

Cluster Expansion of Thermal States using Tensor Networks

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Foreword

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Abstract

Het abstract is maximum één bladzijde en bevat minstens: a) De informatie die werd vermeld op het titelblad (eigen vorm); b) Een summiere beschrijving van het werk (vijftien à twintig regels); c) Eventueel: drie tot vijf goed gekozen trefwoorden die het onderwerp best omschrijven.

Extended Abstract

De extended abstract heeft een standaardlengte van minimaal twee bladzijden, met een maximum van zes bladzijden.

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Chapter 1

Introduction

There is nothing new to be discovered in physics now. All that remains is more and more precise measurement.

Lord Kelvin, 1900

1.1 Introduction

In 2015, there were about 5.6 million known physics papers in literature. At the current rate, this number doubles every 18.7 years [1]. Despite this enormous body of literature, there are a lot of things which are not completely understood. Some examples include a self-consistent theory of quantum gravity, the need for dark energy and matter in cosmology, the arrow of time, the matterantimatter asymmetry. There even is no interpretation of quantum mechanics where everyone agrees upon.

But certainly not all open problems have to do with 'new' physics. In many areas of physics, computing the implications of relatively simple laws becomes exceedingly difficult for many particles. Of historical importance is the three-body problem, describing the trajectory of 3 gravitational bodies such as the earth, moon and sun. The general case is not solved, despite developments over the last 300 years.

In reality, the real challenge is to model the macroscopic properties of quantum many-body system with around 10^{23} particles. Needless to say, this not an easy task at all. Finding good and computable approximations is of primary importance in the fields of quantum chemistry, condensed matter physics, and materials science.

In computational chemistry, the many-body problem is tackled with methods which fall in one of the following categories: (post-) Hartree-Fock methods, density functional theory (DFT) and force-field methods. While they have many

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<u>applications</u>, these methods are not fully able to capture all the properties of the so called strongly correlated matter.

Examples of phase of strongly correlated matter which are not yet understood include high-T superconductors, topological ordered phases, quantum spin liquids [2]. There exist different methods to investigate these exciting materials. A very limited number of models is quantum integrable, meaning they can be solved in a non pertubative way. Also, some properties of models near criticality can be determined exactly with conformal field theory (CFT). But for some systems, we can only simulate the behaviour with numerical techniques. To make progress, new fast and accurate numerical methods are needed, because exact diagonalisation becomes unfeasible for large systems.

Some examples of such numerical techniques, which will not be discussed, are: Dynamical Mean Field (DMFT) / Dynamical Cluster Approximation (DCA), Series expansion, Density Matrix Embedding Theory (DMET), Fixed-node Monte Carlo, Diagrammatic Monte Carlo, Variational Monte Carlo, Functional renormalization group (FRG) and Coupled-cluster methods. [3]

In this thesis, a technique is proposed that builds on the broad field of tensor networks.

1.2 Tensor networks

This is often referred to as the curse of dimensionality. The size of the Hilbert space of quantum states grows exponentially fast. This prevents an effecient description of all possible quantum states.

area law+picture

sign problem monte carlo,

write about tensor networks

Chapter 2

Tensor networks

The aim of this chapter is to give a basic introduction to tensor networks from a computational viewpoint. First the graphical notation for tensors, which is ubiquitous in the field, will be introduced and explained. A incomplete classification of the different kinds of tensor networks will be discussed. Some routine tensor network manipulations are explained. Also, a selection of different algorithms, which are from a computational point of view the reason behind the success of tensor networks, are presented very briefly.

2.1 Tensor networks

2.1.1 Graphical notation

Before explaing tensor networks, some graphical notation should be introduced. This really is a way to conveniently write vectors, matrices an in general tensors without the need to introduce many labels. A tensor T is represented by a circle with a number of external legs, according to the number of external indices. Connected legs are summed. Some examples are shown in table 2.1. Every leg which is connected to multiple tensors, is contracted.

Table 2.1: Caption

conventional	Einstein	tensor notation
\vec{x}	x_{α}	
M	$M_{lphaeta}$	$-\!$
$ec{x}\cdotec{y}$	$x_{\alpha}y_{\alpha}$	x— y



Figure 2.1: Caption

2.1.2 Representing a quantum state

Tensor network come in many shapes and forms. Tensor networks are really used to represent a tensor with many legs. A general quantum state with N sites can be described in a given basis $|i\rangle$ in the following way:

$$|\Psi\rangle = \sum_{i_1 i_2 \cdots i_n} C^{i_1 i_2 \cdots i_n} |i_1\rangle \otimes |i_2\rangle \otimes \cdots \otimes |i_n\rangle$$
 (2.1)

Here the tensor C holds all the information of the quantum state. The graphical representation can be seen in fig. 2.1.

This requires an exponential number d^n of coefficients C where d is the dimensions of basis $|i\rangle$. In order to make the problem tractable, the following form is proposed as wave function:

$$C^{i_1 i_2 \cdots i_n} = C^{1 i_1}_{\alpha_1} C^{2 i_2}_{\alpha_1 \alpha_2} \cdots C^{n i_n}_{\alpha_{n-1}}$$
 (2.2)

Where summation over shared indices is implied. It is always possible to find such a representation by means of matrix decomposition (see section 2.2). The summation over α_i are called a virtual bond and their dimension is denoted by χ . At this point, this is not yet an improvement because the bond dimension needs to be exponentially large in order to represent the tensor C exactly.

Explicit translational invariance is given by tensor $C_{\alpha\beta}^i$ that don't depend on the location. The chain is closed by a matrix M which contains the boundary condiditions. Setting $\alpha_n = \alpha_0$. We can now write this as a Trace over matrix products:

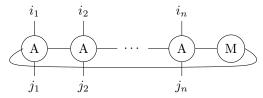
$$|\Psi\rangle = \text{Tr}\left(C^{i_1}C^{i_2}\cdots C^{i_n}M\right)|i_1\rangle \otimes |i_2\rangle \otimes \cdots \otimes |i_n\rangle \tag{2.4}$$

2.1.3 Classification

2.1.3.1 MPS

2.1.3.1.1 MPO Similarly, a Matrix Product Operator (MPO) is of the following form:

$$\hat{O} = \sum Tr(A^{i_1j_1}A^{i_2j_2}\cdots A^{i_nj_n}M)
\times |i_1\rangle \langle j_1| \otimes |i_2\rangle \langle j_2| \otimes \cdots \otimes |i_n\rangle \langle j_n|$$
(2.6)



The matrix M contains the boundary conditions of the operator. Many Hamiltonians can be represented by an MPO. For ins

2.1.3.1.2 PEPS Exact contraction is hashtag P-Hard:

No exact canonical form

2.1.3.1.3 PEPO

2.1.3.1.4 Others MERA,TTN,

2.2 Tensor network manipulations

This section serves as an introduction of tensor network manipulations. The overview mainly focusses on MPS/MPO networks, but most of the oprations translate to the 2D case.

The MPS's are processed by transforming the tensor into a matrix, performing some matrix calculations and casting it back into its original form. In this way, the standard methods from linear algebra can be used. This section gives some examples how this is done in practice:

2.2.1 Basics

2.2.1.1 Grouping legs

One of the most basic manipulations is to group some legs of a tensor into one leg:

$$T^{i_{1}i_{2}j_{1}j_{2}} = T$$

$$T$$

$$j_{1} \qquad j_{1}$$

$$\cong (i_{1}j_{1}) \qquad T$$

$$= T^{(i_{1}j_{1})(i_{2}j_{2})}$$

$$(2.7)$$

The dimension of the new leg is the product of the dimension of the individual legs. Contracting 2 merged legs with 2 merged legs is exactly the same as contracting them separately. The both The 4 leg tensor and matrix contain exactly the same information. Manipulating this in memory requires both permute and reshape commands. This requires some time, the internal representation of the matrix changes.

