
Problem Set 1

Due: Tuesday January 24, 2017, in class.

By $a\|b$ we denote the concatenation of strings $a, b \in \{0, 1\}^*$. (For example $010\|01 = 01001$.) The time to compute a blockcipher E , denoted T_E , is the time to compute E or E^{-1} . All times are worst case.

Problem 1 [60 points] Let $E: \{0, 1\}^k \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ be a blockcipher. Define $F: \{0, 1\}^{k+n} \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ as follows:

Alg $F(K_1\|K_2, M)$
 $C \leftarrow E(K_1, M \oplus K_2)$
Return C

Above, $K_1 \in \{0, 1\}^k$ and $K_2, M \in \{0, 1\}^n$.

- (a) **[8 points]** Is F a blockcipher? Answer YES or NO and prove your answer correct.
 - (b) **[8 points]** How much time is taken by a 3-query exhaustive key search attack on F ? Your answer should be a function of T_E, k, n .
 - (c) **[18 points]** Present in pseudocode a 1-query adversary A_1 that has advantage $\text{Adv}_F^{\text{kr}}(A_1) = 1$ and running time $\mathcal{O}(T_E + k + n)$.
 - (d) **[26 points]** Present in pseudocode a 3-query adversary A_3 that has advantage $\text{Adv}_F^{\text{kr}}(A_3) = 1$ and running time $\mathcal{O}(2^k \cdot (T_E + k + n))$.
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Problem 2 [40 points] Let $E_1: \{0, 1\}^k \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ and $E_2: \{0, 1\}^k \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ be blockciphers. Define $E: \{0, 1\}^k \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ by

Alg $E(K, M)$
 $C_1 \leftarrow E_1(K, M)$; $C_2 \leftarrow E_2(K, C_1)$
Return C_2

Above, $K \in \{0, 1\}^k$ and $M \in \{0, 1\}^n$.

- (a) [10 points] Prove that E is a blockcipher.
- (b) [30 points] Let blockcipher $E_1: \{0, 1\}^k \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ be given. Specify in pseudocode the following:
- A blockcipher $E_2: \{0, 1\}^k \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ —of your choice, allowed to depend on E_1 —and its inverse $E_2^{-1}: \{0, 1\}^k \times \{0, 1\}^n \rightarrow \{0, 1\}^n$ such that the time to compute E_2 is $\mathcal{O}(T_{E_1} + k + n)$.
 - A 1-query adversary A having running time $\mathcal{O}(T_{E_1} + k + n)$ and achieving advantage $\mathbf{Adv}_E^{\text{kr}}(A) = 1$ against E , where E is defined as above based on the given E_1 and your E_2 .
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