

## Class Information

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- Eighth homework has been posted.
- Second project extension: **Monday, April 21, 11:59pm.**
- Sakai submission will be set up soon.

## Review: Dependence — Basics

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### Theorem

Any reordering transformation that preserves every dependence (i.e., visits first the source, and then the sink of the dependence) in a program preserves the meaning of that program.

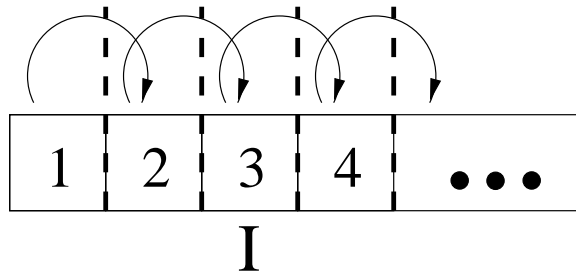
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Note: Dependence starts with the notion of a sequential execution, i.e., starts with a sequential program.

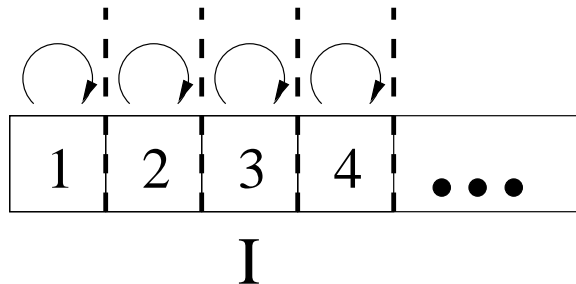
# Review: Dependence Analysis for Array References

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```
do I = 1, 100
  A(I) =
    = A(I-1)
enddo
```



```
do I = 1, 100
  A(I) =
    = A(I)
enddo
```



A **loop-independent** dependence exists regardless of the loop structure. The source and sink of the dependence occur on the same loop iteration.

A **loop-carried** dependence is induced by the iterations of a loop. The source and sink of the dependence occur on different loop iterations.

*Loop-carried dependences can inhibit parallelization and loop transformations*

# Dependence Testing

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Given

```
do  $i_1 = L_1, U_1$   
  ...  
    do  $i_n = L_n, U_n$   
       $S_1 \quad A(f_1(i_1, \dots, i_n), \dots, f_m(i_1, \dots, i_n)) = \dots$   
       $S_2 \quad \dots = A(g_1(i_1, \dots, i_n), \dots, g_m(i_1, \dots, i_n))$ 
```

A *dependence* between statement  $S_1$  and  $S_2$ , denoted  $S_1 \delta S_2$ , indicates that  $S_1$ , the *source*, must be executed before  $S_2$ , the *sink* on some iteration of the nest.

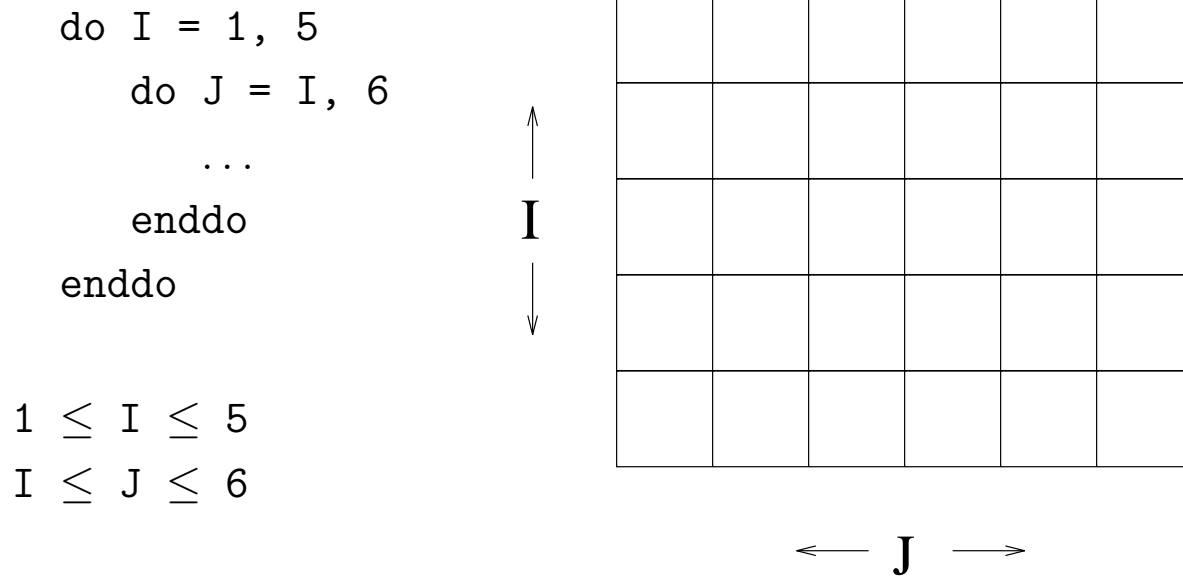
Let  $\alpha$  &  $\beta$  be a vector of  $n$  integers within the ranges of the lower and upper bounds of the  $n$  loops.

Does  $\exists \alpha \leq \beta$ , s.t.

$$f_k(\alpha) = g_k(\beta) \quad \forall k, 1 \leq k \leq m ?$$

# Iteration Space

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- lexicographical (sequential) order for the above iteration space is

$(1,1), (1,2), \dots, (1,6)$   
 $(2,2), (2,3), \dots (2,6)$   
 $\dots$   
 $(5,5), (5,6)$

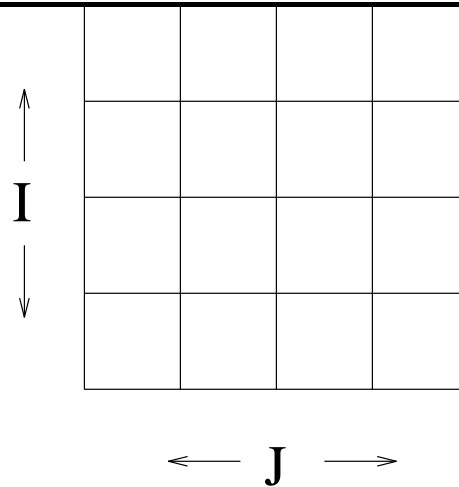
- given  $I = (i_1, \dots, i_n)$  and  $I' = (i'_1, \dots, i'_n)$ ,

$I < I'$  iff

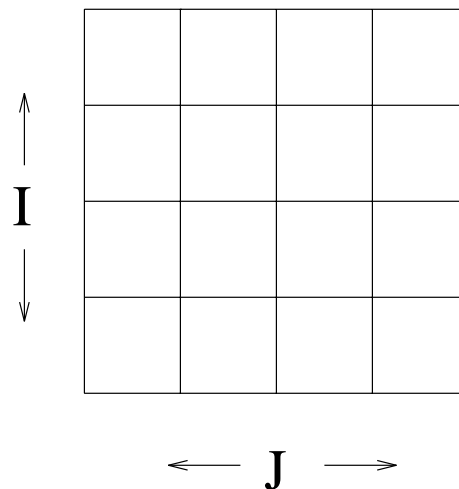
$(i_1, i_2, \dots, i_k) = (i'_1, i'_2, \dots, i'_k) \ \& \ i_{k+1} < i'_{k+1}$

## Distance & Direction Vectors