2.2.1.2 decomposition

The grouping above can be applied to decompose a tensor into 2 tensor with matrix techniques. An example, which will be needed later on, is give here.

The indices U,V and W represent blocks indices. Step 2 reshapes and groups the indices on to one index on the left and one on the right. The dimension of

this index is the product of the separate dimensions. Step 3 decomposes the matrix into a product of 2 matrices. The last step transforms the indices back to separate legs.

For an exact representation, the bond dimension of virtual level v is:

$$\dim v = \min(\dim u, \dim v) + 2\dim i \tag{2.9}$$

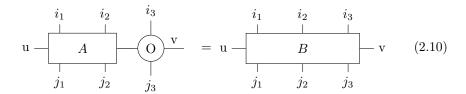
Many matrix decompositions exist. Some useful examples here are SVD decomposition, eigenvalue decomposition, QR, \dots .

2.2.1.3 virtual levels

In the previous example, the levels were indicate with a block index or virtual level. The idea is to create seperate the contraction into blocks. This is completely analogous to matrix block multiplication. This will be a more natural form to represent the algorithm. Of course, one can easily switch between block representation and the full one.

2.2.1.4 inverse

Suppose we want to find a MPO O for given tensors A and B such that the following holds:



Again, the indices can be taken toghether in the following way: $\alpha = (ui_1j_1i_2j_2)$ and $\beta = (i_3j_3v)$:

$$A_{\alpha\gamma}O_{\gamma\beta} = B_{\alpha\beta} \tag{2.11}$$

This a a standard matrix equation and can hence be solved with linear algebra packages. Note that it is not necessary to calculate A^{-1} to obtain the solution. Linear solver are generally much faster. As this is one of the core problems to solve both in 1D and 2D, this will be discussed in detail in section 5.1.

2.2.1.5 contraction order

contraction order

2.2.1.6 Gauge freedom

gauge

2.2.1.7 truncation

svd truncation

2.2.2 MPS algoritms

2.2.2.1 cononical form

schmidt decomp,

2.2.2.2 DMRG

2.2.2.3 Expectation values

Suppose that the there is an MPO representation of $e^{-\beta \hat{H}}$ A and that the mpo representation for X Y is localised over n sites, then the expactation value is given by:

$$\langle X \rangle = \frac{1}{\langle X \rangle} = \frac{1}{\langle$$

In the thermodynamic limit there are an infinity number of A to the left and the right. This can be simulated by taking the left and right fixed points of the traced MPO A corresponding to the largest eigenvector λ .

$$G_l - A \longrightarrow = \lambda G_l \longrightarrow (2.13)$$

Equation eq. (2.12) can now be easily callulated:

$$\langle X \rangle = \frac{1}{\langle X \rangle} - \frac{1}{\langle X \rangle} - \frac{1}{\langle X \rangle} = \frac{1}{\langle$$

2.3 Tensor network algorithms

2.4 MPS algorithms

This section, and the following one, will introduce some different tensor network algorithms. The gaol is to provide an intuitive expanation how these algorithm work. For rigorous derivations and mathematical details, other sources can be read.

2.4.1 Canonical form

A translation invariant MPS has the following form:

It can be easily seen that inserting XX^{-1} on each bond doesn't change the contracted tensor.

$$A_l = \qquad \qquad = X \qquad \qquad X^{-1} \tag{2.17}$$

fix baseline

$$A_r = \qquad \qquad = Y \qquad \qquad Y^{-1} \tag{2.18}$$

and

Then eq. (2.16) can be written as follows:

$$(2.20)$$

Introducing one more tensor:

At the moment the matrices X and Y are not yet defined. To bring an MPS A in its unique canonical form, the following choice is made

more on MPS

2.4.1.1 DMRG

2.5 2D tensor network contraction

PEPS contraction concerns the following problem:

Contract an infinite lattice of identical tensors, with some irregularities on a small patch. For example, the one patch can be used to calculate an expectation value of a certain observable.

2.5.1 Overview methods

- 2.5.1.1 Real-space renormalization-group methods
- 2.5.1.2 corner transfer matrix methods
- 2.5.1.3 Boundary methods
- 2.5.1.3.1 density-matrix renormalization group
- 2.5.1.3.2 time-evolving block decimation [4]

2.5.2 Vumps

The porpuse of this section is to give some intuition on the variational uniform Matrix Product State (VUMPS) algorithm, which will be used later on in this thesis.

The goal is to find a MPS layer for the MPO such that:

$$\approx \qquad (2.24)$$

This is very similar to section 2.2.2.3. The expression holds approximately, because the MPS on the left hand side has a larger bond dimension than on the right hand side.

2.5.2.1 The equations

Suppose there are tensor which fullfil the conditions stated below:

$$Gr = \lambda \qquad Gr \qquad (2.25b)$$

These eigentensor equations are solved in pratice in a slightly different manner:

$$\begin{array}{c|c}
\hline
Gr \\
\hline
\end{array} = \lambda \quad \begin{array}{c}
\hline
Gr \\
\hline
\end{array}$$
(2.26)

Wher the Ar tensor was connected from below on both sides and eq. (2.22) was used on the right hand side. The blocks in eq. (2.25) form a zipper from the left and right respectively. Each application brings down one of the MPS tensors in the following way.

The left and right zipper can now move towards each other, until they meet at the center. In order to become the same MPS as before the application to the MPO (eq. (2.24)), one more condition is needed:

$$Gl \qquad Gr \qquad = \lambda \qquad (2.28)$$

This completely determines the problem, but one more equation is used to solve the problem. Combining on of the equations eq. (2.25), the definition of A_c

eq. (2.28) and one of eq. (2.22) gives C:

$$Gl \qquad Gr \qquad = \lambda \qquad (2.29)$$

fix gap

Then eq. (2.24) is solved:

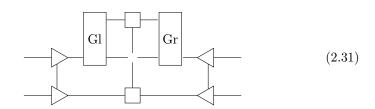
check lambdas in literature

Contracting a 2D tensor network is thus reduced to solving the eq. (2.21), eq. (2.25), eq. (2.28) and eq. (2.29) simultaneously. Inspection of the equations show that the following cycle needs to be solved:

- $A_c, C \to A_l, A_r \text{ eq. } (2.21)$
- $A_l, A_r \to G_l, G_r \text{ eq. } (2.25)$
- $G_l, G_r \to Ac, C$ eq. (2.28) and eq. (2.29)

The calculated environment can now be used to solve the original problem. Due to symmetry, the same MPS can be applied from below. eq. (2.29) now

becomes:



2.5.2.2 The derivation

While the above derivation is reasonable, it is not very rigorous. The algorithm finds its origins in tangent space methods, such as explained in [5].

Not every state in the many body Hillbert space can be represented by an MPS. For a

By carefully constructing the tangent space and making use of the available gauge freedom, a compact expression can be found for the tangent space projector \mathcal{P}_A . This projects a state from the many body Hillbert space to the tangent space of an MPS A. For an optimal MPS approximation A, the error made in the approximation should be orthogonal to the tangent space. For eq. (2.24) this means that the application of the projection of the error on the tangent space should be zero, i.e. the MPS is at a variational minimum. [4]

Applying the projector \mathcal{P}_A to eq. (2.24) exactly gives rise to the equations stated earlier.

Chapter 3

Strongly correlated matter

3.1 Phases and Criticality

3.1.1 Phases of matter

An important area of research is the study of the different phases of (quantum) matter. A phase is a state of matter in which the macroscopic physical properties of the substance are uniform on a macroscopic length scale. These phase can be measured by thermodynamic function, i.e. by function of a few macroscopic parameters. [6]. More precisely, for a given phase the properties vary as an analytic function of the macroscopic variables.

