

RISC-V External Debug Support  
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# Preface

**Warning! This draft specification will change before being accepted as standard, so implementations made to this draft specification will likely not conform to the future standard.**

## Acknowledgments

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# Chapter 1

## Introduction

When a design progresses from simulation to hardware implementation, a user’s control and understanding of the system’s current state drops dramatically. To help bring up and debug low level software and hardware, it is critical to have good debugging support built into the hardware. When a robust OS is running on a core, software can handle many debugging tasks. However, in many scenarios, hardware support is essential.

This document outlines a standard architecture for external debug support on RISC-V platforms. This architecture allows a variety of implementations and tradeoffs, which is complementary to the wide range of RISC-V implementations. At the same time, this specification defines common interfaces to allow debugging tools and components to target a variety of platforms based on the RISC-V ISA.

System designers may choose to add additional hardware debug support, but this specification defines a standard interface for common functionality.

### 1.1 Terminology

A *platform* is a single integrated circuit consisting of one or more *components*. Some components may be RISC-V cores, while others may have a different function. Typically they will all be connected to a single system bus. A single RISC-V core contains one or more hardware threads, called *harts*.

#### 1.1.1 Context

This document is written to work with:

1. The RISC-V Instruction Set Manual, Volume I: User-Level ISA, Document Version 2.2
2. The RISC-V Instruction Set Manual, Volume II: Privileged Architecture, Version 1.10

## 1.2 About This Document

### 1.2.1 Structure

This document contains two parts. The main part of the document is the specification, which is given in the numbered sections. The second part of the document is a set of appendices. The information in the appendix is intended to clarify and provide examples, but is not part of the actual specification.

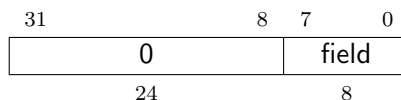
### 1.2.2 Register Definition Format

All register definitions in this document follow the format shown below. A simple graphic shows which fields are in the register. The upper and lower bit indices are shown to the top left and top right of each field. The total number of bits in the field are shown below it.

After the graphic follows a table which for each field lists its name, description, allowed accesses, and reset value. The allowed accesses are listed in Table 1.2.

Names of registers and their fields are hyperlinks to their definition, and are indexed on page 74.

#### 1.2.2.1 Long Name (shortname, at 0x123)



Field	Description	Access	Reset
field	Description of what this field is used for.	R/W	15

Table 1.2: Register Access Abbreviations

R	Read-only.
R/W	Read/Write.
R/W0	Read/Write. Only writing 0 has an effect.
R/W1	Read/Write. Only writing 1 has an effect.
R/W1C	Read/Write. For each bit in the field, writing 1 clears that bit. Writing 0 has no effect.
W	Write-only. When read this field returns 0.
W1	Write-only. Only writing 1 has an effect.



## 1.3 Background

There are several use cases for dedicated debugging hardware, both internal to a CPU core and with an external connection. This specification addresses the use cases listed below. Implementations can choose not to implement every feature, which means some use cases might not be supported.

- Debugging low-level software in the absence of an OS or other software.
- Debugging issues in the OS itself.
- Bootstrapping a system to test, configure, and program components before there is any executable code path in the system.
- Accessing hardware on a system without a working CPU.

In addition, even without a hardware debugging interface, architectural support in a RISC-V CPU can aid software debugging and performance analysis by allowing hardware triggers and breakpoints. This specification aims to define common resources which can be used for different cases.

When debugging software, this specification distinguishes between two forms of external debugging. The first is *halt mode* debugging, where an external debugger halts some or all components of a platform and inspects their state while they are in stasis. The debugger can read and/or modify state, then direct the hardware to execute a single instruction, or continue to run freely.

The second is *run mode* debugging. In this mode a software debug agent runs on a component (eg. triggered by a timer interrupt or breakpoint on a RISC-V core) which transfers data to or from the debugger without halting the component, only briefly interrupting its program flow. This functionality is essential if the component is controlling some real-time system (like a hard drive) where long timing delays could lead to physical damage. This requires additional software support (both on the system as well as on the debugger), and efficient communication channels between the component and the debugger.

## 1.4 Supported Features

The debug interface described in this specification supports the following features:

1. All hart registers (including CSRs) can be read/written.
2. Memory can be accessed either from the hart's point of view, through the system bus directly, or both.
3. RV32, RV64, and future RV128 are all supported.
4. Any hart in the platform can be independently debugged.

5. A debugger can discover almost<sup>1</sup> everything it needs to know itself, without user configuration.
6. Each hart can be debugged from the very first instruction executed.
7. A RISC-V hart can be halted when a software breakpoint instruction is executed.
8. Hardware single-step can execute one instruction at a time.
9. Debug functionality is independent of the debug transport used.
10. The debugger does not need to know anything about the microarchitecture of the harts it is debugging.
11. Arbitrary subsets of harts can be halted and resumed simultaneously. (Optional)
12. Arbitrary instructions can be executed on a halted hart. That means no new debug functionality is needed when a core has additional or custom instructions or state, as long as there exist programs that can move that state into GPRs. (Optional)
13. Registers can be accessed without halting. (Optional)
14. A running hart can be directed to execute a short sequence of instructions, with little overhead. (Optional)
15. A system bus master allows memory access without involving any hart. (Optional)
16. A RISC-V hart can be halted when a trigger matches the PC, read/write address/data, or an instruction opcode. (Optional)

This document does not suggest a strategy or implementation for hardware test, debugging or error detection techniques. Scan, BIST, etc. are out of scope of this specification, but this specification does not intend to limit their use in RISC-V systems.

It is possible to debug code that uses software threads, but there is no special debug support for it.

---

<sup>1</sup>Notable exceptions include information about the memory map and peripherals.

## Chapter 2

# System Overview

Figure 2.1 shows the main components of External Debug Support. Blocks shown in dotted lines are optional.

The user interacts with the Debug Host (eg. laptop), which is running a debugger (eg. gdb). The debugger communicates with a Debug Translator (eg. OpenOCD, which may include a hardware driver) to communicate with Debug Transport Hardware (eg. Olimex USB-JTAG adapter). The Debug Transport Hardware connects the Debug Host to the Platform's Debug Transport Module (DTM). The DTM provides access to one or more Debug Modules (DMs) using the Debug Module Interface (DMI).

Each hart in the platform is controlled by exactly one DM. Harts may be heterogeneous. There is no further limit on the hart-DM mapping, but usually all harts in a single core are controlled by the same DM. In most platforms there will only be one DM that controls all the harts in the platform.

DMs provide run control to their harts in the platform. Abstract commands provide access to GPRs. Additional registers are accessible through abstract commands or by writing programs to the optional Program Buffer.

The Program Buffer allows the debugger to execute arbitrary instructions on a hart. This mechanism can be used to access memory. An optional system bus access block allows memory accesses without using a RISC-V hart to perform the access.

Each RISC-V hart may implement a Trigger Module. When trigger conditions are met, harts will halt and inform the debug module that they have halted.



Figure 2.1: RISC-V Debug System Overview

## Chapter 3

# Debug Module (DM)

The Debug Module implements a translation interface between abstract debug operations and their specific implementation. It might support the following operations:

1. Give the debugger necessary information about the implementation. (Required)
2. Allow any individual hart to be halted and resumed. (Required)
3. Provide status on which harts are halted. (Required)
4. Provide abstract read and write access to a halted hart's GPRs. (Required)
5. Provide access to a reset signal that allows debugging from the very first instruction after reset. (Required)
6. Provide a mechanism to allow debugging harts immediately out of reset (regardless of the reset cause). (Optional)
7. Provide abstract access to other hart registers. (Optional)
8. Provide a Program Buffer to force the hart to execute arbitrary instructions. (Optional)
9. Allow multiple harts to be halted, resumed, and/or reset at the same time. (Optional)
10. Allow memory access from a hart's point of view. (Optional)
11. Allow direct System Bus Access. (Optional)

In order to implement memory access, a target must implement at least one of Program Buffer, System Bus Access, or Abstract Access Memory command mechanisms.

Targets that wish to implement access to all registers must implement either the Program Buffer or abstract access to all registers.

A single DM can debug up to  $2^{20}$  harts.

### 3.1 Debug Module Interface (DMI)

Debug Modules are slaves to a bus called the Debug Module Interface (DMI). The master of the bus is the Debug Transport Module(s). The Debug Module Interface can be a trivial bus with one master and one slave, or use a more full-featured bus like TileLink or the AMBA Advanced Peripheral Bus. The details are left to the system designer.

The DMI uses between 7 and 32 address bits. It supports read and write operations. The bottom of the address space is used for the first (and usually only) DM. Extra space can be used for custom debug devices, other cores, additional DMs, etc. If there are additional DMs on this DMI, the base address of the next DM in the DMI address space is given in `nextdm`.

The Debug Module is controlled via register accesses to its DMI address space.

## 3.2 Reset Control

The Debug Module controls a global reset signal, `ndmreset` (non-debug module reset), which can reset, or hold in reset, every component in the platform, except for the Debug Module and Debug Transport Modules. Exactly what is affected by this reset is implementation dependent, as long as it is possible to debug programs from the first instruction executed. The Debug Module's own state and registers should only be reset at power-up and while `dmactive` in `dmcontrol` is 0. The halt state of harts should be maintained across system reset provided that `dmactive` is 1, although trigger CSRs may be cleared.

Due to clock and power domain crossing issues, it may not be possible to perform arbitrary DMI accesses across system reset. While `ndmreset` or any external reset is asserted, the only supported DM operation is accessing `dmcontrol`. The behavior of other accesses is undefined.

There is no requirement on the duration of the assertion of `ndmreset`. The implementation must ensure that a write of `ndmreset` to 1 followed by a write of `ndmreset` to 0 triggers system reset. The system may take an arbitrarily long time to come out of reset, as reported by `allunavail`, `anyunavail`, or other implementation specific indicators.

When harts have been reset, they must set a sticky `havereset` state bit. The conceptual `havereset` state bits can be read for selected harts in `anyhavereset` and `allhavereset` in `dmstatus`. These bits must be set regardless of the cause of the reset. The `havereset` bits for the selected harts can be cleared by writing 1 to `ackhavereset` in `dmcontrol`. The `havereset` bits may or may not be cleared when `dmactive` is low.

## 3.3 Selecting Harts

Up to  $2^{20}$  harts can be connected to a single DM. The debugger selects a hart, and then subsequent halt, resume, reset, and debugging commands are specific to that hart.

To enumerate all the harts, a debugger must first determine `HARTSELLEN` by writing all ones to `hartsel` (assuming the maximum size) and reading back the value to see which bits were actually set. Then it selects each hart starting from 0 until either `anynonexistent` in `dmstatus` is 1, or the highest index (depending on `HARTSELLEN`) is reached.

The debugger can discover the mapping between hart indices and `mhartid` by using the interface to read `mhartid`, or by reading the system's Device Tree.

### 3.3.1 Selecting a Single Hart

All debug modules must support selecting a single hart. The debugger can select a hart by writing its index to [hartsel](#). Hart indexes start at 0 and are contiguous until the final index.

### 3.3.2 Selecting Multiple Harts

Debug Modules may implement a Hart Array Mask register to allow selecting multiple harts at once. The Nth bit in the Hart Array Mask register applies to the hart with index N. If the bit is 1 then the hart is selected. Usually a DM will have a Hart Array Mask register exactly wide enough to select all the harts it supports, but it's allowed to tie any of these bits to 0.

The debugger can set bits in the hart array mask register using [hawindowssel](#) and [hawindow](#), then apply actions to all selected harts by setting [hasel](#). If this feature is supported, multiple harts can be halted, resumed, and reset simultaneously.

Only the actions initiated by [dmcontrol](#) can apply to multiple harts at once, Abstract Commands apply only to the hart selected by [hartsel](#).

## 3.4 Run Control

For every hart, the Debug Module contains 4 conceptual bits of state: halt request, resume request, halt-on-reset request, and hart reset. (The hart reset and halt-on-reset request bits are optional.) These bits all reset to 0. A debugger can write them for the currently selected harts through [haltreq](#), [resumereq](#), [setresethaltreq](#)/[clrresethaltreq](#) and [hartreset](#) in [dmcontrol](#). In addition the DM receives halted, running, and resume ack signals from each hart.

When a running hart receives a halt request, it responds by halting and asserting its halted signal. The halted signals of all selected harts are reflected in the [allhalted](#) and [anyhalted](#) bits. [haltreq](#) is ignored by halted harts.

When a halted hart receives a resume request, it responds by resuming, clearing its halted signal, and asserting its running signal and resume ack signals. The resume ack signal is lowered when the resume request is deasserted. These status signals of all selected harts are reflected in [allresumeack](#), [anyresumeack](#), [allrunning](#), and [anyrunning](#). [resumereq](#) is ignored by running harts.

When halt or resume is requested, a hart must respond in less than one second, unless it is unavailable. (How this is implemented is not further specified. A few clock cycles will be a more typical latency).

The DM can implement optional halt-on-reset bits for each hart, which it indicates by setting [hasresethaltreq](#) to 1. This means the DM implements the [setresethaltreq](#) and [clrresethaltreq](#) bits. Writing 1 to [setresethaltreq](#) sets the halt-on-reset request bit for each selected hart. When a hart's halt-on-reset request bit is set, the hart will immediately enter debug mode on the next deassertion of its reset. This is true regardless of whether or not the hart's reset is caused by the DM's [ndmreset](#)/[hartreset](#) signals. The hart's halt-on-reset request bit remains set until cleared by the

debugger writing 1 to `clrresethaltreq` while the hart is selected, or by debug module reset.

## 3.5 Abstract Commands

The DM supports a set of abstract commands, most of which are optional. Depending on the implementation, the debugger may be able to perform some abstract commands even when the selected hart is not halted. Debuggers can only determine which abstract commands are supported by a given hart in a given state by attempting them and then looking at `cmderr` in `abstractcs` to see if they were successful. Commands may be supported with some options set, but not with other options set. If a command has unsupported options set, the DM must set `cmderr` to 2 (not supported).

---

*Example: Every system must support the Access Register command, but may not support accessing CSRs. If the debugger requests to read a CSR in that case, the command will return "not supported."*

---

Debuggers execute abstract commands by writing them to `command`. Debuggers can determine whether an abstract command is complete by reading `busy` in `abstractcs`. If the command takes arguments, the debugger must write them to the `data` registers before writing to `command`. If a command returns results, the Debug Module must ensure they are placed in the `data` registers before `busy` is cleared. Which `data` registers are used for the arguments is described in Table 3.1. In all cases the least-significant word is placed in the lowest-numbered `data` register. The argument width depends on the command being executed, and is MXLEN where not explicitly specified.

---

*The Abstract Command interface is designed to allow a debugger to write commands as fast as possible, and then later check whether they completed without error. In the common case the debugger will be much slower than the target and commands succeed, which allows for maximum throughput. If there is a failure, the interface ensures that no commands execute after the failing one. To discover which command failed, the debugger has to look at the state of the DM (eg. contents of `command`) or hart (eg. contents of a register modified by a Program Buffer program) to determine which one failed.*

---

Table 3.1: Use of Data Registers

Argument Width	arg0/return value	arg1	arg2
32	<code>data0</code>	<code>data1</code>	<code>data2</code>
64	<code>data0</code> , <code>data1</code>	<code>data2</code> , <code>data3</code>	<code>data4</code> , <code>data5</code>
128	<code>data0</code> – <code>data3</code>	<code>data4</code> – <code>data7</code>	<code>data8</code> – <code>data11</code>

### 3.5.1 Abstract Command Listing

This section describes each of the different abstract commands and how their fields should be interpreted when they are written to `command`.

Each abstract command is a 32-bit value. The top 8 bits contain `cmdtype` which determines the kind of command. Table 3.2 lists all commands.



Table 3.2: Meaning of `cmdtype`

<code>cmdtype</code>	Command	Page
0	Access Register Command	11
1	Quick Access	12
2	Access Memory Command	13

### 3.5.1.1 Access Register

This command gives the debugger access to CPU registers and program buffer. It performs the following sequence of operations:

1. Copy data from the register specified by `regno` into the `arg0` region of `data`, if `write` is clear and `transfer` is set.
2. Copy data from the `arg0` region of `data` into the register specified by `regno`, if `write` is set and `transfer` is set.
3. Execute the Program Buffer, if `postexec` is set.

If any of these operations fail, `cmderr` is set and none of the remaining steps are executed. An implementation may detect an upcoming failure early, and fail the overall command before it reaches the step that would cause failure.

Debug Modules must implement this command and must support read and write access to all GPRs when the selected hart is halted. Debug Modules may optionally support accessing other registers, or accessing registers when the hart is running. If this command is supported for a register while the hart is running, it must also be supported for a register while the hart is halted. Each individual register (aside from GPRs) may be supported differently across read, write, and halt status.

---

*The encoding of `aarsize` was chosen to match `sbaccess` in `sbc`.*

31	24	23	22	20	19	18	17	16	15	0
<code>cmdtype</code>	0	<code>aarsize</code>	0	<code>postexec</code>	<code>transfer</code>	<code>write</code>	<code>regno</code>			
8	1	3	1	1	1	1	16			

Field	Description
<code>cmdtype</code>	This is 0 to indicate Access Register Command.
<code>aarsize</code>	<p>2: Access the lowest 32 bits of the register.  3: Access the lowest 64 bits of the register.  4: Access the lowest 128 bits of the register.</p> <p>If <code>aarsize</code> specifies a size larger than the register's actual size, then the access must fail. If a register is accessible, then reads of <code>aarsize</code> less than or equal to the register's actual size must be supported.</p> <p>This field controls the Argument Width as referenced in Table 3.1.</p>

*Continued on next page*

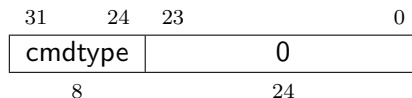
Field	Description
postexec	When 1, execute the program in the Program Buffer exactly once after performing the transfer, if any.
transfer	0: Don't do the operation specified by <code>write</code> . 1: Do the operation specified by <code>write</code> . This bit can be used to just execute the Program Buffer without having to worry about placing valid values into <code>aarsize</code> or <code>regno</code> .
write	When <code>transfer</code> is set: 0: Copy data from the specified register into <code>arg0</code> portion of <code>data</code> . 1: Copy data from <code>arg0</code> portion of <code>data</code> into the specified register.
regno	Number of the register to access, as described in Table 3.6. <code>dpc</code> may be used as an alias for PC if this command is supported on a non-halted hart.

### 3.5.1.2 Quick Access

Perform the following sequence of operations:

1. If the hart is halted, the command sets `cmderr` to `halt/resume` and does not continue.
2. Halt the hart. If the hart halts for some other reason (e.g. breakpoint), the command sets `cmderr` to `halt/resume` and does not continue.
3. Execute the Program Buffer. If an exception occurs, `cmderr` is set to `exception` and the program buffer execution ends, but the quick access command continues.
4. Resume the hart.

Implementing this command is optional.



Field	Description
cmdtype	This is 1 to indicate Quick Access command.

### 3.5.1.3 Access Memory

This command lets the debugger perform memory accesses, with the exact same memory view and permissions as the selected hart has. This includes access to hart-local memory-mapped registers, etc. The command performs the following sequence of operations:

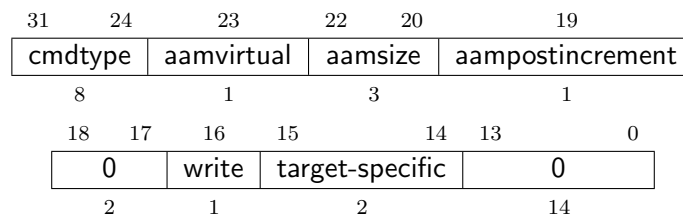
1. Copy data from the memory location specified in **arg1** into the **arg0** portion of **data**, if **write** is clear.
2. Copy data from the **arg0** portion of **data** into the memory location specified in **arg1**, if **write** is set.
3. Increment **arg1**, if **aampostincrement** is set.

