

CSC 1052 – Algorithms & Data Structures II: Avoiding Recursion

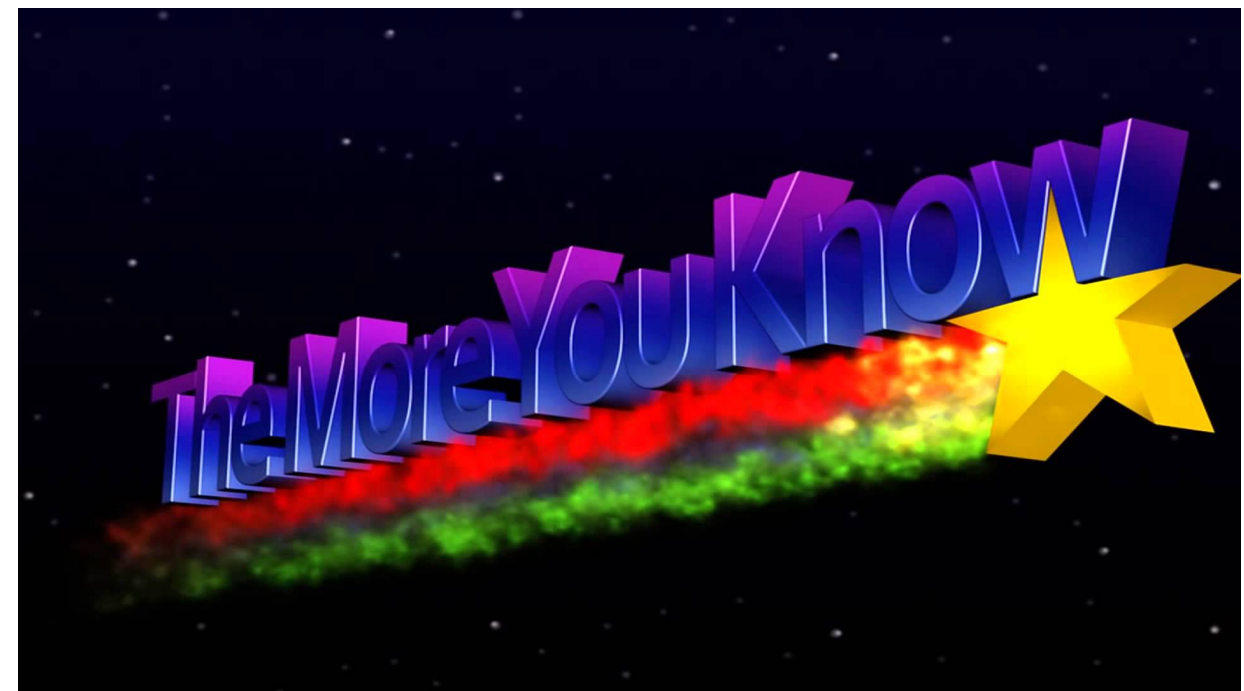
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Project I graded

- Grades are posted on Blackboard
- Breakdown:
 - Submitted code and write up: 60
 - Submitted compiling code: 70
 - Submitted compiling code that returns the correct song order and plays: 80
 - Passed additional method testing: 100

Lessons learned

- “Simple concepts are not simple to implement”
- “Writing out functionality was helpful
- “Testing was difficult”



Recap

- Recursion is an obvious solution for many mathematical problems
- Recursion may be able to solve some problems in more subtle ways
 - Towers of Hanoi
- Recursion can be used as a tool on its own
 - Generating recursive structures in fractals

Is recursion right for me?

- Recall: the hidden cost of recursion is the added space required to store each call
- Some languages do not support recursion at all
- Some applications are simplified but needlessly inefficient when using recursion



How memory is allocated

- Statically vs Dynamically
- Static allocation:
 - Program is compiled into machine code
 - Additional space is reserved along with machine instructions for variables

Static



vs.

Dynamic



Static Code

```
public class Kids
{
    private static int countKids(int girlCount, int boyCount)
    {
        int totalKids;
        . . .
    }

    public static void main(String[] args)
    {
        int numGirls; int numBoys; int numChildren;
        . . .
    }
}
```

Static Variable Allocation

- Space for the main() method
 - Space for variables
 - Code to initialize variables
 - Jump statement
 - Print statement code
 - Exit
- Space for the countKids() method
 - Space for variables
 - Method code
 - Return statement

What's wrong?