Interesting physics happens at the boundary between 2 or more distinc phases. The phase transitions were classified by Ehrenfest [7], who looked at the free energy across the phase boundary. If the free energy shows a discontinuity, it is called first order (or discontinuous) phase transition. Similarly, if the derivitive shows a discontinuity, it is called second order (or continuous). Higher order phase transitions are possible, and there are even examples of infinite order transitions, such as the BKT transition.

3.1.2 symmetry breaking

Sometimes, but not always, a phase transition is related to spontaneous symmetry breaking. A state $|\Psi\rangle$ is said to be symmetric under a unitary transformation U if the state only changes by a phase factor: $\hat{U}|\Psi\rangle = e^{i\phi}|\Psi\rangle$. A hamiltonian posesses a symmetry if it commutes with U: [H,U]=0 [8]. A remarkable fact is that many ground states are not invariant under a symmetry U of the hamiltonian.

For phase transitions associated with a broken symmetry, one can define an order parameter. This parameter evalutes to 0 for the symmetric phase, but not for the sponteous broken phase.

In continuous or second-order phase transitions the order parameter increases continuously from zero as the critical temperature is traversed. The

entropy also changes continuously. On the other hand, the correlation length and related energy scales diverge at the critical temperature. In fact, at the critical temperature of a second-order phase transition, scale invariance systems become scale-invariant, in the sense that physical properties no longer depend on the length (or energy) scale at which they are probed. Many symmetry-breaking phase transitions are second-order, with the onsets of superfluidity, (anti)ferromagnetism and many phases of liquid crystals as famous examples.[8]

3.1.3 Universality

Universality looks at the behaviour of the system near a continuous phase transition. These can be discribed well by so called power laws. For classical phase transitions (driven by temperature) near critical temperature T_c , observables a_i depend in the following way on the reduced temperature $t = \frac{T - T_c}{T_c}$: $a_i(t) \sim t^{\alpha_i}$. One would expect that the set of critical exponents α_i depends on the precise form of the hamiltonian of the system, but it turns out these exponents can be captured by a limited numer of universality classes. This means that the physics near criticality is completely understood once it is understood for one member of the class.

3.1.4 Critical exponents for spin systems

The following table defines some of the critical exponents for the Ising system.

Symbol	name
m	magnetisation
ξ	correlation length
g	external field
t	reduced temperature
d	dimension

The 2 point correlation function is defined as $f(x,y) = \langle m(x)m(y)\rangle - \langle m(x)\rangle \langle m(y)\rangle$. At larger distances this decays exponentially fast (see ??) $f(x,y) = e^{-\frac{|x-y|}{\xi}}$, where ξ is the correlation length.

for the ordered phase, the following relations hold: $m |t|^{\beta}$, $\xi(t) \approx |t|^{-\nu}$. At the critical temperature near a quantum phase transition $m \approx |g - g_c|^{\frac{1}{\delta}}$.

3.1.5 Finite size scaling

Finite size scaling was originally introduced in order to extract critical exponents from monte carlo simulations.

3.1.6 CFT

3.1.7 Quantum phase transitions

A traditional 2nd order phase transition is driven by a change in temperature. Quantum phase transitions on the other hand happen at zero temperature under influence of another parameter g of the model. At finite temperature, 2 things can happen: either there is a line connecting a classical 2nd order phase transition to the quantum phase transition, or the phase transition disappears at finite temperature [9].

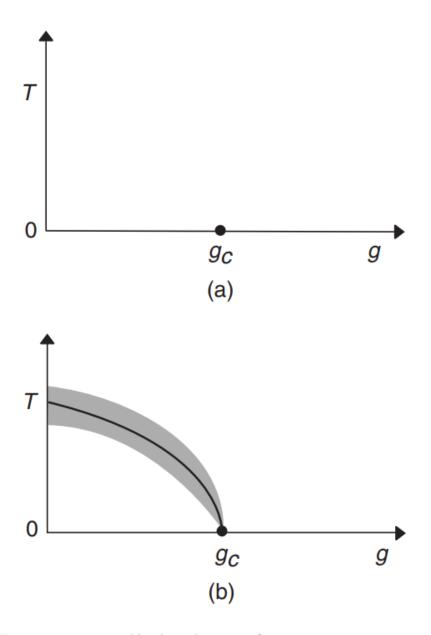


Figure 3.1: Two possible phase diagrams of a system near a quantum phase transition. In both cases there is a quantum critical point at $g=g_c$ and T=0. In (b), there is a line of T>0 second-order phase transitions terminating at the quantum critical point. The theory of phase transitions in classical systems driven by thermal fluctuations can be applied within the shaded region of (b). Figure and caption taken from [9].

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3.1.8 Quantum to classical mapping

Quantum to classical mapping

3.2 Models

The goal of numerical techniques is to simulate the physics of real world systems. These are, to some extend, captured by different models. Models are a simplified mathematical description that captures some relevant physics of more complicated systems. This section introduces some specific models, their relevance and some properties. These models will be used later to benchmark the developed tensor network expansion.

3.2.1 Ising model

The prototypical example of a model in the field of strongly correlated matter is the Ising model. It was first introduced 1925 by Ernest Ising, as a model to capture ferromagnetism. He proved that for a linear chain, there is no phase transition at finite temperature. He wrongly concluded that this would also be the case in higher dimensions, but it turned out t be one of the deepest and far-reaching problems in 20th century [10].

The Ising model, in essence, assigns an energy contribution to neighbouring spins. These spins sit on a fixed position on a chain (1D) or lattice (2D/3D/...). In classical Ising, the operators in the Hamiltonian all commute with each other. An energy is assign between neighbouring spins and possible and energy for alignment with an external magnitic field in the same direction. In quantum Ising model, a transversal field is added. Often, the particles on the grid are spin 1/2 particles, but of course other particles are possible.

Many generalizations exist for the Ising model.

q potts,...

3.2.1.1 Classical Ising

The classical ising model is given by the following hamiltonian:

$$H = -J\left(\sum_{\langle ij \rangle} \sigma_i \sigma_j + h \sum_i \sigma_i\right) \tag{3.1}$$

where $\langle ij \rangle$ runs over all neighbouring lattice sites. The possible values of σ depends on the spin dimension. For spin 1/2 lattices $\sigma \in -1, +1$. g encodes the interaction strength of the external magnetic field.

The sign J determines the low temperature ground state. A positive J will tend to allign all neighbouring spins at low temperature. This is often called ferromagnetic, because all the aligned spins cause a macroscopic magnitisation. On the other hand, a negative J causes neighbouring spins to have an opposite sign.

Depending on the sign of the longitudinal field h, the spins tend to align or anti align with this external field. This lifts the degeneracy of the groundstate.

3.2.1.1.1 1D The classical 1D model was solved analytically by Lens.

3.2.1.1.2 2D In 2D, it becomes important to define the lattice. Here, and in the simulations, we will consider a square lattice. This model was famously solved by Lars Onsager in 1944, by using the transfer matrix method. In 2 dimensions, the Ising model has a phase transition at finite temperature. The critical temperature is $T_c = \frac{2J}{T \ln(\sqrt{2}+1)}$.

Only the h=0 case the is solved analytically. For higher dimensions, no analytical solution is known. For these cases, we need to use numerical techniques if we want to understand the behaviour of these models.

On different lattices, interesting thins can happen. For instance, the ground-state of an antiferromagnet on a triangular lattice is not obvious to determine. The spins tend to antiallign, but at least 2 of 3 spins on the corner of a triangle have to align. Remarkably, as will be explained in section 3.1, the physics at the phase transition does remain invariant when the lattice is changed.