If any of these operations fail, **cmderr** is set and none of the remaining steps are executed. An access may only fail if the hart, running M mode code, might encounter that same failure when it attempts the same access. An implementation may detect an upcoming failure early, and fail the overall command before it reaches the step that would cause failure.

Debug Modules may optionally implement this command and may support read and write access to memory locations when the selected hart is running or halted. If this command supports memory accesses while the hart is running, it must also support memory accesses while the hart is halted.

---

The encoding of **aamsize** was chosen to match **sbaccess** in **sbc**s.



Field	Description
cmdtype	This is 2 to indicate Access Memory Command.
aamvirtual	An implementation does not have to implement both virtual and physical accesses, but it must fail accesses that it doesn't support. 0: Addresses are physical (to the hart they are performed on). 1: Addresses are virtual, and translated the way the would be from M mode, with <b>mprv</b> set.

*Continued on next page*

Field	Description
<b>aamsize</b>	0: Access the lowest 8 bits of the memory location. 1: Access the lowest 16 bits of the memory location. 2: Access the lowest 32 bits of the memory location. 3: Access the lowest 64 bits of the memory location. 4: Access the lowest 128 bits of the memory location. This field controls the Argument Width as referenced in Table 3.1.
<b>aampostincrement</b>	After a memory access has completed, if this bit is 1, increment <b>arg1</b> (which contains the address used) by the number of bytes encoded in <b>aamsize</b> .
<b>write</b>	0: Copy data from the memory location specified in <b>arg1</b> into <b>arg0</b> portion of <b>data</b> . 1: Copy data from <b>arg0</b> portion of <b>data</b> into the memory location specified in <b>arg1</b> .
<b>target-specific</b>	These bits are reserved for target-specific uses.

Table 3.6: Abstract Register Numbers

0x0000 – 0x0fff	CSRs. The “PC” can be accessed here through <b>dpc</b> .
0x1000 – 0x101f	GPRs
0x1020 – 0x103f	Floating point registers
0xc000 – 0xffff	Reserved for non-standard extensions and internal use.

## 3.6 Program Buffer

To support executing arbitrary instructions on a halted hart, a Debug Module can include a Program Buffer that a debugger can write small programs to. Systems that support all necessary functionality using abstract commands only may choose to omit the Program Buffer.

A debugger can write a small program to the Program Buffer, and then execute it exactly once with the Access Register Abstract Command, setting the **postexec** bit in **command**. The debugger can write whatever program it likes (including jumps out of the Program Buffer), but the program must end with **ebreak** or **c.ebreak**. To save hardware, an implementation may support an implied **ebreak** that is executed when a hart runs off the end of the Program Buffer. This is indicated in **impebreak**. With this feature, a Program Buffer of just 2 32-bit words can offer efficient debugging.

If **progbufsize** is 1, **impebreak** must be 1. It is possible that the Program Buffer can hold only one 32- or 16-bit instruction, so the debugger must only write a single instruction in this case, regardless

of its size. This instruction can be a 32-bit instruction, or a compressed instruction in the lower 16 bits accompanied by a compressed **nop** in the upper 16 bits.

---

*The slightly inconsistent behavior with a Program Buffer of size 1 is to accommodate hardware designs that prefer to stuff instructions directly into the pipeline when halted, instead of having the Program Buffer exist in the address space somewhere.*

---

If the debugger executes a program that does not terminate with an **ebreak** instruction, the hart will remain in Debug Mode until it is reset.

While these programs are executed, the hart does not leave Debug Mode (see Section 4.1). If an exception is encountered during execution of the Program Buffer, no more instructions are executed, the hart remains in Debug Mode, and **cmderr** is set to 3 (**exception error**). If the debugger executes a program that doesn't terminate, then it loses control of the hart.

Executing the Program Buffer may clobber **dpc**. If that is the case, it must be possible to read/write **dpc** using an abstract command with **postexec** not set. The debugger must attempt to save **dpc** between halting and executing a Program Buffer, and then restore **dpc** before leaving Debug Mode.

---

*Allowing Program Buffer execution to clobber **dpc** allows for direct implementations that don't have a separate PC register, and do need to use the PC when executing the Program Buffer.*

---

The Program Buffer may be implemented as RAM which is accessible to the hart. A debugger can determine if this is the case by executing small programs that attempt to write and read back relative to **pc** while executing from the Program Buffer. If so, the debugger has more flexibility in what it can do with the program buffer.

### 3.7 Overview of States

Figure 3.1 shows a conceptual view of the states passed through by a hart during run/halt debugging as influenced by the different fields of **dmcontrol**, **abstractcs**, **abstractauto**, and **command**.

### 3.8 System Bus Access

A debugger can access memory from a hart's point of view using a Program Buffer or the Abstract Access Memory command. (Both these features are optional.) A Debug Module may also include a System Bus Access block to provide memory access without involving a hart, regardless of whether Program Buffer is implemented. The System Bus Access block uses physical addresses.

The System Bus Access block may support 8-, 16-, 32-, 64-, and 128-bit accesses. Table 3.7 shows which bits in **sbddata** are used for each access size.

Depending on the microarchitecture, data accessed through System Bus Access may not always be coherent with that observed by each hart. It is up to the debugger to enforce coherency if the implementation does not. This specification does not define a standard way to do this, as it is implementation/platform specific. Possibilities may include writing to special memory-mapped locations, or executing special instructions via the Program Buffer.

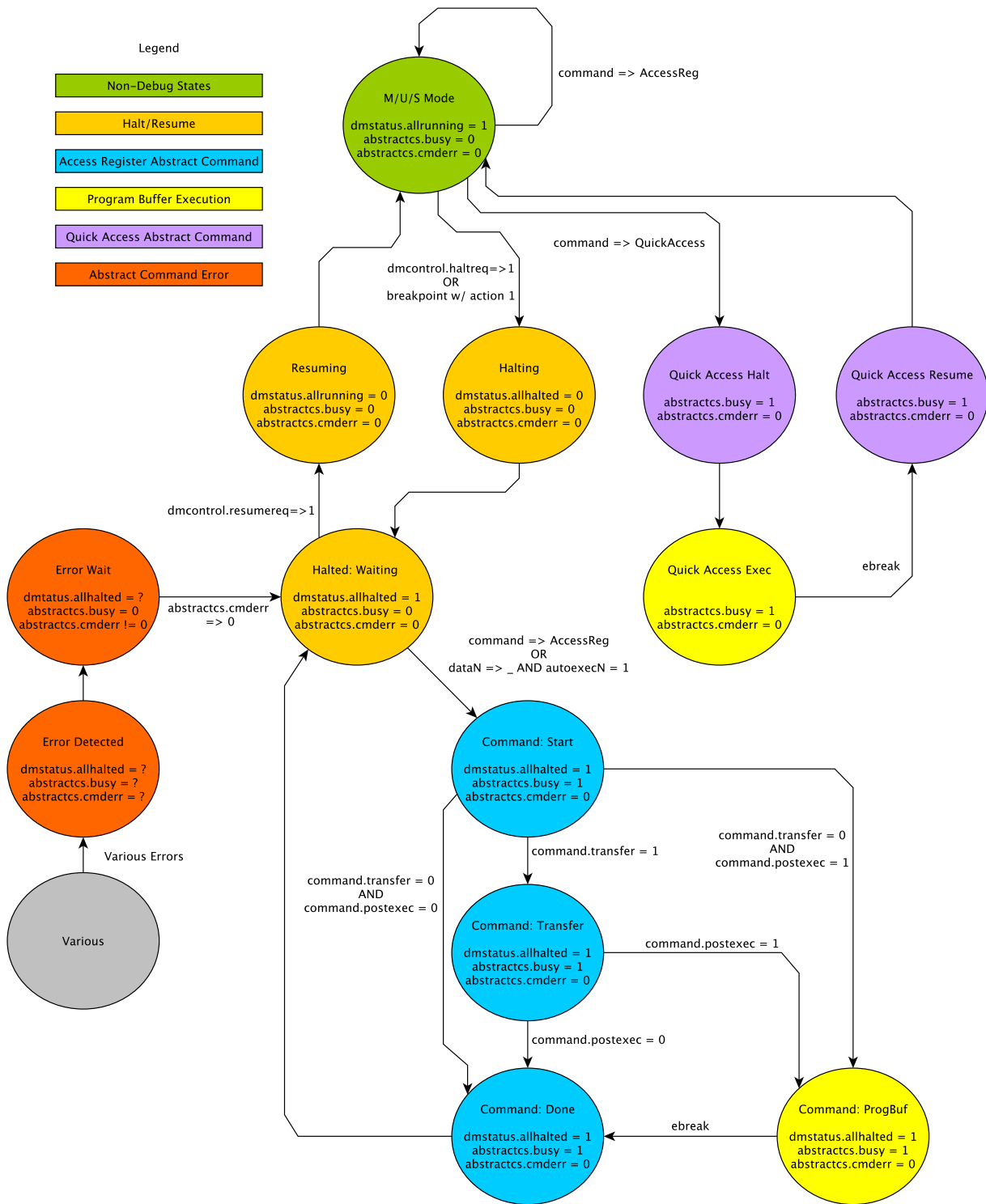


Figure 3.1: Run/Halt Debug State Machine. As only a small amount of state is visible to the debugger, the states and transitions are conceptual.

Table 3.7: System Bus Data Bits

Access Size	Data Bits
8	<a href="#">sbddata0</a> bits 7:0
16	<a href="#">sbddata0</a> bits 15:0
32	<a href="#">sbddata0</a>
64	<a href="#">sbddata1</a> , <a href="#">sbddata0</a>
128	<a href="#">sbddata3</a> , <a href="#">sbddata2</a> , <a href="#">sbddata1</a> , <a href="#">sbddata0</a>

---

*Implementing a System Bus Access block has several benefits even when a Debug Module also implements a Program Buffer. First, it is possible to access memory in a running system with minimal impact. Second, it may improve performance when accessing memory. Third, it may provide access to devices that a hart does not have access to.*

### 3.9 Quick Access

Depending on the task it is performing, some harts can only be halted very briefly. There are several mechanisms that allow accessing resources in such a running system with a minimal impact on the running hart.

First, an implementation may allow some abstract commands to execute without halting the hart.

Second, the Quick Access abstract command can be used to halt a hart, quickly execute the contents of the Program Buffer, and let the hart run again. Combined with instructions that allow Program Buffer code to access the `data` registers, as described in [3.11.3](#), this can be used to quickly perform a memory or register access. For some systems this will be too intrusive, but many systems that can't be halted can bear an occasional hiccup of a hundred or less cycles.

Third, if the System Bus Access block is implemented, it can be used while a hart is running to access system memory.

### 3.10 Security

To protect intellectual property it may be desirable to lock access to the Debug Module. To allow access during a manufacturing process and not afterwards, a reasonable solution could be to add a fuse bit to the Debug Module that can be used to be permanently disable it. Since this is technology specific, it is not further addressed in this spec.

Another option is to allow the DM to be unlocked only by users who have an access key. Between [authenticated](#), [authbusy](#), and [authdata](#) arbitrarily complex authentication mechanism can be supported. When [authenticated](#) is clear, the DM must not interact with the rest of the platform, nor expose details about the harts connected to the DM. All DM registers should read 0, while writes should be ignored, with the following mandatory exceptions:

1. [authenticated](#) in [dmstatus](#) is readable.

2. `authbusy` in `dmstatus` is readable.
3. `version` in `dmstatus` is readable.
4. `dmactive` in `dmcontrol` is readable and writable.
5. `authdata` is readable and writable.

### 3.11 Debug Module Registers

The registers described in this section are accessed over the DMI bus. Each DM has a base address (which is 0 for the first DM). The register addresses below are offsets from this base address.

When read, unimplemented Debug Module DMI Registers return 0. Writing them has no effect.

For each register it is possible to determine that it is implemented by reading it and getting a non-zero value (eg. `sbc`), or by checking bits in another register (eg. `progbu`).



Table 3.8: Debug Module Debug Bus Registers

Address	Name	Page
0x04	Abstract Data 0	<a href="#">28</a>
0x0f	Abstract Data 11	
0x10	Debug Module Control	<a href="#">21</a>
0x11	Debug Module Status	<a href="#">19</a>
0x12	Hart Info	<a href="#">24</a>
0x13	Halt Summary 1	<a href="#">30</a>
0x14	Hart Array Window Select	<a href="#">25</a>
0x15	Hart Array Window	<a href="#">25</a>
0x16	Abstract Control and Status	<a href="#">25</a>
0x17	Abstract Command	<a href="#">26</a>
0x18	Abstract Command Autoexec	<a href="#">27</a>
0x19	Device Tree Addr 0	<a href="#">28</a>
0x1a	Device Tree Addr 1	
0x1b	Device Tree Addr 2	
0x1c	Device Tree Addr 3	
0x1d	Next Debug Module	<a href="#">28</a>
0x20	Program Buffer 0	<a href="#">29</a>
0x2f	Program Buffer 15	
0x30	Authentication Data	<a href="#">29</a>
0x34	Halt Summary 2	<a href="#">30</a>
0x35	Halt Summary 3	<a href="#">30</a>
0x37	System Bus Address 127:96	<a href="#">31</a>
0x38	System Bus Access Control and Status	<a href="#">31</a>
0x39	System Bus Address 31:0	<a href="#">33</a>
0x3a	System Bus Address 63:32	<a href="#">33</a>
0x3b	System Bus Address 95:64	<a href="#">34</a>
0x3c	System Bus Data 31:0	<a href="#">34</a>
0x3d	System Bus Data 63:32	<a href="#">35</a>
0x3e	System Bus Data 95:64	<a href="#">35</a>
0x3f	System Bus Data 127:96	<a href="#">36</a>
0x40	Halt Summary 0	<a href="#">29</a>

### 3.11.1 Debug Module Status (dmstatus, at 0x11)

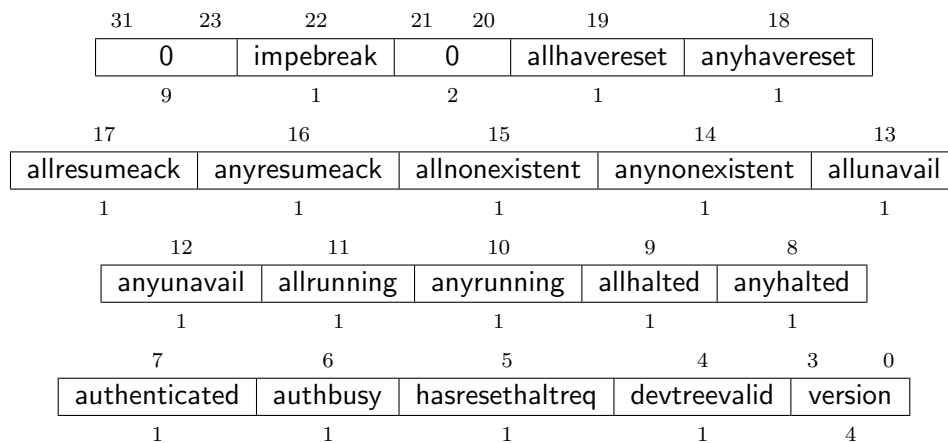
The address of this register will not change in the future, because it contains [version](#). It has changed from version 0.11 of this spec.

This register reports status for the overall Debug Module as well as the currently selected harts, as defined in [hasel](#).

Harts are nonexistent if they will never be part of this system, no matter how long a user waits. Eg. in a simple single-hart system only one hart exists, and all others are nonexistent. Debuggers may assume that a system has no harts with indexes higher than the first nonexistent one.

Harts are unavailable if they might exist/become available at a later time, or if there are other harts with higher indexes than this one. Eg. in a multi-hart system some might temporarily be powered down, or a system might support hot-swapping harts. Systems with very large number of harts may permanently disable some during manufacturing, leaving holes in the otherwise continuous hart index space. In order to let the debugger discover all harts, they must show up as unavailable even if there is no chance of them ever becoming available.

This entire register is read-only.



Field	Description	Access	Reset
impebreak	If 1, then there is an implicit <b>ebreak</b> instruction at the non-existent word immediately after the Program Buffer. This saves the debugger from having to write the <b>ebreak</b> itself, and allows the Program Buffer to be one word smaller. This must be 1 when <a href="#">progbuFSIZE</a> is 1.	R	Preset
allhavereset	This field is 1 when all currently selected harts have been reset but the reset has not been acknowledged.	R	-
anyhavereset	This field is 1 when any currently selected hart has been reset but the reset has not been acknowledged.	R	-
allresumeack	This field is 1 when all currently selected harts have acknowledged the previous resume request.	R	-
anyresumeack	This field is 1 when any currently selected hart has acknowledged the previous resume request.	R	-
allnonexistent	This field is 1 when all currently selected harts do not exist in this system.	R	-
anynonexistent	This field is 1 when any currently selected hart does not exist in this system.	R	-
allunavail	This field is 1 when all currently selected harts are unavailable.	R	-

*Continued on next page*

Field	Description	Access	Reset
anyunavail	This field is 1 when any currently selected hart is unavailable.	R	-
allrunning	This field is 1 when all currently selected harts are running.	R	-
anyrunning	This field is 1 when any currently selected hart is running.	R	-
allhalted	This field is 1 when all currently selected harts are halted.	R	-
anyhalted	This field is 1 when any currently selected hart is halted.	R	-
authenticated	0 when authentication is required before using the DM. 1 when the authentication check has passed. On components that don't implement authentication, this bit must be preset as 1.	R	Preset
authbusy	0: The authentication module is ready to process the next read/write to <a href="#">authdata</a> . 1: The authentication module is busy. Accessing <a href="#">authdata</a> results in unspecified behavior. <a href="#">authbusy</a> only becomes set in immediate response to an access to <a href="#">authdata</a> .	R	0
hasresethaltreq	1 if this Debug Module supports halt-on-reset functionality controllable by the <a href="#">setresethaltreq</a> and <a href="#">clrresethaltreq</a> bits. 0 otherwise.	R	Preset
devtreevalid	0: <a href="#">devtreeaddr0</a> – <a href="#">devtreeaddr3</a> hold information which is not relevant to the Device Tree. 1: <a href="#">devtreeaddr0</a> – <a href="#">devtreeaddr3</a> registers hold the address of the Device Tree.	R	Preset
version	0: There is no Debug Module present. 1: There is a Debug Module and it conforms to version 0.11 of this specification. 2: There is a Debug Module and it conforms to version 0.13 of this specification. 15: There is a Debug Module but it does not conform to any available version of this spec.	R	2

### 3.11.2 Debug Module Control ([dmcontrol](#), at 0x10)

This register controls the overall Debug Module as well as the currently selected harts, as defined in [hasel](#).

Throughout this document we refer to [hartsel](#), which is [hartselhi](#) combined with [hartsello](#). While the spec allows for 20 [hartsel](#) bits, an implementation may choose to implement fewer than that. The actual width of [hartsel](#) is called [HARTSELLEN](#). It must be at least 0 and at most 20. A debugger

should discover HARTSELLEN by writing all ones to [hartsel](#) (assuming the maximum size) and reading back the value to see which bits were actually set.

