198:314 Principles of Programming Languages

Course Goals

- To gain understanding of the basic structure of programming languages:
 - Data types, control structures, naming conventions,...
- To study different language paradigms:

 - To ensure an appropriate language is chosen for a task
- To know the principles underlying all programming languages:
 - To make learning new programming languages easier
 - To enable full use of a programming language
 - To understand the implementation challenges of different programning constructs / features

Programming languages are $\underline{tools} \Rightarrow understand$ how to design or use them

Course Information

Prerequisites (summary):

- CS 205 (Introduction to Discrete Structures)
- CS 211 (Computer Architecture)

Important facts:

staff: Prof. Ulrich Kremer, Hans Christian Woithe,

Liu Liu, Ardash Yoga

lectures: Tu/Fri noon-1:20pm, TIL-254 (Tillett Hall)

recitations: attendance mandatory, start this Friday

section 1, Tu 1:55-2:50pm, LCB 110

section 2, Tu 5:15-6:10pm, TIL-103A

section 3, Fri 1:55-2:50pm, TIL-251

books on: Science & Engineering Resource Center

reserve (SEC), Busch Campus

Basis for grades (subject to changes):

10% homework / recitation

25% mid-term exam

35% final exam (cumulative)

30% three major programming projects

- The textbook for this course is "Programming Language Pragmatics by Michael L. Scott, 3nd Edition, Morgan Kaufmann (Elsevier), 2009.
- Additional (recommended) texts: see course web page

Course material is available on our class website at

www.cs.rutgers.edu/courses/314/classes/spring_2014_kremer

In addition, there is a news group (sakai web page) sakai.rutgers.edu. All questions regarding homeworks and projects should be posted here. YOU SHOULD READ THE SAKAI NEWS GROUP AND LOOK AT THE HOME PAGE AT LEAST EVERY OTHER DAY

All programming will be done on the **ilab cluster**. Get yourself an **ilab** account (see link on bottom of our 314 website). Learn to do the normal things — edit, compile, . . .

Academic Integrity (see our web page)

- \rightarrow read-protect your directories and files (ilab)
- → no group projects
- \rightarrow will use MOSS for detecting software plagiarism

14 weeks, no "make-up" work after the end of the course. If there is a problem, let me know immediately.

IMPORTANT INFORMATION \Rightarrow will be posted on 314 web page and/or on sakai forums!

- Failure to take a scheduled exam
- Grading of homeworks and projects
- Instructions of how to submit programming projects
- Partial credit for late project submissions

Email TAs or me:

- Subject line has to start with 314:, e.g., 314: Question about my midterm exam
- No project and homework questions; post them on the sakai discussion forums;

Special permission numbers

- Put your name (and email address!) on the list. Indicate whether you prefer any particular section.
- Come talk to me after class.
- Will try to accommodate everyone.

I use overhead transparencies

- I try to moderate my speed
- You need to say STOP!
- all transparencies are on the Web (PDF)
- you should still take some notes

I'll tell you where we are in the book

- I don't lecture directly from the book
- You need to read the book
- Going to the recitations is mandatory
- I strongly recommend coming to the lectures

What is the Purpose of a Programming Language?

A programming language is ...

a set of conventions for communicating an algorithm. Horowitz

Purposes:

- specifying algorithm and data structures
- communicating algorithms among people
- establishing correctness (allow reasoning)

This is a C program that uses two one-dimensional arrays **a** and **b** of size **SIZE**. The arrays are initialized, and then a sum reduction is performed. The size of the arrays and the result of the sum reduction is printed out.

example.c

```
#include <stdio.h>

#define SIZE 100
int main() {
   int a[SIZE], b[SIZE];
   int i, sum;

for (i=0; i<SIZE; i++) {
    a[i] = 1;
    b[i] = 2;
   }
   sum = 0;
   for (i=0; i<SIZE; i++)
      sum = sum + a[i] + b[i];

   printf("for two arrays of size %d, sum = %d\n", SIZE, sum);
}</pre>
```

