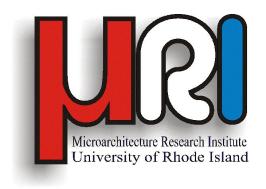
verifying simulated program correspondence to actual execution through function coverage

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outline



- problem
- simulation approaches
- verification approaches
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- results
 - using reference inputs
 - using reduced inputs
- summary

problem



- we want to perform microarchitectural simulation on a program
- microarchitectural simulation takes a long time!
- we generally cannot simulate the program to completion
- how do we verify that our simulated program execution is characteristic of an actual program execution done in its entirety?

simulation approaches



simulation approaches

- run the program from the start but only for a limited number of instructions
- skip some number of instructions at the start and then run the full microarchitectural ("performance") simulation for a limited number of instructions
- use a different program input for the simulation as compared with what would be used normally
- alternately run the skipping simulator with the microarchitectural simulator
- in any of these cases, how do we verify correspondence to actual program execution?

verification approaches



- compare statistics from the actual execution with those from the simulation
 - types of instructions execution
 - percent memory operations
 - percent control-flow instructions
 - other
- what about when all (or most all) of the above are the same?
- try function coverage!

function coverage



- what is it?
 - the percentage of instructions that the program executes in each of its subroutines
- how do we determine it?

```
gather statistical (in time)⇒ monitor(3c)⇒ profil(2)
```

gather actual (instruction counts)

```
⇒ Pixie
⇒ ATOM
```

analyze

```
\Rightarrow \operatorname{prof}(1)
\Rightarrow \operatorname{gprof}(1)
```

- in simulation, we can also determine it using instructions counts
 - but we need a FAST algorithm! (data structures?)

implementation overview



at load time

- read symbol table from the program object file
- record in DB the subroutine names and entry addresses
- sort all subroutines by entry address for quick access later
- on SGI OS determine subroutine length from successive entry points

at run time

- for each instruction: search DB for which function the instruction address is in
 - ⇒use cache of most recent subroutine encountered for extra speed
 - ⇒use binary search for speed -- complexity O(log N)
- increment the associated subroutine counter

after execution

sort subroutines by hit counts

example (BZIP2)



Pixie (whole program)

function % instructions generateMTFValues 21.5% qSort3 15.1% sendMTFValues 13.7% sortlt 11.4% fullGtU 9.6% getRLEpair 8.1% 6.7% simpleSort

IoadAndRLEsource

fgetc

simulation (500 M instructions)

function	% instructions
generateMTFValues	18.8%
sendMTFValues	13.3%
qSort3	10.0%
fullGtU	8.6%
sortlt	7.2%
getRLEpair	6.9%
simpleSort	6.4%
bsW	3.3%
fgetc	2.9%

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3.7% 3.7%

example (COMPRESS)



Pixie (whole program)

simulation (500 M instructions, reference input)

function	% instructions
compress	37.1%
decompress	14.4%
output	12.3%
getcode	10.7%
putbyte	9.0%
getbyte	6.6%
readbytes	4.8%
getranchar	4.3%
fgetc	0.4%

function	% instructions
getranchar	67.7%
ran2	9.0%
fill_text_buffer	6.7%
compress	0.0%
output	0.0%
decompress	0.0%
getcode	0.0%
rindex	0.0%
onintr	0.0%

example (COMPRESS)



Pixie (whole program)

simulation (500 M instructions, reduced input)

function	% instructions
compress	37.1%
decompress	14.4%
output	12.3%
getcode	10.7%
putbyte	9.0%
getbyte	6.6%
readbytes	4.8%
getranchar	4.3%
fgetc	0.4%

function	% instructions
compress	24.5%
getranchar	14.7%
output	10.1%
decompress	9.6%
getcode	7.1%
putbyte	5.9%
getbyte	4.9%
readbytes	2.9%
ran2	2.0%

summary



- shown how function coverage can be used to verify corresponding simulation characteristics with actual program execution
- shown how function coverage is calculated (fast) in simulation environment
- used function coverage analysis to change the input on the COMPRESS program to get closer to the actual whole-program execution characteristics -- saved us on our Micro'02 submission!