31	30	29	28	27	26	25	16
haltreq	resumereq	hartreset	ackhavereset	0	hasel	hartsello	
1	1	1	1	1	1	10	
15	6	5	4	3	2	1	0
hartselhi	0	setresethaltreq	clrresethaltreq	ndmreset	dmactive		
10	2	1	1	1	1	1	

Field	Description	Access	Reset
haltreq	Writes the halt request bit for all currently selected harts. When set to 1, each selected hart will halt if it is not currently halted. Writing 1 or 0 has no effect on a hart which is already halted, but the bit must be cleared to 0 before the hart is resumed. Writes apply to the new value of <a href="#">hartsel</a> and <a href="#">hasel</a> .	W	-
resumereq	Writes the resume request bit for all currently selected harts. When set to 1, each selected hart will resume if it is currently halted. The resume request bit is ignored while the halt request bit is set. Writes apply to the new value of <a href="#">hartsel</a> and <a href="#">hasel</a> .	W	-
hartreset	This optional field writes the reset bit for all the currently selected harts. To perform a reset the debugger writes 1, and then writes 0 to deassert the reset signal. If this feature is not implemented, the bit always stays 0, so after writing 1 the debugger can read the register back to see if the feature is supported. Writes apply to the new value of <a href="#">hartsel</a> and <a href="#">hasel</a> .	R/W	0
ackhavereset	Writing 1 to this bit clears the <b>havereset</b> bits for any selected harts. Writes apply to the new value of <a href="#">hartsel</a> and <a href="#">hasel</a> .	W	-
hasel	Selects the definition of currently selected harts. 0: There is a single currently selected hart, that selected by <a href="#">hartsel</a> . 1: There may be multiple currently selected harts – that selected by <a href="#">hartsel</a> , plus those selected by the hart array mask register. An implementation which does not implement the hart array mask register must tie this field to 0. A debugger which wishes to use the hart array mask register feature should set this bit and read back to see if the functionality is supported.	R/W	0

*Continued on next page*

Field	Description	Access	Reset
hartsello	The low 10 bits of <a href="#">hartsel</a> : the DM-specific index of the hart to select. This hart is always part of the currently selected harts.	R/W	0
hartselhi	The high 10 bits of <a href="#">hartsel</a> : the DM-specific index of the hart to select. This hart is always part of the currently selected harts.	R/W	0
setresethaltreq	This optional field writes the halt-on-reset request bit for all currently selected harts. When set to 1, each selected hart will halt upon the next deassertion of its reset. The halt-on-reset request bit is not automatically cleared. The debugger must write to <a href="#">clrresethaltreq</a> to clear it. Writes apply to the new value of <a href="#">hartsel</a> and <a href="#">hasel</a> . If <a href="#">hasresethaltreq</a> is 0, this field is not implemented.	W	0
clrresethaltreq	This optional field clears the halt-on-reset request bit for all currently selected harts. Writes apply to the new value of <a href="#">hartsel</a> and <a href="#">hasel</a> .	W	0
ndmreset	This bit controls the reset signal from the DM to the rest of the system. The signal should reset every part of the system, including every hart, except for the DM and any logic required to access the DM. To perform a system reset the debugger writes 1, and then writes 0 to deassert the reset.	R/W	0
dmactive	This bit serves as a reset signal for the Debug Module itself. 0: The module's state, including authentication mechanism, takes its reset values (the <a href="#">dmactive</a> bit is the only bit which can be written to something other than its reset value). 1: The module functions normally. No other mechanism should exist that may result in resetting the Debug Module after power up, including the platform's system reset or Debug Transport reset signals. A debugger may pulse this bit low to get the Debug Module into a known state. Implementations may use this bit to aid debugging, for example by preventing the Debug Module from being power gated while debugging is active.	R/W	0

### 3.11.3 Hart Info (hartinfo, at 0x12)

This register gives information about the hart currently selected by [hartsel](#).

This register is optional. If it is not present it should read all-zero.

If this register is included, the debugger can do more with the Program Buffer by writing programs which explicitly access the **data** and/or **dscratch** registers.

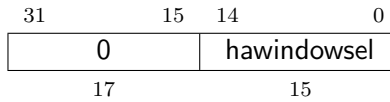
This entire register is read-only.

31	24	23	20	19	17	16	15	12	11	0
0	nscratch				0	dataaccess	datasize	dataaddr		
8	4				3	1	4	12		

Field	Description	Access	Reset
nscratch	Number of <b>dscratch</b> registers available for the debugger to use during program buffer execution, starting from <a href="#">dscratch0</a> . The debugger can make no assumptions about the contents of these registers between commands.	R	Preset
dataaccess	0: The <b>data</b> registers are shadowed in the hart by CSR registers. Each CSR register is MXLEN bits in size, and corresponds to a single argument, per Table 3.1. 1: The <b>data</b> registers are shadowed in the hart's memory map. Each register takes up 4 bytes in the memory map.	R	Preset
datasize	If <a href="#">dataaccess</a> is 0: Number of CSR registers dedicated to shadowing the <b>data</b> registers. If <a href="#">dataaccess</a> is 1: Number of 32-bit words in the memory map dedicated to shadowing the <b>data</b> registers. Since there are at most 12 <b>data</b> registers, the value in this register must be 12 or smaller.	R	Preset
dataaddr	If <a href="#">dataaccess</a> is 0: The number of the first CSR dedicated to shadowing the <b>data</b> registers. If <a href="#">dataaccess</a> is 1: Signed address of RAM where the <b>data</b> registers are shadowed, to be used to access relative to <b>zero</b> .	R	Preset

### 3.11.4 Hart Array Window Select (`hawindowssel`, at 0x14)

This register selects which of the 32-bit portion of the hart array mask register (see Section 3.3.2) is accessible in `hawindow`.

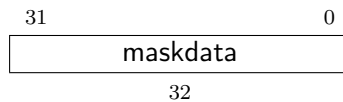


Field	Description	Access	Reset
hawindowssel	The high bits of this field may be tied to 0, depending on how large the array mask register is. Eg. on a system with 48 harts only bit 0 of this field may actually be writable.	R/W	0

### 3.11.5 Hart Array Window (`hawindow`, at 0x15)

This register provides R/W access to a 32-bit portion of the hart array mask register (see Section 3.3.2). The position of the window is determined by `hawindowssel`. I.e. bit 0 refers to hart `hawindowssel * 32`, while bit 31 refers to hart `hawindowssel * 32 + 31`.

Since some bits in the hart array mask register may be constant 0, some bits in this register may be constant 0, depending on the current value of `hawindowssel`.

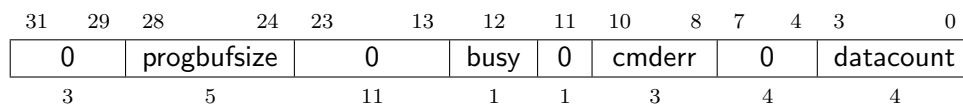


### 3.11.6 Abstract Control and Status (`abstractcs`, at 0x16)

Writing this register while an abstract command is executing causes `cmderr` to be set to 1 (busy) if it is 0.

---

`datacount` must be at least 1 to support RV32 harts, 2 to support RV64 harts, or 4 to support RV128 harts.



Field	Description	Access	Reset
<b>progbufsize</b>	Size of the Program Buffer, in 32-bit words. Valid sizes are 0 - 16.	R	Preset
<b>busy</b>	1: An abstract command is currently being executed. This bit is set as soon as <b>command</b> is written, and is not cleared until that command has completed.	R	0
<b>cmderr</b>	Gets set if an abstract command fails. The bits in this field remain set until they are cleared by writing 1 to them. No abstract command is started until the value is reset to 0. This field only contains a valid value if <b>busy</b> is 0. 0 (none): No error. 1 (busy): An abstract command was executing while <b>command</b> , <b>abstractcs</b> , <b>abstractauto</b> was written, or when one of the <b>data</b> or <b>progbuf</b> registers was read or written. This status is only written if <b>cmderr</b> contains 0. 2 (not supported): The requested command is not supported, regardless of whether the hart is running or not. 3 (exception): An exception occurred while executing the command (eg. while executing the Program Buffer). 4 (halt/resume): The abstract command couldn't execute because the hart wasn't in the required state (running/halted). 5 (bus): The abstract command failed due to a bus error (eg. alignment, access size, timeout). 7 (other): The command failed for another reason.	R/W1C	0
<b>datacount</b>	Number of <b>data</b> registers that are implemented as part of the abstract command interface. Valid sizes are 1 – 12.	R	Preset

### 3.11.7 Abstract Command (**command**, at 0x17)

Writes to this register cause the corresponding abstract command to be executed.

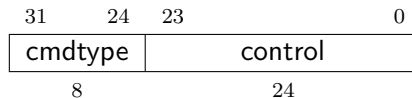
Writing this register while an abstract command is executing causes **cmderr** to be set to 1 (busy) if it is 0.

If **cmderr** is non-zero, writes to this register are ignored.

---



`cmderr` inhibits starting a new command to accommodate debuggers that, for performance reasons, send several commands to be executed in a row without checking `cmderr` in between. They can safely do so and check `cmderr` at the end without worrying that one command failed but then a later command (which might have depended on the previous one succeeding) passed.

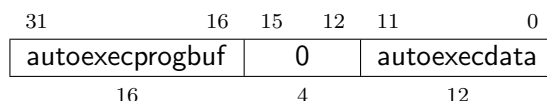


Field	Description	Access	Reset
cmdtype	The type determines the overall functionality of this abstract command.	W	0
control	This field is interpreted in a command-specific manner, described for each abstract command.	W	0

### 3.11.8 Abstract Command Autoexec (abstractauto, at 0x18)

This register is optional. Including it allows more efficient burst accesses. A debugger can detect whether it is support by setting bits and reading them back.

Writing this register while an abstract command is executing causes `cmderr` to be set to 1 (busy) if it is 0.



Field	Description	Access	Reset
autoexecprogbuf	When a bit in this field is 1, read or write accesses to the corresponding <code>progbuf</code> word cause the command in <code>command</code> to be executed again.	R/W	0
autoexecdata	When a bit in this field is 1, read or write accesses to the corresponding <code>data</code> word cause the command in <code>command</code> to be executed again.	R/W	0

### 3.11.9 Device Tree Addr 0 (devtreeaddr0, at 0x19)

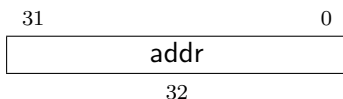
When `devtreevalid` is set, reading this register returns bits 31:0 of the Device Tree address. Reading the other `devtreeaddr` registers returns the upper bits of the address.

When system bus mastering is implemented, this must be an address that can be used with the System Bus Access module. Otherwise, this must be an address that can be used to access the Device Tree from the hart with ID 0.

If `devtreevalid` is 0, then the `devtreeaddr` registers hold identifier information which is not further specified in this document.

The Device Tree itself is described in the RISC-V Privileged Specification.

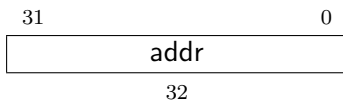
This entire register is read-only.



### 3.11.10 Next Debug Module (nextdm, at 0x1d)

If there is more than one DM accessible on this DMI, this register contains the base address of the next one in the chain, or 0 if this is the last one in the chain.

This entire register is read-only.



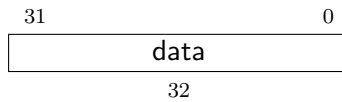
### 3.11.11 Abstract Data 0 (data0, at 0x04)

`data0` through `data11` are basic read/write registers that may be read or changed by abstract commands. `datacount` indicates how many of them are implemented, starting at `sbddata0`, counting up. Table 3.1 shows how abstract commands use these registers.

Accessing these registers while an abstract command is executing causes `cmderr` to be set to 1 (busy) if it is 0.

Attempts to write them while `busy` is set does not change their value.

The values in these registers may not be preserved after an abstract command is executed. The only guarantees on their contents are the ones offered by the command in question. If the command fails, no assumptions can be made about the contents of these registers.

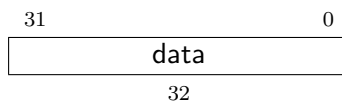


### 3.11.12 Program Buffer 0 (progbuf0, at 0x20)

`progbuf0` through `progbuf15` provide read/write access to the optional program buffer. `progbufsize` indicates how many of them are implemented starting at `progbuf0`, counting up.

Accessing these registers while an abstract command is executing causes `cmderr` to be set to 1 (busy) if it is 0.

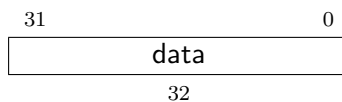
Attempts to write them while `busy` is set does not change their value.



### 3.11.13 Authentication Data (authdata, at 0x30)

This register serves as a 32-bit serial port to the authentication module.

When `authbusy` is clear, the debugger can communicate with the authentication module by reading or writing this register. There is no separate mechanism to signal overflow/underflow.

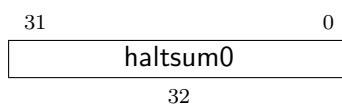


### 3.11.14 Halt Summary 0 (haltsum0, at 0x40)

Each bit in this read-only register indicates whether one specific hart is halted or not. Unavailable/nonexistent harts are not considered to be halted.

The LSB reflects the halt status of hart `{hartsel[19:5],5'h0}`, and the MSB reflects halt status of hart `{hartsel[19:5],5'h1f}`.

This entire register is read-only.



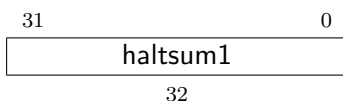
### 3.11.15 Halt Summary 1 (haltsum1, at 0x13)

Each bit in this read-only register indicates whether any of a group of harts is halted or not. Unavailable/nonexistent harts are not considered to be halted.

This register may not be present in systems with fewer than 33 harts.

The LSB reflects the halt status of harts {hartsel[19:10],10'h0} through {hartsel[19:10],10'h1f}. The MSB reflects the halt status of harts {hartsel[19:10],10'h3e0} through {hartsel[19:10],10'h3ff}.

This entire register is read-only.



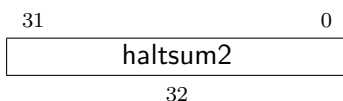
### 3.11.16 Halt Summary 2 (haltsum2, at 0x34)

Each bit in this read-only register indicates whether any of a group of harts is halted or not. Unavailable/nonexistent harts are not considered to be halted.

This register may not be present in systems with fewer than 1025 harts.

The LSB reflects the halt status of harts {hartsel[19:15],15'h0} through {hartsel[19:15],15'h3ff}. The MSB reflects the halt status of harts {hartsel[19:15],15'h7c00} through {hartsel[19:15],15'h7fff}.

This entire register is read-only.



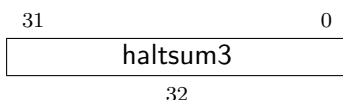
### 3.11.17 Halt Summary 3 (haltsum3, at 0x35)

Each bit in this read-only register indicates whether any of a group of harts is halted or not. Unavailable/nonexistent harts are not considered to be halted.

This register may not be present in systems with fewer than 32769 harts.

The LSB reflects the halt status of harts 20'h0 through 20'h7fff. The MSB reflects the halt status of harts 20'hf8000 through 20'hffff.

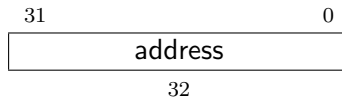
This entire register is read-only.



### 3.11.18 System Bus Address 127:96 (sbaddress3, at 0x37)

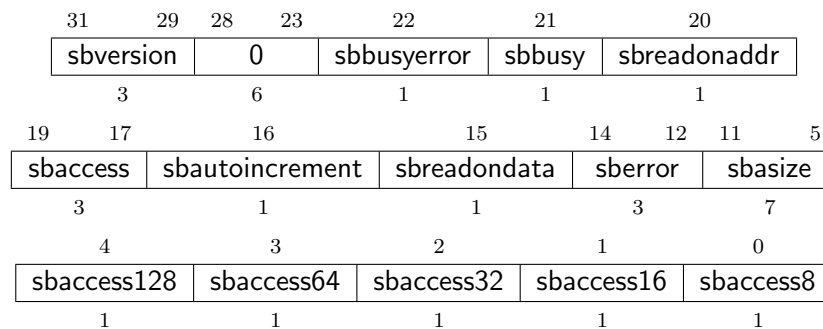
If **sbasize** is less than 97, then this register is not present.

When the system bus master is busy, writes to this register will set **sbbusyerror** and don't do anything else.



Field	Description	Access	Reset
address	Accesses bits 127:96 of the physical address in <b>sbaddress</b> (if the system address bus is that wide).	R/W	0

### 3.11.19 System Bus Access Control and Status (sbcs, at 0x38)



Field	Description	Access	Reset
sbversion	0: The System Bus interface conforms to mainline drafts of this spec older than 1 January, 2018. 1: The System Bus interface conforms to this version of the spec. Other values are reserved for future versions.	R	1
sbbusyerror	Set when the debugger attempts to read data while a read is in progress, or when the debugger initiates a new access while one is already in progress (while <b>sbbusy</b> is set). It remains set until it's explicitly cleared by the debugger. While this field is non-zero, no more system bus accesses can be initiated by the Debug Module.	R/W1C	0

*Continued on next page*

Field	Description	Access	Reset
sbbusy	When 1, indicates the system bus master is busy. (Whether the system bus itself is busy is related, but not the same thing.) This bit goes high immediately when a read or write is requested for any reason, and does not go low until the access is fully completed. Writes to <b>sbcs</b> while <b>sbbusy</b> is high result in undefined behavior. A debugger must not write to <b>sbcs</b> until it reads <b>sbbusy</b> as 0.	R	0
sbreadonaddr	When 1, every write to <b>sbaddress0</b> automatically triggers a system bus read at the new address.	R/W	0
sbaccess	Select the access size to use for system bus accesses. 0: 8-bit 1: 16-bit 2: 32-bit 3: 64-bit 4: 128-bit If <b>sbaccess</b> has an unsupported value when the DM starts a bus access, the access is not performed and <b>sberror</b> is set to 3.	R/W	2
sbautoincrement	When 1, <b>sbaddress</b> is incremented by the access size (in bytes) selected in <b>sbaccess</b> after every system bus access.	R/W	0
sbreadondata	When 1, every read from <b>sbdata0</b> automatically triggers a system bus read at the (possibly auto-incremented) address.	R/W	0
sberror	When the Debug Module's system bus master causes a bus error, this field gets set. The bits in this field remain set until they are cleared by writing 1 to them. While this field is non-zero, no more system bus accesses can be initiated by the Debug Module. An implementation may report "Other" (7) for any error condition. 0: There was no bus error. 1: There was a timeout. 2: A bad address was accessed. 3: There was an alignment error. 4: An access of unsupported size was requested. 7: Other.	R/W1C	0
sbsize	Width of system bus addresses in bits. (0 indicates there is no bus access support.)	R	Preset
sbaccess128	1 when 128-bit system bus accesses are supported.	R	Preset
sbaccess64	1 when 64-bit system bus accesses are supported.	R	Preset

*Continued on next page*

Field	Description	Access	Reset
sbaccess32	1 when 32-bit system bus accesses are supported.	R	Preset
sbaccess16	1 when 16-bit system bus accesses are supported.	R	Preset
sbaccess8	1 when 8-bit system bus accesses are supported.	R	Preset

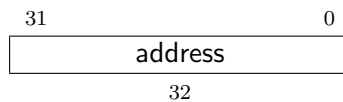
### 3.11.20 System Bus Address 31:0 (sbaddress0, at 0x39)

If `sbasize` is 0, then this register is not present.

When the system bus master is busy, writes to this register will set `sbbusyerror` and don't do anything else.

If `sberror` is 0, `sbbusyerror` is 0, and `sbreadonaddr` is set then writes to this register start the following:

1. Set `sbbusy`.
2. Perform a bus read from the new value of `sbaddress`.
3. If the read succeeded and `sbautoincrement` is set, increment `sbaddress`.
4. Clear `sbbusy`.

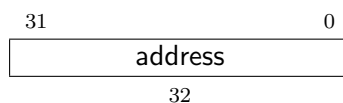


Field	Description	Access	Reset
address	Accesses bits 31:0 of the physical address in <code>sbaddress</code> .	R/W	0

### 3.11.21 System Bus Address 63:32 (sbaddress1, at 0x3a)

If `sbasize` is less than 33, then this register is not present.

When the system bus master is busy, writes to this register will set `sbbusyerror` and don't do anything else.

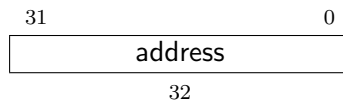


Field	Description	Access	Reset
<b>address</b>	Accesses bits 63:32 of the physical address in <b>sbaddress</b> (if the system address bus is that wide).	R/W	0

### 3.11.22 System Bus Address 95:64 (sbaddress2, at 0x3b)

If **sbasize** is less than 65, then this register is not present.

When the system bus master is busy, writes to this register will set **sbbusyerror** and don't do anything else.