3.2.1.2 Quantum Ising

As we all know, the real world behaves, certainly at small length and time scales, quantum mechanically. Therefore, it is important to understand how the quantum Ising model differs from the classical model. In the quantum Ising model, the operators no longer commute with each other. An example is the transversal Ising model given by the following hamiltonian:

$$\hat{H} = -J \left(\sum_{\langle ij \rangle} \sigma_i^x \sigma_j^x + g \sum_i \sigma_i^z \right)$$
 (3.2)

In the case that g = 0, this is the classical Ising model (in the h = 0 case).

3.2.1.2.1 1D Different to the classical case, the 1D model already contains a phase transition.

3.2.1.2.2 2D

3.2.2 Heisenberg

The heisenberg model is given by:

$$\hat{H} = -\left(\sum_{\langle ij \rangle} J_x \sigma_i^z \sigma_j^z + J_y \sigma_i^y \sigma_j^y + J_z \sigma_i^z \sigma_j^z + h \sum_i \sigma_i^z\right)$$
(3.3)

These models have different names depending on the values of J_{α} with $\alpha = x, y, z$. $J_x = J_y \neq J_z = \Delta$ is called the XXZ model.

3.2.3 Random

It's also possible to construct random hamiltonians.

in basis: hermitian H

3.3 Operator exponentials

While it is often possible to find exact MPO representation to represent a wide class of hamiltonians , it is much harder to do the same for exponentiated operators. These operators play an important role: they act as time evolution operators for quantum systems $|\Psi(t)\rangle = \exp\left(-i\hat{H}t\right)|\Psi(0)\rangle$. A very similar operator governs the partition function in statistical mechanics: the probability of finding a system at inverse temperature $\beta = \frac{1}{T}$ in a microstate i is given by $p_i \exp\left\{-\beta \hat{H}_i\right\}$. This is often called "imaginary" time, due to the substitution $\beta = it$. The ability to calculate these operators is essential for understanding the dynamics of a given quantum model, and making contact with real world obsevations of these systems at finite temperature.

find citation and examples

explain link CFT

3.3.1 Statistical mechanics

The physics of a system in thermodynamical equilibrium can be derived from it's partition function Z. The classical formula generalises to a density matrix ρ as follows:

$$Z = \sum_{n} e^{-\beta E_{n}}$$

$$= \sum_{n} \left\langle n \left| e^{-\beta \hat{H}} \right| n \right\rangle$$

$$= \operatorname{Tr}\left(e^{-\beta \hat{H}}\right)$$
(3.4)

The first line is the partition function for clasical discrete systems. The index n runs of all possible microstates. It is known that the propability to find the system in a given microstates is given by:

$$p_i = \frac{\sum e^{-\beta E_i}}{Z} \tag{3.5}$$

An useful quantity is the density matrix ρ .

$$\rho = \sum_{j} p_{i} |\Psi_{j}\rangle \langle \Psi_{j}|$$

$$= \sum_{j} \frac{e^{-\beta \hat{H}}}{Z} |\Psi_{j}\rangle \langle \Psi_{j}|$$
(3.6)

Whith this notation, the partition function Z and ensemble average of an operator \hat{X} are given by:

$$Z = \text{Tr}(\rho)$$

$$\langle X \rangle = \text{Tr}(\rho \hat{X})$$
(3.7)

3.3.2 Time evolutions

LSZ theorema herbekijken

In quantum field theory, , calculation of n-point correlation functions is extremely important to understand a given field theory.

3.3.2.1 ground state

One practical way of finding the ground state is cooling an intitial stata down very small T.

3.3.3 Tensor network methods

In the following section I will give a very short review of the current tensor network methods to simulate real or imaginary time evolution. This overview is mainly based on the review paper [11].

[11]

3.3.3.1 Approximations to $\hat{U}(\delta)$

TEBD, MPO $W^{I,II}$

3.3.3.2 global Krylov method

3.3.3.3 MPS-local methods

local Krylov TDVP

Chapter 4

Construction Cluster expansion

4.1 Introduction

4.1.1 Notation

In the following, the external legs and virtual level 0 will be omitted:

$$\begin{array}{c|c}
i \\
\hline
0 \\
j \\
\end{array} = \bigcirc$$
(4.1)

$$\begin{array}{c|cccc}
i_1 & i_2 \\
\hline
0 & 1 & 0 \\
\hline
j_1 & j_2
\end{array} = \begin{array}{c}
1 \\
\hline
\end{array} (4.2)$$

This hamiltonian consists of 1 and 2 site operators. Of course more general

hamiltonians can also be captured.

$$\hat{H} = \left(\sum_{\langle ij \rangle} H_2^i H_2^j + \sum_i H_1^i\right) \tag{4.4}$$

The same notation will be used to denote the hamiltonin evaluated on the given geometry:

$$H \left(\bigcirc \longrightarrow \bigcirc \right) = H_1 \otimes 1 \otimes 1$$

$$+1 \otimes H_1 \otimes 1$$

$$+1 \otimes 1 \otimes H_1$$

$$+H_2 \otimes H_2 \otimes 1$$

$$+1 \otimes H_2 \otimes H_2$$

$$(4.5)$$

4.1.2 Idea

This chapter shows the main construction of dissertation. A cluster expansion is used to approximate $e^{\hat{H}}$ for every possible geometry. The goal is to make a MPO/PEPO which captures the tensor exponential in the thermodynamic limit.

This cluster expansions introduced in [12]. The main idea is to make an extensive expansion by adding blocks which solve the model exactly on a local patch. Crucially, the expansion is not in the inverse temperature β but in the size of the patches. The local patches are separated by a virtual level 0 bond.

To make this somewhat more precise, the first steps of the expansion are shown here. The smallest patch, i.e. 1 site, encodes the exponential of that hamiltonian.

$$\bigcirc = \exp(-\beta H(\bigcirc)) \tag{4.6}$$

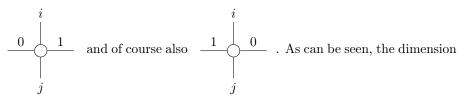
If there were no 2 site interaction, this already captures the full diagonilsation. Of course, such a model wouldn't be useful. The next step is to introduce 2 site interactions, where the one site interactions previously introduced interaction are subtracted from the diagonalised hamiltonian.

$$\underbrace{\begin{array}{c}
1\\
\bigcirc = \exp{-\beta H}(\bigcirc \longrightarrow \bigcirc)\\
-\bigcirc \longrightarrow \bigcirc
\end{array}} \tag{4.7}$$

At this stage, all seperated networks with maximally 2 connected sites in a row are diagonalised exactly. Notice that here, 2 new blocks are introduced:

symmetry, speed

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of sublevel 1 needs to be d^2 , with d the dimension of physical level. Although different possible constructions already differ in the next step, one more step is added to make te construction and notation clear.

It is clear that the right-hand side of eq. (4.8) can also be ommited, as it is just evaluating the exponentiated hamiltonian on the same geometry as the left hand side and substructing all possible contractions of the blocks which were added previously.

4.1.3 Preview

In the following sections some possible configurations in 1D and 2D will be discussed. At this point, the focus is on the construction and its bond dimension. The results will be discussed later in chapter 6. T

4.2 Construction MPO

4.2.1 Type A

This type was originally proposed in [12]. The first few blocks in the expansion are the following:

The following types of blocks appear in the cluster expansion

n-1.

The O^{nn} block is in defined for a chain with an odd number of sites. The contraction of O^{nm} and O^{mn} is defined by for a chain with even order. The decomposition is defined up to a gauge transformation.