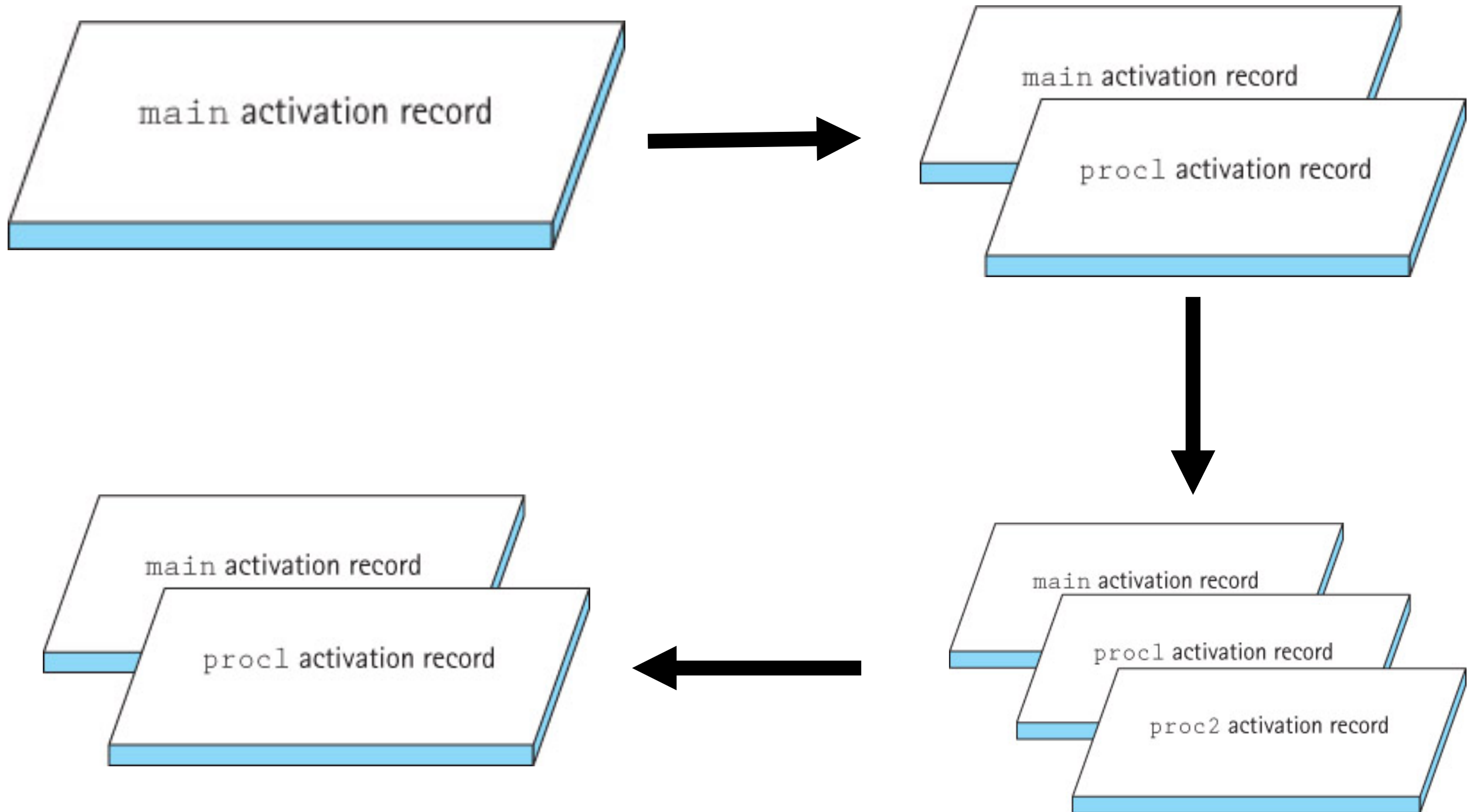
- Space allocated for each method is fixed in size
- Assumption: method will be called once and return before being called again
- Recursion breaks this assumption, so *purely* statically allocated languages don't allow recursion