Field	Description	Access	Reset
<b>address</b>	Accesses bits 95:64 of the physical address in <b>sbaddress</b> (if the system address bus is that wide).	R/W	0

### 3.11.23 System Bus Data 31:0 (sbddata0, at 0x3c)

If all of the **sbaccess** bits in **sbcs** are 0, then this register is not present.

Any successful system bus read updates **sbddata**. If the width of the read access is less than the width of **sbddata**, the contents of the remaining high bits may take on any value.

If **sberror** or **sbbusyerror** both aren't 0 then accesses do nothing.

If the bus master is busy then accesses set **sbbusyerror**, and don't do anything else.

Writes to this register start the following:

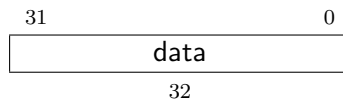
1. Set **sbbusy**.
2. Perform a bus write of the new value of **sbddata** to **sbaddress**.
3. If the write succeeded and **sbautoincrement** is set, increment **sbaddress**.
4. Clear **sbbusy**.



Reads from this register start the following:

1. “Return” the data.
2. Set `sbbusy`.
3. If `sbautoincrement` is set, increment `sbaddress`.
4. If `sbreadondata` is set, perform another system bus read.
5. Clear `sbbusy`.

Only `sbdata0` has this behavior. The other `sbdata` registers have no side effects. On systems that have buses wider than 32 bits, a debugger should access `sbdata0` after accessing the other `sbdata` registers.

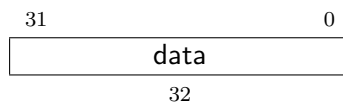


Field	Description	Access	Reset
<code>data</code>	Accesses bits 31:0 of <code>sbdata</code> .	R/W	0

### 3.11.24 System Bus Data 63:32 (`sbdata1`, at 0x3d)

If `sbaccess64` and `sbaccess128` are 0, then this register is not present.

If the bus master is busy then accesses set `sbbusyerror`, and don't do anything else.

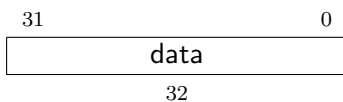


Field	Description	Access	Reset
<code>data</code>	Accesses bits 63:32 of <code>sbdata</code> (if the system bus is that wide).	R/W	0

### 3.11.25 System Bus Data 95:64 (`sbdata2`, at 0x3e)

This register only exists if `sbaccess128` is 1.

If the bus master is busy then accesses set [sbbusyerror](#), and don't do anything else.

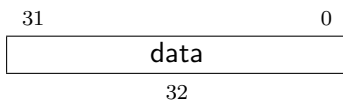


Field	Description	Access	Reset
<b>data</b>	Accesses bits 95:64 of <b>sbddata</b> (if the system bus is that wide).	R/W	0

### 3.11.26 System Bus Data 127:96 (sbddata3, at 0x3f)

This register only exists if [sbaccess128](#) is 1.

If the bus master is busy then accesses set [sbbusyerror](#), and don't do anything else.



Field	Description	Access	Reset
<b>data</b>	Accesses bits 127:96 of <b>sbddata</b> (if the system bus is that wide).	R/W	0

## Chapter 4

# RISC-V Debug

Modifications to the RISC-V core to support debug are kept to a minimum. There is a special execution mode (Debug Mode) and a few extra CSRs. The DM takes care of the rest.

### 4.1 Debug Mode

Debug Mode is a special processor mode used only when a hart is halted for external debugging. How Debug Mode is implemented is not specified here.

When executing code from the Program Buffer, the hart stays in Debug Mode and the following apply:

1. All operations are executed at machine mode privilege level, except that `mprv` in `mstatus` may be ignored according to `mprven`.

---

*In general, the debugger is expected to be able to simulate all the effects of `mprv`. The exception is the case of Sv32 systems, which need `mprv` functionality in order to access 34-bit physical addresses. Other systems are likely to tie `mprven` to 0.*

2. All interrupts (including NMI) are masked.
3. Exceptions don't update any registers. That includes `cause`, `epc`, `tval`, `dpc`, and `mstatus`. They do end execution of the Program Buffer.
4. No action is taken if a trigger matches.
5. Trace is disabled.
6. Counters may be stopped, depending on `stopcount` in `dcsr`.
7. Timers may be stopped, depending on `stoptime` in `dcsr`.
8. The `wfi` instruction acts as a `nop`.
9. Almost all instructions that change the privilege level have undefined behavior. This includes `ecall`, `mret`, `hret`, `sret`, and `uret`. (To change the privilege level, the debugger can write `prv` in `dcsr`). The only exception is `ebreak`. When that is executed in Debug Mode, it halts the hart again but without updating `dpc` or `dcsr`.
10. Completing Program Buffer execution is considered output for the purpose of `fence` instructions.

## 4.2 Load-Reserved/Store-Conditional Instructions

The reservation registered by an `lr` instruction on a memory address may be lost when entering Debug Mode or while in Debug Mode. This means that there may be no forward progress if Debug Mode is entered between `lr` and `sc` pairs.

---

*This is a behavior that debug users must be aware of. If they have a breakpoint set between a `lr` and `sc` pair, or are stepping through such code, the `sc` may never succeed. Fortunately in general use there will be very few instructions in such a sequence, and anybody debugging it will quickly notice that the reservation is not occurring. The solution in that case is to set a breakpoint on the first instruction after the `sc` and run to it. A higher level debugger may choose to automate this.*

---

## 4.3 Single Step

A debugger can cause a halted hart to execute a single instruction and then re-enter Debug Mode by setting `step` before setting `resumereq`.

If executing or fetching that instruction causes an exception, Debug Mode is re-entered immediately after the PC is changed to the exception handler and the appropriate `tval` and `cause` registers are updated.

If executing or fetching the instruction causes a trigger to fire, Debug Mode is re-entered immediately after that trigger has fired. In that case `cause` is set to 2 (trigger) instead of 4 (single step). Whether the instruction is executed or not depends on the specific configuration of the trigger.

If the instruction that is executed causes the PC to change to an address where an instruction fetch causes an exception, that exception does not occur until the next time the hart is resumed. Similarly, a trigger at the new address does not fire until the hart actually attempts to execute that instruction.

## 4.4 Reset

If the halt signal (driven by the hart's halt request bit in the Debug Module) is asserted when a hart comes out of reset, the hart must enter Debug Mode before executing any instructions, but after performing any initialization that would usually happen before the first instruction is executed.

## 4.5 dret Instruction

To return from Debug Mode, a new instruction is defined: `dret`. It has an encoding of 0x7b200073. On harts which support this instruction, executing `dret` in Debug Mode changes `pc` to the value stored in `dpc`. The current privilege level is changed to that specified by `prv` in `dcsr`. The hart is no longer in debug mode.

Executing `dret` outside of Debug Mode causes an illegal instruction exception.

It is not necessary for the debugger to know whether an implementation supports `dret`, as the Debug Module will ensure that it is executed if necessary. It is defined in this specification only to reserve the opcode and allow for reusable Debug Module implementations.

## 4.6 Core Debug Registers

The supported Core Debug Registers must be implemented for each hart that can be debugged. They are CSRs, accessible using the RISC-V `csr` opcodes and optionally also using abstract debug commands.

These registers are only accessible from Debug Mode.

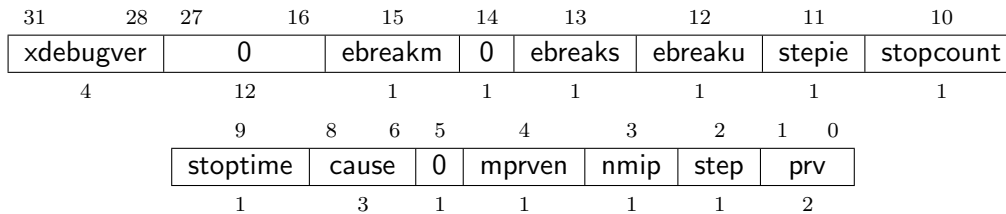
Table 4.1: Core Debug Registers

Address	Name	Page
0x7b0	Debug Control and Status	<a href="#">39</a>
0x7b1	Debug PC	<a href="#">41</a>
0x7b2	Debug Scratch Register 0	
0x7b3	Debug Scratch Register 1	

### 4.6.1 Debug Control and Status (`dcsr`, at 0x7b0)

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*cause* priorities are assigned such that the least predictable events have the highest priority.



Field	Description	Access	Reset
xdebugver	0: There is no external debug support. 4: External debug support exists as it is described in this document. 15: There is external debug support, but it does not conform to any available version of this spec.	R	Preset
ebreakm	When 1, <code>ebreak</code> instructions in Machine Mode enter Debug Mode.	R/W	0
ebreaks	When 1, <code>ebreak</code> instructions in Supervisor Mode enter Debug Mode.	R/W	0

*Continued on next page*

Field	Description	Access	Reset
<b>ebreaku</b>	When 1, <b>ebreak</b> instructions in User/Application Mode enter Debug Mode.	R/W	0
<b>stepie</b>	0: Interrupts are disabled during single stepping. 1: Interrupts are enabled during single stepping. Implementations may hard wire this bit to 0. The debugger must read back the value it writes to check whether the feature is supported. If not supported, interrupt behavior can be emulated by the debugger.	R/W	0
<b>stopcount</b>	0: Increment counters as usual. 1: Don't increment any counters while in Debug Mode or on <b>ebreak</b> instructions that cause entry into Debug Mode. These counters include the <b>cycle</b> and <b>instret</b> CSRs. This is preferred for most debugging scenarios. An implementation may choose not to support writing to this bit. The debugger must read back the value it writes to check whether the feature is supported.	R/W	Preset
<b>stoptime</b>	0: Increment timers as usual. 1: Don't increment any hart-local timers while in Debug Mode. An implementation may choose not to support writing to this bit. The debugger must read back the value it writes to check whether the feature is supported.	R/W	Preset
<b>cause</b>	Explains why Debug Mode was entered. When there are multiple reasons to enter Debug Mode in a single cycle, hardware should set <b>cause</b> to the cause with the highest priority. 1: An <b>ebreak</b> instruction was executed. (priority 3) 2: The Trigger Module caused a breakpoint exception. (priority 4, highest) 3: The debugger requested entry to Debug Mode. (priority 2) 4: The hart single stepped because <b>step</b> was set. (priority 1, lowest) Other values are reserved for future use.	R	0
<b>mprven</b>	When 1, <b>mprv</b> in <b>mstatus</b> takes effect during debug mode. When 0, it is ignored during debug mode. Implementing this bit is optional. If not implemented it should be tied to 0.	R/W	0

*Continued on next page*

Field	Description	Access	Reset
<b>nmip</b>	When set, there is a Non-Maskable-Interrupt (NMI) pending for the hart. Since an NMI can indicate a hardware error condition, reliable debugging may no longer be possible once this bit becomes set. This is implementation-dependent.	R	0
<b>step</b>	When set and not in Debug Mode, the hart will only execute a single instruction and then enter Debug Mode. If the instruction does not complete due to an exception, the hart will immediately enter Debug Mode before executing the trap handler, with appropriate exception registers set.	R/W	0
<b>prv</b>	Contains the privilege level the hart was operating in when Debug Mode was entered. The encoding is described in Table 4.5. A debugger can change this value to change the hart's privilege level when exiting Debug Mode. Not all privilege levels are supported on all harts. If the encoding written is not supported or the debugger is not allowed to change to it, the hart may change to any supported privilege level.	R/W	3

#### 4.6.2 Debug PC (dpc, at 0x7b1)

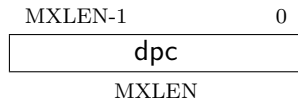
Upon entry to debug mode, **dpc** is updated with the virtual address of the next instruction to be executed. The behavior is described in more detail in Table 4.3.

Table 4.3: Virtual address in DPC upon Debug Mode Entry

Cause	Virtual Address in DPC
<b>ebreak</b>	Address of the <b>ebreak</b> instruction
single step	Address of the instruction that would be executed next if no debugging was going on. Ie. <b>pc</b> + 4 for 32-bit instructions that don't change program flow, the destination PC on taken jumps/branches, etc.
trigger module	If <b>timing</b> is 0, the address of the instruction which caused the trigger to fire. If <b>timing</b> is 1, the address of the next instruction to be executed at the time that debug mode was entered.
halt request	Address of the next instruction to be executed at the time that debug mode was entered

When resuming, the hart's PC is updated to the virtual address stored in **dpc**. A debugger may

write `dpc` to change where the hart resumes.



### 4.6.3 Debug Scratch Register 0 (dscratch0, at 0x7b2)

Optional scratch register that can be used by implementations that need it. A debugger must not write to this register unless `hartinfo` explicitly mentions it (the Debug Module may use this register internally).

### 4.6.4 Debug Scratch Register 1 (dscratch1, at 0x7b3)

Optional scratch register that can be used by implementations that need it. A debugger must not write to this register unless `hartinfo` explicitly mentions it (the Debug Module may use this register internally).

## 4.7 Virtual Debug Registers

Virtual debug registers are a requirement on the debugger SW/interface, not on the Core designer.

Users of the debugger shouldn't need to know about the core debug registers, but may want to change things affected by them. A virtual register is one that doesn't exist directly in the hardware, but that the debugger exposes as if it does.

Table 4.4: Virtual Core Debug Registers

Address	Name	Page
virtual	Privilege Level	42

### 4.7.1 Privilege Level (priv, at virtual)

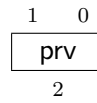
User can read this register to inspect the privilege level that the hart was running in when the hart halted. User can write this register to change the privilege level that the hart will run in when it resumes.

This register contains `prv` from `dcsr`, but in a place that the user is expected to access. The user should not access `dcsr` directly, because doing so might interfere with the debugger.



Table 4.5: Privilege Level Encoding

Encoding	Privilege Level
0	User/Application
1	Supervisor
3	Machine



Field	Description	Access	Reset
<b>prv</b>	Contains the privilege level the hart was operating in when Debug Mode was entered. The encoding is described in Table 4.5, and matches the privilege level encoding from the RISC-V Privileged ISA Specification. A user can write this value to change the hart's privilege level when exiting Debug Mode.	R/W	0



## Chapter 5

# Trigger Module

Triggers can cause a breakpoint exception, entry into Debug Mode, or a trace action without having to execute a special instruction. This makes them invaluable when debugging code from ROM. They can trigger on execution of instructions at a given memory address, or on the address/data in loads/stores. These are all features that can be useful without having the Debug Module present, so the Trigger Module is broken out as a separate piece that can be implemented separately.

Each trigger may support a variety of features. A debugger can build a list of all triggers and their features as follows:

1. Write 0 to `tselect`.
2. Read back `tselect` and check that it contains the written value. If not, exit the loop.
3. Read `tinfo`.
4. If that caused an exception, the debugger must read `tdata1` to discover the type. (If `type` is 0, this trigger doesn't exist. Exit the loop.)
5. If `info` is 1, this trigger doesn't exist. Exit the loop.
6. Otherwise, the selected trigger supports the types discovered in `info`.
7. Repeat, incrementing the value in `tselect`.

---

*There are two ways to check whether a given trigger is the last one to support these implementations:*

1. *When no hardware triggers are implemented at all, all related registers return 0. The algorithm above terminates when checking `type`.*
2. *When 2 triggers are implemented, `tselect` is just a single bit that selects one of the two. When the debugger writes 2, it reads back as 0 which terminates the enumeration.*

### 5.1 Trigger Registers

These registers are CSRs, accessible using the RISC-V `csr` opcodes and optionally also using abstract debug commands.

Most trigger functionality is optional. All `tdata` registers follow write-any-read-legal semantics. If a debugger writes an unsupported configuration, the register will read back a value that is supported (which may simply be a disabled trigger). This means that a debugger must always read back values it writes to `tdata` registers, unless it already knows already what is supported. Writes to one `tdata` register may not modify the contents of other `tdata` registers, nor the configuration of any trigger besides the one that is currently selected.

The trigger registers are only accessible in machine and Debug Mode to prevent untrusted user code from causing entry into Debug Mode without the OS's permission.

Table 5.1: `action` encoding

Value	Description
0	Raise a breakpoint exception. (Used when software wants to use the trigger module without an external debugger attached.)
1	Enter Debug Mode. (Only supported when the trigger's <code>dmode</code> is 1.)
2	Start tracing.
3	Stop tracing.
4	Emit trace data for this match. If it is a data access match, emit appropriate Load/Store Address/Data. If it is an instruction execution, emit its PC.
other	Reserved for future use.

Table 5.2: Trigger Registers

Address	Name	Page
0x7a0	Trigger Select	<a href="#">46</a>
0x7a1	Trigger Data 1	<a href="#">47</a>
0x7a1	Match Control	<a href="#">49</a>
0x7a1	Instruction Count	<a href="#">52</a>
0x7a1	Interrupt Trigger	<a href="#">53</a>
0x7a1	Exception Trigger	<a href="#">53</a>
0x7a2	Trigger Data 2	<a href="#">48</a>
0x7a3	Trigger Data 3	<a href="#">48</a>
0x7a4	Trigger Info	<a href="#">48</a>

### 5.1.1 Trigger Select (`tselect`, at 0x7a0)

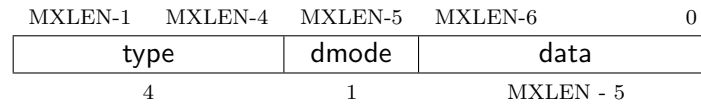
This register determines which trigger is accessible through the other trigger registers. The set of accessible triggers must start at 0, and be contiguous.

Writes of values greater than or equal to the number of supported triggers may result in a different value in this register than what was written. To verify that what they wrote is a valid index, debuggers can read back the value and check that `tselect` holds what they wrote.

Since triggers can be used both by Debug Mode and M Mode, the debugger must restore this register if it modifies it.



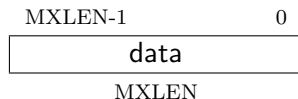
### 5.1.2 Trigger Data 1 (tdata1, at 0x7a1)



Field	Description	Access	Reset
type	0: There is no trigger at this <a href="#">tselect</a> . 1: The trigger is a legacy SiFive address match trigger. These should not be implemented and aren't further documented here. 2: The trigger is an address/data match trigger. The remaining bits in this register act as described in <a href="#">mcontrol</a> . 3: The trigger is an instruction count trigger. The remaining bits in this register act as described in <a href="#">icount</a> . 4: The trigger is an interrupt trigger. The remaining bits in this register act as described in <a href="#">itrigger</a> . 5: The trigger is an exception trigger. The remaining bits in this register act as described in <a href="#">etrigger</a> . 15: This trigger exists (so enumeration shouldn't terminate), but is not currently available. Other values are reserved for future use.	R/W	Preset
dmode	0: Both Debug and M Mode can write the <b>tdata</b> registers at the selected <a href="#">tselect</a> . 1: Only Debug Mode can write the <b>tdata</b> registers at the selected <a href="#">tselect</a> . Writes from other modes are ignored. This bit is only writable from Debug Mode.	R/W	0
data	Trigger-specific data.	R/W	Preset

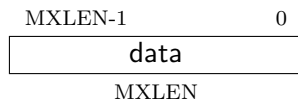
### 5.1.3 Trigger Data 2 (tdata2, at 0x7a2)

Trigger-specific data.



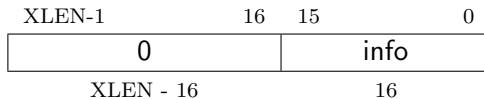
### 5.1.4 Trigger Data 3 (tdata3, at 0x7a3)

Trigger-specific data.



### 5.1.5 Trigger Info (tinfo, at 0x7a4)

This entire register is read-only.



Field	Description	Access	Reset
info	<p>One bit for each possible <a href="#">type</a> enumerated in <a href="#">tdata1</a>. Bit N corresponds to type N. If the bit is set, then that type is supported by the currently selected trigger.</p> <p>If the currently selected trigger doesn't exist, this field contains 1.</p> <p>If <a href="#">type</a> is not writable, this register may be unimplemented, in which case reading it causes an illegal instruction exception. In this case the debugger can read the only supported type from <a href="#">tdata1</a>.</p>	R	Preset

### 5.1.6 Match Control (mcontrol, at 0x7a1)

This register is accessible as `tdata1` when `type` is 2.

