

CS314 Spring 14

HW #3 Sample Solution

1.1: LL(1) property -- whenever the grammar allows a choice of multiple productions for a single non-terminal symbol, the FIRST+ sets of the right hand sides of the productions for that non-terminal symbol have to be pairwise disjoint. This allows a deterministic selection among the rules using only a single input look ahead symbol.

Any nonterminal with only one production rule cannot break the LL(1) property. The nonterminals with multiple productions are $\langle \text{morestmts} \rangle$, $\langle \text{stmt} \rangle$, $\langle \text{expr} \rangle$, $\langle \text{variable} \rangle$, and, $\langle \text{digit} \rangle$.

$\langle \text{digit} \rangle$'s FIRST sets of its productions' right hand sides are pairwise disjoint:
 $\{0\} \cap \{1\} = \{\}$, $\{0\} \cap \{2\} = \{\}$, ...

$\langle \text{variables} \rangle$'s FIRST sets of its productions right hand sides are pairwise disjoint: They are $\{a\}$, $\{b\}$, and $\{c\}$.

$\langle \text{expr} \rangle$'s FIRST sets of its productions right hand sides are pairwise disjoint: They are $\{+\}$, $\{*\}$, $\{a, b, c\}$ and $\{0, 1, 2\}$.

$\langle \text{stmts} \rangle$'s FIRST sets of its productions right hand sides are pairwise disjoint: They are $\{a, b, c\}$, $\{\text{if}\}$, $\{\text{while}\}$ and $\{\text{begin}\}$.

The tricky part is to show that we can make a unique decision for the right hand sides for the two productions of the nonterminal $\langle \text{morestmts} \rangle$. $\text{FIRST}(\langle \text{stmtlist} \rangle) = \{;\}$, so that's not a problem. However, since $\langle \text{morestmts} \rangle ::= \epsilon$, we need to look at the FOLLOW set of $\langle \text{morestmts} \rangle$, i.e., we have to compute FIRST+. Due to the rule

$\langle \text{stmtlist} \rangle ::= \langle \text{stmt} \rangle \langle \text{morestmts} \rangle$, every symbol that is in $\text{FOLLOW}(\langle \text{stmtlist} \rangle)$ has to be added to $\text{FOLLOW}(\langle \text{morestmts} \rangle)$. Due to the rule

$\langle \text{block} \rangle ::= \text{begin } \langle \text{stmtlist} \rangle \text{ end}$, $\text{FOLLOW}(\langle \text{stmtlist} \rangle)$ has to contain symbol "end". In fact, $\text{FOLLOW}(\langle \text{stmtlist} \rangle) = \{\text{end}\}$. Therefore, $\text{FOLLOW}(\langle \text{morestmts} \rangle) = \{\text{end}\}$ as well.

So, $\text{FIRST}^+(\epsilon) = \{\text{end}\}$ for the right hand side ϵ for nonterminal symbol $\langle \text{morestmts} \rangle$. Since the sets $\{;\}$ and $\{\text{end}\}$ are pairwise disjoint, we can make a deterministic decision to pick a single rule for nonterminal $\langle \text{morestmts} \rangle$.

Therefore, our grammar is LL(1).

1.2:

All empty cells are error states.

Parse table is listed in two parts.

Part 1:

NTVT	program	begin	end	;	if	then	else	while	do
<program>	program <block>								
<block>		begin <stmtlist> end							
<stmtlist>		<stmt> <morestmts>			<stmt> <morestmts>			<stmt> <morestmts>	
<morestmts>				; <stmtlist>					
<stmt>		<block>			<ifstmt>			<whilestmt>	
<assign>									
<ifstmt>					if <testexpr> then <stmt> else <stmt>				
<whilestmt>								while <testexpr> > do <stmt>	
<testexpr>									
<expr>									
<variable>									
<digit>									

Part 2:

NTVT	<=	+	*	;	a	b	c	0	1	2	end	EOF
<program>					<stmt> <morestmts>							
<block>												
<stmtlist>												

<morestmts>							<i>epsilon</i>	
<stmt>					<assign>			
<assign>					<variable> = <expr>			
<ifstmt>								
<whilestmt>								
<testexpr>					<variable> <= <expr>			
<expr>		+ <expr> <expr>	* <expr> <expr>		<variable>	<digit>		
<variable>					a	b	c	
<digit>						0	1	2

1.3 / 1.4:

Assume global variable token exists and that it is initialized before entering the function representing the nonterminal program. Also assume the procedure next_token() exists.

One possible solution for 1.4 uses global integer variables to keep track of the number of assignments (#asgn) and numbers of references (#refs).

Added code for problem 1.4 is shown in bold face.

```

program
{
    #asgn := 0;
    #refs := 0;
    switch token
    {
        case "program":
            token := next_token();
            call block();
            if "." != token
            {
                error; die;
            }
            token := next_token();
            if eof != token
            {
                error; die;
            }
            break;
        default:
            error; die;
    }
    print( %d assignments, % references), #asgn, #refs) ;
}

block

```

```

{
    switch token
    {
        case "begin":
            token := next_token();
            call stmtlist();
            if "end" != token
            {
                error; die;
            }
            token := next_token();
            break;
        default:
            error; die;
    }
}

stmtlist
{
    switch token
    {
        case a: case b: case c: case "if": case "while": case "begin":
            call stmt();
            call morestmts();
            break;
        default:
            error; die;
    }
}

morestmts
{
    switch token
    {
        case "(":
            token := next_token();
            call stmtlist();
            break;
        case "end": // this is how you handle epsilon productions
            break;
        default:
            error; die;
    }
}

stmt
{
    switch token
    {
        case a: case b: case c:
            call assign();
            break;
    }
}

```

```

        case "if":
            call ifstmt();
            break;
        case "while":
            call whilestmt();
            break;
        case "begin":
            call block();
            break;
        default:
            error; die;
    }
}

assign
{
    switch token
    {
        case a: case b: case c:
            #asgn = #asgn + 1;
            call variable();
            if "=" != token
            {
                error; die;
            }
            token := next_token();
            call expr();
            break;
        default:
            error; die;
    }
}

ifstmt
{
    switch token
    {
        case "if":
            token := next_token();
            call testexpr();
            if "then" != token
            {
                error; die;
            }
            token := next_token();
            call stmt();
            if "else" != token
            {
                error; die;
            }
            token := next_token();
            call stmt();

```

```

        break;
    default:
        error; die;
    }
}

```

```

whilestmt
{
    switch token
    {
        case "while":
            token := next_token();
            call testexpr();
            if "do" != token
            {
                error; die;
            }
            token := next_token();
            call stmt();
            break;
        default:
            error; die;
    }
}

```

```

testexpr
{
    switch token
    {
        case a: case b: case c:
            call variable();
            if "<=" != token
            {
                error; die;
            }
            token := next_token();
            call expr();
            break;
        default:
            error; die;
    }
}

```

```

expr
{
    switch token
    {
        case +: case *:
            token := next_token();
            call expr();
            call expr();
            break;
    }
}

```

```

        case a: case b: case c:
            call variable();
            break;
        case 0: case 1: case 2:
            call digit();
            break;
        default:
            error; die;
    }
}

variable
{
    switch token
    {
        case a: case b: case c:
            #refs = #refs + 1;
            token := next_token();
            break;
        default:
            error; die;
    }
}

digit
{
    switch token
    {
        case 0: case 1: case 2:
            token := next_token();
            break;
        default:
            error; die;
    }
}

```
