

Q5.1 Base case: height = 0 \Rightarrow Nodes = 1 $\geq 2 \cdot 0 + 1$ True

Q5.2 Assume that for all full binary trees with height $\leq h$, we have
 $\# \text{nodes} \geq 2(\text{height}) + 1$

Q5.3 Let T be a nonbinary subtree with height $h+1$. Then

T has a root and two subtrees, L, R , with height l, r .

Since the max height of the subtree is h , we have

$\max(l, r) = h$. Since L, R have height at most h , we

know that the $\# \text{nodes}(L) \geq 2l+1$, $\# \text{nodes}(R) \geq 2r+1$

$\therefore \# \text{nodes}(T) = 1 + \# \text{nodes}(L) + \# \text{nodes}(R) \geq 1 + (2l+1) + (2r+1) = 2(l+r) + 3$

However we have $\max(l, r) = h \Rightarrow l+r \geq h$

$\therefore \# \text{nodes}(T) \geq 2(h) + 3 \geq 2(h+1) + 1 = 2 \cdot \text{Height}(T) + 1$ and we are done.