Address and data trigger implementation are heavily dependent on how the processor core is implemented. To accommodate various implementations, execute, load, and store address/data triggers may fire at whatever point in time is most convenient for the implementation. The debugger may request specific timings as described in [timing](#). Table 5.5 suggests timings for the best user experience.

Table 5.5: Suggested Breakpoint Timings

Match Type	Suggested Trigger Timing
Execute Address	Before
Execute Instruction	Before
Execute Address+Instruction	Before
Load Address	Before
Load Data	After
Load Address+Data	After
Store Address	Before
Store Data	Before
Store Address+Data	Before

This trigger type may be limited to address comparisons (`select` is always 0) only. If that is the case, then `tdata2` must be able to hold all valid virtual addresses but it need not be capable of holding other values.

MXLEN-1	MXLEN-4	MXLEN-5	MXLEN-6	MXLEN-11	MXLEN-12	21	20	19					
type		dmode	maskmax			0		hit	select				
4		1		6			MXLEN - 32		1	1			
18		17	12	11	10	7	6	5	4	3	2	1	0
timing		action		chain	match		m	0	s	u	execute	store	load
1		6		1	4		1	1	1	1	1	1	1

Field	Description	Access	Reset
maskmax	Specifies the largest naturally aligned powers-of-two (NAPOT) range supported by the hardware when <code>match</code> is 1. The value is the logarithm base 2 of the number of bytes in that range. A value of 0 indicates that only exact value matches are supported (one byte range). A value of 63 corresponds to the maximum NAPOT range, which is $2^{63}$ bytes in size.	R	Preset

*Continued on next page*

Field	Description	Access	Reset
hit	If this optional bit is implemented, the hardware sets it when this trigger matches. The trigger's user can set or clear it at any time. The trigger's user can use this bit to determine which trigger(s) matched. If the bit is not implemented, it is always 0 and writing it has no effect.	R/W	0
select	0: Perform a match on the virtual address. 1: Perform a match on the data value loaded/stored, or the instruction executed.	R/W	0
timing	0: The action for this trigger will be taken just before the instruction that triggered it is executed, but after all preceding instructions are committed. 1: The action for this trigger will be taken after the instruction that triggered it is executed. It should be taken before the next instruction is executed, but it is better to implement triggers and not implement that suggestion than to not implement them at all. Most hardware will only implement one timing or the other, possibly dependent on <a href="#">select</a> , <a href="#">execute</a> , <a href="#">load</a> , and <a href="#">store</a> . This bit primarily exists for the hardware to communicate to the debugger what will happen. Hardware may implement the bit fully writable, in which case the debugger has a little more control. Data load triggers with <a href="#">timing</a> of 0 will result in the same load happening again when the debugger lets the hart run. For data load triggers, debuggers must first attempt to set the breakpoint with <a href="#">timing</a> of 1. A chain of triggers that don't all have the same <a href="#">timing</a> value will never fire (unless consecutive instructions match the appropriate triggers).	R/W	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	R/W	0

*Continued on next page*



Field	Description	Access	Reset
chain	<p>0: When this trigger matches, the configured action is taken.</p> <p>1: While this trigger does not match, it prevents the trigger with the next index from matching.</p> <p>A trigger chain starts on the first trigger with <b>chain</b> = 1 after a trigger with <b>chain</b> = 0, or simply on the first trigger if that has <b>chain</b> = 1. It ends on the first trigger after that that has <b>chain</b> = 0. This final trigger is part of the chain. The action on that final trigger will be taken if and only if all the triggers in the chain match at the same time. Because <b>chain</b> affects the next trigger, hardware must zero it in writes to <b>mcontrol</b> that set <b>dmode</b> to 0 if the next trigger has <b>dmode</b> of 1. In addition hardware should ignore writes to <b>mcontrol</b> that set <b>dmode</b> to 1 if the previous trigger has both <b>dmode</b> of 0 and <b>chain</b> of 1. Debuggers must avoid the latter case by checking <b>chain</b> on the previous trigger if they're writing <b>mcontrol</b>.</p> <p>Implementations that wish to limit the maximum length of a trigger chain (eg. to meet timing requirements) may do so by zeroing <b>chain</b> in writes to <b>mcontrol</b> that would make the chain too long.</p>	R/W	0
match	<p>0: Matches when the value equals <b>tdata2</b>.</p> <p>1: Matches when the top M bits of the value match the top M bits of <b>tdata2</b>. M is MXLEN-1 minus the index of the least-significant bit containing 0 in <b>tdata2</b>.</p> <p>2: Matches when the value is greater than (unsigned) or equal to <b>tdata2</b>.</p> <p>3: Matches when the value is less than (unsigned) <b>tdata2</b>.</p> <p>4: Matches when the lower half of the value equals the lower half of <b>tdata2</b> after the lower half of the value is ANDed with the upper half of <b>tdata2</b>.</p> <p>5: Matches when the upper half of the value equals the lower half of <b>tdata2</b> after the upper half of the value is ANDed with the upper half of <b>tdata2</b>.</p> <p>Other values are reserved for future use.</p>	R/W	0
m	When set, enable this trigger in M mode.	R/W	0
s	When set, enable this trigger in S mode.	R/W	0
u	When set, enable this trigger in U mode.	R/W	0
execute	When set, the trigger fires on the virtual address or opcode of an instruction that is executed.	R/W	0

*Continued on next page*

Field	Description	Access	Reset
store	When set, the trigger fires on the virtual address or data of a store.	R/W	0
load	When set, the trigger fires on the virtual address or data of a load.	R/W	0

### 5.1.7 Instruction Count (icount, at 0x7a1)

This register is accessible as `tdata1` when `type` is 3.

*This trigger type is intended to be used as a single step that's useful both for external debuggers and for software monitor programs. For that case it is not necessary to support `count` greater than 1. The only two combinations of the mode bits that are useful in those scenarios are `u` by itself, or `m`, `s`, and `u` all set.*

*If the hardware limits `count` to 1, and changes mode bits instead of decrementing `count`, this register can be implemented with just 2 bits. One for `u`, and one for `m` and `s` tied together. If only the external debugger or only a software monitor needs to be supported, a single bit is enough.*

MXLEN-1	MXLEN-4	MXLEN-5	MXLEN-6	25	24	23	10	9	8	7	6	5	0
type	dmode	0	hit	count	m	0	s	u	action				
4	1	MXLEN - 30	1	14	1	1	1	1	6				

Field	Description	Access	Reset
hit	If this optional bit is implemented, the hardware sets it when this trigger matches. The trigger's user can set or clear it at any time. The trigger's user can use this bit to determine which trigger(s) matched. If the bit is not implemented, it is always 0 and writing it has no effect.	R/W	0
count	When count is decremented to 0, the trigger fires. Instead of changing <code>count</code> from 1 to 0, it is also acceptable for hardware to clear <code>m</code> , <code>s</code> , and <code>u</code> . This allows <code>count</code> to be hard-wired to 1 if this register just exists for single step.	R/W	1
m	When set, every instruction completed or exception taken in M mode decrements <code>count</code> by 1.	R/W	0
s	When set, every instruction completed or exception taken in S mode decrements <code>count</code> by 1.	R/W	0
u	When set, every instruction completed or exception taken in U mode decrements <code>count</code> by 1.	R/W	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	R/W	0

### 5.1.8 Interrupt Trigger (ittrigger, at 0x7a1)

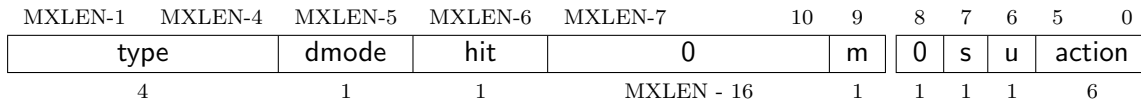
This register is accessible as `tdata1` when `type` is 4.

This trigger may fire on any of the interrupts configurable in `mie` (described in the Privileged Spec). The interrupts to fire on are configured by setting the same bit in `tdata2` as would be set in `mie` to enable the interrupt.

Hardware may only support a subset of interrupts for this trigger. A debugger must read back `tdata2` after writing it to confirm the requested functionality is actually supported.

The trigger only fires if the hart takes a trap because of the interrupt. (Eg. it does not fire when a timer interrupt occurs but that interrupt is not enabled in `mie`.)

When the trigger fires, all CSRs are updated as defined by the Privileged Spec, and the requested action is taken just before the first instruction of the interrupt/exception handler is executed.



Field	Description	Access	Reset
hit	If this optional bit is implemented, the hardware sets it when this trigger matches. The trigger's user can set or clear it at any time. The trigger's user can use this bit to determine which trigger(s) matched. If the bit is not implemented, it is always 0 and writing it has no effect.	R/W	0
m	When set, enable this trigger for interrupts that are taken from M mode.	R/W	0
s	When set, enable this trigger for interrupts that are taken from S mode.	R/W	0
u	When set, enable this trigger for interrupts that are taken from U mode.	R/W	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	R/W	0

### 5.1.9 Exception Trigger (etrigger, at 0x7a1)

This register is accessible as `tdata1` when `type` is 5.

This trigger may fire on up to MXLEN of the Exception Codes defined in `mcause` (described in the Privileged Spec, with Interrupt=0). Those causes are configured by writing the corresponding bit in `tdata2`. (Eg. to trap on an illegal instruction, the debugger sets bit 2 in `tdata2`.)

Hardware may support only a subset of exceptions. A debugger must read back `tdata2` after writing it to confirm the requested functionality is actually supported.

When the trigger fires, all CSRs are updated as defined by the Privileged Spec, and the requested action is taken just before the first instruction of the interrupt/exception handler is executed.

MXLEN-1	MXLEN-4	MXLEN-5	MXLEN-6	MXLEN-7	10	9	8	7	6	5	0
type	dmode	hit	0	m	0	s	u	action			
4	1	1	MXLEN - 16	1	1	1	1	6			

Field	Description	Access	Reset
hit	If this optional bit is implemented, the hardware sets it when this trigger matches. The trigger's user can set or clear it at any time. The trigger's user can use this bit to determine which trigger(s) matched. If the bit is not implemented, it is always 0 and writing it has no effect.	R/W	0
m	When set, enable this trigger for exceptions that are taken from M mode.	R/W	0
s	When set, enable this trigger for exceptions that are taken from S mode.	R/W	0
u	When set, enable this trigger for exceptions that are taken from U mode.	R/W	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	R/W	0

## Chapter 6

# Debug Transport Module (DTM)

Debug Transport Modules provide access to the DM over one or more transports (eg. JTAG or USB).

There may be multiple DTMs in a single platform. Ideally every component that communicates with the outside world includes a DTM, allowing a platform to be debugged through every transport it supports. For instance a USB component could include a DTM. This would trivially allow any platform to be debugged over USB. All that is required is that the USB module already in use also has access to the Debug Module Interface.

Using multiple DTMs at the same time is not supported. It is left to the user to ensure this does not happen.

This specification defines a JTAG DTM in Section 6.1. Additional DTMs may be added in future versions of this specification.

## 6.1 JTAG Debug Transport Module

This Debug Transport Module is based around a normal JTAG Test Access Port (TAP). The JTAG TAP allows access to arbitrary JTAG registers by first selecting one using the JTAG instruction register (IR), and then accessing it through the JTAG data register (DR).

### 6.1.1 JTAG Background

JTAG refers to IEEE Std 1149.1-2013. It is a standard that defines test logic that can be included in an integrated circuit to test the interconnections between integrated circuits, test the integrated circuit itself, and observe or modify circuit activity during the components normal operation. This specification uses the latter functionality. The JTAG standard defines a Test Access Port (TAP) that can be used to read and write a few custom registers, which can be used to communicate with debug hardware in a component.

### 6.1.2 JTAG DTM Registers

JTAG TAPs used as a DTM must have an IR of at least 5 bits. When the TAP is reset, IR must default to 00001, selecting the IDCODE instruction. A full list of JTAG registers along with their encoding is in Table 6.1. If the IR actually has more than 5 bits, then the encodings in Table 6.1 should be extended with 0's in their most significant bits. The only regular JTAG registers a debugger might use are BYPASS and IDCODE, but this specification leaves IR space for many other standard JTAG instructions. Unimplemented instructions must select the BYPASS register.

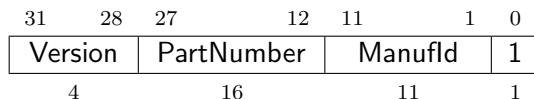
Table 6.1: JTAG DTM TAP Registers

Address	Name	Description	Page
0x00	BYPASS	JTAG recommends this encoding	57 58
0x01	IDCODE	JTAG recommends this encoding	
0x10	DTM Control and Status	For Debugging	
0x11	Debug Module Interface Access	For Debugging	
0x12	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x13	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x14	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x15	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x16	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x17	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x1f	BYPASS	JTAG requires this encoding	

### 6.1.3 IDCODE (at 0x01)

This register is selected (in IR) when the TAP state machine is reset. Its definition is exactly as defined in IEEE Std 1149.1-2013.

This entire register is read-only.



Field	Description	Access	Reset
Version	Identifies the release version of this part.	R	Preset
PartNumber	Identifies the designer's part number of this part.	R	Preset
Manufld	Identifies the designer/manufacturer of this part. Bits 6:0 must be bits 6:0 of the designer/manufacturer's Identification Code as assigned by JEDEC Standard JEP106. Bits 10:7 contain the modulo-16 count of the number of continuation characters (0x7f) in that same Identification Code.	R	Preset

### 6.1.4 DTM Control and Status (dtmcs, at 0x10)

The size of this register will remain constant in future versions so that a debugger can always determine the version of the DTM.

31	18	17	16	15	14	12	11	10	9	4	3	0
0	dmihardreset	dmireset	0	idle	dmistat	abits	version					
14	1	1	1	3	2	6	4					

Field	Description	Access	Reset
dmihardreset	Writing 1 to this bit does a hard reset of the DTM, causing the DTM to forget about any outstanding DMI transactions. In general this should only be used when the Debugger has reason to expect that the outstanding DMI transaction will never complete (e.g. a reset condition caused an inflight DMI transaction to be cancelled).	W1	0
dmireset	Writing 1 to this bit clears the sticky error state and allows the DTM to retry or complete the previous transaction.	W1	0
idle	This is a hint to the debugger of the minimum number of cycles a debugger should spend in Run-Test/Idle after every DMI scan to avoid a ‘busy’ return code (dmistat of 3). A debugger must still check dmistat when necessary. 0: It is not necessary to enter Run-Test/Idle at all. 1: Enter Run-Test/Idle and leave it immediately. 2: Enter Run-Test/Idle and stay there for 1 cycle before leaving. And so on.	R	Preset
dmistat	0: No error. 1: Reserved. Interpret the same as 2. 2: An operation failed (resulted in op of 2). 3: An operation was attempted while a DMI access was still in progress (resulted in op of 3).	R	0
abits	The size of address in dmi.	R	Preset
version	0: Version described in spec version 0.11. 1: Version described in spec version 0.13 (and later?), which reduces the DMI data width to 32 bits. 15: Version not described in any available version of this spec.	R	1

### 6.1.5 Debug Module Interface Access (dmi, at 0x11)

This register allows access to the Debug Module Interface (DMI).

In Update-DR, the DTM starts the operation specified in **op** unless the current status reported in **op** is sticky.

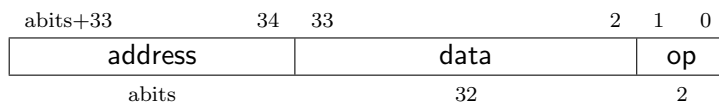
In Capture-DR, the DTM updates **data** with the result from that operation, updating **op** if the current **op** isn't sticky.

See Section B.1 and Table B.1 for examples of how this is used.

---

*The still-in-progress status is sticky to accommodate debuggers that batch together a number of scans, which must all be executed or stop as soon as there's a problem.*

*For instance a series of scans may write a Debug Program and execute it. If one of the writes fails but the execution continues, then the Debug Program may hang or have other unexpected side effects.*



Field	Description	Access	Reset
address	Address used for DMI access. In Update-DR this value is used to access the DM over the DMI.	R/W	0
data	The data to send to the DM over the DMI during Update-DR, and the data returned from the DM as a result of the previous operation.	R/W	0

*Continued on next page*



Field	Description	Access	Reset
op	<p>When the debugger writes this field, it has the following meaning:</p> <p>0: Ignore <a href="#">data</a> and <a href="#">address</a>. (nop)</p> <p>Don't send anything over the DMI during Update-DR. This operation should never result in a busy or error response. The address and data reported in the following Capture-DR are undefined.</p> <p>1: Read from <a href="#">address</a>. (read)</p> <p>2: Write <a href="#">data</a> to <a href="#">address</a>. (write)</p> <p>3: Reserved.</p> <p>When the debugger reads this field, it means the following:</p> <p>0: The previous operation completed successfully.</p> <p>1: Reserved.</p> <p>2: A previous operation failed. The data scanned into <a href="#">dmi</a> in this access will be ignored. This status is sticky and can be cleared by writing <a href="#">dmireset</a> in <a href="#">dtmcs</a>.</p> <p>This indicates that the DM itself responded with an error. Note: there are no specified cases in which the DM would respond with an error, and DMI is not required to support returning errors.</p> <p>3: An operation was attempted while a DMI request is still in progress. The data scanned into <a href="#">dmi</a> in this access will be ignored. This status is sticky and can be cleared by writing <a href="#">dmireset</a> in <a href="#">dtmcs</a>. If a debugger sees this status, it needs to give the target more TCK edges between Update-DR and Capture-DR. The simplest way to do that is to add extra transitions in Run-Test/Idle.</p> <p>(The DTM, DM, and/or component may be in different clock domains, so synchronization may be required. Some relatively fixed number of TCK ticks may be needed for the request to reach the DM, complete, and for the response to be synchronized back into the TCK domain.)</p>	R/W	2

### 6.1.6 BYPASS (at 0x1f)

1-bit register that has no effect. It is used when a debugger does not want to communicate with this TAP.

This entire register is read-only.

0  
0  
 1

### 6.1.7 Recommended JTAG Connector

To make it easy to acquire debug hardware, this spec recommends a connector that is compatible with the Atmel AVR JTAG Connector, as described below.

The connector is a .05”-spaced, gold-plated male header with .016” thick hardened copper or beryllium bronze square posts (SAMTEC FTSH-105 or equivalent). Female connectors are compatible 20 $\mu$ m gold connectors.

Viewing the male header from above (the pins pointing at your eye), a target’s connector looks as it does in Table 6.5. The function of each pin is described in Table 6.6.

Table 6.5: JTAG Connector Diagram

TCK	1	2	GND
TDO	3	4	VCC
TMS	5	6	(SRST <sub>n</sub> )
(NC)	7	8	(TRST <sub>n</sub> )
TDI	9	10	GND

Target connectors may be shrouded. In that case the key slot should be next to pin 5. Female headers should have a matching key.

Debug adapters should be tagged or marked with their isolation voltage threshold (i.e. unisolated, 250V, etc.).

All debug adapter pins other than GND should be current-limited to 20mA.

Table 6.6: JTAG Connector Pinout

1	TCK	JTAG TCK signal, driven by the debug adapter. This pin must be clearly marked in both male and female headers.
5	TMS	JTAG TMS signal, driven by debug adapter.
9	TDI	JTAG TDI signal, driven by the debug adapter.
3	TDO	JTAG TDO signal, driven by the target.
8	TRSTn	Test Reset (optional, only used by some devices. Used to reset the JTAG TAP Controller).
4	VCC	Reference voltage for logic high. A debug adapter may attempt to draw up to 20mA from this pin to power itself, but a target is not obligated to provide that power.
2, 10	GND	Target ground.
6	SRSTn	Active-low reset signal, driven by the debug adapter. Asserting reset should reset any RISC-V cores as well as any other peripherals on the PCB. It should not reset the debug logic. Although connecting this pin is optional, it is recommended as it allows the debugger to hold the target device in a reset state, which may be essential to debug some scenarios. If not implemented in a target, this pin must not be connected.



