



Lab 07

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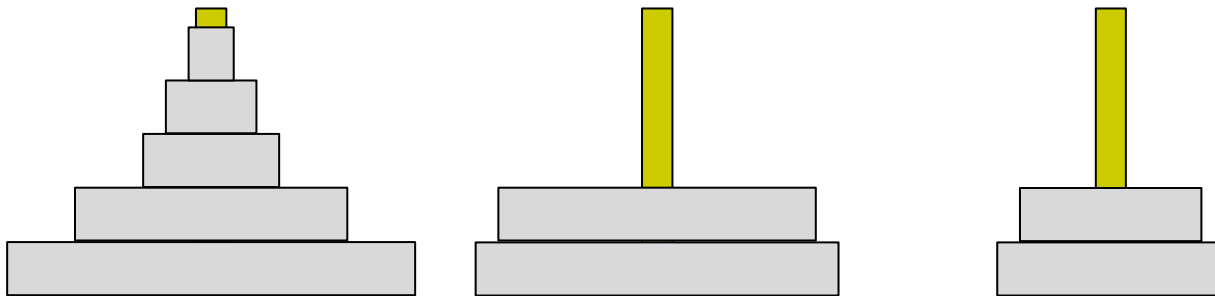
Torino - Italy

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Tower of Hanoi

- The tower of Hanoi consists of 3 poles.
- Each pole contains a certain number of discs (possibly zero), arranged in descending order of size (the largest disc is the lowest).



Tower of Hanoi: implementation

- Each pole can be represented across a stack.
- The elements of the stack indicate the size of the discs contained in the corresponding pole.
- The 3 stacks are defined in a READWRITE area: the 3 stack pointers point to different memory locations within the data area.
- Freedom of choice is left to manage the 3 stacks as full descending or empty ascending.

Exercise 1

- Write a `fillStack` subroutine to initialize the stack associated to a pole.
- The subroutine receives two parameters:
 - the stack pointer associated to the pole
 - the address to a READONLY area containing a sequence of constants.
- The subroutine updates the two parameters:
 - the stack pointer points to the last entered data
 - the address in the READONLY area is subsequent to the last constant entered.

Exercise 1 (cont.)

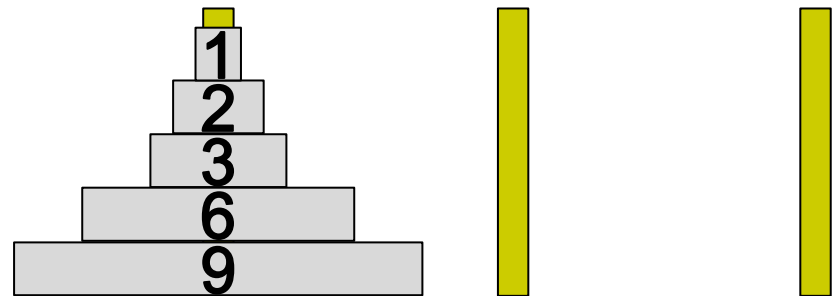
- The parameters are passed to the subroutine through the main stack (SP).
- The insertion ends when one of the following two conditions occurs:
 - the constant to be entered is 0
 - the constant to be inserted is greater than the last element on the stack.

Exercise 1: example

- `r1` is the stack pointer associated with the first pole
- The sequence of constants is
DCD 9, 6, 3, 2, 1, 8, 7, 0, 5, 4, 0
- `r0` contains the address of the first constant (9).
- After calling `fillStack` the situation is

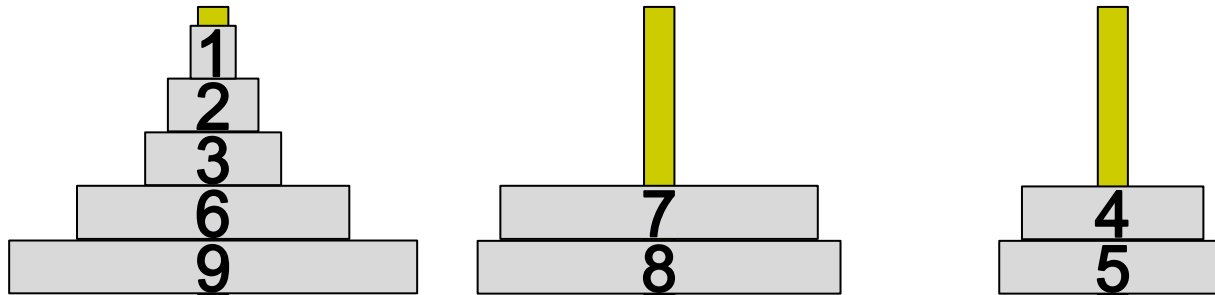
- $r1_{\text{new}} = r1_{\text{old}} \pm 20$

- $r0_{\text{new}} = r0_{\text{old}} + 20$



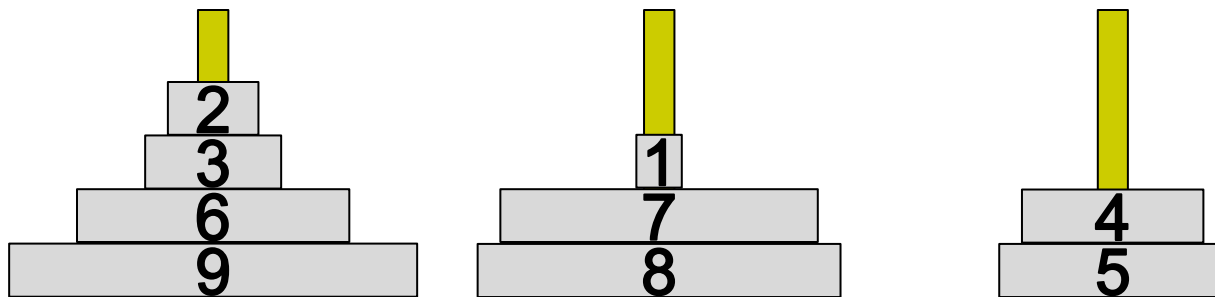
Exercise 1: example

- With three consecutive calls to `fillStack` (one for each pole) the tower of Hanoi becomes:



Tower of Hanoi: single movement

- It is possible to move a disc from one pole to another if one of two conditions occurs:
 - the new pole is empty
 - the upper disc in the new pole has a larger size than the disc to be moved.
- Example: a disc is moved from pole 1 to pole 2



Exercise 2

- Write a `move1` subroutine that handles moving a disk.
- `move1` receives 3 parameters across the stack
 - stack pointer of the starting pole
 - stack pointer of the destination pole
 - space for the return value
- `move1` returns 1 if it was possible to move the disk, 0 otherwise. Stack pointers are updated if the disk has been moved.

Exercise 2: warning

- To pass parameters to `move1`, it is preferable to use 3 `PUSH` instead of just one.
- Example: let `r1` be the stack pointer of the starting pole, `r2` the stack pointer of the destination pole, `r0` contain the return value.
- ```
PUSH {r1, r2, r0}
BL move1
POP {r1, r2, r0}
```

moves from pole `r2` to pole `r1`, because the registers are reordered as `{r2, r1, r0}`

## Exercise 2: warning

- To ensure the movement from pole  $r1$  to pole  $r2$ , the following code can be used:

```
PUSH { r1 }
```

```
PUSH { r2 }
```

```
PUSH { r0 }
```

```
BL move1
```

```
POP { r0 }
```

```
POP { r2 }
```

```
POP { r1 }
```

# Tower of Hanoi: multiple movements

- It is possible to move  $N$  discs from pole  $X$  to pole  $Y$  using a recursive procedure:
  - Move the first  $N-1$  discs from  $X$  to  $Z$
  - Move disk  $N$  from  $X$  to  $Y$
  - Move  $N-1$  discs from  $Z$  to  $Y$

## Exercise 3

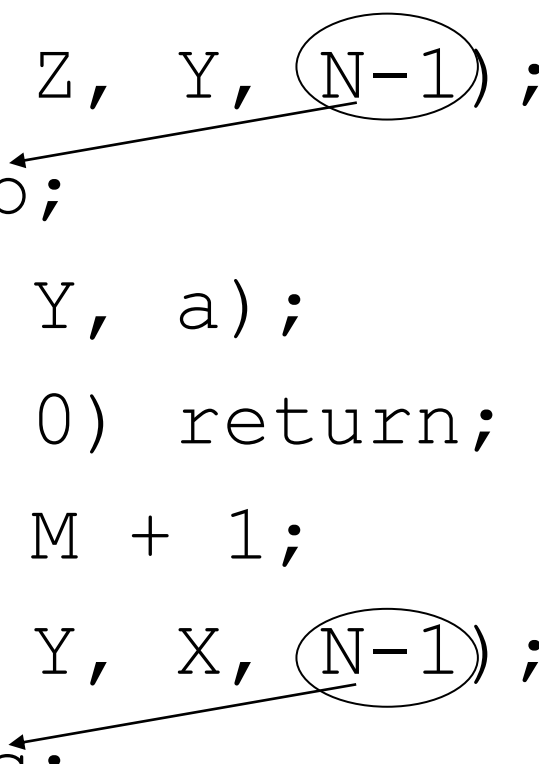
- Write a `moveN` subroutine that manages the movement of N disks.
- `moveN` receives across the stack:
  - the stack pointer of the starting pole X
  - the stack pointer of the destination pole Y
  - the stack pointer of the auxiliary pole Z
  - the number of discs to move
- `moveN` updates the stack pointers and replaces the fourth parameter with the number of movements performed.

# Pseudocode of moveN(X, Y, Z, N)

```
/* M is the number of mov. made */
M = 0;
if (N == 1)
{
 move1(X, Y, a);
/* a is the return value: 0-1 */
 M = M + a;
}
```

# Pseudocode of moveN(X, Y, Z, N)

```
else {
 moveN(X, Z, Y, N-1);
 M = M + b;
 move1(X, Y, a);
 if (a == 0) return;
 else M = M + 1;
 moveN(Z, Y, X, N-1);
 M = M + c;
}
```



The diagram illustrates the recursive calls and return values in the pseudocode. Two arrows point from the circled 'N-1' arguments in the recursive calls to the variables 'b' and 'c' in the assignment statements. The first arrow points from the 'N-1' in 'moveN(X, Z, Y, N-1);' to the 'b' in 'M = M + b;'. The second arrow points from the 'N-1' in 'moveN(Z, Y, X, N-1);' to the 'c' in 'M = M + c;'. This indicates that the recursive calls return values that are added to 'M'.

## Exercise 3: suggestions

- Since each recursive call adds its own parameters to the stack, it is suggested to increase the stack size by changing the value of `Stack_Size`.
- As a final check, considering a tower of Hanoi with  $N$  discs in one pole and 0 in the other two, note that the number of movements required to move the  $N$  discs into another pole is equal to  $2^N - 1$ .