- **4.2.1.0.1 Dimension** In this scheme, virtual level n has dimension d^n . Of course, this dimension can be lowered if some error is allowed for the longest chain.
- **4.2.1.0.2** discussion Type A can form long chains, which where not explicitly optimised for. The question arise whether this will results in accurate results for cyclic systems or not.

4.2.2 Type B

Type B only contains blocks of the following form; O^{mn} and O^{n0} . The first few blocks are:

$$\begin{array}{c|c}
i_n & i_{n+1} \\
\hline
\mathbf{m} & 0 & 0 \\
j_n & j_{n+1}
\end{array} = U^n \Sigma V^{\dagger} \tag{4.11}$$

The following split is made: $O^{mn} \cong U^n$ and $O^{n0} \cong \Sigma V^{\dagger}$. In this way the left inverse exists and doesn't need any calculation: $O^{mn} = U^{\dagger}$.

4.2.2.0.1 dimension From the construction the bond dimension grows from the left to the right. For the last step, there are only d^2 non zero singular values. Each steps adds d^2 to the dimension. For the last step, only d^2 non zero singular values need to be keeped. With the following natation:

$$\begin{array}{ccc}
 & i & \\
 & & \\
\hline
 & & \\
 & j & \\
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The bond dimension of lower virtual levels can be reduced if we can solve the following equations simultaneously:

Then the MPO doesn't change if there are matrices $A^{\prime n},\,A^{\prime n+1}$ and $B^{\prime n}$ such that

$$S = A^{m}A^{n} = A'^{m}A'^{n}$$

$$T = A^{m}B^{n} = A'^{m}B'^{n}$$
(4.13)

Such matrices with optimal bond dimension can be found with generalised SVD. Generalised SVD decomposes 2 matrices as follows:

$$S^{\dagger} = (U\Sigma_1)Q^{\dagger}$$

$$T^{\dagger} = (V\Sigma_2)Q^{\dagger}$$
(4.14)

The new bond dimension is the dim $n' = d^2 \cdot \min(\dim n - 1, \dim(n+1))$. This is higher than the dimension of type A.

4.2.2.0.2 Discussion The bond dimension is larger than type A, but the long chains from A are absent. The left inverse is always well defined and doesn't need any computation, because hermitian matrix U can be inverted easily. One major drawback is that for long chains, the virtula bonds are very large before they can be shrunk with the gsvd procedure.

4.2.3 Type C

primed virtual levels

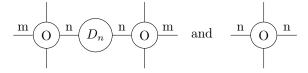
This type implements the same strict type as Type B, but in a different way. No calculation is involved, except the calculation of the the exponentiated hamiltonian to certain order. The following kind of MPO strings are allowed:

and so forth. All but one MPO elements are chosen to be the identity matrix. The middle one is the exponentiated hamiltonian with reshaped legs.

4.2.3.0.1 discussion As can be expected from the construction, the bond dimension grows very fast. This type is just as precise as Type B.

4.2.4 Type D

This type uses a different setup which tries to capture the best of both Type A and B. Type could handle long range correlation better because of the introuction of O^{nn} , but the inverse was not necessarily well defined. Type B had well conditioned inverses, but performed in most of the cases worse. The block appearing in type D are as follows:



Similar to type A,

Matrix D_n is the singular value diagonal matrix devided by a normalisation factor ϕ . Both U and V are multiplied by $\sqrt{\phi}$.

4.2.4.0.1 discussion It's not completely clear what the values of ϕ should be. If ϕ is to large, large chains are not surpressed. If phi is to small, the O^{nn} blocks will become large and hence the chain will diverge again. A reasonable value is the sum of the singular values.

other things could be tried here, WIP

4.2.4.0.2 matrisation The cost of this type lies in the fact that it has no compact way of casting it to a matrix. The following works, but has quite a large dimension:

O_{00}	O_{01}		$-2O_{01}$		O_{01}		$O_{01}D_1^{1/2}$	
$O_{00} \\ O_{10}$		O_{12}		$-2O_{12}$		O_{12}	_	
	O_{21}							
O_{10}								
	O_{21}							
O_{10}					O_{11}			
	O_{21}					O_{22}		
$D_1^{1/2}O_{10}$								$D_1^{-1/2}O_{12}D_2^{1/2}$
1							$D_2^{1/2}O_{21}D_1^{-1/2}$	fix this
					I		1 - 2 - 21 - 1	

4.2.5 Type E

Again, this is a strict variant which needs exactly twice the bond dimension of type A. The idea is to split every chain in a left and a right part. For the left part, the numbers increase while right part they decrease. This construction carries over well to higher dimensions. The first few blocks are:

The construction is very similar to type A

4.2.6 Type F

The idea behind this type is very similar to type D. The blocks look as follows:

The blocks $O^{n-1,n}$ and $O^{n,n-1}$ unitary matrices from the svd decomposition, scaled by the largest singular value. The blocks $O^{n-1,n'}$ and $O^{n',n-1}$ are then used to actually solve the problem for the given chain. In the next step, $O^{n,n}$ is added as usual. The idea here is once again to keep this block small, in order to not cuase any divergences.

4.2.7 Conclusion

Many 1D constructions are discussed here. They can be roughly divided in 2 groups. B, C and E are strict variants, meaning only the explicitly constructed blocks will appear in the final expansion. As a consequence, they have exactly the same predictive power. While B has some advantageous properties such as its final bond dimension and well-defined inverses, type E will be used in the results chapter due to its simplicity and scalability. It also generalises well to 2D, in contrast to type B.

The second category are the unstrict types. Type A has the lowest possible bond dimension to exactly represent a chain of a given length. One hurdle to be overcome are the badly conditioned inverses, when implemented naively. Types D and F try to remedy this. D scales very badly with the maximum number of sites, and has a construction which doesn't fit in with the simple diagrams. The construction was only implemented in 1D code due to this. This type won't be reported, and has a similar performance to type F.

In short, type A, E and F will be reported in section 6.2.

4.3 Construction PEPO

While there where some initeresting choices in the 1D construction, the number of possibilities in 2D is virtually limitless. The focus will mainly be to generalise type A to 2D. As can be expected, the construction starts of quite similar:

$$\bigcirc \qquad (4.19)$$

From 3 blocks onwards, the number of extra blocks start to increase:

Of course, besides linear blocks also T and + shaped need to be constucted to fully capture the model.

$$\begin{array}{c|c}
 & (4.22) \\
 & 1 \\
 & 1 \\
 & 0
\end{array}$$

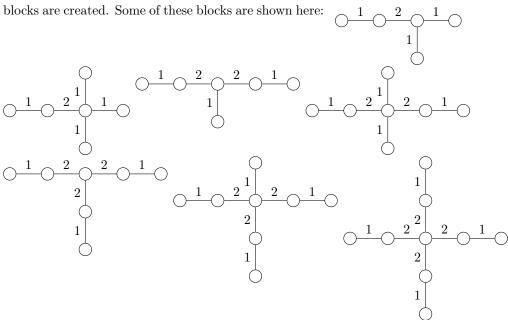
$$\begin{array}{c|c}
 & (4.23) \\
 & 1 \\
 & 1 \\
 & 1
\end{array}$$

Here, and in the following, care has to be taken in which order the blocks are made. In general, every block which fits in the given block needs to be constructed earlier. The construction continues by introducing a 2 block as before:

$$\bigcirc \frac{1}{2} \bigcirc \frac{2}{1} \bigcirc \Rightarrow \tag{4.24}$$

Again, 2 blocks are introduced. For other variations of the linear chain, only one block needs to be solved

Due to the way they are constructed, the error for every linear chain of a given lenght will be de same as in the 1D case. Once again, all possible T and +



For each block show, there are still multiple permutations of the legs possible. It is clear that a completely automated solver is needed to construct all these different blocks. From here on the construction generalises easily to higher block numbers, and to higher dimensions.

