

Chapter 4: Graphs

Graph theory

See Alberto Montresor theory here: <http://disi.unitn.it/~montreso/sp/slides/06-grafi.pdf> (<http://disi.unitn.it/~montreso/sp/slides/06-grafi.pdf>)

See [Graphs on the book](https://interactivepython.org/runestone/static/pythonds/Graphs/toctree.html) (<https://interactivepython.org/runestone/static/pythonds/Graphs/toctree.html>)

In particular, see :

- [Vocabulary and definitions](https://interactivepython.org/runestone/static/pythonds/Graphs/VocabularyandDefinitions.html)
(<https://interactivepython.org/runestone/static/pythonds/Graphs/VocabularyandDefinitions.html>)

To keep it short, a graph is a set of vertices linked by edges.

Directed graphs

In this worksheet we are going to use so called Directed Graphs (DiGraph for brevity), that is graphs that have *directed* edges: each edge can be pictured as an arrow linking source node *a* to target node *b*. With such an arrow, you can go from *a* to *b* but you cannot go from *b* to *a* unless there is another edge in the reverse direction.

- A DiGraph for us can also have no edges or no vertices at all.
- A vertex for us can be anything, a string like 'abc', the number 3, etc
- In our model, edges simply link vertices and have no weights
- The DiGraph is represented as an adjacency list, mapping each vertex to the vertices it is linked to.

QUESTION: DiGraph model is thus good for dense or sparse graphs?

Serious graphs

In this worksheet we follow the Do-It-Yourself methodology and create graph classes from scratch for didactical purposes. Of course, in Python world you have already nice libraries entirely devoted to graphs like [networkx](https://networkx.github.io/) (<https://networkx.github.io/>), you can also use them for visualizing graphs. If you have huge graphs to process you might consider big data tools like [Spark GraphX](http://spark.apache.org/graphx/) (<http://spark.apache.org/graphx/>) which is programmable in Python.

0) Code skeleton

First off, download [the Python skeleton \(graphs.py\)](#) to modify. Solutions are in a [separate file \(graphs_solution.py\)](#).

1) Building graphs

IMPORTANT: All the functions until `1.8 has_edge()` excluded are already provided and you don't need to implement them !

1.1) Building basics

Let's look at the constructor `__init__` and `add_vertex`. They are already provided and you don't need to implement it:

```

class DiGraph:
    def __init__(self):
        # The class just holds the dictionary _edges: as keys it has the vertices, and
        # to each vertex associates a list with the vertices it is linked to.

        self._edges = {}

    def add_vertex(self, vertex):
        """ Adds vertex to the DiGraph. A vertex can be any object.

        If the vertex already exist, does nothing.
        """
        if vertex not in self._edges:
            self._edges[vertex] = []

```

You will see that inside it just initializes `_edges`. So the only way to create a `DiGraph` is with a call like

In [4]:

```
g = DiGraph()
```

`DiGraph` provides an `__str__` method to have a nice printout:

In [5]:

```
print g
```

```
DiGraph()
```

You can add then vertices to the graph like so:

In [6]:

```
g.add_vertex('a')
g.add_vertex('b')
g.add_vertex('c')
```

In [7]:

```
print g
```

```
a: []
b: []
c: []
```

Adding a vertex twice does nothing:

In [8]:

```
g.add_vertex('a')
print g
```

```
a: []
b: []
c: []
```

Once you added the vertices, you can start adding directed edges among them with the method `add_edge`:

```
def add_edge(self, vertex1, vertex2):
    """ Adds an edge to the graph, from vertex1 to vertex2

    If vertexes don't exist, raises an Exception.
    If there is already such an edge, exits silently.
    """

    if not vertex1 in self._edges:
        raise Exception("Couldn't find source vertex:" + str(vertex1))

    if not vertex2 in self._edges:
        raise Exception("Couldn't find target vertex:" + str(vertex2))

    if not vertex2 in self._edges[vertex1]:
        self._edges[vertex1].append(vertex2)
```

In [9]:

```
g.add_edge('a', 'c')
print g
```

```
a: ['c']
b: []
c: []
```

In [10]:

```
g.add_edge('a', 'b')
print g
```

```
a: ['c', 'b']
b: []
c: []
```

Adding an edge twice makes no difference:

In [11]:

```
g.add_edge('a', 'b')
print g
```

```
a: ['c', 'b']
b: []
c: []
```

Notice a `DiGraph` can have self-loops too (also called *caps*):

In [12]:

```
g.add_edge('b', 'b')
print g
```

```
a: ['c', 'b']
b: ['b']
c: []
```

1.2) dig()