Compiler: $gcc - O3 - S example.c \Rightarrow example.s$

```
.file "example.c"
.version "01.01"
gcc2_compiled.:
.section .rodata.str1.32, "aMS", @progbits,1
.align 32
.LCO:
.string "for two arrays of size %d, sum = %d\n"
.text
.align 4
.globl main
.type main, @function
main:
pushl %ebp
movl %esp, %ebp
xorl %eax, %eax
subl $808, %esp
mov1 $99, %edx
.p2align 2
.L21:
movl $1, -408(%ebp,%eax)
movl $2, -808(%ebp,%eax)
addl $4, %eax
decl %edx
jns .L21
xorl %ecx, %ecx
xorl %eax, %eax
movl $99, %edx
.p2align 2
.L26:
addl -408(%ebp, %eax), %ecx
addl -808(%ebp, %eax), %ecx
addl $4, %eax
decl %edx
jns .L26
pushl %eax
pushl %ecx
pushl $100
pushl $.LCO
call printf
addl $16, %esp
leave
ret
.Lfe1:
.size main, .Lfe1-main
.ident "GCC: (GNU) 2.96 20000731 (Red Hat Linux 7.3 2.96-112)"
```

gcc -o example.o -O3 example.c; strip example.o; objdump -d example.o

```
objdump: example.o: No symbols
             file format elf32-sparc
Disassembly of section .text:
00010444 <.text>:
  10444: bc 10 20 00 clr %fp
  10448: e0 03 a0 40 ld [ %sp + 0x40 ], %10
  1044c: a2 03 a0 44 add %sp, 0x44, %11
  10450: 9c 23 a0 20 sub %sp, 0x20, %sp
  10454: 80 90 00 01 tst %g1
  10458: 02 80 00 04 be 0x10468
  1045c: 90 10 00 01 mov %g1, %o0
  10460: 40 00 40 c4 call 0x20770
  10464: 01 00 00 00 nop
  10468: 11 00 00 41 sethi %hi(0x10400), %o0
  1046c: 90 12 22 d8 or %o0, 0x2d8, %o0 ! 0x106d8
  10470: 40 00 40 c0 call 0x20770
  10474: 01 00 00 00 nop
  10478: 40 00 00 91 call 0x106bc
  1047c: 01 00 00 00 nop
  10480: 90 10 00 10 mov %10, %00
  10484: 92 10 00 11 mov
                          %11, %o1
  10488: 95 2c 20 02 sll %10, 2, %o2
  1048c: 94 02 a0 04 add %o2, 4, %o2
  10490: 94 04 40 0a add %11, %o2, %o2
  10494: 17 00 00 82 sethi %hi(0x20800), %o3
  10498: 96 12 e0 a8 or %o3, 0xa8, %o3 ! 0x208a8
  1049c: d4 22 c0 00 st %o2, [ %o3 ]
  104a0: 40 00 00 4e call 0x105d8
  104a4: 01 00 00 00 nop
  104a8: 40 00 40 b5 call 0x2077c
  104ac: 01 00 00 00 nop
   104b0: 40 00 40 b6 call 0x20788
   104b4: 01 00 00 00 nop
  104b8: 81 c3 e0 08 retl
  104bc: ae 03 c0 17 add %o7, %17, %17
  104c0: 9d e3 bf 90 save %sp, -112, %sp
  104c4: 11 00 00 00 sethi %hi(0), %o0
  104c8: 2f 00 00 40 sethi %hi(0x10000), %17
  104cc: 7f ff ff fb call 0x104b8
  104d0: ae 05 e2 54 add %17, 0x254, %17 ! 0x10254
  104d4: 90 12 20 0c or %00, 0xc, %00
  104d8: d2 05 c0 08 ld [ %17 + %00 ], %o1
   104dc: d4 02 40 00 ld [ %o1 ], %o2
  104e0: 80 a2 a0 00 cmp %o2, 0
  104e4: 12 80 00 23 bne 0x10570
  104e8: 11 00 00 00 sethi %hi(0), %o0
  104ec: 90 12 20 10 or %o0, 0x10, %o0 ! 0x10
  104f0: d4 05 c0 08 ld [ %17 + %00 ], %02
