

gcc x86 Assembly Quick Reference ("Cheat Sheet")

Instructions		
Mnemonic	Purpose	Examples
<code>mov <i>src,dest</i></code>	Move data between registers, load immediate data into registers, move data between registers and memory.	<code>mov \$4,%eax # Load constant into eax</code> <code>mov %eax,%ebx # Copy eax into ebx</code> <code>mov %ebx,123 # Copy ebx to memory address 123</code>
<code>push <i>src</i></code>	Insert a value onto the stack. Useful for passing arguments, saving registers, etc.	<code>push %ebp</code>
<code>pop <i>dest</i></code>	Remove topmost value from the stack. Equivalent to " <code>mov (%esp),<i>dest</i>; add \$4,%esp</code> "	<code>pop %ebp</code>
<code>call <i>func</i></code>	Push the address of the next instruction and start executing <i>func</i> .	<code>call print_int</code>
<code>ret</code>	Pop the return program counter, and jump there. Ends a subroutine.	<code>ret</code>
<code>add <i>src,dest</i></code>	<i>dest=dest+src</i>	<code>add %ebx,%eax # Add ebx to eax</code>
<code>mul <i>src</i></code>	Multiply <i>eax</i> and <i>src</i> as unsigned integers, and put the result in <i>eax</i> . High 32 bits of product go into <i>eax</i> .	<code>mul %ebx #Multiply eax by ebx</code>
<code>jmp <i>label</i></code>	Goto the instruction <i>label</i> :. Skips anything else in the way.	<code>jmp post_mem</code> <code>mov %eax,0 # Write to NULL!</code> <code>post_mem: # OK here...</code>
<code>cmp <i>a,b</i></code>	Compare two values. Sets flags that are used by the conditional jumps (below). WARNING: compare is relative to *last* argument, so " <code>jl</code> " jumps if <i>b<a</i> !	<code>cmp \$10,%eax</code>
<code>jl <i>label</i></code>	Goto <i>label</i> if previous comparison came out as less-than. Other conditionals available are: <code>jle</code> (<code><=</code>), <code>jeq</code> (<code>==</code>), <code>jge</code> (<code>>=</code>), <code>jg</code> (<code>></code>), <code>jne</code> (<code>!=</code>), and many others.	<code>jl loop_start # Jump if eax<10</code>

Stack Frame

(example without %ebp or local variables)

Contents	off esp
caller's variables	12(%esp)
Argument 2	8(%esp)
Argument 1	4(%esp)
Caller Return Address	0(%esp)

```

my_sub: # Returns first argument
    mov 4(%esp), %eax
    ret
    
```

(example when using %ebp and two local variables)

Contents	off ebp	off esp
caller's variables	16(%ebp)	24(%esp)
Argument 2	12(%ebp)	20(%esp)
Argument 1	8(%ebp)	16(%esp)
Caller Return Address	4(%ebp)	12(%esp)
Saved ebp	0(%ebp)	8(%esp)
Local variable 1	-4(%ebp)	4(%esp)
Local variable 2	-8(%ebp)	0(%esp)

```

my_sub2: # Returns first argument
    push %ebp # Prologue
    mov %esp, %ebp
    mov 8(%ebp), %eax
    mov %ebp, %esp # Epilogue
    pop %ebp
    
```

ret

Constants, Registers, Memory

Constants MUST be preceeded with "\$". "\$12" means decimal 12; "\$0xF0" is hex. "\$some_function" is the address of the first instruction of the function. WARNING: a bare "12", "0xF0", or "some_function" dereferences the expression like it was a pointer!

Registers MUST be preceeded with "%". "%eax" means register eax.

Memory access (use register as pointer): "(%esp)". Same as C "`*esp`".

Memory access with offset (use register + offset as pointer): "`4(%esp)`". Same as C "`*(esp+4)`".

Memory access with scaled index (register + another register * scale): "`(%eax, %ebx, 4)`". Same as C "`*(eax+ebx*4)`".

Registers

%esp is the stack pointer
%ebp is the stack frame pointer
Return value in %eax
Arguments are on the stack
Free for use (no save needed):
 %eax, %ebx, %ecx, %edx
Must be saved:
 %esp, %ebp, %esi, %edi

Common Errors

Segfault on innocent-looking code.

- Do you need to add "\$" in front of a constant?
- Did you clean up the stack properly?

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The Intel [Software Developer's Manuals](#) are incredibly long, boring, and complete--they give all the nitty-gritty details. [Volume 1](#) lists the processor registers in Section 3.4.1. [Volume 2](#) lists all the x86 instructions in Section 3.2. [Volume 3](#) gives the performance monitoring registers. For Linux, the [System V ABI](#) gives the calling convention on page 39. Also see the Intel [hall of fame](#) for historical info. [Sandpile.org](#) has a good opcode table.

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