`dig()` is a shortcut to build graphs, it is already provided and you don't need to implement it. **USE IT ONLY WHEN TESTING, NOT IN THE DiGraph CLASS CODE !!!!**

With no parameter prints the empty graph:

In [13]:

```
print dig()
```

DiGraph()

To build more complex graphs, provide pairs source vertex / target vertices list like in the following examples:

In [14]:

```
print dig('a', ['b', 'c'])
```

```
a: ['b', 'c']
b: []
c: []
```

In [15]:

```
print dig('a', ['b', 'c'],
          'b', ['b'],
          'c', ['a'])
```

```
a: ['b', 'c']
b: ['b']
c: ['a']
```

1.3) Equality

Graphs for us are equal irrespectively of the order in which elements in adjacency lists are specified. So for example these two graphs will be considered equal:

In [16]:

```
dig('a', ['c', 'b']) == dig('a', ['b', 'c'])
```

Out[16]:

True

1.4) Basic querying

There are some provided methods to query the DiGraph: `adj`, `vertices`, `is_empty`

1.5) adj

To obtain the edges, you can use the method `adj(self, vertex)`. It is already provided and you don't need to implement it:

```
def adj(self, vertex):
    """ Returns the vertices adjacent to vertex.

    NOTE: vertices are returned in a NEW list.
    Modifying the list will have NO effect on the graph!
    """
    if not vertex in self._edges:
        raise Exception("Couldn't find a vertex " + str(vertex))

    return self._edges[vertex][:]
```

In [17]:

```
lst = dig('a', ['b', 'c'],
          'b', ['c']).adj('a')
print lst

['b', 'c']
```

Let's check we actually get back a new list (so modifying the old one won't change the graph):

In [18]:

```
lst.append('d')
print lst

['b', 'c', 'd']
```

In [19]:

```
print g.adj('a')

['c', 'b']
```

NOTE: This technique of giving back copies is also called *defensive copying*: it prevents users from modifying the internal data structures of a class instance in an uncontrolled manner. For example, if we allowed them direct access to the internal `verteces` list, they could add duplicate edges, which we don't allow in our model. If instead we only allow users to add edges by calling `add_edge`, we are sure the constraints for our model will always remain satisfied.

1.6) `is_empty()`

We can check if a `DiGraph` is empty. It is already provided and you don't need to implement it:

```
def is_empty(self):
    """ A DiGraph for us is empty if it has no verteces and no edges """

    return len(self._edges) == 0
```

In [20]:

```
print dig().is_empty()
```

True

In [21]:

```
print dig('a', []).is_empty()
```

False

1.7) `verteces()`

To obtain the `verteces`, you can use the function `verteces`. (NOTE for Italians: method is called `verteces`, with two **es** !!!). It is already provided and you don't need to implement it:

```
def verteces(self):
    """ Returns a set of the graph verteces. Verteces can be any object. """

    # Note dict keys() return a list, not a set. Bleah.
    # See http://stackoverflow.com/questions/13886129/why-does-pythons-dict-keys-ret
    return set(self._edges.keys())
```

In [22]:

```
g = dig('a', ['c', 'b'],
        'b', ['c'])
print g.verteces()

set(['a', 'c', 'b'])
```

Notice it returns a *set*, as verteces are stored as keys in a dictionary, so they are not supposed to be in any particular order. When you print the whole graph you see them vertically ordered though, for clarity purposes:

In [23]:

```
print g
```

```
a: ['c', 'b']
b: ['c']
c: []
```

Verteces in the edges list are instead stored and displayed in the order in which they were inserted.

1.8) has_edge

Enough for talking! Implement this method in DiGraph:

```
def has_edge(self, source, target):
    """ Returns True if there is an edge between source vertex and target vertex.
        Otherwise returns False.

        If either source, target or both verteces don't exist raises an Exception.
    """

    raise Exception("TODO IMPLEMENT ME!")
```

1.9) full_graph

Implement this function **outside** the class definition. It is **not** a method of DiGraph !

```
def full_graph(verteces):
    """ Returns a DiGraph which is a full graph with provided verteces list.

        In a full graph all verteces link to all other verteces (including themselves!).
    """

    raise Exception("TODO IMPLEMENT ME!")
```

1.10) dag

Implement this function **outside** the class definition. It is **not** a method of DiGraph !

```
def dag(verteces):
    """ Returns a DiGraph which is DAG (Directed Acyclic Graph) made out of provided ver
        teces list

        Provided list is intended to be in topological order.
        NOTE: a DAG is ACYCLIC, so caps (self-loops) are not allowed !!
    """

    raise Exception("TODO IMPLEMENT ME!")
```