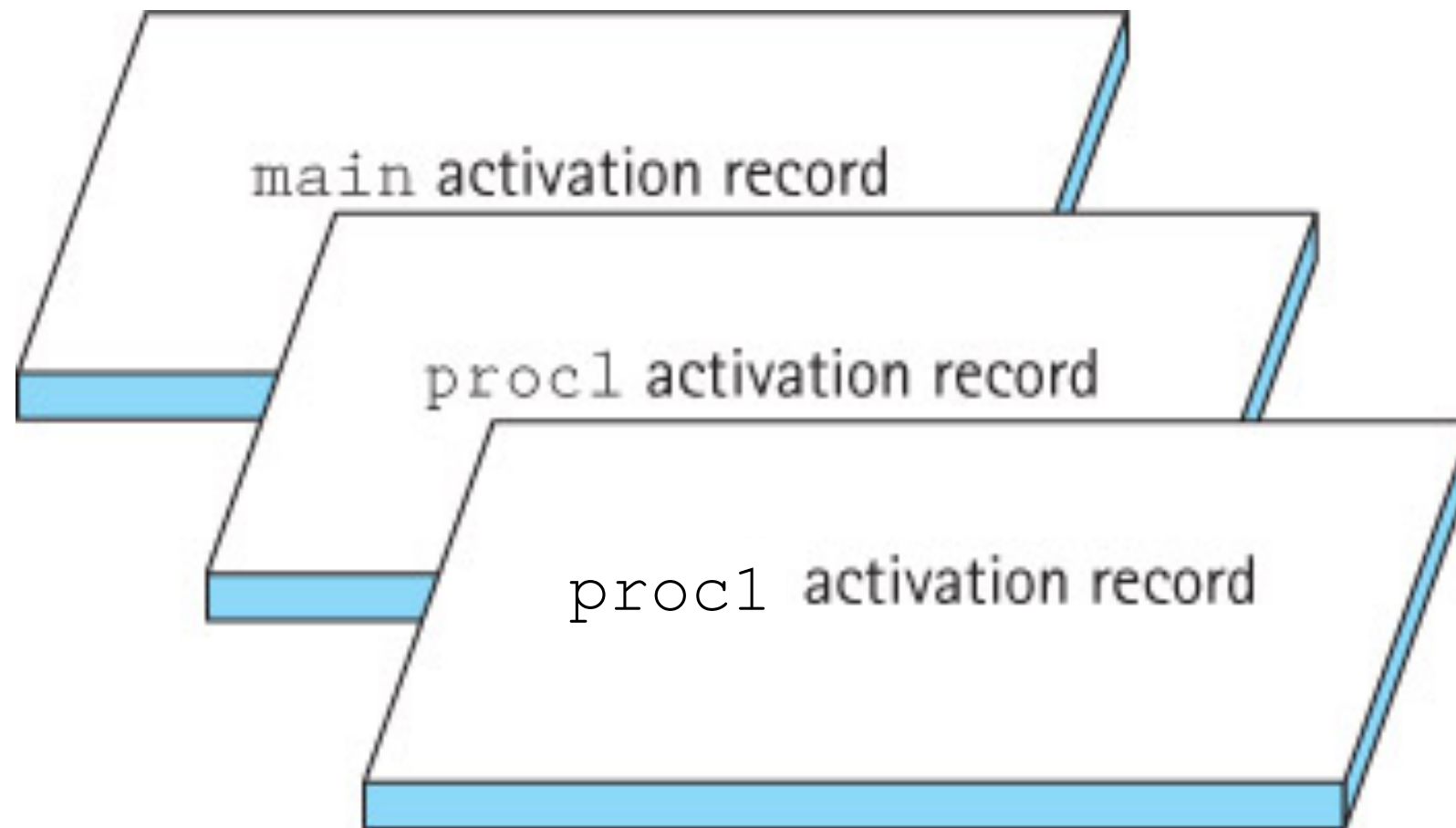
Dynamic allocation

- Method storage is allocated when the method is called
- Parameters, variables, and return address are nested in an activation record
- Activation records are stored on a stack to ensure they are handled and return in the correct order

Example Call Stack



Example Recursive Stack



Call Stack

- Stack ensures that all required return statements are handled in the correct order
- Depth of recursion refers to the depth of the call stack required
- Creating, storing, and handling activation records is the major reduction to efficiency in recursion



When to avoid recursion

- Some recursive functions have easily replicable functionality in a loop
- Some recursive functions can be emulated by explicitly using a stack
- Some recursive functions unnecessarily repeat computation



Recursion Overhead

- Traditional recursion calls the recursive function, performs an operation, and returns the result
- Note the nested structure requires all of the activation records to be maintained
- Requires stack space $O(n)$ for n recursive calls

Example Traditional Call

```
public static int factorial(int n)
{
    if (n == 0)
        return 1;
    else
        return (n * factorial(n - 1));
}
```


Tail Recursion

- Makes the recursive call last, passing a running tally into the call parameters
- After the base case is reached, call frames simply return in sequence
- Smart compilers account for this and deallocate activation records that aren't needed
- Requires stack space $O(1)$ if optimized



Example Tail Recursive Call

```
public static int factorial(int n,  
                             int current)  
{  
    if (n == 0)  
        return current;  
    else  
        return factorial(n - 1, current*n);  
}
```

