Operating systems and concurrency B04

David Kendall

Northumbria University

Introduction

- Towards an operating system: CMSIS and Mbed
- Multi-tasking operating system services
- μ C/OS-II (uC/OS-II)
- Task management in uC/OS-II

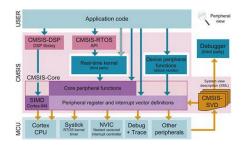
Standard names for the microcontroller registers

```
#include <stdbool b>
#include < stdint.h>
#include <LPC407x 8x 177x 8x.h>
#define LED1PIN
                   (1UL << 18)
void delay(uint32_t ms);
int main() {
  LPC IOCON->P1 18 &= ~0x1FUL;
  LPC GPIO1->DIR |= LED1PIN;
  while (true) {
    LPC GPIO1->SET = LED1PIN;
    delay(1000):
    LPC GPIO1->CLR = LED1PIN;
    delay (1000);
```

- Header file
 LPC407x_8x_177x_8x.h
 provided by the
 manufacturer of the
 microcontroller (NXP) to
 give a set of standard
 names for the
 microcontroller registers
- No need to write #defines after looking up the addresses yourself.
- Header file written in standard ANSI C -CMSIS-complaint portable between tool providers.

CMSIS

- Cortex Microcontroller Software Interface Standard
- CMSIS-Core
 - Standard set of names provided by ARM for accessing the microprocessor
 - Standard startup files for bringing up the CPU out of reset:
 - configure system clocks
 - define the vector table
 - run the main function
- Improves software portability and reusability across different microcontrollers and toolchains

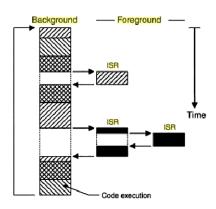


Martin, T. The Designer's Guide to the Cortex-M Processor

Family: A Tutorial Approach, Newnes, 2013

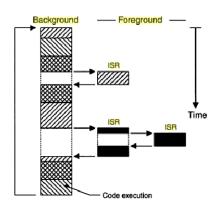
Foreground/background tasks

- Simple multitasking
- Super loop calls functions for computation (background)
- Interrupt service routines (ISRs) handle asynchronous events (foreground)



Foreground/background tasks

- Simple multitasking
- Super loop calls functions for computation (background)
- Interrupt service routines (ISRs) handle asynchronous events (foreground)



Problem

Model with only one computation task is not flexible enough

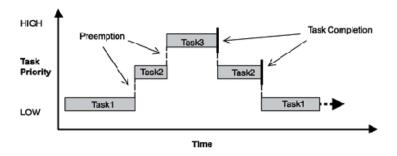
Multitasking

- Easier to structure the application as a set of concurrent tasks rather than as a single program or as foreground/background tasks: each task is responsible for some well-defined part of the system's overall behaviour
- But only the illusion of concurrency the OS switches quickly between tasks, executing some instructions from one task before moving on to another task
- switching from one text to another requires a context switch
- deciding which task to switch to requires a scheduling algorithm

Scheduling

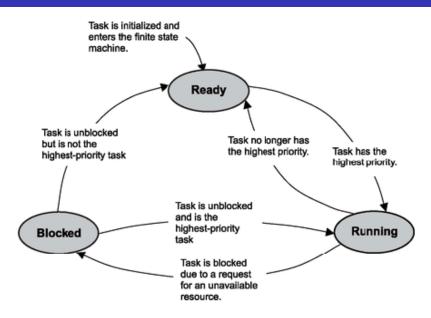
- Deciding when one task should stop executing and which one should begin next is a scheduling problem
- Two main approaches to scheduling:
 - Preemptive scheduling
 - Task is forced to yield the CPU
 - Round robin
 - Priority-based
 - Non-preemptive (cooperative) scheduling
 - Task voluntarily yields the CPU and signals the next task to begin

Fixed-priority preemptive scheduling



• We focus on fixed-priority premptive scheduling

A task and its execution states



uC/OS-II: A small operating system

- Main features:
 - Multi-tasking
 - Preemptive
- Other features:
 - Predictable
 - Robust and reliable
 - Standards-compliant
 - Portable
 - Scalable
 - Source code available

uC/OS-II Services

- Task management
- Delay management

uC/OS-II Services

- Task management
- Delay management
- Semaphores
- Mutual exclusion semaphores
- Event flags
- Message mailboxes
- Message queues
- Memory management
- Timers
- Miscellaneous

uC/OS-II Services

- Task management
- Delay management
- Semaphores
- Mutual exclusion semaphores
- Event flags
- Message mailboxes
- Message queues
- Memory management
- Timers
- Miscellaneous

See uC/OS-II Quick Reference

Tasks behaviour

- The behaviour of a task is defined by a C function that:
 - never terminates
 - blocks repeatedly

Example of task behaviour definition

```
static void appTaskLED2(void *pdata) {
   while (true) {
     OSTimeDlyHMSM(0,0,0,500);
     gpioPinToggle(&pin[LED2]);
   }
}
```