# Appendix A

## Hardware Implementations

Below are two possible implementations. A designer could choose one, mix and match, or come up with their own design.

### A.1 Abstract Command Based

Halting happens by stalling the hart execution pipeline.

Muxes on the register file(s) allow for accessing GPRs and CSRs using the Access Register abstract command.

Memory is accessed using the Abstract Access Memory command or through System Bus Access.

### A.2 Execution Based

This implementation only implements the Access Register abstract command for GPRs on a halted hart, and relies on the Program Buffer for all other operations.

This method uses the hart’s existing pipeline and ability to execute from arbitrary memory locations to avoid modifications to a hart’s datapath. When the halt request bit is set, the Debug Module raises a special interrupt to the selected hart(s). This interrupt causes each hart to enter Debug Mode and jump to a defined memory region that is serviced by the DM. When taking this exception, `pc` is saved to `dpc` and `cause` is updated in `dcsr`.

The code in the Debug Module causes the hart to execute a “park loop”. In the park loop the hart writes its `mhartid` to a memory location within the Debug Module to indicate that it is halted. To allow the DM to individually control one out of several halted harts, each hart polls for flags in a DM-controlled memory location to determine whether the debugger wants it to execute the Program Buffer or perform a resume.

To execute an abstract command, the DM first populates some internal words of program buffer

according to `command`. When `transfer` is set, the DM populates these words with `lw <gpr>, 0x400(zero)` or `sw 0x400(zero), <gpr>`. 64- and 128-bit accesses use `ld/sd` and `lq/sq` respectively. If `transfer` is not set, the DM populates these instructions as `nops`. If `execute` is set, execution continues to the debugger-controlled Program Buffer, otherwise the DM causes a `ebreak` to execute immediately.

When `ebreak` is executed (indicating the end of the Program Buffer code) the hart returns to its park loop. If an exception is encountered, the hart jumps to a defined debug exception address within the Debug Module. The code at that address causes the hart to write to an address in the Debug Module which indicates exception. This address is considered I/O for `fence` instructions (see #10 on page 37). Then the hart jumps back to the park loop. The DM infers from the write that there was an exception, and sets `cmderr` appropriately.

To resume execution, the debug module sets a flag which causes the hart to execute a `dret`. When `dret` is executed, `pc` is restored from `dpc` and normal execution resumes at the privilege set by `prv`.

`data0` etc. are mapped into regular memory at an address relative to `zero` with only a 12-bit `imm`. The exact address is an implementation detail that a debugger must not rely on. For example, the `data` registers might be mapped to `0x400`.

For additional flexibility, `progbuf0`, etc. are mapped into regular memory immediately preceding `data0`, in order to form a contiguous region of memory which can be used for either program execution or data transfer.

## Appendix B

# Debugger Implementation

This section details how an external debugger might use the described debug interface to perform some common operations on RISC-V cores using the JTAG DTM described in Appendix ???. All these examples assume a 32-bit core but it should be easy to adapt the examples to 64- or 128-bit cores.

To keep the examples readable, they all assume that everything succeeds, and that they complete faster than the debugger can perform the next access. This will be the case in a typical JTAG setup. However, the debugger must always check the sticky error status bits after performing a sequence of actions. If it sees any that are set, then it should attempt the same actions again, possibly while adding in some delay, or explicit checks for status bits.

### B.1 Debug Module Interface Access

To read an arbitrary Debug Module register, select `dmi`, and scan in a value with `op` set to 1, and `address` set to the desired register address. In Update-DR the operation will start, and in Capture-DR its results will be captured into `data`. If the operation didn't complete in time, `op` will be 3 and the value in `data` must be ignored. The busy condition must be cleared by writing `dmireset` in `dtmcs`, and then the second scan must be performed again. This process must be repeated until `op` returns 0. In later operations the debugger should allow for more time between Capture-DR and Update-DR.

To write an arbitrary Debug Bus register, select `dmi`, and scan in a value with `op` set to 2, and `address` and `data` set to the desired register address and data respectively. From then on everything happens exactly as with a read, except that a write is performed instead of the read.

It should almost never be necessary to scan IR, avoiding a big part of the inefficiency in typical JTAG use.

## B.2 Checking for Halted Harts

A user will want to know as quickly as possible when a hart is halted (eg. due to a breakpoint). To efficiently determine which harts are halted when there are many harts, the debugger uses the `haltsum` registers. Assuming the maximum number of harts exist, first it checks `haltsum3`. For each bit set there, it writes `hartsel`, and checks `haltsum2`. This process repeats through `haltsum1` and `haltsum0`. Depending on how many harts exist, the process should start at one of the lower `haltsum` registers.

## B.3 Halting

To halt one or more harts, the debugger selects them, sets `haltreq`, and then waits for `allhalted` to indicate the harts are halted before clearing `haltreq` to 0.

## B.4 Running

First, the debugger should restore any registers that it has overwritten. Then it can let the selected harts run by setting `resumereq`. Once `allresumeack` is set, the debugger knows the hart has resumed, and it can clear `resumereq`. Note that harts might halt very quickly after resuming (e.g. by hitting a software breakpoint) so the debugger cannot use `allhalted/anyhalted` to check whether the hart resumed.

## B.5 Single Step

Using the hardware single step feature is almost the same as regular running. The debugger just sets `step` in `dcsr` before letting the hart run. The hart behaves exactly as in the running case, except that interrupts may be disabled (depending on `stepie`) and it only fetches and executes a single instruction before re-entering Debug Mode.

## B.6 Accessing Registers

### B.6.1 Using Abstract Command

Read `s0` using abstract command:

Op	Address	Value	Comment
Write	<code>command</code>	<code>aarsize = 2, transfer, 0x1008</code>	Read <code>s0</code>
Read	<code>data0</code>	-	Returns value that was in <code>s0</code>

Write `mstatus` using abstract command:



Op	Address	Value	Comment
Write	<code>data0</code>	new value	
Write	<code>command</code>	<code>aarsize = 2, transfer, write, 0x300</code>	Write <code>mstatus</code>

### B.6.2 Using Program Buffer

Abstract commands are used to exchange data with GPRs. Using this mechanism, other registers can be accessed by moving their value into/out of GPRs.

Write `mstatus` using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>csrw s0, MSTATUS</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>data0</code>	new value	
Write	<code>command</code>	<code>aarsize = 2, postexec, transfer, write, 0x1008</code>	Write <code>s0</code> , then execute program buffer

Read `f1` using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>fmv.x.s s0, f1</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>command</code>	<code>postexec</code>	Execute program buffer
Write	<code>command</code>	<code>transfer 0x1008</code>	read <code>s0</code>
Read	<code>data0</code>	-	Returns the value that was in <code>f1</code>

## B.7 Reading Memory

### B.7.1 Using System Bus Access

With system bus access, addresses are physical system bus addresses.

Read a word from memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbc</code>	<code>sbaccess = 2, sbreadonaddr</code>	Setup
Write	<code>sbaddress0</code>	address	
Read	<code>sbdata0</code>	-	Value read from memory

Read block of memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbcs</code>	<code>sbaccess = 2, sbreadonaddr, sbreadondata, sbautoincrement</code>	Turn on autoread and autoincrement
Write	<code>sbaddress0</code>	address	Writing address triggers read and increment
Read	<code>sbddata0</code>	-	Value read from memory
Read	<code>sbddata0</code>	-	Next value read from memory
...	...	...	...
Write	<code>sbcs</code>	0	Disable autoread
Read	<code>sbddata0</code>	-	Get last value read from memory.

### B.7.2 Using Program Buffer

Through the Program Buffer, the hart performs the memory accesses. Addresses are physical or virtual (depending on `mprven` and other system configuration).

Read a word from memory using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>lw s0, 0(s0)</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>data0</code>	address	
Write	<code>command</code>	<code>write, postexec, 0x1008</code>	Write <code>s0</code> , then execute program buffer
Write	<code>command</code>	<code>0x1008</code>	Read <code>s0</code>
Read	<code>data0</code>	-	Value read from memory

Read block of memory using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>lw s1, 0(s0)</code>	
Write	<code>progbuf1</code>	<code>addi s0, s0, 4</code>	
Write	<code>progbuf2</code>	<code>ebreak</code>	
Write	<code>data0</code>	address	
Write	<code>command</code>	<code>write, postexec, 0x1008</code>	Write <code>s0</code> , then execute program buffer
Write	<code>command</code>	<code>postexec, 0x1009</code>	Read <code>s1</code> , then execute program buffer
Write	<code>abstractauto</code>	<code>autoexecdata [0]</code>	Set <code>autoexecdata [0]</code>
Read	<code>data0</code>	-	Get value read from memory, then execute program buffer
Read	<code>data0</code>	-	Get next value read from memory, then execute program buffer
...	...	...	...
Write	<code>abstractauto</code>	0	Clear <code>autoexecdata [0]</code>
Read	<code>data0</code>	-	Get last value read from memory.

TODO: Table B.1 shows the scans involved in reading a single word using this method.

Table B.1: Memory Read Timeline

	JTAG State	Activity
TODO	TODO	TODO

### B.7.3 Using Abstract Memory Access

Abstract memory accesses act as if they are performed by the hart, although the actual implementation may differ.

Read a word from memory using abstract memory access:

Op	Address	Value	Comment
Write	<code>data1</code>	address	
Write	<code>command</code>	<code>cmdtype=2</code> , <code>aamsize =2</code>	
Read	<code>data0</code>	-	Value read from memory

Read block of memory using abstract memory access:

Op	Address	Value	Comment
Write	<code>abstractauto</code>	1	Re-execute the command when <code>data0</code> is accessed
Write	<code>data1</code>	address	
Write	<code>command</code>	<code>cmdtype=2</code> , <code>aamsize =2</code> , <code>aampostincrement =1</code>	
Read	<code>data0</code>	-	Read value, and trigger reading of next address
...	...	...	...
Write	<code>abstractauto</code>	0	Disable auto-exec
Read	<code>data0</code>	-	Get last value read from memory.

## B.8 Writing Memory

### B.8.1 Using System Bus Access

With system bus access, addresses are physical system bus addresses.

Write a word to memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbaddress0</code>	address	
Write	<code>sbddata0</code>	value	

Write block of memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbc</code>	<code>sbaccess = 2, sbautoincrement</code>	Turn on autoincrement
Write	<code>sbaddress0</code>	address	
Write	<code>sbddata0</code>	value0	
Write	<code>sbddata0</code>	value1	
...	...	...	...
Write	<code>sbddata0</code>	valueN	

### B.8.2 Using Program Buffer

Through the Program Buffer, the hart performs the memory accesses. Addresses are physical or virtual (depending on `mprven` and other system configuration).

Write a word to memory using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>sw s1, 0(s0)</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>data0</code>	value	
Write	<code>command</code>	<code>write, 0x1008</code>	Write <code>s0</code>
Write	<code>data0</code>	address	
Write	<code>command</code>	<code>write, postexec, 0x1009</code>	Write <code>s1</code> , then execute program buffer

Write block of memory using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>sw s1, 0(s0)</code>	
Write	<code>progbuf1</code>	<code>addi s0, s0, 4</code>	
Write	<code>progbuf2</code>	<code>ebreak</code>	
Write	<code>data0</code>	address	
Write	<code>command</code>	<code>write, 0x1008</code>	Write <code>s0</code>
Write	<code>data0</code>	value0	
Write	<code>command</code>	<code>write, postexec, 0x1009</code>	Write <code>s1</code> , then execute program buffer
Write	<code>abstractauto</code>	<code>autoexecdata [0]</code>	Set <code>autoexecdata [0]</code>
Write	<code>data0</code>	value1	
...	...	...	...
Write	<code>data0</code>	valueN	
Write	<code>abstractauto</code>	0	Clear <code>autoexecdata [0]</code>

### B.8.3 Using Abstract Memory Access

Abstract memory accesses act as if they are performed by the hart, although the actual implementation may differ.

Write a word to memory using abstract memory access:

Op	Address	Value	Comment
Write	<b>data1</b>	address	
Write	<b>data0</b>	value	
Write	<b>command</b>	cmdtype=2, <b>aamsize</b> =2, write=1	

Write block of memory using abstract memory access:

Op	Address	Value	Comment
Write	<b>data1</b>	address	
Write	<b>data0</b>	value0	
Write	<b>command</b>	cmdtype=2, <b>aamsize</b> =2, write=1, <b>aampostincrement</b> =1	
Write	<b>abstractauto</b>	1	Re-execute the command when <b>data0</b> is accessed
Write	<b>data0</b>	value1	
Write	<b>data0</b>	value2	
...	...	...	...
Write	<b>data0</b>	valueN	
Write	<b>abstractauto</b>	0	Disable auto-exec

## B.9 Triggers

A debugger can use hardware triggers to halt a hart when a certain event occurs. Below are some examples, but as there is no requirement on the number of features of the triggers implemented by a hart, these examples may not be applicable to all implementations. When a debugger wants to set a trigger, it writes the desired configuration, and then reads back to see if that configuration is supported.

Enter Debug Mode just before the instruction at 0x80001234 is executed, to be used as an instruction breakpoint in ROM:

<b>tdata1</b>	0x105c	action=1, match=0, m=1, s=1, u=1, execute=1
<b>tdata2</b>	0x80001234	address

Enter Debug Mode right after the value at 0x80007f80 is read:

<b>tdata1</b>	0x4159	timing=1, action=1, match=0, m=1, s=1, u=1, load=1
<b>tdata2</b>	0x80007f80	address

Enter Debug Mode right before a write to an address between 0x80007c80 and 0x80007cef (inclusive):

<b>tdata1</b> 0	0x195a	action=1, chain=1, match=2, m=1, s=1, u=1, store=1
<b>tdata2</b> 0	0x80007c80	start address (inclusive)
<b>tdata1</b> 1	0x115a	action=1, match=2, m=1, s=1, u=1, store=1
<b>tdata2</b> 1	0x80007cf0	end address (exclusive)

Enter Debug Mode right before a write to an address between 0x81230000 and 0x8123ffff (inclusive):

<b>tdata1</b>	0x10da	action=1, match=1, m=1, s=1, u=1, store=1
<b>tdata2</b>	0x81237fff	16 bits to match exactly, then 0, then all ones.

Enter Debug Mode right after a read from an address between 0x86753090 and 0x8675309f or between 0x96753090 and 0x9675309f (inclusive):

<b>tdata1</b> 0	0x41a59	timing=1, action=1, chain=1, match=4, m=1, s=1, u=1, load=1
<b>tdata2</b> 0	0xffff03090	Mask for low half, then match for low half
<b>tdata1</b> 1	0x412d9	timing=1, action=1, match=5, m=1, s=1, u=1, load=1
<b>tdata2</b> 1	0x7fff8675	Mask for high half, then match for high half

## B.10 Handling Exceptions

Generally the debugger can avoid exceptions by being careful with the programs it writes. Sometimes they are unavoidable though, eg. if the user asks to access memory or a CSR that is not implemented. A typical debugger will not know enough about the platform to know what's going to happen, and must attempt the access to determine the outcome.

When an exception occurs while executing the Program Buffer, **cmderr** becomes set. The debugger can check this field to see whether a program encountered an exception. If there was an exception, it's left to the debugger to know what must have caused it.

## B.11 Quick Access

There are a variety of instructions to transfer data between GPRs and the **data** registers. They are either loads/stores or CSR reads/writes. The specific addresses also vary. This is all specified in **hartinfo**. The examples here use the pseudo-op **transfer dest, src** to represent all these options.

Halt the hart for a minimum amount of time to perform a single memory write:

Op	Address	Value	Comment
Write	<a href="#">progbuf0</a>	transfer arg2, s0	Save s0
Write	<a href="#">progbuf1</a>	transfer s0, arg0	Read first argument (address)
Write	<a href="#">progbuf2</a>	transfer arg0, s1	Save s1
Write	<a href="#">progbuf3</a>	transfer s1, arg1	Read second argument (data)
Write	<a href="#">progbuf4</a>	sw s1, 0(s0)	
Write	<a href="#">progbuf5</a>	transfer s1, arg0	Restore s1
Write	<a href="#">progbuf6</a>	transfer s0, arg2	Restore s0
Write	<a href="#">progbuf7</a>	ebreak	
Write	<a href="#">data0</a>	address	
Write	<a href="#">data1</a>	data	
Write	<a href="#">command</a>	0x10000000	Perform quick access

This shows an example of setting the [m](#) bit in [mcontrol](#) to enable a hardware breakpoint in M mode. Similar quick access instructions could have been used previously to configure the trigger that is being enabled here:

Op	Address	Value	Comment
Write	<a href="#">progbuf0</a>	transfer arg0, s0	Save s0
Write	<a href="#">progbuf1</a>	li s0, (1 << 6)	Form the mask for <a href="#">m</a> bit
Write	<a href="#">progbuf2</a>	csrrs x0, <a href="#">tdata1</a> , s0	Apply the mask to <a href="#">mcontrol</a>
Write	<a href="#">progbuf3</a>	transfer s0, arg2	Restore s0
Write	<a href="#">progbuf4</a>	ebreak	
Write	<a href="#">command</a>	0x10000000	Perform quick access

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# Appendix C

## Change Log

Revision	Date	Author(s)	Description
38fde94	2018-07-09	Tim Newsome	datacount cannot be 0 (#286)
800ca8d	2018-07-06	Tim Newsome	Clarifications requested by Jeremy Bennett (#280)
b363afa	2018-07-06	Tim Newsome	Add missing .tex file to dependencies. (#302)
93340e4	2018-07-06	Tim Newsome	Clarify that trigger registers are WARL. (#306)
95af58a	2018-07-06	Tim Newsome	Force the register-address in place. (#304)
d83039d	2018-07-06	Tim Newsome	\Fcause priority numbers: higher means higher (#307)
921c6a3	2018-07-03	Tim Newsome	Completing progbuf exec is I/O for fence insts. (#305)
99e01fa	2018-06-27	Tim Newsome	Add target-specific bits to abstract access memory. (#295)
4a0152d	2018-06-19	Tim Newsome	Only write busy to \Fcmderr if \Fcmderr is 0. (#296)
b0dc615	2018-06-16	Tim Newsome	Rebuild the PDF.
90873eb	2018-06-16	Tim Newsome	Fix typo in abstract access memory examples. (#297)
5fe8e08	2018-06-16	Tim Newsome	dret is a section, not a subsection of reset (#294)
abfd8a0	2018-06-14	Tim Newsome	Revert "Only write busy to \Fcmderr if \Fcmderr is 0."
7c66968	2018-06-14	Tim Newsome	Only write busy to \Fcmderr if \Fcmderr is 0.
0f28f27	2018-06-08	Tim Newsome	Abstract memory (#283)
7c840dd	2018-06-08	Tim Newsome	Specify an Exception Trigger (#266)
9d0d8af	2018-06-06	Tim Newsome	Clarify what address space these registers are in (#281)
a7f293d	2018-06-03	Tim Newsome	Add missing dependency to Makefile (#285)
37893aa	2018-05-30	Tim Newsome	Make trigger types writable. (#279)
6730cc0	2018-05-29	Tim Newsome	Explain priority assignment rationale. (#277)
b6d5d66	2018-05-25	Tim Newsome	Prevent M mode triggers affecting D mode ones (#282)
08ee84f	2018-05-22	Tim Newsome	Reading tselect doesn't guarantee a valid trigger. (#271)

6dfe375	2018-04-18	Megan Wachs	Debug Module should be capitalized
dac2120	2018-04-11	Megan Wachs	resethaltreq: Proposal for forcing a hart into debug mode out of reset
3b6442f	2018-05-16	Tim Newsome	tdata2 need only hold valid addresses if select=0 (#278)
68501cb	2018-04-26	mwachs5	mprven: Add a bit to enable MPRV to take effect in debug mode
9fcabe0	2018-05-03	Megan Wachs	Appendix: correct and clarify what debugger vs DM does
30773fd	2018-05-03	Tim Newsome	Debuggers must not write sbcs while sbbusy is set (#270)
50d8cd8	2018-05-03	Megan Wachs	Remove merge commits from the changelog
3b7a296	2018-05-02	Tim Newsome	Fix typo.
b26072b	2018-05-02	Tim Newsome	Explain that 1 in hart array mask means selected
41f6026	2018-05-02	Megan Wachs	Examples: Give an example of CSR access with Quick Access (#268)
675bb14	2018-05-01	Tim Newsome	Replace XLEN with MXLEN. #257
848cca1	2018-04-30	Megan Wachs	Overview Diagram: increase number of Progbuf words (#267)
a719ee6	2018-04-25	Megan Wachs	fix misspelled name
097c701	2018-04-23	Tim Newsome	Fix typo.
01dabd5	2018-04-23	Tim Newsome	Incorporate review feedback.
ca7a9d0	2018-04-18	Tim Newsome	Add trigger examples for match types 1, 4, and 5
cd5a15c	2018-04-16	Tim Newsome	Give a few trigger examples.
4375927	2018-04-12	Tim Newsome	Clarify that maskmax applies only to NAPOT trigger
acadfe9	2018-04-13	Megan Wachs	NMI: debugging may not be possible if an NMI happens
8fb190c	2018-04-12	Tim Newsome	Another attempt at SBA errors.
714c5d1	2018-04-11	Megan Wachs	Core Debug: all interrupts are masked includes NMI
56fbd9d	2018-04-11	Megan Wachs	DCSR: add nmip bit to indicate NMI is pending
ffe3c2	2018-04-10	Tim Newsome	Clarify SBA unsupport access size error.
b4006ac	2018-04-10	Tim Newsome	Clarify high bits of sbdata in narrow reads.
4ca83dd	2018-03-28	Tim Newsome	Clarify progbuf=1 some more
3b62243	2018-03-26	Tim Newsome	Clarify debugger requirements when progbufsize=1
ffba4d0	2018-03-26	Tim Newsome	Explain why progbufsize=1 is special
6b88905	2018-03-19	Megan Wachs	haltsum1: correct its address to be BWC and not overlap with ABSTRACTAUTO
2382e2e	2018-03-06	Megan Wachs	Correct some inaccuracies in the chisel generated files
3e88e11	2018-03-06	Megan Wachs	travis: add 'make chisel' target to regression
32cbb9b	2018-03-19	Tim Newsome	Nonexistent/unavailable harts are not halted.
f8a7bb7	2018-03-19	Tim Newsome	More clarification.
e21ae4c	2018-03-16	Tim Newsome	Allow any bit in hart array mask to be tied to 0
efb7e45	2018-03-15	Tim Newsome	Change dcsr.prv reset value to 3
f19946b	2018-03-15	Tim Newsome	Clarify hart array mask register size.
ddec145	2018-03-14	Tim Newsome	Be more precise about core vs hart
4e5f4ad	2018-03-14	Tim Newsome	Review feedback.
8ac9273	2018-03-14	Tim Newsome	Be more precise about processor vs hart
83c9774	2018-03-14	Tim Newsome	Clarify abstract command errors.

