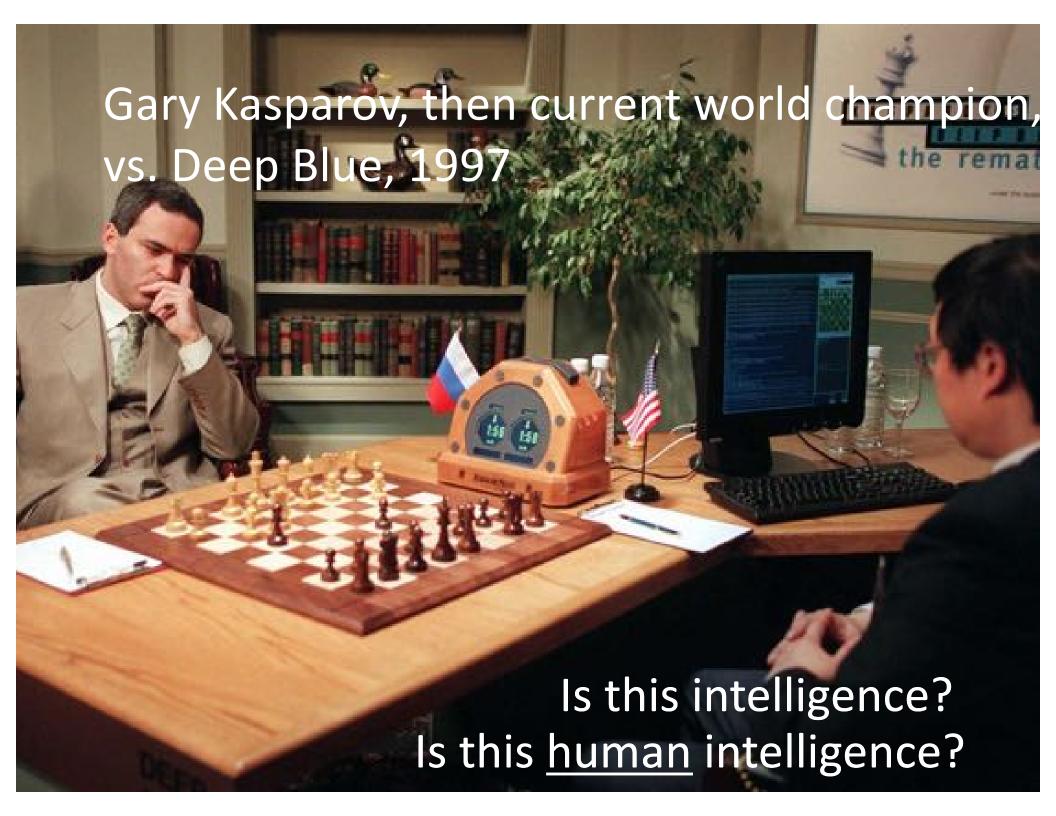
Progressive (Iterative) deepening search

- Optimal DFS method (redundant use of space, but doesn't lose on arbitrary depth d trees)
- Earlier searches tend to improve the commonly used heuristics, so that a more accurate estimate of the score of various nodes at the final depth search can occur
- Because early iterations use small values for d
 they execute extremely quickly. This allows the
 algorithm to supply early indications of the result
 almost immediately, followed by refinements as d
 increases

Deep Blue, IBM, circa 1997.
≈ 10⁸ static evaluations per sec

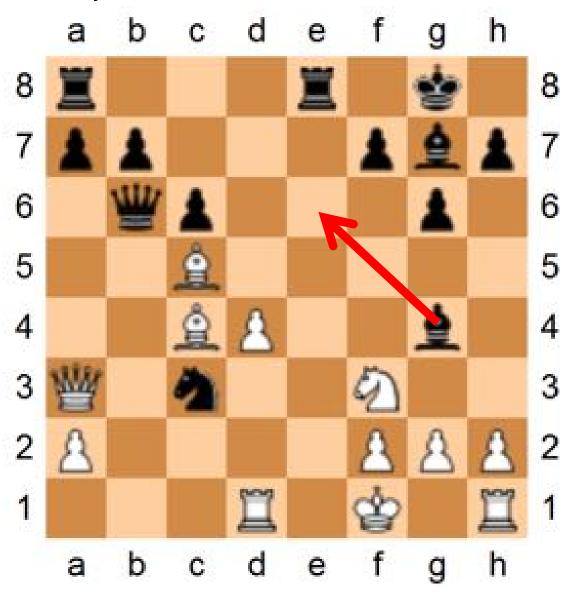


Deep Blue = minimax + alphabeta + progressive deepening+ parallel computing + opening book + end game + secret sauce = uneven tree development



Donald Byrne vs Bobby Fischer (black) 1956, move 17, Queen sacrifice





Gold star ideas today

Dead horse principle, aka "AWP": α-β search
 Martial arts principle – use adversary's strength against them: progressive deepening
 Anytime algorithms: progressive deepening
 Simple ≠ Trivial: sometimes, bulldozers work