```
do I = 1, N
  do J = 1, N
    S1 A(I,J) = A(I,J-1)
  enddo
enddo
```



```
do I = 1, N
  do J = 1, N
    S2 A(I,J) = A(I-1,J-1)
    S3 B(I,J) = B(I-1,J+1)
  enddo
enddo
```



**Distance Vector** = number of iterations between accesses to the same location

**Direction Vector** = direction in iteration space ( $=$ ,  $<$ ,  $>$ )

	distance vector	direction vector
$S_1 \delta S_1$	(0,1)	( $=$ , $<$ )
$S_2 \delta S_2$	(1,1)	( $<$ , $<$ )
$S_3 \delta S_3$	(1,-1)	( $<$ , $>$ )

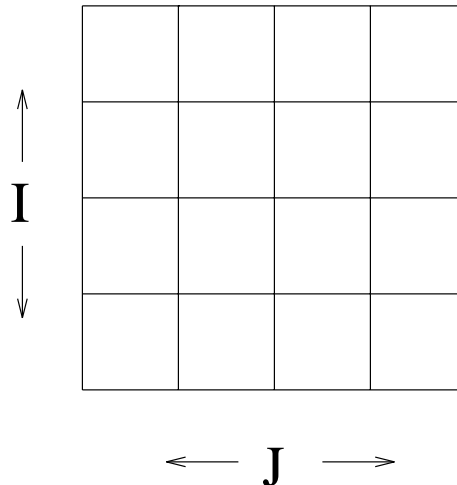
# Which Loops are Parallel?

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```
do I = 1, N
  do J = 1, N
    S1 A(I,J) = A(I,J-1)
```

```
do I = 1, N
  do J = 1, N
    S2 A(I,J) = A(I-1,J-1)
```

```
do I = 1, N
  do J = 1, N
    S3 B(I,J) = B(I-1,J+1)
```



- a dependence  $D = (d_1, \dots, d_k)$  is *carried* at *level*  $i$ , if  $d_i$  is the first nonzero element of the distance/direction vector
- a loop  $l_i$  is *parallel*, if  $\nexists$  a dependence  $D_j$  carried at level  $i$

	distance vector	direction vector
$\forall D_j$	$d_1, \dots, d_{i-1} > 0$	$d_1, \dots, d_{i-1} = "<"$
OR	$d_1, \dots, d_i = 0$	$d_1, \dots, d_i = "="$