4ebc177	2018-03-14	Tim Newsome	hawindowssel can be smaller, depends on # of harts
11e1b5c	2018-03-14	Tim Newsome	Split future ideas section into a notes doc
bafeeaa	2018-03-13	Tim Newsome	Rebuild PDF
6a85d53	2018-03-13	Tim Newsome	Incorporate review feedback.
f213315	2018-03-09	Tim Newsome	Clarify user responsibilities when debugging lr/sc
3641305	2018-03-09	Tim Newsome	Remove implemented features from Future Ideas.
1135bf3	2018-03-06	Tim Newsome	Incorporate feedback.
8f35e7e	2018-03-05	Megan Wachs	gt_1024: Clarify that some registers may not be present for small numbers of harts
683ae37	2018-02-14	Megan Wachs	hartsum-;haltsum
ee51758	2018-02-14	Megan Wachs	Modification of ; 1024 hart proposal that maintains backwards compatibility
370d222	2018-03-05	Tim Newsome	Rephrase description of hit bit.
eee5e0c	2018-03-05	Tim Newsome	Clarify multiple DMs/harts
4d5acef	2018-02-28	Tim Newsome	Clarify what happens when \Fauthenticated is clear
6a0c9ec	2018-02-27	Tim Newsome	Move hit bit per review feedback.
097bd8e	2018-02-21	Tim Newsome	Fix link to pre-built pdf
d21774b	2018-02-21	Omer Faruk IR-MAK	Python interpreter to be used should default to Python2
a8c10cf	2018-02-20	Tim Newsome	Incorporate review feedback.
a0f947c	2018-02-20	Tim Newsome	Make trigger hit bit optional.
77e4634	2018-02-08	Tim Newsome	Add hit bit to hardware triggers.
140390a	2018-02-05	Tim Newsome	Better wording.
e35b1ff	2018-02-05	Tim Newsome	Move Reg Access Abbrev table after sample register
e887433	2018-02-05	Tim Newsome	Use longtable instead of xtabular.
5c84437	2018-01-31	Tim Newsome	Abstract Command data usage depends on the command
3d508ea	2018-01-25	Tim Newsome	HARTSELBITS-;HARTSELLEN and other feedback
eb653f7	2018-01-24	Tim Newsome	Be explicit about the size of \Fhartsel.
822bd81	2018-01-24	Tim Newsome	Revert incrementing version number.
4c755af	2018-01-24	Tim Newsome	\Fsbbusyerror also inhibits new accesses.
457413d	2018-01-24	Tim Newsome	Update how to enumerate all harts.
2180801	2018-01-18	Tim Newsome	Fix ambiguity in busy error reporting.
3140efa	2018-01-09	Tim Newsome	Re-apply e698a5001aa4583d31dde484d78f4f10e4e3148f
390daa7	2018-01-18	mwachs5	. No need to list out all the consecutive registers. sbaddress: Only writes to address will actually cause an error. Reads while busy are permitted.
5c820f3	2018-01-18	Megan Wachs	Remove reference to "caches"
4533648	2018-01-18	Megan Wachs	correct access spelling
d37c1ac	2018-01-16	Tim Newsome	Fix table column overruns by going full manual
e9100ea	2018-01-16	Tim Newsome	Correct when sbbusy error is set for being busy.
c029cc7	2018-01-16	Tim Newsome	Complete partial sentence.
494338a	2018-01-15	Tim Newsome	Add clarifications about error handling.
e14c34e	2018-01-15	Tim Newsome	Incorporate review feedback.
68720e5	2018-01-15	Tim Newsome	Remove H bits from triggers.
b8eb62a	2018-01-15	Tim Newsome	Clarify when sbaccess is checked for validity
8b50d29	2018-01-12	Tim Newsome	Add \Fsbbusy, to avoid race clearing \Fserror

50b1b41	2018-01-12	Tim Newsome	Clarify: writes to \Rsbdata0 write the new data
7f26759	2018-01-12	Tim Newsome	Clarify exactly which bits are used for SB access.
47a019c	2018-01-11	Tim Newsome	Fix typo.
a49d6ad	2018-01-11	Tim Newsome	sbreadonaddr is R/W
42195c2	2018-01-11	Tim Newsome	Fix cut-and-paste error.
6c95235	2018-01-11	Tim Newsome	Add sbaddress3, for future proofing.
e3345ea	2018-01-11	Tim Newsome	Incorporate review feedback.
6da48f8	2018-01-11	Tim Newsome	Remove dmerr.
e99c092	2018-01-10	Tim Newsome	Add system bus version field.
a6aa531	2018-01-10	Tim Newsome	Talk about all data and progbuf regs in first reg
af272db	2018-01-09	Megan Wachs	Update dret font
3d579d8	2018-01-09	Tim Newsome	Explicitly list data[1-10] and progbuf[1-15]
c6481ae	2018-01-09	Tim Newsome	Revert "Explicitly list data[1-10] and progbuf[1-15]"
e698a50	2018-01-09	Tim Newsome	Explicitly list data[1-10] and progbuf[1-15]
e547ed5	2018-01-09	Tim Newsome	Clarify that we deal in physical addresses only.
b377b89	2018-01-09	Tim Newsome	Revert "Clarify that we deal in physical addresses only."
f7da066	2018-01-09	Tim Newsome	Clarify that we deal in physical addresses only.
99a1599	2018-01-09	Tim Newsome	Clarify that \Fdatasize contains at most 12.
ae6e88a	2018-01-09	mwachs5	dret: Legal only in Debug Mode
18f392d	2017-11-24	Tim Newsome	Get rid of sbsingleread in favor of sbreadonaddr
5754a3b	2018-01-05	Megan Wachs	Use a different word than "clobbered"
aca7e0b	2018-01-03	Megan Wachs	Add missing "to"s to abstractauto description
d59ddf3	2018-01-03	Megan Wachs	Correct plurality of halted harts in haltsum
57c53ed	2017-12-22	Tim Newsome	Put parens around all macros that need it.
7ded846	2017-12-18	Tim Newsome	Refer to existing hart instead of "valid"
68b8ac8	2017-12-15	Tim Newsome	Make \Fhaltsel WARL.
6a72f45	2017-12-18	Tim Newsome	Mark this as a draft, which it is.
dd8d871	2017-12-18	Tim Newsome	Properly deal with \ chars in the changelog.
42f920c	2017-12-18	Tim Newsome	Deal with \ chars in the changelog.
b13891c	2017-12-15	Tim Newsome	Revert "Make \Fhaltsel WARL."
26d76a0	2017-12-15	Tim Newsome	Make \Fhaltsel WARL.
afda8d7	2017-11-28	mwachs5	update PDF
134d310	2017-11-28	Megan Wachs	Correct compressed version of ebreak
caa1258	2017-11-27	Megan Wachs	badaddr -i tval (Priv Spec 1.9 -i 1.9.1)
32b0f08	2017-11-22	Tim Newsome	Incorporate feedback.
2f7aa54	2017-11-22	Tim Newsome	Simplify, and explain trigger behavior.
3e5887f	2017-11-21	Tim Newsome	Clarify some single step corner cases.
f4b9ae2	2017-11-21	Tim Newsome	Make ackhavereset write-only. (#178)
efe3dc8	2017-11-21	Tim Newsome	Make hartreset R/W (#177)
ce1b359	2017-11-17	Megan Wachs	Reset clarifications (#172)
852a70d	2017-11-16	Megan Wachs	icount: remove warning (#173)
363348f	2017-11-16	Tim Newsome	Explain cache coherency wrt to system bus access (#171)
26ea898	2017-11-15	Tim Newsome	Refer to ISA and priv docs.
ffc8c62	2017-11-03	Tim Newsome	Mention the index in "about this doc"
a4257ef	2017-11-02	Tim Newsome	Add an index to the document.
f5f45a5	2017-10-30	Megan Wachs	Add 'has reset' status and control (#168)

46f3f54	2017-10-25	Tim Newsome	Incorporate review feedback.
104247f	2017-10-24	Megan Wachs	Update README.md
6dd5c80	2017-10-24	Megan Wachs	Update README.md
cb1a847	2017-10-24	Megan Wachs	Add a note to the README about the built PDF
e00625f	2017-10-18	Tim Newsome	Include pdf.
c23e729	2017-10-18	Tim Newsome	Clarify more.
83f9faf	2017-10-11	Tim Newsome	Clarify what \Fimpebreak does.
78082b5	2017-10-11	Tim Newsome	Mention \Fimpebreak in Program Buffer description.
0378324	2017-10-11	mwachs5	Add legend and update some transitions on the Abstract Command State Machine diagram
fa2b600	2017-10-11	Megan Wachs	add missing period
0610630	2017-10-11	Megan Wachs	Just do simple hmode -i dmode replacement
16e11f3	2017-10-11	Tim Newsome	Remove hmode reference, to fix build.
84b9a6a	2017-10-11	Tim Newsome	Add \Fimpebreak, to support of implicit ebreak.
cc90b77	2017-10-11	mwachs5	Remove reference to 'H' mode from the figure
cc6a9de	2017-10-11	Megan Wachs	Change old reference to 'hmode' to 'dmode'
ea2877d	2017-10-10	Tim Newsome	Move how-to-debug into the relevant section.
486ecc6	2017-10-05	Tim Newsome	Refuse unsupported bus accesses.
6ca221d	2017-10-05	Tim Newsome	haltreq, resumereq, hartreset are per-hart bits
d4118ab	2017-09-30	Tim Newsome	ndmreset can't reset logic required to access DM.
c6bd8d1	2017-09-29	Tim Newsome	and -i or
58c2441	2017-09-29	Tim Newsome	Mention \Fstepie in Single Step
94c5f78	2017-09-29	Tim Newsome	Clarify ndmreset.
12810b4	2017-09-29	Tim Newsome	Clarify that sbaddress is physical.
5862fdf	2017-09-29	Tim Newsome	Unify M mode and mprv comment.
aea1bd5	2017-09-29	Tim Newsome	Define behavior when haltreq and resumereq are set
146b348	2017-09-28	Megan Wachs	remove superfluous 'an'
a5d16c4	2017-09-28	Megan Wachs	remove superfluous 'a'
052a8ab	2017-09-28	Tim Newsome	Clarify that a debugger can lose hart control.
cc52cff	2017-09-28	Tim Newsome	Add \Fdmerr.
25685eb	2017-09-28	Tim Newsome	Explain that bus master or progbuf is required.
f75ee7d	2017-09-28	Tim Newsome	Clarify debugger can discover "almost" everything
71e6788	2017-09-27	Tim Newsome	Remove description of manual stepping.
9aea347	2017-09-27	Tim Newsome	Move Running/Single Step near Halting.
2090d9b	2017-09-27	Tim Newsome	data0 should be sbdata0 in this table.
5858cfe	2017-09-27	Tim Newsome	Clarify why \Rpriv exists.
bc3c2aa	2017-09-27	Tim Newsome	Mention where priv encoding comes from.
ef77cc4	2017-09-27	Tim Newsome	One more attempt to clarify DPC after single step.
80a288e	2017-09-27	Tim Newsome	Clarify instret not incrementing on ebreak.
c163d22	2017-09-20	Tim Newsome	Remove ebreakh.
9971075	2017-09-20	Tim Newsome	Clarify we're talking about privilege
3fbe495	2017-09-20	Tim Newsome	Clarify that we're talking about *implementation*
3684854	2017-09-20	Tim Newsome	Use steps environment in sbdata0.
d4eda18	2017-09-20	Tim Newsome	Explain that only sbdata0 has side effects.
ae781c6	2017-09-20	Tim Newsome	Don't refer to internal system bus registers.
875922e	2017-09-20	Tim Newsome	Explain sbdata0 being stale a bit more.
cd44fd5	2017-09-20	Tim Newsome	Clarify autoread
194484b	2017-09-20	Tim Newsome	Clarify hawindow.

02f1aac	2017-09-20	Tim Newsome	Clarify that \Fdataaddr is relative to \Rzero.
0e9b6ae	2017-09-20	Tim Newsome	Clarify nonexistent vs unavailable.
b55ff41	2017-09-20	Tim Newsome	Fix devtreevalid.
2eccb86	2017-09-20	Tim Newsome	Explicitly state which registers are read-only.
4af505c	2017-09-20	Tim Newsome	Show section numbers for registers.
cbd5573	2017-09-20	Tim Newsome	Thank Nikhil
19c206f	2017-09-20	Tim Newsome	Clarify how to determine whether progbuf is RAM
0651f7d	2017-09-20	Tim Newsome	Explain what happens if ebreak is missing.
e889dae	2017-09-20	Tim Newsome	Move figure of states into its own section.
cff7b80	2017-09-20	Tim Newsome	Explain when \Ftransfer might be used.
6b2ee61	2017-09-20	Tim Newsome	Explain where \Fsize encoding came from.
c9f3b73	2017-09-14	Tim Newsome	Fix typo.
4b25400	2017-09-13	Tim Newsome	Mention dpc in CSRs abstract register numbers.
c3ee426	2017-09-13	Tim Newsome	Move abstract regno table closer to its reference.
111b9a3	2017-09-13	Tim Newsome	cycle -i operation
994afdc	2017-09-13	Tim Newsome	Account for multiple selected harts.
aa4a297	2017-09-13	Tim Newsome	Halt Control -i Run Control
e97c821	2017-09-13	Tim Newsome	continuous -i contiguous
97f73ff	2017-09-13	Tim Newsome	Clarify ndmreset behavior.
6078220	2017-09-13	Tim Newsome	Explain ndmreset
a3d4f30	2017-09-13	Tim Newsome	Describe ‘halt region’
272b3d9	2017-09-13	Tim Newsome	Clarify accessing unimplemented DM DMI regs
3e91f1b	2017-09-13	Tim Newsome	Clarify either Prog Buf or Sys Bus Acc is required
e8a6145	2017-09-13	Tim Newsome	Clarify CSR access; remove serial port
ce20766	2017-09-13	Tim Newsome	Remove section referencing itself.
1195a61	2017-09-18	Tim Newsome	Generate constants to be unsigned for clang.
8967b0a	2017-08-16	Megan Wachs	Compressed instructions are c.foo, not foo.c
b5698a9	2017-08-16	Megan Wachs	clarify progbufsize description
d221bab	2017-08-16	Megan Wachs	Remove progbufsize enums from register description
0498102	2017-08-16	Megan Wachs	appendix: Use standard assembly format for sw
4456d99	2017-08-09	Tim Newsome	Rename progsiz to progbufsize.
55d5b66	2017-08-09	Tim Newsome	Clarify that trigger comparisons are unsigned.
21e35ef	2017-08-09	Tim Newsome	Configuration String -i Device Tree
f044f45	2017-08-02	Tim Newsome	Don’t require a target to provide 25mA on VCC.
c883943	2017-08-02	Tim Newsome	Add table of Abstract Command Types
985a3df	2017-08-02	Tim Newsome	Fix and speed up build.
95b9108	2017-08-02	mwachs5	DTM: Clarify that there are no cases when DMI would actually return an error.
9c9e0c0	2017-08-02	mwachs5	SystemBus: No longer returns error. So DMI has no ‘error’ return code.
5ba18f9	2017-07-27	Tim Newsome	Fix more typos.
dbc65bf	2017-07-26	Tim Newsome	Fix typos.
bba0ad9	2017-07-26	Tim Newsome	Tighten up introduction lists.
e22d5eb	2017-07-26	Tim Newsome	Add version constants for ”not compatible”.
c79038e	2017-07-26	Tim Newsome	Small clarification.
9df0411	2017-07-21	Tim Newsome	Incorporate review feedback.
d67419c	2017-07-21	Tim Newsome	Clarify dpc contents.
9f50c05	2017-07-11	Tim Newsome	Use LL instead of L for 64-bit constant suffix.