  104f4: d2 02 80 00 ld [ %o2 ], %o1
  104f8: d0 02 40 00 ld [ %o1 ], %o0
  104fc: 80 a2 20 00 cmp %o0, 0
  10500: 02 80 00 0f be 0x1053c
  10504: 11 00 00 00 sethi %hi(0), %o0
   10508: a0 10 00 0a mov %o2, %10
  1050c: d0 04 00 00 ld [ %10 ], %00
  10510: 90 02 20 04 add %00, 4, %00
  10514: d0 24 00 00 st %o0, [ %10 ]
  10518: d2 02 3f fc ld [ %o0 + -4 ], %o1
  1051c: 9f c2 40 00 call %o1
  10520: 01 00 00 00 nop
  10524: d0 04 00 00 ld [ %10 ], %o0
  10528: d2 02 00 00 ld [ %o0 ], %o1
  1052c: 80 a2 60 00 cmp %o1, 0
   10530: 12 bf ff f9 bne 0x10514
   10534: 90 02 20 04 add %00, 4, %00
  10538: 11 00 00 00 sethi %hi(0), %o0
  1053c: 90 12 20 1c or %o0, 0x1c, %o0 ! 0x1c
  10540: d2 05 c0 08 ld [ %17 + %o0 ], %o1
  10544: 80 a2 60 00 cmp %o1, 0
  10548: 02 80 00 05 be 0x1055c
```

```
1054c: 13 00 00 00 sethi %hi(0), %o1
10550: 92 12 60 08 or %o1, 8, %o1 ! 0x8
10554: 40 00 40 90 call 0x20794
10558: d0 05 c0 09 ld [ %17 + %o1 ], %o0
1055c: 11 00 00 00 sethi %hi(0), %o0
10560: 90 12 20 0c or %00, 0xc, %00 ! 0xc
10564: d4 05 c0 08 ld [ %17 + %00 ], %o2
10568: 92 10 20 01 mov 1, %o1
1056c: d2 22 80 00 st %o1, [ %o2 ]
10570: 81 c7 e0 08 ret
10574: 81 e8 00 00 restore
10578: 9d e3 bf 90 save %sp, -112, %sp
1057c: 81 c7 e0 08 ret
10580: 81 e8 00 00 restore
10584: 9d e3 bf 90 save %sp, -112, %sp
10588: 11 00 00 00 sethi %hi(0), %o0
1058c: 2f 00 00 40 sethi %hi(0x10000), %17
10590: 7f ff ff ca call 0x104b8
10594: ae 05 e1 90 add %17, 0x190, %17 ! 0x10190
10598: 90 12 20 18 or %00, 0x18, %00
1059c: d2 05 c0 08 ld [ %17 + %00 ], %o1
105a0: 80 a2 60 00 cmp %o1, 0
105a4: 02 80 00 08 be 0x105c4
105a8: 13 00 00 00 sethi %hi(0), %o1
105ac: 92 12 60 08 or %o1, 8, %o1 ! 0x8
105b0: 15 00 00 00 sethi %hi(0), %o2
105b4: d0 05 c0 09 ld [ %17 + %o1 ], %o0
105b8: 94 12 a0 04 or %o2, 4, %o2
105bc: 40 00 40 79 call 0x207a0
105c0: d2 05 c0 0a ld [ %17 + %o2 ], %o1
105c4: 81 c7 e0 08 ret
105c8: 81 e8 00 00 restore
105cc: 9d e3 bf 90 save %sp, -112, %sp
105d0: 81 c7 e0 08 ret
105d4: 81 e8 00 00 restore
105d8: 9d e3 bc 70 save %sp, -912, %sp
105dc: 92 07 be 60 add %fp, -416, %o1
105e0: 94 07 bc d0 add %fp, -816, %o2
105e4: 86 10 00 09 mov %o1, %g3
105e8: 84 10 00 0a mov %o2, %g2
105ec: 9a 10 20 01 mov 1, %o5
105f0: 98 10 20 02 mov 2, %o4
105f4: 90 10 20 00 clr %o0
105f8: 96 10 20 63 mov 0x63, %o3
105fc: da 22 00 03 st %o5, [ %o0 + %g3 ]
10600: d8 22 00 02 st %o4, [ %o0 + %g2 ]
10604: 96 82 ff ff addcc %o3, -1, %o3
10608: 1c bf ff fd bpos 0x105fc
1060c: 90 02 20 04 add %00, 4, %00
10610: 9a 10 00 0a mov %o2, %o5
10614: 84 10 00 09 mov %o1, %g2