Tasks: other requirements

Priority

- Used for fixed-priority pre-emptive scheduling
- a number between 0 and OS_LOWEST_PRIO
- low number ⇒ high priority
- high number ⇒ low priority
- OS reserves priorities 0 to 3 and OS_LOWEST_PRIO 3 to OS_LOWEST_PRIO
- Advice: define an enumeration of task priority constants, starting at priority level 4.
- Example

```
enum {
   APP_TASK_BUTTONS_PRIO = 4,
   APP_TASK_LED1_PRIO,
   APP_TASK_LED2_PRIO
}:
```

Tasks: other requirements

Stack

- Each task needs its own data area (stack) for storing
 - context
 - local variables
- Example stack definition

```
enum {APP_TASK_LED2_STK_SIZE = 256};
static OS_STK appTaskLED2Stk[APP_TASK_LED2_STK_SIZE];
```

Tasks: other requirements

Stack

- Each task needs its own data area (stack) for storing
 - context
 - local variables
- Example stack definition

```
enum {APP_TASK_LED2_STK_SIZE = 256};
static OS_STK appTaskLED2Stk[APP_TASK_LED2_STK_SIZE];
```

User data

- Optionally tasks can be given access to user data when they are created
- We will not use this feature in this module
- Advice: always specify this as (void *) 0 when creating a task

Task creation

A task is created using the OS function

```
INT8U OSTaskCreate(
    void (*task)(void *pdata), /* function for the task */
    void *pdata, /* user data for function */
    OS_STK *ptos, /* pointer to top of stack */
    INT8U priority /* task priority */
);
```

Task creation

A task is created using the OS function

```
INT8U OSTaskCreate(
    void (*task)(void *pdata), /* function for the task */
    void *pdata, /* user data for function */
    OS_STK *ptos, /* pointer to top of stack */
    INT8U priority /* task priority */
);
```

Example

Task delay

- Often, a task will block itself by explicitly asking the OS to delay it for some period of time
- void OSTimeDly(INT16U ticks);
- Causes a context switch if ticks is between 1 and 65535
- If ticks is 0, OSTimeDly () returns immediately to caller
- On context switch uC/OS-II executes the next highest priority task
- Task that called OSTimeDly() will be made ready to run when the specified number of ticks elapses - actually runs when it becomes the highest priority ready task
- Resolution of the delay is between 0 and 1 tick
- Another task can cancel the delay by calling OSTimeDlyResume()

Task delay

- OSTimeDly () specifies delay in terms of a number of ticks
- Use OSTimeDlyHMSM() to specify delay in terms of Hours, Minutes, Seconds and Milliseconds
- Otherwise OSTimeDlyHMSM() behaves as OSTimeDly()

Complete example

```
#include <stdbool b>
#include <ucos ii.h>
#include "gpioPin.h"
typedef enum
 APP TASK_LED1_PRIO = 4,
 APP TASK LED2 PRIO
} taskPriorities t:
typedef enum {
 APP TASK LED1 STK SIZE = 256.
 APP TASK LED2 STK SIZE = 256
} stackSizes t:
static OS STK appTaskLED1Stk[APP TASK LED1 STK SIZE];
static OS STK appTaskLED2Stk[APP TASK LED2 STK SIZE];
static void appTaskLED1(void *pdata);
static void appTaskLED2(void *pdata);
typedef enum {
LED1 = 0,
LFD2
} deviceNames t:
```

gpioPin t pin[2];

Complete example

```
int main() {
 /* Initialise the GPIO pins */
 gpioPinInit(&pin[LED1], 1, 18, OUTPUT PIN);
 gpioPinInit(&pin[LED2], 0, 13, OUTPUT PIN);
 /* Initialise the OS */
 OSInit();
 /* Create the tasks */
 OSTaskCreate (appTaskLED1,
               (void *)0,
               (OS STK *)&appTaskLED1Stk[APP TASK LED1 STK SIZE - 1],
               APP TASK LED1 PRIO);
 OSTaskCreate(appTaskLED2.
               (void *)0,
               (OS STK *)&appTaskLED2Stk[APP TASK LED2 STK SIZE - 1],
               APP TASK LED2 PRIO):
 /* Start the OS */
 OSStart():
 /* Should never arrive here */
 return 0:
```

Complete example

```
static void appTaskLED1(void *pdata) {
  /* Start the OS ticker — must be done in the highest priority task */
  SysTick_Config(SystemCoreClock / OS_TICKS_PER_SEC);

/* Task main loop */
  while (true) {
     gpioPinToggle(&pin[LED1]);
     OSTimeDlyHMSM(0,0,0,500);
  }
}

static void appTaskLED2(void *pdata) {
  while (true) {
     gpioPinToggle(&pin[LED2]);
     OSTimeDlyHMSM(0,0,0,500);
  }
}
```

Acknowledgements

- Qing Li and Caroline Yao, Real-time concepts for embedded systems, CMP, 2003
- Jean Labrosse, MicroC/OS-II: The Real-Time Kernel, CMP, 2002