23fd24a	2017-07-10	Megan Wachs	Cleaning up whitespaces
c5ab04c	2017-07-10	Megan Wachs	Update abstract_commands.xml
6e8cdf1	2017-07-10	Megan Wachs	Update abstract_commands.xml
cf6e3f2	2017-07-10	Megan Wachs	clarify DCSR.cause
79ffbb9	2017-07-10	Megan Wachs	Clarify implications of CSR read, write, halt
013e191	2017-07-10	Megan Wachs	Clarify when you would get error halt/resume
231e457	2017-07-10	Megan Wachs	Quick Access error clarification
c54c2f2	2017-07-03	mwachs5	serial: add the XML file, not the TEX file
ac77477	2017-07-03	mwachs5	serial: Fix compile errors after moving serial port to appendix
6defcb8	2017-07-03	mwachs5	serial: Move serial ports out of main spec and into Future Work appendix
a28f639	2017-06-30	mwachs5	remove trace dependencies from Makefile
52a122b	2017-06-30	mwachs5	remove trace section
d9e166b	2017-06-30	mwachs5	remove trace registers
7caf4e5	2017-06-30	mwachs5	remove trace appendix
4688988	2017-06-29	mwachs5	DCSR: define a 'stepie' bit which may be hard-wired to 0.
9a0492c	2017-06-13	Megan Wachs	Add missing period and some other small text edits
13ccdbf	2017-06-13	Megan Wachs	fix typo in ProgBuf register macro
b01f989	2017-06-13	mwachs5	implementations: be a bit more concrete about the one example implementation we have.
a7b5f83	2017-06-13	mwachs5	jtagdtm: Move it out of the appendix as it is really part of the specification
87aceb0	2017-06-13	Megan Wachs	remove "spontaneous"
50b9950	2017-06-13	Megan Wachs	Forward reference for anynonexistent
adea3e2	2017-06-13	Megan Wachs	More clarifications on dret
1b8dd0e	2017-06-13	Megan Wachs	Define DRET instruction
617da4c	2017-06-08	Megan Wachs	Update description of R/W1C
de2c56b	2017-06-08	Megan Wachs	Clarify that DCSR is also not updated on ebreak
efa615d	2017-06-07	Tim Newsome	Increase xdebugver field size to 4 bits. (#92)
a0e147a	2017-06-07	Tim Newsome	Address some review comments.
89ffe50	2017-06-06	mwachs5	NDMRESET: Clarify what it may and may not do
1932da0	2017-06-06	mwachs5	DPC: Clarifications on its meaning
6470fdb	2017-06-06	mwachs5	ABSTRACTCS: Correct inconsistency on the number of data words.
3ca82b4	2017-06-06	Megan Wachs	More corrections for R vs R/W1C on SERCS
9705fb8	2017-06-06	Megan Wachs	Correct a bunch of W0 registers
1347371	2017-06-05	Tim Newsome	Add intdisable to dcsr.
989c60d	2017-06-05	Tim Newsome	Fix language. We can only halt harts, not cores.
517a08b	2017-06-05	Tim Newsome	Incorporate review feedback.
802be28	2017-06-05	Tim Newsome	Clarify/fix Quick Access example.
b8cc523	2017-06-02	Tim Newsome	Add included tex files as dependencies. (#78)
15f864a	2017-06-01	Tim Newsome	Language cleanups, consistency and typo fixes.
4ecae86	2017-06-01	Tim Newsome	Add page numbers to list-of-register tables.
59b3e4a	2017-05-19	Megan Wachs	Setting up a Travis regression to check for build errors (#72)

124bf44	2017-05-17	mwachs5	Debug Module: CMDERR is Write-1-to clear, not R/W0
bb6c7f0	2017-05-17	mwachs5	SW Registers file should be XML, not TEX
d360358	2017-05-10	Megan Wachs (Temporary Acct.)	Remove virtual register from core_registers.xml
bfc64fb	2017-05-10	Megan Wachs (Temporary Acct.)	Add missing sw_registers.tex file
0512f5d	2017-05-06	mwachs5	Move virtual 'prv' register to a separate section to make it more clear it is not a real register.
6b3c9d7	2017-05-06	mwachs5	Clarify haltreq/resumereq/resumack
0a487eb	2017-04-26	mwachs5	jtag: Change specified JTAG pinout from Coretex to AVR, to provide for TRSTn option.
93cdfaf	2017-04-26	mwachs5	DM : Clarify that DATA/PROGBUF can't be written while busy.
ef98f23	2017-04-19	mwachs5	jtag: Make it clear that a NOP is really a NOP.
a6f8efa	2017-04-17	mwachs5	single_step: Exceptions count as the 'step' completion.
bf11e9e	2017-04-17	mwachs5	resumeack: fix some LaTeX cross references
4afa081	2017-04-11	mwachs5	halt/resumereq: Clarify what setting them to 0 or 1 does
297a39b	2017-04-06	mwachs5	fix chisel build
082c499	2017-04-06	mwachs5	Rename resumed to resumeack, and add more text about what these bits mean.
909d617	2017-04-06	mwachs5	Correct some cross references after removing all the multiply listed registers
dd09914	2017-04-06	mwachs5	Add 'resumedall' and 'resumedany' bits to avoid race condition on about to resume and just halted
feb88fc	2017-04-05	mwachs5	JTAG DTM: Clarify that leading bits are 0 for more than 5-bit IR
75b96ea	2017-04-04	mwachs5	use renamed dm_registers file
9f3ec7e	2017-04-04	mwachs5	debugger_implementation: remove some old TODO and commentary.
45dd5b5	2017-04-04	mwachs5	Don't list out every single DM register for those that are just indexed versions
b8b3aa2	2017-04-04	mwachs5	remove core-side register definitions from Debug Module. Rename dm1 to dm
d979a13	2017-04-04	mwachs5	remove core-side serial port specification, as these should look like implementation-specific devices with appropriate drivers.
b56870b	2017-04-04	mwachs5	Remove the wording about 'debug exception', as it is called breakpoint exception in the RISC-V Spec.
1e9347d	2017-04-03	mwachs5	Add description of hasel
0dda84d	2017-04-03	mwachs5	JTAG DTM: Clean up TAP register descriptions
82ccde5	2017-04-03	mwachs5	JTAG DTM: Add a hard DMI bit which cancels the outstanding DMI transaction
bd2a3d1	2017-04-03	mwachs5	remove preexec

02c733a	2017-04-03	mwachs5	remove preexec from Abstract State diagram.
1e271d6	2017-04-03	mwachs5	Update Debugger implementation for DMI register access, and fix tex compile issues.
155dda4	2017-04-03	mwachs5	Rewrite HW Implementation examples to describe a pure abstract command approach, and to not rely on harts executing every instruction which is fetched from the Debug Module
556c2be	2017-04-03	mwachs5	minor wording edits about RISC-V core registers
523c64a	2017-04-03	mwachs5	Edits to the Debug Module section.
b9a371f	2017-04-03	mwachs5	add missing trace.tex file.
58b2396	2017-04-03	mwachs5	Re-order the JTAG DTM Sections
a8827e2	2017-04-03	mwachs5	Edits to the System Overview.
c5417ce	2017-04-03	mwachs5	add more sections as separate files.
287d5c6	2017-04-03	mwachs5	moving more files to separate tex files.
9e873f4	2017-04-03	mwachs5	move trigger info into separate file.
2c89a86	2017-04-03	mwachs5	move risc-v core debug info into separate file.
e676491	2017-04-03	mwachs5	Move System Overview to separate file
03df6ee	2017-04-03	mwachs5	Move Debug Module description to a separate file.
5faa430	2017-04-03	mwachs5	add back in JTAG DTM in appendix
7b28b11	2017-04-03	mwachs5	Move jtag DTM to appendix. Move some text to commentary.
cc183ba	2017-04-03	mwachs5	move introduction to a separate file. Comment out reading order.
f727d14	2017-04-03	mwachs5	Use Chapters vs Sections. Needs reorganization.
815951d	2017-04-03	mwachs5	Formatting updates. Make this look more like the RISC-V specs. Need to use chapter vs. section
69ffaf8	2017-03-31	mwachs5	Move XML files into a subdirectory.
b276384	2017-03-31	mwachs5	Remove debug_rom.S
112bbac	2017-03-31	mwachs5	figures: reorganize the figures into directories.
1e5c068	2017-03-27	Megan Wachs	Add LICENSE
fc17730	2017-03-22	Po-wei Huang	Change some halt mode into debug mode.
8ccf029	2017-03-22	Po-wei Huang	All halt mode changed to debug mode to synchronize with the priv spec.
f143d9e	2017-03-21	mwachs5	Correct duplicated progbuf register names
0797ec1	2017-03-17	mwachs5	autoexec: make autoexec bits match the number of data words there really are.
8e76d93	2017-03-17	mwachs5	dm1_registers: move a few more things around. Reduce abstract data words back to 12.
f8bf292	2017-03-17	mwachs5	dm1_registers: resolve some address conflicts and inconsistencies
a74dff9	2017-03-17	mwachs5	access_register: some small bit changes
2e6b0ca	2017-03-15	mwachs5	config string: Fix LaTeX compile errors.
f83260a	2017-03-10	mwachs5	Abstract Commands: clarify that 32-bit reads should always work. This allows reading MISA.
6f9347a	2017-03-10	mwachs5	Config String: change the Abstract Command to DMI registers. Allow the same registers to be used for unspecified identifier information.

4ea10ff	2017-03-10	mwachs5	abstract: Make autoexec apply to all data and progbuf words. Make a seperate register which is optional.
5008436	2017-03-10	mwachs5	abstract: Allow up to 16 progbuf and/or data words. Inform debugger about dscratch registers available for its use.
aaa13e5	2017-03-06	mwachs5	Command: use the name 'cmdtype' not 'type' to allow easier auto-generation of Scala code.
e9bb72c	2017-03-06	mwachs5	Hart Array: Add registers for hart array.
5d17a35	2017-03-06	mwachs5	DM: Move addresses around for better seperation of functionalities in HW
25ccaa8	2017-03-06	mwachs5	CONTROL: Rename control and status registers to ___CS for consistency and to accurately reflect their functionality.
45cf6c2	2017-03-06	mwachs5	Errors: fix up the bit assignments in SERSTATUS with the addition of error bit.
38cb5a0	2017-03-06	mwachs5	Errors: Make errors write-1-to-clear.
b436d77	2017-03-03	mwachs5	triggers: Clarify that matches are against virtual addresses.
793bb85	2017-03-03	mwachs5	triggers: Add suggested timings for best user experience.
2669866	2017-03-03	mwachs5	stoptime/stopcycle: Make their functionality match their name. Allow any reset value.
c85a1cf	2017-03-01	mwachs5	config_string: Simplify the Config String Address abstract command.
a303a6b	2017-03-02	Megan Wachs	Update README.md
92a4923	2017-03-01	mwachs5	serial: tweak addresses.
b09f460	2017-03-01	mwachs5	serial: tweak addresses.
6477837	2017-03-01	mwachs5	chisel: tweaks to class names.
be83e3e	2017-02-28	Tim Newsome	Clarify stoptime, stopcycle.
c17c17c	2017-02-27	Tim Newsome	Abstract command that returns config string addr.
096dfbc	2017-02-27	Tim Newsome	Acknowledge Alex.
c0253ab	2017-02-24	Tim Newsome	Explain tdata1 type a bit more.
e43ac2e	2017-02-24	Tim Newsome	Clarify how to enumerate triggers again.
c6e3e20	2017-02-23	Tim Newsome	Revert previous commit.
ef770bf	2017-02-23	Tim Newsome	mcontrol and icount mask tdata2, not tdata1.
27806f2	2017-02-23	mwachs5	rename 'type' to 'cmdtype' purely so my auto-generation scripts work.
e46798d	2017-02-22	mwachs5	Add Abstract Commands to automatic chisel
b3bb939	2017-02-21	mwachs5	Generate Chisel headers as well for Debug Module.
c9db98c	2017-02-22	Tim Newsome	Simplify description of op statuses.
bda39cc	2017-02-22	mwachs5	Add explicit type field to Abstract Command.
f83a1ca	2017-02-22	mwachs5	Finish up replacement of ibuf-¿progbuf
9666e51	2017-02-22	mwachs5	IBUF-¿PROGBUF
5308ecd	2017-02-22	mwachs5	Remove last references to "Instruction Supply"
f6ebde9	2017-02-22	Tim Newsome	Move authentication to a serial protocol.
0f079c8	2017-02-22	Tim Newsome	Reserve bit for per-hart reset.
f2c93ac	2017-02-22	Tim Newsome	Clarify that dmactive resets authentication.

f5e7b1c	2017-02-22	Alex Bradbury	Clarify that the halt state of all harts is maintained through reset
3dfe8fd	2017-02-22	Tim Newsome	More Debug Mode -> Halt Mode.
d29fc1f	2017-02-22	Tim Newsome	Debug Mode -> Halt Mode
55d6030	2017-02-21	Tim Newsome	Generate debug_defines.h as part of normal make
b0e6a7f	2017-02-21	Tim Newsome	Minor clarifications.
0f9885c	2017-02-20	Tim Newsome	Various clarifications.
0802d5a	2017-02-15	mwachs5	Use consistent 'Control and Status' naming for CS registers.
5accc7d	2017-02-15	Tim Newsome	Change all the "other" JTAG IRs to just reserved.
bcbd7da	2017-02-15	mwachs5	sm_diagram: Show using resumereq bit to resume.
18f6e55	2017-02-14	Tim Newsome	Introduce resumereq command, similar to haltreq.
4b62c40	2017-02-14	mwachs5	SystemBus: Clean up some formatting and error specification notes.
bc97723	2017-02-14	mwachs5	quick-access: Update SM Diagram for Quick Access
d27066e	2017-02-14	Tim Newsome	Clarify haltreq bit.
6f8ec43	2017-02-14	Tim Newsome	Always generate long constants when required.
c6ac6bc	2017-02-13	Tim Newsome	Include field descriptions in C header file.
b849213	2017-02-13	Tim Newsome	Fix the build.
1cf8033	2017-02-12	mwachs5	jtag: More clarifications
6203bd6	2017-02-12	Megan Wachs	Update requirements- W GPRs Required
f2b43a7	2017-02-12	Megan Wachs	Remove double 'the'
2c64ef1	2017-02-12	Megan Wachs	Remove comma
f84abce	2017-02-12	Megan Wachs	Whitespace edits and address come comments
23c2648	2017-02-11	mwachs5	jtag-dtm: ask for clarification on TAP sharing.
7020d23	2017-02-11	mwachs5	jtag-dtm: Clarifications, DBUS->DMI
292d49c	2017-02-11	Megan Wachs	fix indentation
b879b86	2017-02-11	Megan Wachs	Add missing period
bbe0521	2017-02-11	mwachs5	Make comments on program buffer size match the address map.
4ceaa37	2017-02-11	mwachs5	Flesh out and edit the introduction/background Add a description of use cases this spec has in mind, and what it doesn't cover.
cbf89d6	2017-02-11	Tim Newsome	Rewrite Quick Access.
170bff1	2017-02-10	Megan Wachs	Allow size 4 for the program buffer
c911e6e	2017-02-10	Tim Newsome	Clarify use of dmactive.
2ca296f	2017-02-09	Tim Newsome	Reserve command register space for custom use.
e49666e	2017-02-09	Tim Newsome	Clarify hart index change per Megan's comments.
84865e9	2017-02-09	Tim Newsome	Add header prefix for abstract commands.
2434f4f	2017-02-09	Tim Newsome	Select harts by index instead of hart ID.
7bf112a	2017-02-09	Tim Newsome	Generate correct headers for <32-bit registers.
7f0f09a	2017-02-08	Tim Newsome	Reset dbus status to "failure" to avoid confusion.
8b1c6f0	2017-02-08	Megan Wachs	Fix line wrap issue
345c33f	2017-02-08	Megan Wachs	Call out "arg0" specifically.
9f080f5	2017-02-08	Megan Wachs	Clarify "arguments" to commands
259badd	2017-02-08	Tim Newsome	Make haltsum/halt registers mandatory.
eb0f1d3	2017-02-07	Tim Newsome	Allow for early abstract command failures.
bb49bd1	2017-02-07	Tim Newsome	Clarify error handling a little.

3fc0a97	2017-02-07	Tim Newsome	Explain when abstract data regs may be clobbered.
c37167e	2017-02-07	Tim Newsome	Fix old language in description of halt registers.
6943c96	2017-02-07	Tim Newsome	Generate more useful C header files from reg defs
98639df	2017-02-05	mwachs5	Include the SM Diagram as a figure. Also some minor capitalization fixes.
a95e4c3	2017-02-05	mwachs5	Update State Machine diagram to show uncertainty of halt bit during auto halt/resume.
ba76744	2017-02-05	Tim Newsome	Combine loabits and hiabits.
02b1d92	2017-02-05	Tim Newsome	DMI can get away with just 6 address bits.
35d6e33	2017-02-05	mwachs5	Update State machine diagram to show BUSY without HALTED
f511b05	2017-02-04	Tim Newsome	Clarify command busy bit.
d0f8961	2017-02-03	mwachs5	Update figures
e18a68d	2017-02-03	Tim Newsome	Clarify prehalt/postresume failure.
ac3e2a9	2017-02-02	Tim Newsome	Clarify abstract command failure behavior.
ce4baee	2017-02-02	Tim Newsome	Add Quick Access section.
0490377	2017-02-02	Tim Newsome	Add prehalt and postresume to reg command.
67515bd	2017-02-02	Tim Newsome	Deal with a few minor TODOs.
96456fc	2017-02-02	Tim Newsome	Turn register names into links.
317cd98	2017-02-02	Tim Newsome	Explain what register access is required.
f3ad2f2	2017-02-01	Tim Newsome	Revert Plain Exception implementation to be simple
a0ad281	2017-02-01	Tim Newsome	execb -i, preexec, execa -i, postexec
1d4a2c3	2017-02-01	Tim Newsome	Limit Program Buffer sizes to 0, 1, 8.
cc40815	2017-02-01	Tim Newsome	Incorporate Po-wei's feedback.
c8b45d6	2017-02-01	Tim Newsome	Clarify how all autoexec bits work.
dbb1deb	2017-02-01	Tim Newsome	Remove stale TODO.
c5f8f59	2017-02-01	Tim Newsome	Explain why cmderr inhibits starting new commands.
5c69194	2017-02-01	Tim Newsome	Fix editing error.
50f7c48	2017-02-01	Tim Newsome	Remove empty hart info register.
781c68e	2017-02-01	Megan Wachs	Update README.md
f46b32e	2017-02-01	mwachs5	Add a diagram of Abstract Command flow.
633bd63	2017-02-01	Tim Newsome	Move Reading Order into About This Document
51ec4d1	2017-02-01	Tim Newsome	Add reading order section.
03d20ad	2017-02-01	Tim Newsome	autoexec0 applies to data0, not inst0.
c302353	2017-01-31	Tim Newsome	Don't rely on hart fetching instructions once.
2558c25	2017-01-31	Tim Newsome	Change how exceptions in Halt Mode are handled.
a36ddce	2017-01-31	Tim Newsome	Add size to abstract register command.
64de458	2017-01-31	Tim Newsome	Detail bus master reads.
c08486f	2017-01-31	Megan Wachs	reset: Add some comments (#5)
1558049	2017-01-30	Tim Newsome	Automate Change Log.
51525a4	2017-01-29	Tim Newsome	Update System Overview
7d39ac0	2017-01-29	Tim Newsome	Update Supported Features.
9e7cbea	2017-01-29	Tim Newsome	Update RISC-V Core section.
515188d	2017-01-29	Tim Newsome	Update Hardware Implementations section.
4b19ed8	2017-01-29	mwachs5	system.bus: be consistent and always call it 'System Bus'. Even if some dislike the name, we should be consistent and clear in the spec.
9ccefc3d	2017-01-29	Tim Newsome	Fleshed out some debugger implementation.

04b9176	2017-01-28	Tim Newsome	Rename debug exception to breakpoint exception.
5ac4ea1	2017-01-27	Tim Newsome	WIP on big update on instruction supply.
2d9c3e2	2017-01-27	Tim Newsome	Reorganize dm registers.
de50ba8	2017-01-27	Tim Newsome	Abstract command support is already addressed.
5085046	2017-01-26	mwachs5	Rename registers and fields like 'access' that were confusingly the same name.
10bbf6f	2017-01-26	Tim Newsome	Fix #2: DM address space table
a05c582	2017-01-26	Tim Newsome	Add debugger inspection as a feature.
4062681	2017-01-24	Tim Newsome	Add publish target.
5c8bb83	2017-01-24	Tim Newsome	Clarify use of data registers.
1504da6	2017-01-24	Tim Newsome	Replace manual date with automatic git hash/date.
997f2a0	2017-01-23	Tim Newsome	Deal with unsupported abstract commands.
cb6f2b8	2017-01-23	Tim Newsome	Renumber registers to prevent duplicates.
8b4db96	2017-01-23	Tim Newsome	Don't print out addresses if they're not provided.
b00cd21	2017-01-23	Tim Newsome	Add an abstract command.
675b556	2017-01-23	Tim Newsome	Reorganize DM bits into functional group regs.
5fc7512	2017-01-23	Tim Newsome	Remove bits 33:32 from sbdata[23].
ceb5d66	2017-01-20	Tim Newsome	Starting point for a comprehensive spec