10618: 94 10 20 00 clr %o2
1061c: 98 10 20 00 clr %o4
10620: 96 10 20 63 mov 0x63, %o3
10624: d0 03 00 02 ld [ %o4 + %g2 ], %o0
10628: 96 82 ff ff addcc %o3. -1, %o3
1062c: d2 03 00 0d ld [ %o4 + %o5 ], %o1
10630: 90 02 80 08 add %o2, %o0, %o0
10634: 94 02 00 09 add %00, %01, %02
10638: 1c bf ff fb bpos 0x10624
1063c: 98 03 20 04 add %o4, 4, %o4
10640: 11 00 00 41 sethi %hi(0x10400), %o0
10644: 90 12 22 f8 or %o0, 0x2f8, %o0 ! 0x106f8
10648: 40 00 40 59 call 0x207ac
1064c: 92 10 20 64 mov 0x64, %o1
10650: 81 c7 e0 08 ret
10654: 81 e8 00 00 restore
10658: 81 c3 e0 08 retl
1065c: ae 03 c0 17 add %o7, %17, %17
10660: 9d e3 bf 90 save %sp, -112, %sp
10664: 11 00 00 00 sethi %hi(0), %o0
10668: 2f 00 00 40 sethi %hi(0x10000), %17
1066c: 7f ff ff fb call 0x10658
10670: ae 05 e0 b4 add %17, 0xb4, %17 ! 0x100b4
10674: 90 12 20 14 or %00, 0x14, %00 10678: d2 05 c0 08 ld [ %17 + %00 ], %o1
1067c: d4 02 7f fc ld [ %o1 + -4 ], %o2
10680: 80 a2 bf ff cmp %o2, -1
10684: 02 80 00 09 be 0x106a8
10688: a0 02 7f fc add %o1, -4, %10
1068c: d0 04 00 00 ld [ %10 ], %00
10690: 9f c2 00 00 call %o0
10694: a0 04 3f fc add %10, -4, %10
```

```
10698: d0 04 00 00 ld [ %10 ], %o0
   1069c: 80 a2 3f ff \mbox{cmp} %o0, -1
   106a0: 12 bf ff fb bne 0x1068c
   106a4: 01 00 00 00 nop
   106a8: 81 c7 e0 08 ret
   106ac: 81 e8 00 00 restore
   106b0: 9d e3 bf 90 save %sp, -112, %sp
   106b4: 81 c7 e0 08 ret
   106b8: 81 e8 00 00 restore
Disassembly of section .init:
000106bc <.init>:
  106bc: 9d e3 bf a0 save %sp, -96, %sp
   106c0: 7f ff ff b1 call 0x10584
   106c4: 01 00 00 00 nop
   106c8: 7f ff ff e6 call 0x10660
   106cc: 01 00 00 00 nop
   106d0: 81 c7 e0 08 ret
   106d4: 81 e8 00 00 restore
Disassembly of section .fini:
000106d8 <.fini>:
  106d8: 9d e3 bf a0 save %sp, -96, %sp
106dc: 7f ff ff 79 call 0x104c0
   106e0: 01 00 00 00 nop
   106e4: 81 c7 e0 08 ret
  106e8: 81 e8 00 00 restore
Disassembly of section .plt:
00020740 <.plt>:
   20770: 03 00 00 30 sethi %hi(0xc000), %g1
   20774: 30 bf ff f3 b,a 0x20740
   20778: 01 00 00 00 nop
   2077c: 03 00 00 3c sethi %hi(0xf000), %g1
   20780: 30 bf ff f0 b,a 0x20740
   20784: 01 00 00 00 nop
   20788: 03 00 00 48 sethi %hi(0x12000), %g1
   2078c: 30 bf ff ed b,a 0x20740
   20790: 01 00 00 00 nop
   20794: 03 00 00 54 sethi %hi(0x15000), %g1
   20798: 30 bf ff ea b,a 0x20740
   2079c: 01 00 00 00 nop
   207a0: 03 00 00 60 sethi %hi(0x18000), %g1
   207a4: 30 bf ff e7 b,a 0x20740
   207a8: 01 00 00 00 nop
   207ac: 03 00 00 6c sethi %hi(0x1b000), %g1
   207b0: 30 bf ff e4 b,a
   207b4: 01 00 00 00 nop
   207b8: 01 00 00 00 nop
```

