

Classical computation on a quantum computer

(Nielsen and Chuang 3.2.5)

A quantum computer can efficiently simulate a classical computer, by compiling everything down to NAND gates

quantum computer \longrightarrow reversible classical computer \longrightarrow general classical computer

This seems to imply that we can run a classical subroutine as part of a quantum computation

But it's not quite as easy as this

Consider only reversible circuits (since unitary circuits are of this form)

Consider a circuit that maps an input z to an output $f(z)$, and also outputs z to ensure reversibility

In general it will also output 'garbage' $g(z)$

$$(z, 0) \longrightarrow (z, f(z), g(z))$$

If the input was not included in the output, the garbage must be required for unitarity. But this is not the case here

So $g(z)$ is just some undeleted remnants of the computation

For a classical computation, this garbage can be ignored or deleted

For a quantum computation, we have superpositions to worry about

$$U |z\rangle = |f(z)\rangle \Rightarrow U \sum_z c_z |z\rangle = \sum_z c_z |f(z)\rangle$$

$$U' |z\rangle \otimes |0\rangle = |f(z)\rangle \otimes |g(z)\rangle \Rightarrow U' \sum_z c_z |z\rangle \otimes |0\rangle = \sum_z c_z |f(z)\rangle \otimes |g(z)\rangle$$

The operations U and U' are fundamentally different. U' entangles the computation to garbage in an ancilla, which could mess up required interference effects

So, can a function $f(z)$ that can be computed efficiently with a classical computer also be computed efficiently with a reversible classical computer that produces no garbage?

Yes! By means of 'uncomputation'

$(z, 0, 0, 0)$ We use 4 registers, one with input, rest set to 0

$\rightarrow (z, f(z), g(z), 0)$ Then do the computation, including garbage output

Then copy $f(x)$ to the fourth register

$\rightarrow (z, f(z), g(z), f(z))$

$\rightarrow (z, 0, 0, f(z))$

Finally invert the computation (but not the copy) to reset the second and third registers

Same complexity, but no garbage

Note: we assumed above that the additional registers were initialized in the zero state

$$(z, 0, 0, 0) \rightarrow (z, 0, 0, f(z))$$

This need not be true in general, they can be initialized in any state

$$(z, a, b, c) \rightarrow (z, 0, 0, f(z))$$

This allows us to have nontrivial a , b and c if we want (pre-stored constants, mathematical convenience)

Coming back to the quantum realm (and ignoring any systems that start and finish in the same state) an efficient classical computation

$$z \rightarrow (f(z), g(z)) \quad [\text{garbage present in general}]$$

Implies an efficient quantum computation

$$|z, 0\rangle \rightarrow |z, f(z)\rangle \quad [\text{no garbage}]$$