Converting to a loop

- Tail recursion can often be converted into a loop
- Instead of passing the running total into a recursive call, maintain an explicit variable
- Loop over the input value until the base case is reached

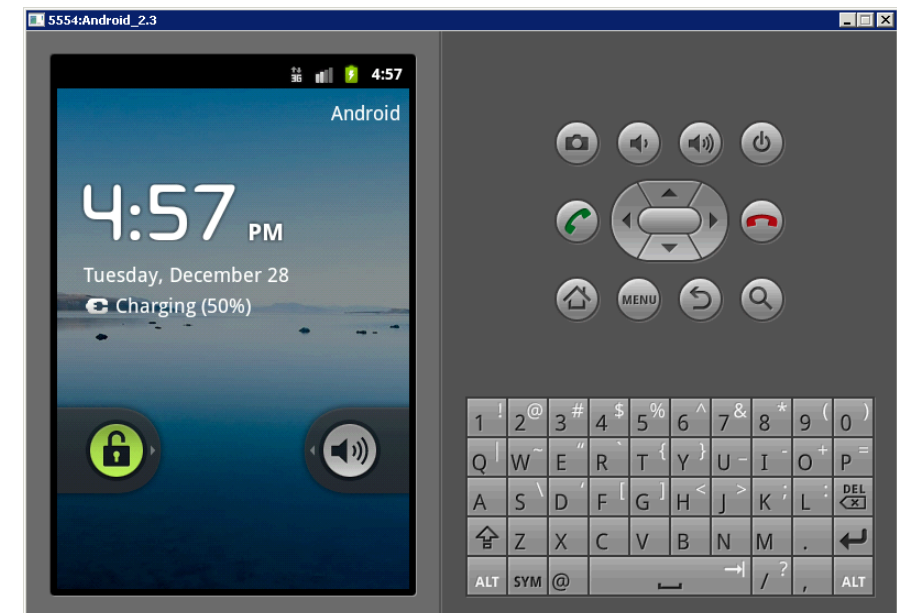


Example Loop Call

```
private static int factorial(int n)
{
    int retValue = 1;    // return value
    while (n != 0)
    {
        retValue = retValue * n;
        n = n - 1;
    }
    return (retValue);
}
```

Stack Emulation

- Recursion works because of the implicit use of a call stack
- Explicit use of a stack in a loop can emulate this behavior
- Requires additional complexity in the implementation code and data structure code
 - We're doing what the OS was doing for us in recursion



Recall recursive reverse print

```
void recPrintList(LLNode<String> listRef)
{
    if (listRef != null)
    {
        recPrintList(listRef.getLink());
        System.out.println(listRef.getInfo());
    }
}
```

Loop and stack version

```
static void iterRevPrintList(LLNode<String> listRef)
// Prints the contents of the listRef linked list to standard output
// in reverse order
{
    StackInterface<String> stack = new LinkedStack<String>();
    while (listRef != null) // put info onto the stack
    {
        stack.push(listRef.getInfo());
        listRef = listRef.getLink();
    }

    // Retrieve references in reverse order and print elements
    while (!stack.isEmpty())
    {
        System.out.println(stack.top());
        stack.pop();
    }
}
```

Efficiency issues: repeated work

- Recursive solutions are often clearer but may repeat calls
- Example: Combinations
- Recall, choosing k different elements from an n element pool



Combinations

```
public static int combo(int n,int k)
{
    if (k == 1)
        return n;
    else if (k == n)
        return 1;
    else
        return combo(n-1,k-1) + combo(n-1,k);
}
```

Example: 4 choose 3

Example: 6 choose 4

Solutions

- Remember previously calculated steps
 - Dynamic Programming
- Closed form solution
 - $N \text{ choose } k = N! / (N-k)!k!$
- Note that the closed form solution (factorial) can be easily implemented with a loop



Practice

- Convert traditional recursion to tail recursion:
 - Exponentiation
 - Summation
 - Multiplication
- Convert the same functions to a loop

Recap

- Variable space can be allocated dynamically or statically
- Dynamic allocation is critical to allow recursion but represents a hidden cost
- Recursion can be eliminated from several inefficient scenarios:
 - Tail recursion or looping
 - Explicit use of a stack
 - Closed form solutions or dynamic programming

Next Time...

- Dale, Joyce, Weems Chapter 4.1-4.2
 - Remember, you need to read it BEFORE you come to class!
- Check the course webpage for practice problems
- Peer Tutors
 - <http://www.csc.villanova.edu/help/>