Need for high-level programming languages for

- Readable, familiar notations
- Machine independence (portability)
- Consistency checks during implementation
- Dealing with scale

The art of programming is the art of organizing complexity. Example: Dijkstra, 1972

However:

• Acceptable loss of efficiency

First FORTRAN compiler built by IBM, in 1957, translated into code as efficient as hand-coded code. *John Backus*

Why Learn More than One Programming Language?

- Each language encourages thinking about a problem in a particular way.
- Each language provides (slightly) different expressiveness & efficiency.
- \Rightarrow The language should match the problem.

Why Learn About Programming Language PRINCIPLES?

A programming language is a **tool**.

Studying the design of a tool leads to:

- Better understanding of its functionality and limitations.
- Increased competence in using it.
- Basis for lots of other work in computer science.

Computational Paradigms

Imperative:

Sequence of state-changing actions.

- Manipulate an abstract machine with:
 - 1. Variables naming memory locations
 - 2. Arithmetic and logical operations
 - 3. Reference, evaluate, assign operations
 - 4. Explicit control flow statements
- Fits the von Neumann architecture closely
- Key operations: Assignment and "Goto"

Functional:

Composition of operations on data.

- No named memory locations
- Value binding through parameter passing
- ullet Key operations: Function application and Function abstraction

Basis in lambda calculus

Computational Paradigms (Cont.)

Logic:

Formal logic specification of problem.

- Programs say what properties the solution must have, not how to find it
- Solutions through reasoning process.
- Key operation: *Unification*

Basis in first order predicate logic

Object-Oriented:

Communication between abstract objects.

- "Objects" collect both the data and the operations
- "Objects" provide data abstraction
- Can be either imperative or functional
- Key operation: Message passing or Method invocation

Computational Paradigms (Cont.)

Event-Driven:

Objects are associated with events

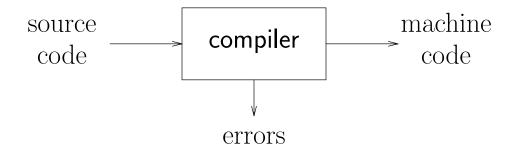
- events are asynchronous
- arrival of an event triggers action
- main applications: GUI, simulations
- Key operation: event handling

Parallel:

Computations and data accesses at the same time

- functional (task/threads) and data parallelism
- different granularities: instruction, loop, or task level
- synchronization: locks, message passing, ...
- Key notions: control and data dependencies