4.3.1 loops

While the blocks above certainly encode a large number of finite size patches, there are still quite some patches need to be encoded. The simplest case is a 1 square loop.

$$(4.26)$$

It is clear that this problem cannot be solved with the techniques from the previous section. The square loop needs a new virtual level α . In general, all the loop levels will be named with greek letters for convenience. The simplest choice for the loop is:

$$\begin{array}{c|c}
\alpha & \alpha \\
\alpha & \alpha \\
\hline
 & \alpha
\end{array}$$
(4.27)

At this points all blocks of order 4 are solved.

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4.3.1.1 Single extensions

The loops need to be connected to the linear blocks. One way to do this is as follows:

$$\begin{array}{c|c}
\alpha & \alpha \\
\alpha & \alpha \\
\hline
1 & \alpha \\
\hline
\end{array}$$
(4.28)

With the given blocks, the following combination is also possible

$$\begin{array}{c|c}
 & \alpha \\
 & \alpha \\
 & \alpha \\
 & 1 \\
 & 1
\end{array}$$
(4.29)

This results in very large errors, and it would require many more blocks to be added in order to counteract this. Luckily, there are many options in 2D. One example is:

$$\begin{array}{c|c}
\alpha \\
\uparrow \\
\uparrow \\
\beta \\
\alpha
\end{array}$$
(4.30)

No other combinations are possible, except a loop with on one corner an extension. Of course, there is no need to stop here, the following blocks can now be constructed easily.

4.3.1.2 Double extensions

It seems as if one of the corner pieces can be used as follows:

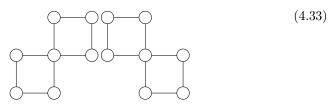
But in other to make a meaningful change in the residual error, the bond dimension of both α and β need to be enlarged significantly. It is more advantageous to introduce yet another level δ , which forms the link between the 2 parts.

As both corner tensors can be optimised at once, the total bond dimension is lower than for the previous suggestion, but still larger than the dimensions of the other loop levels.

$$\begin{array}{c|c}
 & \alpha \\
 & \gamma \\
 & \delta \\
 & 1 \\
 & 0
\end{array}$$
(4.32)

4.3.1.3 Larger loops

One question which comes to mind is where the focus should be for the construction. Making blocks for all possible single loop extensions comes at a increasing cost in total bond dimension. From a physics point of view, the model will be better approximated when smallest non solved patch is solved by introducing new blocks. On the other hand, the already included blocks may cause an error which was not present before the blocks were introduced. One example is a linear chain which closes upon itself, in for instance a 2x3 rectangle. Another example are + blocks which connect upon themselves in the following shape:



One way to solve this, without breaking the rotation pattern of β and γ , is

$$\begin{array}{c|c}
 & \alpha \\
 & \gamma \\
 & \alpha \\
 & \alpha \\
 & \alpha
\end{array}$$

$$\begin{array}{c|c}
 & \alpha \\
 & \alpha \\
 & \alpha
\end{array}$$

$$\begin{array}{c|c}
 & \alpha \\
 & \alpha \\
 & \alpha
\end{array}$$

$$\begin{array}{c|c}
 & \alpha \\
 & \alpha
\end{array}$$

But this block introduces an infinite tiling, corrupting the expansion.

4.3.1.4 failed ideas

Here it starts to get really tricky.

Chapter 5

Framework implementation

An expert is a person who has made all the mistakes that can be made in a very narrow field

Niels Bohr

5.1 Solvers

The construction can be written down quite compactly as done in the previous section, but actually implementing the code in full generality is somewhat more complex. In practice, the 1D and 2D implementation were performed separately. where a part of the code but mainly the ideas and lessons learned from 1D where taken to the 2D code. Of course, the 1D construction is just a subset of the 2D implementation. In fact, the 2D implementation even outperforms the 1D implementation for reasons with will be explained later. The main focus of this chapter will go to the 2D implementation. At the end some particular optimisations in 1D will be highlighted.

Solving the blocks introduced in the previous sections need 2 different approaches. The graphs of the maps can be split in 2 groups. The first one, considers problems where there are no loops and at most one node with more

will be reduced to a standard matrix problem, and solve with matrix (pseudo-) inversion.

The other group, of course, constitutes the nonlinear problems. This includes every problem where a block (or rotated version) occurs more than once, problems which include loops, ...

5.1.1 Linear solver

The linear solver is a general purpose block solver which reduces the problem to a set of linear matrix equations. Linear block consist of a tree structure, where the new block is the root of the tree, and all the branches need to be inverted. Let $I^m = (i_1^1 i_2^1 \cdots i_{n_1}^1)$, then the problem can in general, after some tedious tensor leg bookkeeping, be rewritten in the following form:

$$A_{I_1 I_2 \dots I_n \alpha^1 \alpha^2 \dots \alpha^m} X_{\alpha^1 \alpha^2 \dots \alpha^m j}$$

$$= B_{I_1 I_2 \dots I_n j}$$

$$(5.1)$$

Here i_N^M has the following meaning: M numbers the different legs or branches of the tree, N number of sites of the leg and i numbers the bra and ket states and has dimension d^2 . Hence the bond dimension of $I_n = d^{2n_m}$. The most obvious way to solve this system is by using a linear solver. The problem is that the bond dimension increases very fast: matrix A has dimension $d^2 \sum_m n_m \times d^2 \sum_m n_m$. Although using a linear solver instead of full inversion is considerably faster, this becomes infeasable for very quickly. A second method consist of solving the following sequence of linear problems one leg at a time:

$$A_{I^{1}\alpha^{1}}^{1}X_{\alpha^{1}I^{2}...I^{m}j} = B_{I_{1}I_{2}...I_{n}j}$$

$$A_{I^{2}\alpha^{2}}^{2}X_{\alpha^{1}\alpha^{2}I^{3}...I^{m}j} = B_{\alpha^{1}I_{2}...I_{n}j}$$

$$\vdots$$

$$A_{I^{m}\alpha^{m}}^{m}X_{\alpha^{1}\alpha^{2}...\alpha^{m}j} = B_{\alpha^{1}\alpha^{2}...\alpha^{m-1}I_{m}j}$$
(5.2)

While this method is very quick and scales well, in practice it results in unstable result. This is a result of the potentially ill conditioned inverses inherent to the construction. A pseudo-inverse of the full matrix can be easily obtained and resolves this issue. Solving in a sequetial way, the errors of the pseudo-inverses accumulate. Luckily the problem can be resolved by first performming an SVD decomposion of $A^m = U^m S^m V^{m\dagger}$ matrices, with S diagonal and U and V unitary. All the U^m matrices can be inverted by applying the hermitian transpose to B. The Tensor $S^1 \otimes S^2 \cdots \otimes S^m$ is very sparse and can be inverted at once. The last step consist of inverting all unitary V.

5.1.2 Nonlinear solver

In some cases, the above solver does not return the best possible solution to a given problem. The reason is that it is not able to incorporate symmetries or solve problems where the new blocks appear more than once. A new solver is needed which does not rely on methods from linear algebra, but on more general non-linear leat squares solvers.

In essence, the non-linear least squares solver needs as input a vector with the error values $\vec{f}(\vec{x})$, and if possible also the jocabian, i.e. $J_{I,J} = \frac{\partial f_I}{\partial x_J}$. An

link to right section

5.1. SOLVERS 37

improved point x is chosen by the algoritm, untill some convergence criterium is reached. The implementation uses matlab fsolve routine, which uses Levenberg-Marquardt algoritm under the hood.