1.11) list_graph

Implement this function **outside** the class definition. It is **not** a method of DiGraph !

```
def list_graph(n):
    """ Return a graph of n verteces displaced like a
        monodirectional list: 1 -> 2 -> 3 -> ... -> n

        Each vertex is a number i,  $1 \leq i \leq n$  and has only one edge connecting it
        to the following one in the sequence
        If  $n = 0$ , return the empty graph.
        if  $n < 0$ , raises an Exception.
    """

    raise Exception("TODO IMPLEMENT ME!")
```

1.12) star_graph

Implement this function **outside** the class definition. It is **not** a method of DiGraph !

```
def star_graph(n):
    """ Returns graph which is a star with n nodes

        First node is the center of the star and it is labeled with 1. This node is link
        ed to all the others. For example, for  $n=4$  you would have a graph like this:

            3
            ^
            |
        2 <- 1 -> 4

        If  $n = 0$ , the empty graph is returned
        If  $n < 0$ , raises an Exception
    """

    raise Exception("TODO IMPLEMENT ME!")
```

2) Manipulate graphs

You will now implement some methods to manipulate graphs.

2.1) remove_vertex

```
def remove_vertex(self, vertex):  
    """ Removes the provided vertex and returns it  
  
    If the vertex is not found, raises an Exception.  
    """  
  
    raise Exception("TODO IMPLEMENT ME!")
```

2.2) reverse

```
def reverse(self):  
    """ Reverses the direction of all the edges """  
  
    raise Exception("TODO IMPLEMENT ME!")
```

2.3) has_self_loops

```
def has_self_loops(self):  
    """ Returns True if the graph has any self loop (a.k.a. cap), False otherwise """  
    "  
  
    raise Exception("TODO IMPLEMENT ME !")
```

2.4) remove_self_loops

```
def remove_self_loops(self):  
    """ Removes all of the self-loops edges (a.k.a. caps)  
  
    NOTE: Removes just the edges, not the verteces!  
    """  
  
    raise Exception("TODO IMPLEMENT ME!")
```


3) Query graphs

You can query graphs the "Do it yourself" way with Depth First Search (DFS) or Breadth First Search (BFS).

3.1) Visit and VertexLog

If you noticed, in the skeleton there are two extra classes `Visit` and `VertexLog`. Also, in `DiGraph` the functions `dfs` and `bfs` are already provided. The idea here is that both `dfs` and `bfs` will traverse the graph and report the intermediate results of the visit inside instances of `Visit` and `VertexLog`. At the end of the traversal, they will give back one instance of `Visit`. Maybe when you do exercises on paper it is convenient to write for example the discovery times inside the nodes of your graphs, but when programming writing intermediate results directly in the vertices of the input graph may cause troubles to the users of your methods. So it is better to store such visit logs in separate data structures: basically, `Visit` contains a map that associates to each vertex its `VertexLog`:

```
class Visit:
    """ The visit of a DiGraph visit sequence.

    """

    def __init__(self):
        """ Creates a Visit """

        self._logs = {}
```

In `VertexLog` you can put the intermediate info like i.e. `discovery_time`, or parents of the node if you are interested in building a tree.

```
class VertexLog:
    """ Represents the visit log a single DiGraph vertex

    This class is very simple and doesn't even have getters methods.

    You can just access fields by using the dot:

        print vertex_log.discovery_time

    and set them directly:

        vertex_log.finish_time = 5

    If you want, an instances you can set your own fields:

        vertex_log.my_own_field = "whatever"

    """

    def __init__(self, vertex):
        self.vertex = vertex
        self.discovery_time = -1
        self.finish_time = -1
        self.parent = None
```

Let's make a simple example:

In [24]:

```
g = dig('a', ['a','b', 'c'],
        'b', ['c'],
        'd', ['e'])
print g.dfs('a')
```

```
Visit:
[ { 'discovery_time': 1, 'finish_time': 6, 'parent': None, 'vertex': 'a'},
  { 'discovery_time': 2, 'finish_time': 5, 'parent': 'a', 'vertex': 'b'},
  { 'discovery_time': 3, 'finish_time': 4, 'parent': 'b', 'vertex': 'c'}]
```

Notice we started from 'a', so by default unreachable nodes like d and e were not displayed. Let's try a bfs:

In [25]:

```
print g.bfs('a')
```

Visit:

```
[ { 'discovery_time': 1, 'finish_time': -1, 'parent': None, 'vertex': 'a'},
  { 'discovery_time': 2, 'finish_time': -1, 'parent': 'a', 'vertex': 'b'},
  { 'discovery_time': 3, 'finish_time': -1, 'parent': 'a', 'vertex': 'c'}]
```

Predictably, results are different, you can see it by the parent fields. Note how the `finish_time` here is always `-1` because it is less meaningful to calculate it for a 'bfs'.