Compilers



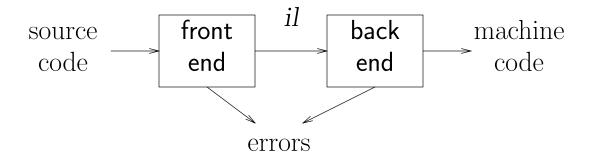
Implications:

- recognize legal (and illegal) programs
- generate correct code
- manage storage of all variables and code
- need format for object (or assembly) code

Big step up from assembler – higher level notations

Traditional two pass compiler

Pass: reading and writing entire program



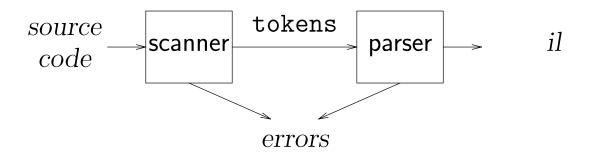
Implications:

- intermediate language (il)
- front end maps legal code into il
- back end maps il onto target machine
- simplify retargeting
- allows multiple front ends
- multiple passes \Rightarrow better code

Front end is **O**(n)

Back end is NP-Complete

Front end



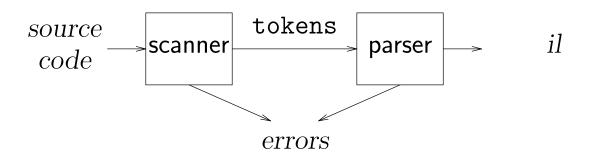
Parser: syntax & semantic analyzer, il code generator (syntax-directed translator)

Front End Responsibilities:

- recognize legal programs
- report errors
- produce il
- preliminary storage map
- shape the code for the back end

Much of front end construction can be automated

Scanner



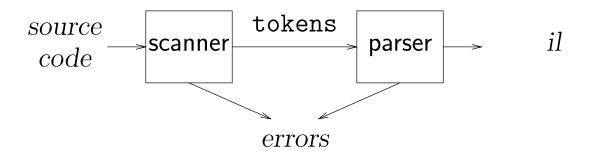
Scanner

• maps characters into tokens – the basic unit of syntax

$$x = x + y;$$
becomes
 $< id, x > = < id, x > + < id, y > ;$

- character string for a token is a lexeme
- typical tokens: number, id, +, -, *, /, do, end
- eliminates white space (tabs, blanks, comments)
- a key issue is speed⇒ use specialized recognizer (lex)

Parser



Parser:

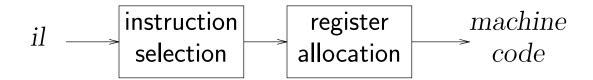
- recognize context-free syntax (Context Free Grammars)
- guide context-sensitive analysis
- \bullet construct il(s)
- produce meaningful error messages
- attempt error correction

Parser generators mechanize much of the work

Example il: Abstract syntax tree (AST)

Compilers often use an abstract syntax tree.

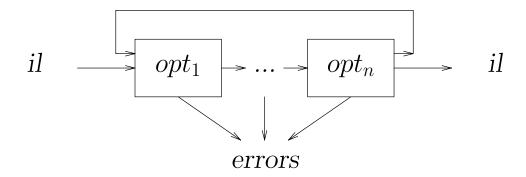
Back end



Responsibilities

- translate il into target machine code
- choose instructions for each il operation
- decide what to keep in registers at each point

Optimizer (middle end)



Modern optimizers are usually built as a set of passes.

Typical passes

- discover & propagate constant values
- reduction of operator strength
- common subexpression elimination
- redundant computation elimination
- encode an idiom in some powerful instruction
- move computation to less frequently executed place (e.g., out of loops)

Things to Do

Things to do for next lecture:

- read Scott: Chapter 1 (covers today's lecture)
- read Scott: Chapters 2.1 and 2.2; ALSU: Chapters 3.1 3.4
- get an ilab account
- learn to read Sakai news group

Recitations will start this Friday (section 3).