5.1.2.0.1 automatic differentiation With some care, the jacobian can be calculated for a general tensor network in an automated way. Suppose we want to differentiate the contracted tensor $T^{i_1\cdots i_n}$ with respect to one of the PEPO blocks $x_n = O^{i_n}_{\alpha\beta\gamma\delta}$. Denote $I = (i_1\cdots i_n)$ and $J = (i_m\alpha\beta\gamma\delta)$, and this block only occurs once. Then $J_{IJ} = \frac{\partial T^{i_1\cdots i_n}}{\partial O^{i_m}_{\alpha\beta\gamma\delta}} = T^{i_1\cdots i_n}_{i_m\alpha\beta\gamma\delta}\delta^{i_n}_{i_m}$ amounts to contracting the network with the tensor x_m removed, and treating the non contracted indices as external ones. If a tensor appears in multiple places, the sum of these contributions has to be taken into account.

source

5.1.2.0.2 Symmetry The non-linear solver can handle rotated and permuted blocks. For instance, a simple loop (square) can be solved by rotating one tensor $T^I_{\alpha\alpha00}$ 4 times, once for every corner. Another example is the following decomposition: $X^I_{\alpha}X^J_{\alpha}=T^{IJ}$.

5.1.3 Sequential linear solver

While from the previous section it seems all non-linear problems need to be solved with the non-linear solver, this is in fact not the case. For instance the following corner block can perfectly be solved with the linear solver.

$$\begin{array}{c|c}
\alpha \\
\gamma \\
\beta \\
\alpha
\end{array}$$
(5.3)

The loop is treated exactly the same as regular leg. The error is of the same magnitude as the non-linear solver, but performs much faster, and is therefore the solver of choice. For problems with multiple tensors, which may be a rotated version of each other such as in a 2x3 rectangle, the linear solver can be useful.

$$\begin{array}{c|c}
\hline
z & y \\
\hline
y & z
\end{array}$$
(5.4)

The idea is to solve each of the tensors with a linear solver. As this is not truly a linear system, the error will not be zero after one pass. But solving the tensors repeatedly lowers the error at each step, giving a iterative procedure. This procedure can be sped up by reusing some parts of the calculations involved in the linear solver. For example, the exponentiated hamiltonian and contraction of all virtual levels that do not involve a given block only need to be performed once.

5.2 Optimasation

5.2.1 Bookkeeping

One important aspect of programming these general solvers is to devise a scheme that keeps track of all the involved tensors and transforms to problem to the form discribed above. In the code, the geometric info is represented by a map. This keeps track of the neighbours for each site, numbers the internal and external legs and a list to perform the contractions.

The framework provides some tools to transform these maps into other maps, for instance by removing 1 site.

5.2.2 Fast contraction

One particular task is to determine all the possible combinations of virtual levels for a given geometry. Simply looping over all possible combinations scales as n^m , with the number of virtual levels and m the number of internal legs. This quickly becomes a bottleneck. This problem can be restated as a PEPS contraction in the following way: for each site make a tensor $T^i_{\alpha\beta\gamma\delta}$ where i encodes all the non-empty combinations of legs $(\alpha\beta\gamma\delta)$. On each site, the right boundary conditions need to be applied to get the right geometry. After setting the boundary conditions, the sparse PEPS network can be contracted and the resulting tensor gives, after decoding, all the possible contractions. Due to its sparsity, this performs quite fast. As an added bonus, removing a tensor from T gives all contractions without this tensor. As both results are sorted list, the subset of contractions containing a given tensor can also be found fast.

5.2.3 Normalisation

For many of the end results, the PEPO cells can be divided by a normalisation factor. Normalising the calculations is important, because $\exp(\hat{H})$ scales exponentially in the number of sites. Luckily, the exponential can be calculated directly in normalised form. Suppose H is the matrisation of the hamiltonian evaluated for a certain geometry. This is a hermitian matrix and can be diagonalised $H = QDQ^{\dagger}$ with Q unitary. Then

$$exp(H_{d^N} - N\alpha I) = Qexp(D - N\log(\alpha)I)Q^{\dagger}$$

$$= Q \begin{bmatrix} exp(D_{11} - N\log(\alpha)) & & & \\ & \ddots & & \\ & & exp(D_{d^Nd^N} - N\log(\alpha)) \end{bmatrix} Q$$

$$= \frac{exp(H_{d^N})}{\alpha^N}$$

$$(5.5)$$

. With I the unit matrix. Next to a global normalisation factor, every block calculation calculates a specific normalisation factor such that the largest eigen-

value of exp(H) is of the order 1.

5.2.4 Internal representation

Two main internal representations are used to construct the given MPO. Either, the MPO is stored as a cell of matrices, or as one big matrix where the blocks are added to during the construction. The output type can be chosen. For some types, sparse matrices are used during the construction. Given that Matlab doesn't support multidimensional matrices by default, this library is used.

expand

5.2.5 Even faster inverses

While the inversion procedure above states how to make use of pseudoinverses, it was not yet clear in the 1D case it was needed. The 1D implementation uses a trick to get all the inverses for free from the SVD decomposition. Take the MPO which corresponds to a unitary matrix:

$$\frac{\delta}{O_n} \stackrel{\beta}{\underset{i}{\nearrow}} \cong U^n_{\alpha(ij\beta)}$$
(5.8)

Then the inverse MPO can be calcuated by taking its Hermitian conjugate and reshaping.

$$\frac{\beta}{O_n^{-1}} \frac{\gamma}{\gamma} \cong U_{(ij\beta)\gamma}^{n\dagger} \tag{5.9}$$

The left inverse looks like

$$2 = \frac{\alpha}{O_n^{-1}} - O_m^{-1} - \cdots - O_1^{-1} = 0$$
 (5.10)

where the physical indices need to be contracted with the corresponding indices of the right hand side.

update

5.2.6 Buffering Results

Some calculations, such as calculating the matrix exponential, take some time. In 1D code, the same calculations were performed over and over again, and hence

a buffer mechanism was written to store these results. In the 2D framework, this not necessary as the solvers only calculate the matrix exponential once and return the blocks together with the made error.

5.3 Calculating phase diagrams

- 5.3.1 Points sampling
- 5.3.2 Storing the information
- 5.3.3 Fitting

5.4 How to use

All the code neede to generate all the results from this dissertation is available on my github page https://github.com/DavidDevoogdt/Thesis_Tensor_Networks. The starting points to explore the code are in the readme file.

5.4.1 Source code structure 1D

The source code for this project can be found on github. The implementation of these types can be found under src/generateMP0.m. In this class the different types of MPO can be constructed. It bundles some helper functions such as contracting a chain or cycle of MPO's or construction of an exponentiated hamiltonian for the given input hamiltonian. Other examples are making L_n^{-1} by sequential invers MPO contractions,...

src/test.m contains the code to create the plots to compare different types
and orders. The other files in the folder are self-explanatory.

5.4.2 Source code structure 2D

source code 2D

5.5 Limitations and outlook

In short, all components are in place in the framework to generate easily and efficiently all the blocks. Since day one, this is a constant work in progress

5.5.1 Size limitation

The main bottleneck is, as expected, calculating the matrix exponential for large systems. The other components are efficient enough to not cause any troubles.

5.5.2 Lattices

The models studied were all on a square lattice. The universal physics does not depend on the details of the model, such as the lattice. It would be beneficial to be able to simulate to these lattices, but here also higher dimensions could be included. In essence all the information of the lattice is contained in the maps generated for each calculation, such as all the connected sites, how to contract them, etcetera. Although undoubtedly many details will need to be changed to use it in practice, the solvers can stay almost the same.

5.5.3 Symmetries

At the moment rotation and permutation symmetries can be included in the construction of the blocks. Internal symmetries are not yet included at the moment. Including them could push the computational boundaries further.