You can extract the logs from the Visit object by calling `logs()`:

In [26]:

```
pp(g.dfs('a').logs())
```

```
[ { 'discovery_time': 1, 'finish_time': 6, 'parent': None, 'vertex': 'a'},
  { 'discovery_time': 2, 'finish_time': 5, 'parent': 'a', 'vertex': 'b'},
  { 'discovery_time': 3, 'finish_time': 4, 'parent': 'b', 'vertex': 'c'}]
```

By default, they are sorted ascending by discovery time. To see them in descending order, use `descendant=False`:

In [27]:

```
pp(g.dfs('a').logs(descendant=True))
```

```
[ { 'discovery_time': 3, 'finish_time': 4, 'parent': 'b', 'vertex': 'c'},
  { 'discovery_time': 2, 'finish_time': 5, 'parent': 'a', 'vertex': 'b'},
  { 'discovery_time': 1, 'finish_time': 6, 'parent': None, 'vertex': 'a'}]
```

To see the last timestamp, use `last_time`:

In [28]:

```
print g.dfs('a').last_time()
```

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3.2) distances()

Try to implement this method of DiGraph:

```
def distances(self, source):
    """
    Returns a dictionary where the keys are vertexes, and each vertex v is associate
    d
    to the *minimal* distance in number of edges required to go from the source
    vertex to vertex v. If node is unreachable, the distance will be -1

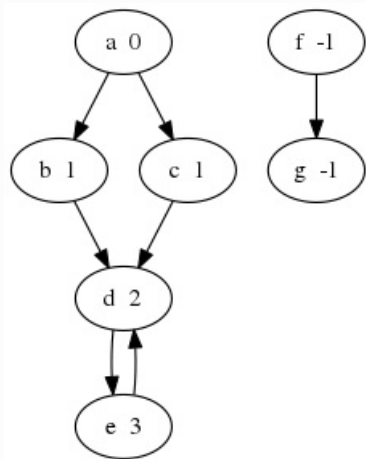
    Source has distance zero from itself
    Vertexes immediately connected to source have distance one.

    if source is not a vertex, raises an Exception

    HINT: to implement this, copy and edit either dfs or bfs. Question: which one ?
    """
```

If you look at the following graph, you can see an example of the distances to associate to each vertex, supposing that the source is a. Note that a itself is at distance zero from itself and also that unreachable nodes like f and g will be at distance `-1` :

Out[29]:



`distances('a')` called on this graph would return a map like this:

```
{
  'a':0,
  'b':1,
  'c':1,
  'd':2,
  'e':3,
  'f':-1,
  'g':-1,
}
```

3.2) Play with dfs and bfs

Create small graphs (like linked lists $a \rightarrow b \rightarrow c$, triangles, mini-full graphs, trees - you can also use the functions you defined to create graphs like `full_graph`, `dag`, `list_graph`, `star_graph`) and try to predict the visit sequence (vertices order, with discovery and finish times) you would have running a dfs or bfs. Then write tests that assert you actually get those sequences when running provided dfs and bfs

3.3) Blow up your computer

Try to call the already implemented function `gen_graphs` with small numbers for n , like 1, 2, 3, 4 Just with 2 we get back a lot of graphs:

```
def gen_graphs(n):
    """ Returns a list with all the possible  $2^{(n^2)}$  graphs of size  $n$ 

        Vertices will be identified with numbers from 1 to  $n$ 
    """
```

In [30]:

```
print gen_graphs(2)
```

```
[
1: []
2: []
,
1: []
2: [2]
,
1: []
2: [1]
,
1: []
2: [1, 2]
,
1: [2]
2: []
,
1: [2]
2: [2]
,
1: [2]
2: [1]
,
1: [2]
2: [1, 2]
,
1: [1]
2: []
,
1: [1]
2: [2]
,
1: [1]
2: [1]
,
1: [1]
2: [1, 2]
,
1: [1, 2]
2: []
,
1: [1, 2]
2: [2]
,
1: [1, 2]
2: [1]
,
1: [1, 2]
2: [1, 2]
]
```

QUESTION: What happens if you call `gen_graphs(10)` ? How many graphs do you get back ?

4) Do cool stuff with theory

- find connected components
- determine if a graph is acyclic
- find node distances

In [32]:

```
from graphs_solution import *  
algolab.run(VisitTest)
```

...

Ran 3 tests in 0.007s

OK

Solution

Solutions are in a [separate file \(graphs_solution.py\)](#).