5.5.4 Code quality

The first goal of writing numerical code is of course that it compiles and that the results are as correct. But this is only the first step. In the 2D framework, the neatly orders the different task is functions to avoid as much code duplication as possible. For instance, there are functions generating and manipulating maps, which depends on the geometry of the problem. Other functions only use the bonds and contractions listed in that map object, and hence another lattice can be used without changing the other components too much.

Another example is

Chapter 6

Results

With four parameters I can fit an elephant, and with five I can make him wiggle his trunk

John von Neumann

6.1 Benchmarking

blabla

6.2 Results 1D

In this section the results of the MPO construction detailed in

6.2.1 Exact tensor matrix exponential

The performance of the MPO construction can be compared with the exact diagonalisation of the hamiltonian for a given number of sites. To obtain a faithful results, the number of sites should be as high as possible. In practice, diagonalisation of large matrices becomes slow and memory consuming. The size grows exponentially in the number of sites: $d^n \times d^n$. A double takes 8 bytes of memory. A Rough estimated of the amount of RAM R needed to store this complex array is:

$$R = d^{2n} \times 16bytes \tag{6.1}$$

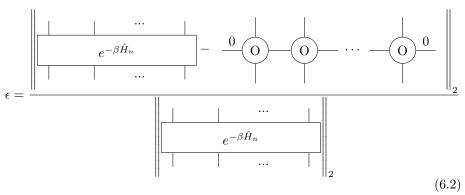
Which means a 14 site chain already takes up more than 4 GB of RAM. The complexity to calculate a matrix exponential scales as $O(n^3)$ [13]. In practice this means that, without any tricks, the matrix exponential can be calculated for 12 sites.

make version for cyclic

trace norm, schatten p norm, ...

6.2.1.1 norms

The schatten 2 norm is used in the following analysis, dentoted by $\|\cdot\|_2$. In the figures the relative error ϵ is reported.



This norm can only be calculated for a finite number of sites. The influence of the number of sites for a linear and cyclic fig. 6.1. As expected, the cyclic norm represents large systems better for the same number of sites. The linear norm keeps increasing with every added site.

Calculating the cyclic norm comes at the extra cost of contracting a cyclic tensor network. In this chapter, the cyclic norm will be given for M=8 sites.

6.2.2 Inversion procedure

6.2.2.1 Full pseudo-inversion

The procedure for inversion of the linear blocks was detailed in section 5.1.1. To demonstrate the need for a pseudoinverse, the error for transversal Ising models is shown in fig. 6.2. The pseudoinversion has 1 parameter σ_0 , the cutoff below which the sinular values are set to zero. A value of 10^{-12} seems optimal in the sense that it doesn't introduce large fluctuations, but still is able to produce good inverses.

6.2.2.1.1 Truncation The original 1D code does not yet have this full inversion procedure. Instead an optimality criterium was devised to truncate the series. Needless to say, even with this criterium, the results were a lot worse at low beta

6.2.3 Models

With these basic definitions and findings out of the way, the different serie expansions can be tested against some different physical models.

6.2. RESULTS 1D 45

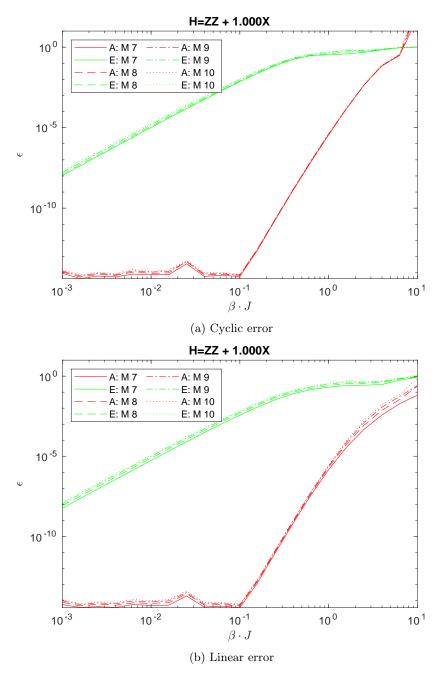


Figure 6.1: Different error measures for 1D transversal Ising model

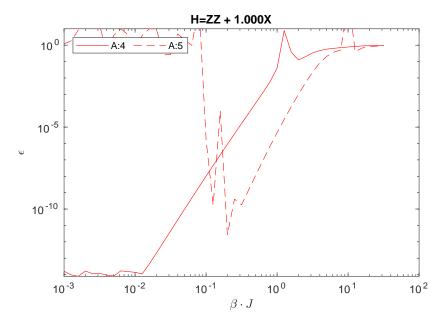


Figure 6.2: Error for transversal Ising, computed with full inversion

6.2.3.1 Ising

The results for the different types are all bundled in fig. 6.3. The virtical axis show the logaritm of the relative error ϵ , horizontal axis the normalised inverse $beta=\frac{J}{T}$. The most surprising finding is the fact that the strict type E performs worse than the others. For low β , it is more than 5 orders of magnitude worse, Taken together with its large bond dimension, this expansion and all other strict variants are seen to be of no use.

Now lets focus on the 2 other variants. Type A clearly outperforms type F for all $\beta < 2$ by quite some margin. Only for large β , type F has the upper hand. While not completely clear in the picture, type A has quite a large error for some orders at $\beta \in (2,10)$, but higher orders seems to solve the problem. The construction of type F requires twice the bond dimension, and hence type A performs clearly better overal

6.2.3.2 Heisenberg

Now we focus on the spin 1/2 Heisenberg model on a chain. The results are again displayd in fig. 6.4. The exact type seems to perform a lot better here, but still has consistently the largest error of the 3 methods for a given order. For low β , type A again performs better than type F. At $\beta \approx 1$ and larger, F starts to improve upon the result of A, and doesn't have an divergent error.

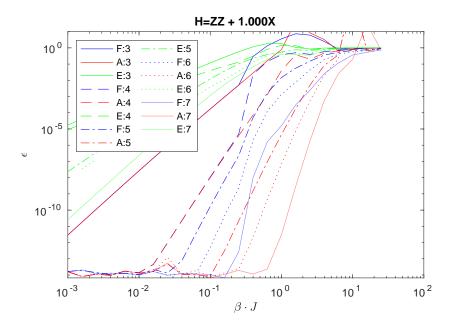


Figure 6.3: Comparison type A, E and F for Transversal Ising.

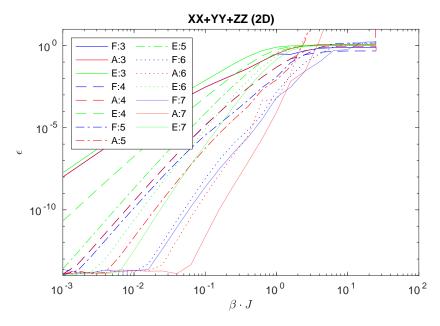


Figure 6.4: Comparison type A, E and F for Heisenberg model.

6.2.3.3 Random

To give a representative overview for random hamiltonians, several simulations were run. The single site and nearest neighbourgh hamiltonians are generated by making hermitian matrices with random real and complex numbers between -1 and 1. In order to compare the different graphs, the engergy scale is set such that the norm of the hamiltonian evaluated on 2 sites is 1.

run random hamiltonian

6.2.4 Conclusion

write this

- strict variant are useless
- full smmetry is not needed
- pseudoinverse neccesary
- Type A is the best

6.3 Results 2D

- 6.3.1 Direct results
- 6.3.1.1 norm
- 6.3.1.2 Ising
- 6.3.1.3 Heisenberg
- 6.3.2 2D transveral Ising model
- 6.3.2.1 Data generation
- 6.3.2.2 classical Ising phase transition
- 6.3.2.3 g=2.5 phase transition
- 6.3.2.4 Towards quantum phase transition

Chapter 7

Conclusion and lookout

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